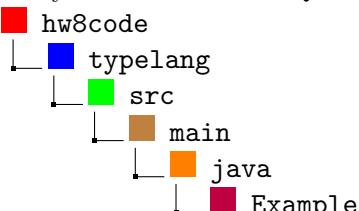


Homework Solutions: TypeLang

Learning Objectives:

1. Understanding and implementing typing rules

Instructions:

- Total points 66 pt
- Early deadline: Apr 14 (Wed) at 11:59 PM; Regular deadline: Apr 16 (Fri) at 11:59 PM (or till TAs start grading the homework)
- Download hw8code.zip from Canvas. Interpreter for Typelang is significantly different compared to previous interpreters:
 - Env in Typelang is generic compared to previous interpreters.
 - Two new files Checker.java and Type.java have been added
 - Type.java defines all the valid types of Typelang.
 - Checker.java defines type checking semantics of all expressions.
 - Typelang.g has changed to add type information in expressions. Please review the changes in file to understand the syntax.
 - Finally Interpreter.java has been changed to add type checking phase before evaluation of Typelang programs.
- Set up the programming project following the instructions in the tutorial from hw2 (similar steps)
- Extend the Typelang interpreter for Q1 - Q5.
- How to submit:
 - Please submit your solutions in one zip file with all the source code files (just zip the complete project's folder).
 - Write your solutions to Q6 - Q7 in a HW8.scm file and store it under your code directory.
 - Submit the zip file to Canvas under Assignments, Homework 8.

Questions:

1. (10 pt) [Implement type rules] Implement the type rules for memory related expressions:

(a) (5 pt) DerefExp: Let a deref expression be (deref e1), where e1 is an expression.

- if e1's type is ErrorT then (deref e1)'s type should be ErrorT
- if e1's type is RefT then (deref e1)'s type should RefT.nestType(). Note that nestType() is method in RefT class.
- otherwise, (deref e1)'s type is ErrorT with message "The dereference expression expect a reference type " + "found " + e1's type + " in " + expression.

Note that you have to add e1's and e2's type and expression in the error message. Examples: \$ (deref (ref : num 45))

45

// no explicit error cases
\$ (deref 45)

Type error: The dereference expression expects a reference type, found num in (deref 45)

(b) (5 pt) AssignExp: Let a set expression be (set! e1 e2), where e1 and e2 are expressions.

- if e1's type is ErrorT then (set! e1 e2)'s type should be ErrorT
- if e1's type is RefT and nestedType of e1 is T then
 - if e2's type is ErrorT then (set! e1 e2)'s type should be ErrorT
 - if e2's type is typeEqual To T then (set! e1 e2)'s type should be e2's type.
 - otherwise (set! e1 e2)'s type is ErrorT with message "The inner type of the reference type is " + nestedType T + " the rhs type is " + e2's type + " in " + expression
- otherwise (set! e1 e2)'s type is ErrorT with message "The lhs of the assignment expression expect a reference type found " + e1's type + " in " + expression.

Note that you have to add e1's and e2's type and expression in the error message. Examples:

\$ (set! (ref : num 0) #t)

Type error: The inner type of the reference type is number the rhs type is bool in (set! (ref 0) #t)

\$ (set! (ref : bool #t) (list : num 1 2 3 4 5 6))

Type error: The inner type of the reference type is bool the rhs type is List<number> in (set! (ref #t) (list 1 2 3 4 5 6))

Sol

- DerefExp

```

1  public Type visit(DerefExp e, Env<Type> env) {
2      Exp exp = e.loc_exp();
3      Type type = (Type)exp.accept(this, env);
4      if (type instanceof ErrorT) { return type; }
5
6      if (type instanceof RefT) {
7          RefT rt = (RefT)type;

```

```

8         return rt.nestType();
9     }
10
11    return new ErrorT("The dereference expression expect a reference type " +
12        "found " + type.toString() + " in " + ts.visit(e, null));
13 }
```

- AssignExp

```

1  public Type visit(AssignExp e, Env<Type> env) {
2      Exp lhs_exp = e.lhs_exp();
3      Type lhsType = (Type)lhs_exp.accept(this, env);
4      if (lhsType instanceof ErrorT) { return lhsType; }
5
6      if (lhsType instanceof RefT) {
7          Exp rhs_exp = e.rhs_exp();
8          Type rhsType = (Type)rhs_exp.accept(this, env);
9          if (rhsType instanceof ErrorT) { return rhsType; }
10
11         RefT rt = (RefT)lhsType;
12         Type nested = rt.nestType();
13
14         if (rhsType.typeEqual(nested)) { return rhsType; }
15
16         return new ErrorT("The inner type of the reference type is " +
17             nested.toString() + " the rhs type is " + rhsType.toString()
18             + " in " + ts.visit(e, null));
19     }
20     return new ErrorT("The lhs of the assignment expression expect a "
21         + "reference type found " + lhsType.toString() + " in " +
22         ts.visit(e, null));
23 }
```

2. (10 pt) [Implement type rules] Implement the type rules for list expressions:

(a) (5 pt) CarExp: Let a car expression be (car e1), where e1 is an expression.

- if e1's type is ErrorT then (car e1)'s type should be ErrorT
- if e1's type is PairT then (car e1)'s type should be the type of the first element of the pair
- otherwise, (car e1)'s type is ErrorT with message "The car expect an expression of type Pair, found" + e1's type + "in" + expression

Note that you have to add e1's type and expression in the error message. See some examples below.

\$ (car 2)

Type error: The car expect an expression of type Pair, found num in (car 2)

\$ (car (car 2))

Type error: The car expect an expression of type Pair, found num in (car 2)

(b) (5 pt) ListExp: Let a list expression be (list : T e1 e2 e3 ... en), where T is type of list and e1, e2, e3 ... en are expressions

- if type of any expression ei, where ei is an expression of element in list at position i, is ErrorT then type of (list : T e1 e2 e3 ... en) is ErrorT.

- if type of any expression e_i , where e_i is an expression of an element of list, is not T then type of $(\text{list} : T \ e_1 \ e_2 \ e_3 \dots \ e_n)$ is ErrorT with message "The " + index + " expression should have type " + T + " found " + Type of e_i + " in " + "expression". where index is the position of expression in list's expression list.
- else type of $(\text{list} : T \ e_1 \ e_2 \ e_3 \dots \ e_n)$ is ListT .

Note that you have to add e_i 's type and expression in the error message. Some examples appear below.

$\$ (\text{list} : \text{bool} \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7)$

Type error: The 0 expression should have type bool, found number in (list 1 2 3 4 5 6 7)

$\$ (\text{list} : \text{num} \ 1 \ 2 \ 3 \ 4 \ 5 \ \#t \ 6 \ 7 \ 8)$

Type error: The 5 expression should have type number, found bool in (list 1 2 3 4 5 #t 6 7 8)

- **carexp**

```

1  public Type visit(CarExp e, Env<Type> env) {
2      Exp exp = e.arg();
3      Type type = (Type)exp.accept(this, env);
4      if (type instanceof ErrorT) { return type; }
5
6      if (type instanceof PairT) {
7          PairT pt = (PairT)type;
8          return pt.fst();
9      }
10
11     return new ErrorT("The car expect an expression of type Pair, found "
12         + type.tostring() + " in " + ts.visit(e, null));
13 }
```

- **listexp**

```

1  public Type visit(ListExp e, Env<Type> env) {
2      List<Exp> elems = e.elems();
3      Type type = e.type();
4
5      int index = 0;
6      for (Exp elem : elems) {
7          Type elemType = (Type)elem.accept(this, env);
8          if (elemType instanceof ErrorT) { return elemType; }
9
10         if (!assignable(type, elemType)) {
11             return new ErrorT("The " + index +
12                 " expression should have type " + type.tostring() +
13                 " found " + elemType.tostring() + " in " +
14                 ts.visit(e, null));
15         }
16         index++;
17     }
18     return new ListT(type);
19 }
```

3. (5 pt) [Implement type rules] Implement the type rules for function calls.

CallExp: Let a call expression be $(\text{ef} \ e_1 \dots \ e_n)$ with type:

- if the type of ef is ErrorT, return ErrorT
- if the type of ef is not FuncT, the type of the call expression is ErrorT, reporting the message "Expect a function type in the call expression, found "+ef's type+"in "+ expression
- if any one of e1, e2, ...en has ErrorT, the call expression has ErrorT
- given that ef has FuncT ($T_1 \dots T_n \rightarrow T_b$), if the actual parameter ei does not have a type T_i , the call expression has ErrorT, reporting the message "The expected type of the " + i + "th actual parameter is " + T_i + ", found " + ei's type +"in "+expression
- otherwise, the type of call expression is T_b

Some examples appear below.

```
1      $(define add: (num num num -> num) (lambda (x: num y: num z: num) (+ x (+ y z
)))))
```

```
$ (add 5 56 #t)
Type error: The expected type of the 2 argument is number found bool in (add 5.0 56.0 #t )
$ (3 4)
Type error: Expect a function type in the call expression, found number in (3 4)
```

Sol.

- CallExp
-

```
1
2  public Type visit(CallExp e, Env<Type> env) {
3      Exp operator = e.operator();
4      List<Exp> operands = e.operands();
5
6      Type type = (Type)operator.accept(this, env);
7      if (type instanceof ErrorT) { return type; }
8
9      String message = "Expect a function type in the call expression, found "
10     + type.toString() + " in ";
11     if (type instanceof FuncT) {
12         FuncT ft = (FuncT)type;
13
14         List<Type> argTypes = ft.argTypes();
15         int size_actuals = operands.size();
16         int size_formals = argTypes.size();
17
18         message = "The number of arguments expected is " + size_formals +
19         " found " + size_actuals + " in ";
20         if (size_actuals == size_formals) {
21             for (int i = 0; i < size_actuals; i++) {
22                 Exp operand = operands.get(i);
23                 Type operand_type = (Type)operand.accept(this,
24                     env);
25
26                 if (operand_type instanceof ErrorT) { return
27                     operand_type; }
```

```

26             if (!assignable(argTypes.get(i), operand_type)) {
27                 return new ErrorT("The expected type of
28                     the " + i +
29                     " argument is " + argTypes.get(i).
30                     toString() +
31                     " found " + operand_type.toString() + "
32                     in " +
33                     ts.visit(e, null));
34             }
35         }
36     }
37     return new ErrorT(message + ts.visit(e, null));
38 }
```

4. (18 pt) [Design and implement type rules] Design and implement the type rules for comparison expressions:

BinaryComparator: Let a BinaryComparator be (binary operator e1 e2), where e1 and e2 are expressions.

- (a) (4 pt) Describe the type rules (see the example type rules or test cases provided in the above questions) to support the comparisons of two numbers (NumT)
- (b) (4 pt) Describe the type rules to support the comparison of two lists (listT)
- (c) (10 pt) Implement the type checking rules for number and list comparisons.

Sol

- (a)
 - if e1's type is ErrorT then (binary operator e1 e2)'s type should be ErrorT
 - if e2's type is ErrorT then (binary operator e1 e2)'s type should be ErrorT
 - if e1's type is not NumT then (binary operator e1 e2)'s type should be ErrorT with message : "The first argument of a binary expression should be num Type, found " + e1's type + " in " + expression.
 - if e2's type is not NumT then (binary operator e1 e2)'s type should be ErrorT with message : "The second argument of a binary expression should be num Type, found " + e2's type + " in " + expression.
 - otherwise (binary operator e1 e2)'s type should be BoolT.
- (b)
 - if e1's type is ErrorT then (binary operator e1 e2)'s type should be ErrorT
 - if e2's type is ErrorT then (binary operator e1 e2)'s type should be ErrorT
 - if e1's type is ListT and e2's type is ListT then (binary operator e1 e2)'s type should be ErrorT with message : "arguments of the binary expression are of List Type, found " + e1's type + " and " + e2's type " in " + expression.
 - otherwise (binary operator e1 e2)'s type should be BoolT.
- (c) BinaryComparator

```

1   private Type visitBinaryComparator(BinaryComparator e, Env<Type> env,
2   String printNode) {
3       Exp first_exp = e.first_exp();
4       Exp second_exp = e.second_exp();
5
6       Type first_type = (Type)first_exp.accept(this, env);
7       if (first_type instanceof ErrorT) { return first_type; }
8
9       Type second_type = (Type)second_exp.accept(this, env);
10      if (second_type instanceof ErrorT) { return second_type; }
11
12      if ((first_type instanceof ListT) & (second_type instanceof ListT)) {
13          return new ErrorT("The arguments of the binary expression "
14          + "are of List Type, found " + first_type.tostring() + " and " +
15          second_type.tostring() +
16          " in " + printNode);
17      }
18
19      if (!(first_type instanceof NumT)) {
20          return new ErrorT("The first argument of a binary expression "
21          + "should be num Type, found " + first_type.tostring() +
22          " in " + printNode);
23      }
24
25      if (!(second_type instanceof NumT)) {
26          return new ErrorT("The second argument of a binary expression "
27          + "should be num Type, found " + second_type.tostring() +
28          " in " + printNode);
29      }
30
31      return BoolT.getInstance();
}

```

5. (9 pt) [Design and implement type rules] Design and implement the type checking rules for condition expressions.

IfExp: Let a IfExp be (if cond then else), where cond, then, else are expressions.

- (a) (4 pt) Describe the type rules to support the condition expressions.
- (b) (5 pt) Implement the type checking rules for condition expressions.

Sol

- (a)
 - if cond's type is ErrorT then (if cond then else)'s type should be ErrorT
 - if cond's type is not BoolT then (if cond then else)'s type should be ErrorT with message: "The condition should have boolean type, found " + cond's type + " in " + expression
 - if then's type is ErrorT then (if cond then else)'s type should be ErrorT
 - if else's type is ErrorT then (if cond then else)'s type should be ErrorT
 - if then's type and else's type are typeEqual then (if cond then else)'s type should be then's type.

- else (if cond then else)'s type should be ErrorT with message: "The then and else expressions should have the same " + "type, then has type " + then's type + " else has type " + else's type+ " in " + expression.
- \$ (if 5 56 67)
Type error: The condition should have boolean type, found number in (if 5 56 67)
\$ (if #t #t 56)
Type error: The then and else expressions should have the same type, then has type bool
else has type number in (if #t #t 56)

(b) IfExp

```

1  public Type visit(IfExp e, Env<Type> env) {
2      Exp cond = e.conditional();
3      Type condType = (Type)cond.accept(this, env);
4      if (condType instanceof ErrorT) { return condType; }
5
6      if (!(condType instanceof BoolT)) {
7          return new ErrorT("The condition should have boolean type, found
8                  " +
9                  condType.toString() + " in " + ts.visit(e, null));
10
11     Type thentype = (Type)e.then_exp().accept(this, env);
12     if (thentype instanceof ErrorT) { return thentype; }
13
14     Type elsetype = (Type)e.else_exp().accept(this, env);
15     if (elsetype instanceof ErrorT) { return elsetype; }
16
17     if (thentype.typeEqual(elsetype)) { return thentype; }
18
19     return new ErrorT("The then and else expressions should have the same "
20             + "type, then has type " + thentype.toString() +
21             " else has type " + elsetype.toString() + " in " +
22             ts.visit(e, null));
23 }
```

6. (8 pt) [Typelang programming] For all the above typing rules (total 8 of them) you implement, write a typelang program for each type rule to test and demonstrate your type check implementation. (You can use typelang.g in hw8code.zip as a reference for the syntax of TypeLang). For each expression, put in comments which type rules the expression is exercising.

For example: \$ (list: num 45 45 56 56 67) // test correct types for list expressions
\$(> 45.0 #t) // test incorrect types for binary comparator expressions

```

1
2  $ (let ((x: bool (deref counter))) (if x 5 6)) // test incorrect types for if
   expressions
3  $ Type error: The declared type of the 0 let variable and the actual type mismatch,
   expect bool found number in (let ( (x (deref counter))) (if x 5.0 6.0) )
4  $ (let ((x: num (deref counter))) (if (< x 4) x 6)) // test correct types for if, let
   and binary comparison expressions
5  $ 0
6  $ (define inc: (-> num) (lambda () (set! counter (+ 1 (deref counter))))) // test
   correct types for set! expressions and lambda expressions
```

```

7 $ (inc 4) // test incorrect actual parameter types for call expressions
8 $ Type error: The number of arguments expected is 0 found 1 in (inc 4.0 )
9 $ (inc) // test correct call expression
10 $ 1
11 $ (define 12: List<num> (list: num 1 2 #t)) // test incorrect type for list element
12 $ Type error: The 2 expression should have type number found bool in (list 1.0 2.0 #t
   )
13 $ (define 12: List<num> (list: num 1 7)) // test correct type for list
14 $(let ((first: num (car 12)) (second: num (car (cdr 12)))) (> first second))
15 // test correct types for car, cdr expressions and binary comparator
16 $ #f

```

7. (6 pt) [TypeLang programming] In HW5, you have written a function `processlists` that takes three arguments `op`, `list1`, `list2`, where `op` is a function that takes two pairs as parameters, and `list1` and `list2` are the two lists of pairs. The return value is the result of applying `op` on each pair of `list1` and `list2`. You have also written functions `common` and `diff` to test `processlists`.

Please refer to the examples below:

```

$ (define list1 (list (cons 1 3) (cons 4 2) (cons 5 6)))
$ (define list2 (list (cons 2 6) (cons 4 2) (cons 1 3)))
$ (processlists common list1 list2)
((-1,-1)(4,2)(-1,-1))
$ ((processlists diff list1 list2)
((1,3)(-1,-1)(5,6))
$ (processlists diff list2 list1)
((2,6)(-1,-1)(1,3))

```

In this problem, you are required to write `common` and `diff` and `processlists` in TypeLang. The types of function arguments should be compatible with the types used in the above examples.

Note:

- (a) You will need to modify your code from HW5 to output a default pair (-1,-1) when `op` is not satisfied instead of empty list (), so that the output type is consistent.
- (b) Assume TypeLang supports type checking for recursive functions as you answer this question.

Sol:

```

1
2
3 (define list1 : List<(num,num)> (list: (num, num) (cons 1 3) (cons 4 2) (cons 5 6)))
4 (define list2 : List<(num,num)> (list: (num, num) (cons 2 6) (cons 4 2) (cons 1 3)))
5
6
7 (define common: ((num, num) (num, num) -> (num, num))
8   (lambda (pair1: (num, num) pair2: (num, num))
9     (if (= (car pair1) (car pair2))
10       (if (= (cdr pair1) (cdr pair2))
11         pair1
12         (cons -1 -1)))

```

```
13                     (cons -1 -1))))  
14  
15  
16 (define diff: ((num, num) (num, num) -> (num, num))  
17     (lambda (pair1: (num, num) pair2: (num, num))  
18         (if (= (car pair1) (car pair2))  
19             (if (= (cdr pair1) (cdr pair2))  
20                 (cons -1 -1)  
21                 pair1)  
22             pair1)))  
23  
24 (define processlists: (((num, num) (num, num) -> (num, num)) List<(num, num)> List<(num, num)> -> List<(num, num)>)  
25     (lambda (op: ((num, num) (num, num) -> (num, num)) lst1: List<(num, num)>  
26         lst2: List<(num, num)>)  
27         (if (null? lst1)  
28             lst2  
29             (if (null? lst2)  
30                 lst1  
                 (cons (op (car lst1) (car lst2))(processlists  
op (cdr lst1) (cdr lst2)))))))  
-----
```