



COMPILER CONSTRUCTION CS F363

Code Optimization

Dr. R. Elakkiya, AP/CS

Text Book

✓Aho, Lam, Sethi and Ullman, "Compilers-Principles, Techniques & Tools", Pearson/Addison-Wesley, Second Edition, 2013.





Text Book Reading:

Chapter 6

Section 6.1 to 6.9

✓ Use algebraic identities to simplify computation

$$\sqrt{x} + 0 = 0 + x = x$$

$$\sqrt{x * 1} = 1 * x = x$$

$$\sqrt{x} - 0 = x$$

$$\sqrt{x/1} = x$$

$$\sqrt{x * 0} = 0 * x = 0$$

✓ Reduction in strength

Expensive	Cheaper
χ^2	x * x
2 * x	x + x
x / 2	x * 0.5

✓ Constant folding:

- ✓ Evaluate expression during compile time, if both operands values of an operation are constants. Replace the constant expressions by its value.
- ✓ Replace variables that have constant values with their values.

$$\sqrt{\text{Ex}}$$
: tmp=5*3+8-12/2 tmp=17





✓ Commutativity and associativity

$$\checkmark x * y = y * x$$

$$\sqrt{a} = b + c$$
;

$$\sqrt{e} = c + d + b$$
;

$$\sqrt{a} = b + c$$

$$\sqrt{t} = c + d$$

$$\sqrt{e} = t + b$$

If t is not needed outside this block, we can change this sequence to

$$a = b + c$$

$$e = a + d$$



- ✓ Common sub-expressions: Instructions that compute a value that has already been computed
- ✓ Types:
 - ✓ Local common sub-expression: common sub-expressions that appear in single basic block
 - ✓ Global common sub-expression: common sub-expressions that appear across basic block
 - ✓ Examples

```
a = 10;

b = a + 1;

c = a + 1;

d = c + a;

a = 10;

c = a + 1;

d = c + a;

Assuming b's value is not needed anymore
```





✓ Consider the BB and the DAG

$$a = b + c$$

$$b = a - d$$

$$c = b + c$$

$$d = a - d$$

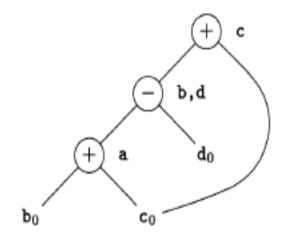
√The optimized BB (asuming b is not live)

$$a = b + c$$

$$d = a - d$$

$$c = d + c$$

b = d (required if b is live)







✓ Consider the BB and the DAG

√The optimized BB (asuming a is live)

- ✓ Dead code
 - ✓ Instructions that compute a value that is never used
 - ✓ Code in a branch that is never taken





- ✓ Copy propagation:
- ✓ After a 'copy' statement, x=y, try to use y as far as possible

$$t=I*4$$

$$a=s+4$$

$$t=I*4$$

$$a=t+4$$

- Code motion: Move the loop-invariant computation outside the loop
- Loop-invariant: An expression that gives the same result, independent of no of times loop is executed.

```
while (i<= limit-2)
{
/* value of limit is not changed
}</pre>
```

- limit-2 is computed every time the loop is executed
- If value of limit is not changed, then we can move limit-2 outside the loop

```
t = limit-2
while (i <= t)
```





Tutorial

- ✓ Consider the following basic block, in which all variables are integers and ** denotes exponentiation.
- ✓ Assume that the only variables that are live at the exit of this block are v and z. In order, apply the following optimizations to this basic block. Show the result of each transformation.
- 1. algebraic simplification
- 2. common sub-expression elimination
- 3. copy propagation
- 4. constant folding
- 5. dead code elimination

a := b + c

z := a ** 2

x := 0 * b

y := b + c

w := y * y

u := x + 3

v := u + w

Is the resulting program optimal? What optimizations, in what order, can you apply to optimize the result



Tutorial

Original BB

a := b + c

$$a := b + c$$

$$x := 0 * b$$
 $x := 0$

$$x := 0$$

$$y := b + c$$

$$y := b + c$$
 $y := b + c$

$$w := y * y$$
 $w := y * y$

$$u := x + 3$$
 $u := x + 3$

$$u := x + 3$$

$$v := u + w$$

$$v := u + w$$
 $v := u + w$

1) Algebraic optimization: 2) Common sub-expression elimination:

$$a := b + c$$

$$z := a * a$$

$$x := 0$$

$$w := y * y$$

$$u := x + 3$$

$$v := u +$$

3) Copy propagation: 4) Constant folding:

$$a := b + c$$

$$x := 0$$

$$w := a * a$$

$$u := 0 + 3$$

$$v := u + w$$

$$a := b + c$$

$$x := 0$$

$$v := u + w$$

5) Dead code elimination:

$$a := b + c$$

$$z := a * a$$

$$u := 3$$

$$v := u + w$$

Tutorial

Final optimized code

a := b + c

z := a * a

v := 3 + z

Thank you

