

## Queue object

*buf*



*lock's wait queue*

*C1*



*lock's blocked queue*



There are two kinds of threads -- consumer threads  $C1, C2, \dots$ , that remove characters from  $buf$ , and producer threads  $P1, P2, \dots$ , that add characters to  $buf$ . For our purposes,  $buf.size()$  returns the number of characters in the buffer.

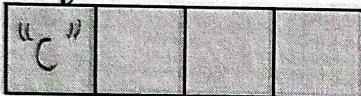
The buffer  $buf$  is initially empty.

1. Consumer  $C1$  enters the synchronized block for the *get* method
2.  $buf.size() == 0$  is true
3.  $wait()$  is executed, placing  $C1$  on the lock's *wait queue*

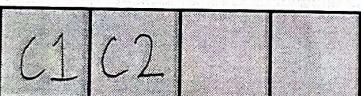
Show the status of  $buf$ , lock's *wait queue* and lock's *blocked queue*.

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1. Consumer 2 ( $C_2$ ) is just about to enter the synchronized block for the *get* method, but has not acquired the lock.
2. Producer  $P_1$  enters the synchronized method *put*, acquires the lock, places the character "c" into *buf*, and calls *notify()*.
3.  $C_1$  is woken up by the notify and must reacquire the lock before proceeding. Thus both  $C_1$  and  $C_2$  are competing for the lock.

Show the status of *buf*, *lock's wait queue* and *lock's blocked queue*.

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*lock's blocked queue*



1. One of  $C_1$  and  $C_2$  is non-deterministically chosen to get the lock. Let's say  $C_2$  gets the lock. It gets to enter the method since  $C_1$  is awake it is put on the *blocked queue*, not back on the *wait queue*.
2.  $C_2$  gets the character and releases the lock which is then acquired by  $C_1$ .

Is there a character in *buf* for  $C_1$  to get? No

What will happen in the program as written?

$C_1$  will be blocked because it is waiting for an item in the

What would have happened if the *while* loop was not buffer.  
in the *get()* method?

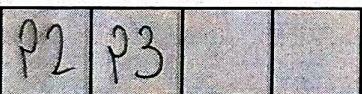
The program would have tried to take an item out  
of the empty buffer which would result in an Exception.

## Queue object

*buf*



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Let's look at a scenario that shows the need for `notifyAll` instead of `notify` in the code.

To make this easy, assume a buffer size of 1. Producer and consumer threads are named as before. *buf* is initially empty.

1. *P1* puts a “c” into the buffer.
2. *P2* attempts a put, checks the while loop and performs a `wait()`
3. *P3* attempts a put, checks the *while* loop and performs a `wait()`

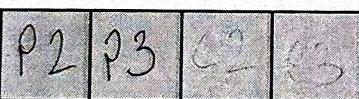
Show the status of *buf*, *lock's wait queue* and *lock's blocked queue*.

## Queue object

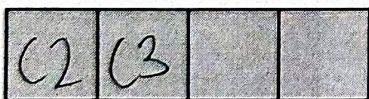
*buf*



*lock's wait queue*



*lock's blocked queue*



4. The following happen at time step 4:

- C1* attempt to get 1 character and enters the *get* method;
- C2* attempts to get 1 character but blocks on entry to the *get* method;
- C3* attempts to get 1 character but blocks on entry to the *get* method;

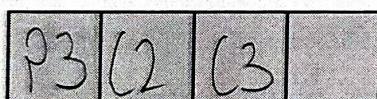
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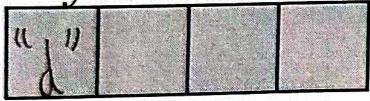
5. The following happen at time step 5.

- a.  $C_1$  is executing the *get* method, gets the character, calls *notify* and exits the method (releasing the lock and giving  $C_2$  and  $C_3$  a chance to acquire it);
- b. The *notify* wakes up  $P_2$
- c. BUT,  $C_2$  enters the method before  $P_2$  can ( $P_2$  must reacquire the lock), so  $P_2$  blocks on entry to the *put* method;
- d.  $C_2$  checks the wait loop, sees there are no more characters in the buffer and so it waits (releasing the lock in the process)
- e.  $C_3$  enters the method after  $C_2$ , but before  $P_2$ , checks the wait loop, sees there are no more characters in the buffer, and so it waits

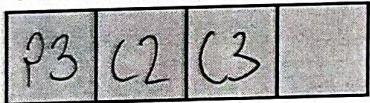
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## Queue object

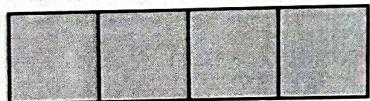
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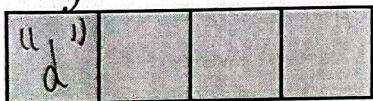
6. The following happen at time step 6.

- Now  $P_3$ ,  $C_2$  and  $C_3$  are all waiting!
- $P_2$  acquires the lock, puts a "d" in the buffer, calls *notify* and exits the method

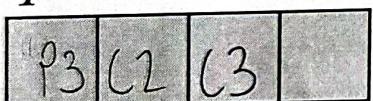
**Show the status of *buf*, lock's *wait queue* and lock's *blocked queue*.**

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7. The following happens at time step 7.

- $P_2$ 's notification wakes up  $P_3$  (any thread can be woken up)
- $P_3$  checks the wait loop condition. There is already a character ("d") in the buffer and so it waits.

Show the status of *buf*, *lock's wait queue* and *lock's blocked queue*.

Is it possible for any thread to be woken up by another notify? Yes, any of the three threads in deadlock can be woken up.

What would have happened if in 6b a `notifyAll()` was called?

A `notifyAll()` call would have awakened all of the threads and one by one each would reacquire the lock and recheck their respective while conditions.