

Ecorrection

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Introduction to Quantum Computing

Exam: IN2381 / Final Exam Date: Tuesday 20th July, 2021

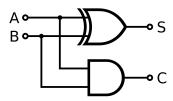
Examiner: Christian Mendl **Time:** 14:15 – 15:45

Working instructions

- This exam consists of 10 pages with a total of 3 problems.
 Please make sure now that you received a complete copy of the exam.
- The total amount of achievable credits in this exam is 60 credits.
- Detaching pages from the exam is prohibited.
- Allowed resources: open book
- Subproblems marked by * can be solved without results of previous subproblems.
- Answers are only accepted if the solution approach is documented. Give a reason for each answer unless explicitly stated otherwise in the respective subproblem.
- Do not write with red or green colors nor use pencils.

Problem 1 (20 credits)

In this problem, you will build a quantum version of a half adder – the basic building block of addition on a classical computer. The most important part of such a circuit is the half-adder:



where A and B are classical bits, $S = A \oplus B$ is the sum modulo two and $C = A \cdot B$ is ordinary multiplication of A and B called the carry. The carry is the part of the summation that adds to the next digit (it is only 1 if both A = 1 and B = 1).



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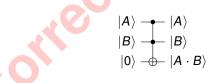
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a) Assume you start in the arbitrary two-qubit state $|AB\rangle$. Provide a quantum gate / series of quantum gates that performs the operation:

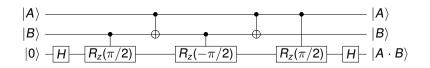
$$\begin{vmatrix} A \rangle - \begin{vmatrix} A \rangle \\ B \rangle - \end{vmatrix} - \begin{vmatrix} A \otimes B \end{vmatrix}$$

$$|A\rangle \xrightarrow{} |A\rangle \\ |B\rangle \xrightarrow{} |A \oplus B\rangle$$

b) The A · B operation can be performed by a Toffoli gate:



Verify that the circuit below performs that operation up to a global phase constant.



Hint: Follow the state of each qubit through the circuit for all 4 possible input basis states $|00\rangle$, $|01\rangle$, $|10\rangle$ and $|11\rangle$ separately.

We account for the activated controlled gates:

Input |00>:

$$|0\rangle \stackrel{}{-} \stackrel{}{H} \stackrel{}{-} \stackrel{}{H} - |0\rangle$$

1 credit

Input |01>:

$$|0\rangle$$
 $-H$ $R_z(\pi/2)$ $R_z(-\pi/2)$ H $|0\rangle$

2 credits

Input |10>:

$$|0\rangle$$
 $-H$ $R_z(-\pi/2)$ $R_z(\pi/2)$ H $|0\rangle$

2 credits

Input |11):

Here we have used that $HR_z(\pi/2)R_z(\pi/2)H = HR_z(\pi)H = -iHZH = -iX$, so there is a global phase constant of -i. 4 credits

c) Build a quantum half-adder using the Toffoli gate and the result from a): i.e. find the circuit that performs the operation:

$$|A\rangle$$
 - $|A\rangle$ - $|A\rangle$ $|B\rangle$ - $|A \oplus B\rangle$ $|A \oplus B\rangle$

You do not need to write out the Toffoli decomposition explicitly.

$$|A\rangle \longrightarrow |A\rangle |B\rangle \longrightarrow |A \oplus B\rangle |0\rangle \longrightarrow |A \cdot B\rangle$$

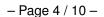
2 credits for Toffoli, 1 for CNOT; -2 if the order is wrong.



d) Given $|A \oplus B\rangle$ and $|A \cdot B\rangle$, is it possible to determine $|A\rangle$ and $|B\rangle$ in all cases? Provide reasoning for your answer.

It is not possible because both $|01\rangle$ and $|10\rangle$ result in $|A \oplus B\rangle = |1\rangle$ and $|A \cdot B\rangle = |0\rangle$. This operation is only reversible with an extra qubit.

4 credits for complete correct answer. 2 credits if this is shown in only one case.



Problem 2 (20 credits)

Consider an ensemble of quantum states $\{p_i, |\psi_i\rangle\}$, where the quantum system is in state $|\psi_i\rangle$ with probability p_i . Recall from the lecture that the density operator ρ of such an ensemble is defined as:

$$\rho = \sum_{i} p_{i} \left| \psi_{i} \right\rangle \left\langle \psi_{i} \right|$$

a) Given the ensemble $\left\{\left(\frac{1}{2},\frac{|0\rangle+|1\rangle}{\sqrt{2}}\right),\left(\frac{1}{2},\frac{|0\rangle+i|1\rangle}{\sqrt{2}}\right)\right\}$, compute ρ and write it in the form:

$$\rho = \frac{1}{2}I + \alpha_x X + \alpha_y Y + \alpha_z Z.$$

What is the connection of $\vec{\alpha} = (\alpha_x, \alpha_y, \alpha_z)$ with the Bloch sphere representation?

Simply by plugging in, expanding and rearranging:

$$\begin{split} \rho &= \frac{1}{2} \frac{\left(\left| 0 \right\rangle + \left| 1 \right\rangle \right)}{\sqrt{2}} \frac{\left(\left\langle 0 \right| + \left\langle 1 \right| \right)}{\sqrt{2}} + \frac{1}{2} \frac{\left(\left| 0 \right\rangle + i \left| 1 \right\rangle \right)}{\sqrt{2}} \frac{\left(\left\langle 0 \right| - i \left\langle 1 \right| \right)}{\sqrt{2}} \\ &= \frac{1}{4} \left(\left| 0 \right\rangle \left\langle 0 \right| + \left| 0 \right\rangle \left\langle 1 \right| + \left| 1 \right\rangle \left\langle 0 \right| + \left| 1 \right\rangle \left\langle 1 \right| \right) + \frac{1}{4} \left(\left| 0 \right\rangle \left\langle 0 \right| - i \left| 0 \right\rangle \left\langle 1 \right| + i \left| 1 \right\rangle \left\langle 0 \right| + \left| 1 \right\rangle \left\langle 1 \right| \right) \\ &= \frac{1}{2} \left(\left| 0 \right\rangle \left\langle 0 \right| + \left| 1 \right\rangle \left\langle 1 \right| \right) + \frac{1}{4} \left(\left| 0 \right\rangle \left\langle 1 \right| + \left| 1 \right\rangle \left\langle 0 \right| \right) + \frac{1}{4} \left(-i \left| 0 \right\rangle \left\langle 1 \right| + i \left| 1 \right\rangle \left\langle 0 \right| \right) \\ &= \frac{1}{2} I + \frac{1}{4} X + \frac{1}{4} Y \end{split}$$

which means that $(\alpha_x, \alpha_y, \alpha_z) = (\frac{1}{4}, \frac{1}{4}, 0)$. Recall that a density operator can be represented as

$$\rho = \frac{I + \vec{r} \cdot \vec{\sigma}}{2},$$

where \vec{r} is the Bloch vector. So $\vec{\alpha} = \frac{\vec{r}}{2}$.

4 credits for correct ρ , 1 credit for ρ in right form, 1 credit for realising this is the Bloch vector over 2

b) Now consider the ensemble

$$\left\{ \left(\frac{1}{2},|0\rangle\right),\left(\frac{1}{2},|1\rangle\right)\right\}$$

and compute its density matrix ρ . Draw a Bloch sphere, clearly labeling $|0\rangle$ and $|1\rangle$, and indicate the position of this ensemble within the sphere.

$$\rho = \frac{1}{2} \left(\left| 0 \right\rangle \left\langle 0 \right| + \left| 1 \right\rangle \left\langle 1 \right| \right) = \frac{1}{2} I$$

This means that the Bloch vector of this ensemble is (0,0,0), i.e., it is located at the origin of the Bloch sphere.

1 credit for correct ρ , 2 credits for correct location of ensemble, 1 credit for correct drawing of a Bloch sphere

0

2 3

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$$U = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \sqrt{1-\lambda} & -i\sqrt{\lambda} & 0 \\ 0 & -i\sqrt{\lambda} & \sqrt{1-\lambda} & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix},$$

where $0 \le \lambda \le 1$, acts on a system of two qubits. The first qubit is initially in an arbitrary state ρ and the second one is initialized at $|0\rangle$. Trace out the second qubit to obtain the two operators E_0 and E_1 which represent the action of U on the first one.

From Exercise 11.2 we recall that

$$\mathcal{E}(\rho) = \sum_{k} E_{k} \rho E_{k}^{\dagger}.$$

$$(E_k)_{\ell,m} = \langle \ell, k | U | m, 0 \rangle$$

2 credits for these or related formulas which show recognition of this being a quantum operation. Therefore,

$$E_0 = \begin{pmatrix} 1 & 0 \\ 0 & \sqrt{1-\lambda} \end{pmatrix}$$
 and $E_1 = \begin{pmatrix} 0 & -i\sqrt{\lambda} \\ 0 & 0 \end{pmatrix}$

Full credits if these two matrices are correct, even if no explanation.

d) Compute the effect of the operators you found on the general density matrix $\rho = \frac{1}{2}I + \alpha_x X + \alpha_y Y + \alpha_z Z$. Interpret their action on the Bloch sphere.

One possible solution is by writing E_0 and E_1 in terms of the Pauli matrices:

$$E_0 = \frac{1}{2}(I+Z) + \frac{\sqrt{1-\lambda}}{2}(I-Z) = \left(\frac{1+\sqrt{1-\lambda}}{2}\right)I + \left(\frac{1-\sqrt{1-\lambda}}{2}\right)Z = aI + bZ$$

$$E_1 = \frac{\sqrt{\lambda}}{2}(-iX + Y) = \frac{-i\sqrt{\lambda}}{2}X + \frac{\sqrt{\lambda}}{2}Y = cX + dY$$

For a general ρ , $\mathcal{E}(\rho) = \frac{1}{2}\mathcal{E}(I) + \alpha_x \mathcal{E}(X) + \alpha_y \mathcal{E}(Y) + \alpha_z \mathcal{E}(Z)$. We compute each term:

$$\mathcal{E}(I) = (aI + bZ)I(aI + bZ) + (cX + dY)I(c^*X + dY)$$

= $a^2I + 2abZ + b^2I + |c|^2I + icdZ - idc^*Z + d^2I = I + \lambda Z$

$$\mathcal{E}(X) = (aI + bZ)X(aI + bZ) + (cX + dY)X(c^*X + dY)$$

= $a^2X - b^2X + |c|^2X - d^2X = \sqrt{1 - \lambda}$

$$\mathcal{E}(Y) = (aI + bZ)Y(aI + bZ) + (cX + dY)Y(c^*X + dY)$$

= $a^2Y - b^2Y - |c|^2Y + d^2X = \sqrt{1 - \lambda}$

$$\mathcal{E}(Z) = (aI + bZ)Z(aI + bZ) + (cX + dY)Z(c^*X + dY)$$

= $a^2Z + 2abI + b^2Z - |c|^2Z - icdI + idc^*I - d^2Z = (1 - \lambda)Z$

4 credits for these or similar Therefore,

$$\mathcal{E}(\rho) = \frac{1}{2}I + \sqrt{1-\lambda}\alpha_x X + \sqrt{1-\lambda}\alpha_y Y + \left(\frac{\lambda}{2} + (1-\lambda)\alpha_z\right)Z,$$

and $\vec{\alpha}' = \left(\sqrt{1-\lambda}\alpha_x, \sqrt{1-\lambda}\alpha_y, (\frac{\lambda}{2}+(1-\lambda)\alpha_z)\right)$. 1 credit for this or similar This is an amplitude damping channel. (This identification is not required.) It shrinks the Bloch sphere (0.5 credits) towards the north pole. (0.5 credits)

Problem 3 (20 credits)

We consider a quantum system of n qubits, and use the notation X_j , Y_j , Z_j to denote that one of the Pauli matrices acts on the jth qubit; e.g., $X_1Z_3 \equiv X \otimes I \otimes Z$ for n = 3.

Conjugation by U refers to the transformation UgU^{\dagger} of a quantum gate g by a unitary operation U. The following table summarizes several conjugation transformations:

Here $S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$ is the phase gate.

a) State the check matrix representation of $g_1, g_2 \in G_4$ given by

$$g_1 = Z \otimes Y \otimes X \otimes X,$$

$$g_2 = X \otimes Z \otimes I \otimes Y.$$

Based on this representation, show that g_1 anti-commutes with g_2 .

The check matrix of (g_1, g_2) is

$$\begin{pmatrix} r(g_1) \\ r(g_2) \end{pmatrix} = \begin{pmatrix} 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix}$$

(3 credits)

Now we recall that elements of the Pauli group either commute or anti-commute. (0.5 credits) According to exercise 12.2 (e), g_1 commutes with g_2 precisely if $r(g_1)\Lambda r(g_2)^T = 0 \mod 2$, with

$$\Lambda = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}.$$

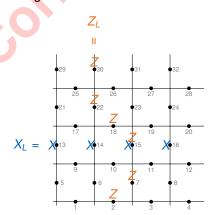
(0.5 credits)

A effectively interchanges the first with the second half of a check row. We obtain

$$r(g_1)\Lambda r(g_2)^T = \begin{pmatrix} 0 & 1 & 1 & 1 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 & 0 & 1 & 1 & 0 & 0 & 1 \end{pmatrix}^T = 3 = 1 \mod 2,$$

i.e., they do not commute, and hence must anti-commute. (1 credit)

b)* Given a square lattice, a qubit is associated with each *edge* of the lattice (dots in the figure below). We define logical Pauli operators X_L and Z_L as tensor products of strings of X and Z operators: $X_L = X_{13}X_{14}X_{15}X_{16}$ and $Z_L = Z_2Z_7Z_{15}Z_{18}Z_{22}Z_{30}$, as visualized in the figure.



Show that X_L and Z_L anti-commute, i.e., $X_L Z_L = -Z_L X_L$. How can one define a logical Y_L operator such that X_L , Y_L , Z_L satisfy the anti-commutation relations of the Pauli-matrices?

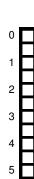


$$\begin{split} X_L Z_L &= Z_2 Z_7 X_{13} X_{14} X_{15} Z_{15} X_{16} Z_{18} Z_{22} Z_{30} \\ &= Z_2 Z_7 X_{13} X_{14} (-Z_{15} X_{15}) X_{16} Z_{18} Z_{22} Z_{30} = -Z_L X_L. \end{split}$$

3 credits

The Pauli-Y matrix can be expressed in terms of X and Z via Y = iXZ. We can use this relation to analogously define $Y_L = iX_LZ_L$. Then Y_L anti-commutes with X_L since X_L commutes with itself and anti-commutes with Z_L , as we have just shown. Likewise, Y_L anti-commutes with Z_L .

2 credits; alternative solutions possible



c)* The subgroup $R = \langle X_1 Y_2 Z_3, Y_1 Y_2 Y_3 \rangle$ of the Pauli group G_3 stabilizes the subspace $V_R = \text{span}\{|\chi_0\rangle, |\chi_1\rangle\}$ with

$$|\chi_0\rangle = \frac{1}{2} \left(|000\rangle + |001\rangle + i|110\rangle - i|111\rangle \right), \qquad |\chi_1\rangle = \frac{1}{2} \left(|010\rangle + |011\rangle - i|100\rangle + i|101\rangle \right).$$

(A proof of this statement is not required here.) Determine the result (eigenvalue) when measuring the operator $Y_1Y_2Y_3$ with respect to the quantum state $(S \otimes H \otimes (SH)) | \chi_0 \rangle$, where S is the phase gate.

In general, an eigenvalue equation $M|\psi\rangle = \lambda |\psi\rangle$ is equivalent to $UMU^{\dagger}U|\psi\rangle = \lambda U|\psi\rangle$ for any unitary matrix U. Setting $U = S \otimes H \otimes (SH)$ here, we search for an operator M such that $UMU^{\dagger} = Y_1Y_2Y_3$, with formal solution

$$M = U^{\dagger}(Y_1 Y_2 Y_3)U = (S^{\dagger} YS) \otimes (HYH) \otimes (HS^{\dagger} YSH).$$

Here we have already used that $H^{\dagger} = H$. From the above conjugation table, we see that $SXS^{\dagger} = Y$, i.e., $S^{\dagger}YS = X$, as well as HYH = -Y and $HS^{\dagger}YSH = HXH = Z$. In summary,

$$M = X \otimes (-Y) \otimes Z = -X_1 Y_2 Z_3$$
.

Since R stabilizes $|\chi_0\rangle$ and $X_1Y_2Z_3\in R$, the state $|\chi_0\rangle$ is an eigenstate of M with eigenvalue (-1). In particular, the measurement result will be (-1) with probability 1.

d)* We consider the two qubit code $C = \text{span}\{|0_L\rangle, |1_L\rangle\}$ with $|0_L\rangle = |00\rangle$ and $|1_L\rangle = |01\rangle$. It is affected by a simultaneous bit flip noise process described by the operation elements $E_0 = \frac{1}{\sqrt{2}}I_4$ and $E_1 = \frac{1}{\sqrt{2}}X \otimes X$, where I_n the $n \times n$ identity matrix. Show that this noise process is error-correctable.

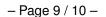
We make use of the quantum error-correction conditions (see lecture). The projector onto C in the present case is $P = |00\rangle \langle 00| + |01\rangle \langle 01|$. The code is error-correctable precisely if

$$PE_k^{\dagger}E_{\ell}P = \alpha_{k\ell}P$$

for all $k, \ell \in \{0, 1\}$ and some Hermitian matrix $(\alpha_{k\ell})$ of complex numbers. We first note that $E_k^{\dagger} E_k = \frac{1}{2} I_4$ for $k \in \{0, 1\}$ since $X^2 = I_2$. Thus the condition is satisfied via $\alpha_{kk} = \frac{1}{2}$ in the case $k = \ell$. We now explicitly compute

$$PE_0^{\dagger}E_1P=PE_1^{\dagger}E_0P=\frac{1}{2}P(X\otimes X)P=\left(\left|00\right\rangle \left\langle 00\right|+\left|01\right\rangle \left\langle 01\right|\right)\left(\left|11\right\rangle \left\langle 00\right|+\left|10\right\rangle \left\langle 01\right|\right)=0,$$

i.e., the condition is satisfied via $\alpha_{01} = \alpha_{10} = 0$ for $k \neq \ell$. In summary, the quantum error-correction conditions hold true for all combinations of $k, \ell \in \{0, 1\}$.



Additional space for solutions-clearly mark the (sub)problem your answers are related to and strike out invalid solutions.

