	the transfer of the transfer o
A -	Big O natation
4 -	Analyzing porformence of orange and objects.
1 -	Problem Solving opproun and pattern
A .	Recursion
-	Searching Algorithm.
•	Bubble Sort, selection sort, insortion sort
	mery sort, quick sort, Roder Sort.
. ^.	
div -	intro to pata structure.
	- Singley linked list, doubly linked hist. - Stocks and gueus
	- Stocks and gueus
	- Binory Sevich trees.
(3/1)	True troversal
A stellar	Binory Heaps
	- Hash tubles.
	- Gropus Troverson
	Dijk stor Pijkstra's olyanithm.
76	PIJE ISHOT PIJESTICES
3,0	And the second of the first of the second of
15 o.)	Live to the same of the same o

Big O natation

Suppose we have to odd the tatol number. possed like if we pass 5 50 it will be 1+2+3+4+5= 15 foctoriof.

Ost was using for loop

furthion (

Const oddupto=(n) => { Cet tuty = 0;

for (+i = 0; i<=n; if+) {

take = feeter + i 1/ Shorthone freque + = i. consola log (tuty)

Edd of to forth (6)

Quil 0,1,3,6,10,19

- Sunction old up to (n) {

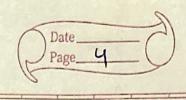
return n+ (n+1)/2;

Both ore Some

what does bottor mean?

· mme redoble?

	To cheek performance we will use buit in
	performance now co frection.
	drawn fort soit the of man in 20/202
-	· vor +1 = Performance. nois ();
	Dd upto1(70000000);
•	vor tz = (rfammer, now ())
	constelly ('time Flogsells & Et2-t1)/10003 suoria's
	A Carpbinson A.
	Some wit oddupto 2 () byth time
	The second of th
	of is 3(+3): 4=21:0=110 x0
	+mi Elopsuls: 0.011.6686.61, silonels.
	med Elopsud2: 0.0000587-4 seconds.
	The state of the s
	Second won one is was foster in the
1	
	Problem with time.
	· Diffrent mothing will richel different Home
	· Some mothine will record defort time.
	· for fost algorithms, Speed missement may nut
	be prosia enaight. 1) to my
	To hondle this Big to notation comosting
	Pricture.
	I posm vota I dols to Med ?
-	1 901 mm, C + 25/4
	Carporter - promove 2245.
	Suldebor som.



If not time, then what?

Rather than counting sciends, which are so vovioble ...

lets count the number of simple operations the Computer has to perform

we con use this

function oddupto (n) { return no (no) 102;

Here we have 3 operation.

35 mple operation rugless the size of the operations of

function odd up to (n) {

Let today = 0;

Tossignat.

for (1 Let i 30; i < = n; i (+ +) {

tatu (+) = i

n odditions one

n assignments.

return tutul:

Tone operation but it depends on a

i's n'= 10 then it's loopings if n = 10 thm it's 10000 gration

So it is n operations
assignment operation also

as nyrous chechen grows of a

R	.'.	n
D	17	U
	U	

Big O Natation is a way to formalize

tert we had the time of I

it allows us to tuk formally obout how the runtime of an algorithm grows as the input grow.

Bry O defination:

if the number is of simple spreations the compative has to do to is eventually less than a constant time & f(n), as n incress.

- . f(n) could be linear (f(n) = n)
- · f(n) could be quodretic (f(n)=12)
- ·f(n) could be constant (f(n)=1)
- . f (n) well be something enount defrent.

we are tulking about uppor board.

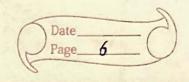
orld up to 1 = o(1) Aways 3 growth on

odoupto 1 - O(n) number of speration

i's (eventually)

bounded by a multiply of n

(soy | lon)



count pord down

function (autopmoddown (n) {

for (let i=0, i <= on; i++) {

Console ly (i)

}

for ((uti-n, i>=0;j--){

(onsul, lay (i))

3

Contapordown (5)
Print all pains

function Print de pairs (n) {

for (let i=0; i<n; i++)? = O(n)

for (let s=0; j<n; j++) s = O(n) * O(n)

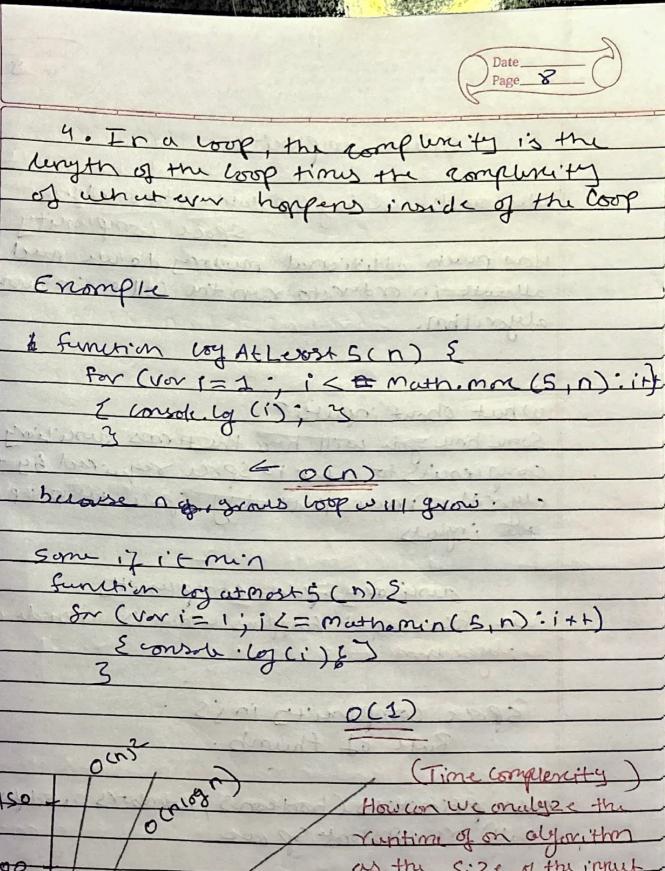
Console leg (i; i)

3

Roma all fors (5)

- Lord Daniel to see to more Cn2) o svert

Simplifying Big O expressions.
Constants Don't matter
O(2n) O(n)
O(2n) $O(n)$ $O(500)$ $O(1)$
$O(13n^2)$ $O(n^2)$
Smaller Terms Done matter
o(n+10) o(n)
0(1000n+50) 0(n). 1.1.
0 (Us+ 2U+8) 0 (Us) .
By O Shorth ands.
3 81111 11 81123.
· Analyzing complimity with big o con
get complicated. (2) vol. 1.
con help
, these rules work & Always work, but
or helpful starting point.
6) vorioble assignment is constant
(2) Allers in all much is constant
or object (by key) is constant



Time complexity

How can we analyze the

runtime of an algorithm

as the 5:2e of the input

increses

octob

octob

so

octob

o

Spale complexity.

we con also use by a putation to onalyzes space complexity

allow much order to run the code in our algorithm.

Some time you will have the turns ownitions space comprenity to refer to space required by the algorithm i mus including space turn up by the inputs

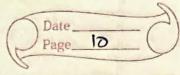
auxilorg Spac compunity
we follow what hoppins what is itisade of
algorithm

Spore complements in is Rule of thumb.

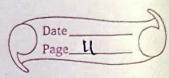
most primitivis Chordeons, numbers, undefined

Strings require our Space (where nisther
string length)

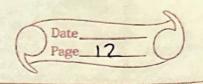
Rywarle types ore growing o(n), which nis the length or (for orrows) or the rumber of keys (for objects)



	An Example.
	Story is the supplemental Second at the order
	function Sun (orr) 2
	let tutul = 0' = one number
	for (luti = 0; i < orr. length; i++) {
2 011	Tuty += or [i]; (or hary total = tatal + corr[i]
2010	y tontinue
	return tutul;
	3
	it man we have Constant Space
	= O(1) SPS(2)
1	Finition double (orv) &
	Let newAvr = [];
	for (1ct: =0; i'c orr. length; i++) {
	newArr. p. sh (2 + or si7):
	on input med it want areas
	return new Arr; a nonumbers.
	3
	dauble ([1,2,3,4])
	at Amino mo (n) sport
1647	marine of more and hillede them prode
1000	in most well that entry with the ment
	(me
	(S) 53 do.
	251-57
	201361-1-10-10-10-10-10-10-10-10-10-10-10-10-
11/2 6 1 1 2	



Common complexities: 0(1), 0(1), 0(12) Sometime by a enpressions involves more compressions ore logorithm! what is log ogoin? log (8) =8 -> 23 =8 (og (volue) = enforent -) zurponent = volue will omit the 2 log = = = coj 2 , 11 1 we ore young to see just general trinel Keele of thems 8,5,13) Habel The log of a number roughtly measured the number of times you can divide their number by 2 (byour you get value that's less than or expect orth = 2(8) : 2 6.25 2 + 2 4 + 2 2 ÷ 2 3.128 120.78125 (og (8) = 3 lgg(25) = 4.64



Logarithm Complemites

Logovitmic time complnity is great.

(ocnosm) 0(1gh) =

ueho cores

Certoin Searching algorithm have logorithmic time complements.

Efficient st sorting objection involves Renovion sometimes involves vogovithmic Spore complementy.

To onalyze the preformance of on algorithm, we use Big O Notation. uses · Big O Notation congives us a night level understanding

of the time or spore complexity of on algorithm · Big a Neutrion does to come shout pruisin, only Shout general trends (linear) quadratic ?, constant? · The timeor spore compaining is mosural by Big O