

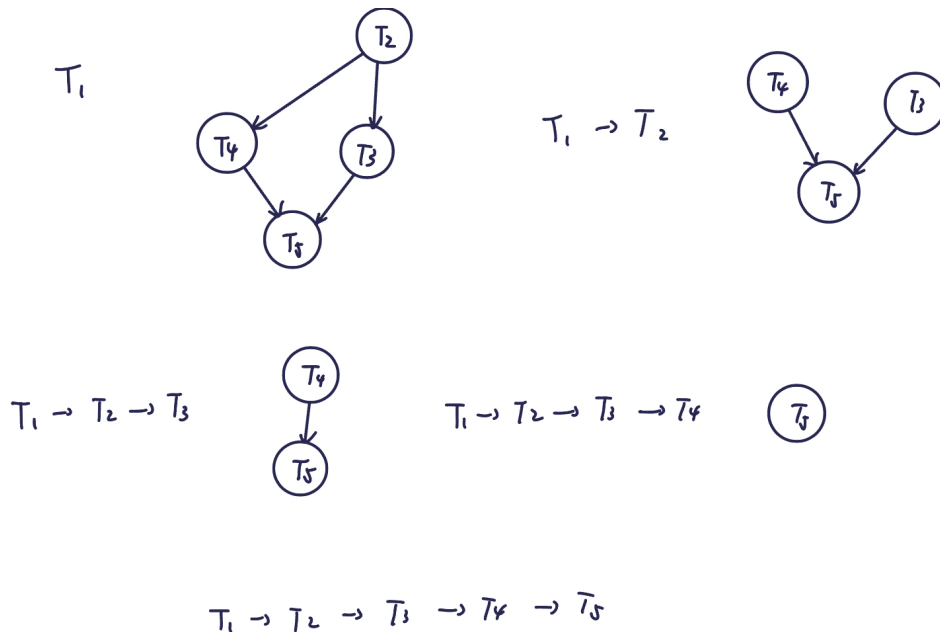
# DB - HW13

## 17.6

Consider the precedence graph of Figure 17.16. Is the corresponding schedule conflict serializable? Explain your answer.

Answer:

There is a serializable schedule corresponding to the Figure 17.16, for we can obtain a possible schedule by doing a topological sort by the way shown in the Figure below.



Clearly,  $T_1 \rightarrow T_2 \rightarrow T_3 \rightarrow T_4 \rightarrow T_5$  is a possible sequence.

## 17.7

What is a cascadeless schedule? Why is cascadelessness of schedules desirable? Are there circumstances under which it would be desirable to allow noncascadeless schedules? Explain your answer.

Answer:

A cascadeless schedule is a schedule where, for each pair of transactions  $T_i$  and  $T_j$  that  $T_j$  reads data previously written by  $T_i$ , the commitment operation of  $T_i$  always appears before the read operation of  $T_j$ .

The cascadelessness of schedule is desirable for, in cascadeless schedules, the failure of a transaction would not cause the aborting of any other transaction.

However, when the probability of failure is very low, the price of the cascading aborts is acceptable, and the concurrency would be increased. Under these circumstances, noncascadeless schedule might be desirable.

## 17.12

**List the ACID properties. Explain the usefulness of each.**

Answer:

- Atomicity

For all the operations of a certain transaction should either be reflected properly in the database, or none of them were, the lack of atomicity would cause the inconsistency in the database.

- Consistency

Execute the transactions non-concurrently preserves the consistency. This property insures that the transactions must transform the database from one state of consistency to another.

- Isolation

Isolation means, when multiple transactions are executed concurrently, for every pair of transactions  $T_i$  and  $T_j$ , it appears to  $T_i$  that either  $T_j$  finished execution before itself started, or  $T_j$  started execution after  $T_i$  finished. Thus, each transaction is unaware of other transactions executing concurrently with it.

Isolation is required for the user view of a transaction system needs it. Besides, it ensures that:

- During the execution of a transaction, the intermediate (and possibly inconsistent) state of the data should not be exposed to all other transactions.
- Two concurrent transactions should not be able to manipulate the same piece of data.

- Durability

Durability means after a transaction is successfully completed, the changes it made to the database would persist perpetual, even when the system fails. It makes the data recoverable.