

Experience with and Plans for Rascal

a DSL for software analysis and transformation

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Supporting Modular, Extensible & Efficient Equation Solving in Rascal



Rascal is a JVM language

A One-Stop Shop
for ...

suitable for

- Meta-programming
 - Software analysis & transformation
 - Compiler construction
 - Applications: software metrics, refactoring, repository analysis, code generation
- Design & implementation Domain-Specific Languages
 - Applications: gaming, questionnaires, banking, tax regulation, laws and legal analysis, ...



Rascal Features

- Sophisticated built-in datatypes: list, set, relation, listrelation, map, datetime, location, ...
- *Immutable* values, but mutable variables (references to immutable values)
- Static types with local type inference
- Pattern matching on all values
- (Higher order) functions using pattern-based dispatch
- Visiting/traversing values
- Syntax definitions and parsing
- Concrete syntax values
- Familiar (Java-like) syntax
- Compiled to JVM byte code
- Java and Eclipse integration
- Command line (REPL)



Rascal Applications

- Compiler for Numerical Simulation Language (Magnolia)
 - Compiler for GPU language
 - Rascal to JVM compiler
 - Software metrics (OSSMETER)
 - PHP analysis (PHP AiR)
 - Java Refactoring
 - Hibernate performance analysis
 - Javascript analysis & transformation
- DSLs for
- Digital Forensics
 - Financial Transactions
 - Game Economics
 - Tax Forms
 - Accountancy
 - Legal rules



Rascal Ecosystem

- Rascal Language
 - Rascal interpreter
 - Eclipse integration
 - Rascal Libraries
 - Data types, statistics, ...
 - Command Line Interface (REPL)
- • Rascal to JVM compiler (uses coroutines internally to implement pattern matching)
- • Compiler-based REPL



Today's Topic: solving equations

- Rascal already provides everything you need for static analysis.
- From day one Rascal has supported a **sol**ve statement to support fixed point equations. Used in many applications with acceptable performance
- Now we want to
 - modularize the solve statement to support large sets of equations
 - use/adapt/develop more efficient solution techniques



Example 1:

Transitive Closure (datalog)

Given a binary relation r , its transitive closure can be defined as follows:

```
trans(e1, e2) :- r(e1, e2).  
trans(e1, e3) :- r(e1, e2), trans(e2, e3).
```


Example 1: Transitive Closure (Rascal **solve**)

```
rel[int,int] trans(rel[int,int] r){  
  
    rel[int,int] t = r;  
  
    solve (t) {  
        t += (t o r);  
    }  
  
    return t;  
}
```

Initialize t to r

Iterate until fixed point
of t is reached

Composition

≡

```
t1 = {};  
while (t != t1) {  
    t1 = t;  
    t += (t o r);  
}
```



Aside

- This is just a very simple example to illustrate **solve**
- Rascal has very good built-in support **for transitive closure, reachability**, and other **relational algebra** operators
- Datalog approach is based on **implications** and a **search procedure**
- Rascal approach is **constructive & programmable**



Example 2:

Dataflow equations (Rascal **soLve**)

```

rel[stat,def] liveVariables(rel[stat,var] DEFS, rel[stat, var] USES, rel[stat,stat] PRED){

  set[stat] STATEMENT = carrier(PRED);
  rel[stat,def] DEF  = definition(DEFS);
  rel[stat,def] USE  = use(USES);

  rel[stat,def] LIN = {};
  rel[stat,def] LOUT = DEF;

  solve (LIN, LOUT) {
    LIN  = { <S, D> | stat S <- STATEMENT,  def D <- USE[S] + (LOUT[S] - (DEF[S])) };
    LOUT = { <S, D> | stat S <- STATEMENT,  stat Succ <- successors(PRED,S),
                                             def D <- LIN[Succ] };
  }
  return LIN;
}

```

Dragon book solution

```

for each block  $B$  do  $in[B] := \emptyset$ ;
while changes to any of the  $in$ 's occur do
  for each block  $B$  do begin
     $out[B] = \bigcup_{\substack{S \text{ a suc-} \\ \text{cessor of } B}} in[S]$ 
     $in[B] := use[B] \cup (out[B] - def[B])$ 
  end

```

Fig. 10.33. Live variable calculation.



Example 3:

Expression Simplification in PHP AiR

Apply available normalization functions to simplify an expression:

```
Expr simplifyExpr(Expr e, loc baseLoc) {  
    e = normalizeConstCase(inlineMagicConstants(e, baseLoc));  
    solve(e) {  
        e = algebraicSimplification(simulateCalls(e));  
    }  
    return e;  
}
```

See: Hills, Klint, & Vinju: Static, Lightweight Includes Resolution for PHP, ASE, 2014

Uses a 4.5 MLOC PHP corpus that has now been extended to 27.5 MLOC
(not yet published work)



General Format of **so**lve

$r_1 = \text{init}_1;$

...

$r_n = \text{init}_n;$

We can also specify an upper bound on the number of iterations here

solve (r_1, \dots, r_n) {

$r_1 = \{ \langle x, y \rangle \mid \langle x, y \rangle \leftarrow r_1, c_1(r_1, \dots, r_n) \}$

...

$r_n = \{ \langle x, y \rangle \mid \langle x, y \rangle \leftarrow r_n, c_n(r_1, \dots, r_n) \}$

}

Assessment of solve

Readability	Good, declarative, high abstraction level
Direction	Solutions can grow or shrink towards a solution
Information Use	Completely open, any visible variable (bound to AST, table, auxiliary relation, ...) may be queried
Complexity	Turing complete, termination not guaranteed (but # of iterations can be restricted), speed in EXP
Safety	Immutable data
Algorithm	Brute force, visit all elements in each iteration
Modularity	Bad, solve imposes a lexical scope for mutually recursive relations



How to achieve Modular, Extensible, Efficient Equation Solving in Rascal?

(A sample of) Related Work

- Paige & Koenig, Finite Differencing of Computable Expressions, 1982.
- Whaley, Avots, Carbin & Lam, Using Datalog with Binary Decision Diagrams for *Program Analysis*, 2005
- Liu & Stoller, Dynamic Programming via Static Incrementalization, 2008.
- O. de Moor et al., .QL Object-oriented Queries made Easy, 2007
- M. Bravenboer and Y Smaragdakis., Exception analysis and points-to analysis: better together. 2009
- T. Veldhuizen, Leapfrog triejoin: A simple worst-case optimal join algorithm, 2012
- M. Arntzenius & N. R.Krishnaswami, Datafun: a Functional Datalog, 2016



Design Considerations

- Profit from the success of Datalog variants in static analysis
- Integrate with Rascal's immutable values, mostly functional semantics and syntactic style
- Support open extension and modularity for solve
- Create a good match with efficient implementation techniques (e.g., finite differencing, magic sets, BDDs, TrieJoin, ...)
- Disclaimer: first, exploratory, ideas!



Steps Towards an open, modular **soLve** statement

```
rel[int,int] trans(rel[int,int] r){  
  
    rel[int,int] t = r;  
  
    solve (t) {  
        t += (t o r);  
    }  
  
    return t;  
}
```

Declare context values.
Together with the variable name
this **identifies** a specific fixed
point computation

```
fix rel[int,int] t(rel[int,int] r);  
fix rel[int,int] t() = r;  
fix rel[int,int] t() += t() o r;
```

Initial value of t

Increments to t

Given the context r solve t

```
rel[int,int] trans(rel[int,int] r) = t(r);
```



A classic PointsTo Analysis in a DataLog (bddbddb)

2.2 Example

Algorithm 1 Context-insensitive points-to analysis with a precomputed call graph, where parameter passing is modeled with assignment statements.

DOMAINS

V	262144	<code>variable.map</code>
H	65536	<code>heap.map</code>
F	16384	<code>field.map</code>

See: Whaley, Avots, Carbin & Lam,
Using Datalog with Binary Decision Diagrams for
Program Analysis, 2005

RELATIONS

input	vP_0	$(variable : V, heap : H)$
input	$store$	$(base : V, field : F, source : V)$
input	$load$	$(base : V, field : F, dest : V)$
input	$assign$	$(dest : V, source : V)$
output	vP	$(variable : V, heap : H)$
output	hP	$(base : H, field : F, target : H)$

RULES

- (1) $vP(v, h) : - vP_0(v, h).$
- (2) $vP(v_1, h) : - assign(v_1, v_2), vP(v_2, h).$
- (3) $hP(h_1, f, h_2) : - store(v_1, f, v_2), vP(v_1, h_1), vP(v_2, h_2).$
- (4) $vP(v_2, h_2) : - load(v_1, f, v_2), vP(v_1, h_1), hP(h_1, f, h_2).$



PointsTo: classic **soLve**

```
alias VP = rel[V variable, H heap];
alias HP = rel[H base, F field, H target];
alias FS = rel[V base, F field, V source];
alias FL = rel[V base, F field, V destination];
alias ASG = rel[V dest, V source];
```

Convenience abbreviations for
some types

```
tuple[VP vP, HP hP] pointsTo1(VP vP0, FS store, FL load, ASG assign){
```

```
vP = vP0;  
hP = {};
```

(1) $vP(v, h) :- vP0(v, h).$

(2) $vP(v1, h) :- assign(v1, v2), vP(v2, h).$

```
solve(vP, hP){
```

(3) $hP(h1, f, h2) :- store(v1, f, v2), vP(v1, h1), vP(v2, h2).$

```
vP += { <v1, h> | <v1, v2> <- assign, <v2, h> <- vP };
```

```
hP += { <h1, f, h2> | <v1, f, v2> <- store, <v1, h1> <- vP, <v2, h2> <- vP };
```

```
vP += { <v2, h2> | <v1, f, v2> <- load, <v1, h1> <- vP, <h1, f, h2> <- hP };
```

```
}
```

```
return <vP, hP>;
```

(4) $vP(v2, h2) :- load(v1, f, v2), vP(v1, h1), hP(h1, f, h2).$

```
}
```

PointsTo: open & modular

```
fix VP vP(VP vP0, FS store, FL load, ASG assign);  
fix HP hP(VP vP0, FS store, FL load, ASG assign);
```

Signature and context
for vP and hP

```
fix VP vP() = vP0; // (1)
```

```
fix HP hP() = {};
```

```
fix VP vP() += { <v1, h> | <v1, v2> <- assign, <v2, h> <- vP() }; // (2)
```

```
fix HP hP() += { <h1, f, h2> | <v1, f, v2> <- store,  
                           <v1, h1> <- vP(), <v2, h2> <- vP() }; // (3)
```

```
fix VP vP() += { <v2, h2> | <v1, f, v2> <- load,  
                           <v1, h1> <- vP(), <h1, f, h2> <- hP() }; // (4)
```

```
tuple[VP vP, HP hP] pointsTo2(VP vP0, FS store, FL load, ASG assign) =  
  <vp(vP0, store, load, assign), hp(vP0, store, load, assign)>;
```



Discussion

- Initial experiment, expect further syntactic/semantic improvements
- We are extending our **capsule** immutable collections library* with bidirectional multi-maps to support incremental computation (in finite differencing style) on binary relations
- Still open: what implementation technique is best suited?

(*) See Steindorfer & Vinju, Optimizing Hash-array Mapped Tries for Fast and Lean Immutable JVM Collections, 2015



You are invited to join!

If you want to learn more about Rascal:

- <http://www.rascal-mpl.org>
- <http://tutor.rascal-mpl.org>
- <http://stackoverflow.com/questions/tagged/rascal>

If you are interested in source code:

<https://github.com/usethesource/rascal>

If you want to give us feedback:

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