Algorithm Template Library

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1 图论

1.1 线段树维护树直径

```
/*LCA 用 ST 表, 总复杂度 O(nlog)*/
/* 每次询问删去两条边后, 剩下 3 棵树的最大直径长度 */
const int N = 2e5 + 5;
int T, n, m;
int len, head[N], ST[20][N];
struct edge{int u, v, w;}ee[N];
int cnt, fa[N], log_2[N], st[N], en[N], dfn[N], dis[N], dep[N], pos[N];
struct edges{int to, next, cost;}e[N];
void add(int u, int v, int w) {
    e[++ len] = (edges){v, head[u], w}, head[u] = len;
   e[++ len] = (edges)\{u, head[v], w\}, head[v] = len;
void dfs1(int u) {
   st[u] = ++ cnt, dfn[cnt] = u;
   for (int v, i = head[u]; i; i = e[i].next) {
       v = e[i].to;
        if (v == fa[u]) continue;
       fa[v] = u, dep[v] = dep[u] + 1;
       dis[v] = dis[u] + e[i].cost, dfs1(v);
   }
   en[u] = cnt;
void dfs2(int u) {
   dfn[++ cnt] = u, pos[u] = cnt;
   for (int v, i = head[u]; i; i = e[i].next) {
       v = e[i].to;
        if (v == fa[u]) continue;
       dfs2(v), dfn[++ cnt] = u;
   }
}
int mmin(int x, int y) {
    if (dep[x] < dep[y]) return x;</pre>
   return y;
}
int lca(int u, int v) {
   static int w;
    if (pos[u] > pos[v]) swap(u, v);
   w = log_2[pos[v] - pos[u] + 1];
   return mmin(ST[w][pos[u]], ST[w][pos[v] - (1 << w) + 1]);
}
int dist(int u, int v) {
    int Lca = lca(u, v);
   return dis[u] + dis[v] - dis[Lca] * 2;
void build() {
   for (int i = 1; i <= cnt; i ++)
       ST[0][i] = dfn[i];
   for (int i = 1; i < 20; i ++)
       for (int j = 1; j \le cnt; j ++)
            if (j + (1 << (i - 1)) > cnt) ST[i][j] = ST[i - 1][j];
```

```
else ST[i][j] = mmin(ST[i - 1][j], ST[i - 1][j + (1 << (i - 1))]);
}
int M;
struct node {
    int 1, r, dis;
}tr[N << 1];</pre>
void update(int o, int o1, int o2) {
    static int d; static node tmp;
    if (tr[o1].dis == -1) {tr[o] = tr[o2]; return;}
    if (tr[o2].dis == -1) {tr[o] = tr[o1]; return;}
    if (tr[o1].dis > tr[o2].dis) tmp = tr[o1];
   else tmp = tr[o2];
   d = dist(tr[o1].1, tr[o2].1);
   if (d > tmp.dis) tmp.l = tr[o1].1, tmp.r = tr[o2].1, tmp.dis = d;
   d = dist(tr[o1].1, tr[o2].r);
    if (d > tmp.dis) tmp.l = tr[o1].l, tmp.r = tr[o2].r, tmp.dis = d;
   d = dist(tr[o1].r, tr[o2].1);
    if (d > tmp.dis) tmp.l = tr[o1].r, tmp.r = tr[o2].l, tmp.dis = d;
   d = dist(tr[o1].r, tr[o2].r);
    if (d > tmp.dis) tmp.l = tr[o1].r, tmp.r = tr[o2].r, tmp.dis = d;
   tr[o] = tmp;
void ask(int s, int t) {
   if (s > t) return;
   for (s += M - 1, t += M + 1; s ^ t ^ 1; s >>= 1, t >>= 1) {
        if (~s&1) update(0, 0, s ^ 1);
        if ( t&1) update(0, 0, t ^ 1);
    }
}
int main() {
    ios::sync_with_stdio(false);
    int u, v, w, ans; log_2[1] = 0;
    for (int i = 2; i <= 200000; i ++)
        if (i == 1 << (log_2[i - 1] + 1))
            log_2[i] = log_2[i - 1] + 1;
        else log_{2}[i] = log_{2}[i - 1];
   for (cin >> T; T --; ) {
        cin >> n >> m, cnt = len = 0;
        for (int i = 1; i <= n; i ++)
            head[i] = 0;
        for (int i = 1; i < n; i ++) {
            cin >> ee[i].u >> ee[i].v >> ee[i].w;
            add(ee[i].u, ee[i].v, ee[i].w);
        }
        dfs1(1);
        for (M = 1; M < n + 2; M <<= 1);
        for (int i = 1; i <= n; i ++)
            tr[i + M].l = tr[i + M].r = dfn[i], tr[i + M].dis = 0;
        for (int i = n + M + 1; i <= (M << 1) + 1; i ++)
            tr[i].dis = -1;
        cnt = 0, dfs2(1), build();
        for (int i = M; i; i --)
            update(i, i << 1, i << 1 | 1);
        for (int i = 1; i < n; i ++)
```

```
if (dep[ee[i].u] > dep[ee[i].v])
                swap(ee[i].u, ee[i].v);
        for (int u, v, i = 1; i <= m; i ++) {
            cin >> u >> v, ans = 0;
            u = ee[u].v, v = ee[v].v, w = lca(u, v);
            if (w == u | | w == v) {
                if (w != u) swap(u, v);
                tr[0].dis = -1, ask(1, st[u] - 1), ask(en[u] + 1, n), and ext{s} = max(ans, tr[0])
                 \rightarrow tr[0].dis);
                tr[0].dis = -1, ask(st[u], st[v] - 1), ask(en[v] + 1, en[u]), ans =

→ max(ans, tr[0].dis);
                tr[0].dis = -1, ask(st[v], en[v]), ans = max(ans, tr[0].dis);
            }
            else {
                if (st[u] > st[v]) swap(u, v);
                tr[0].dis = -1, ask(1, st[u] - 1), ask(en[u] + 1, st[v] - 1), ask(en[v] + 1, st[v] - 1)
                 \rightarrow 1, n), ans = max(ans, tr[0].dis);
                tr[0].dis = -1, ask(st[u], en[u]), ans = max(ans, tr[0].dis);
                tr[0].dis = -1, ask(st[v], en[v]), ans = max(ans, tr[0].dis);
            }
            printf("%d\n", ans);
        }
    }
    return 0;
}
     有向图判断两个点能否到达
```

```
/* 时间 O(n*m/32) 空间 O(n*n/blk) blk 随便取 */
const int N = 1e5 + 2;
const int BLK = 5000;
int n, k, p;
int d[N];
vector <int> e[N], f[N], ck[N];
int top, sta[N], in[N];
int cnt, dfn[N], low[N], vis[N];
int sum, bel[N];
bitset <BLK> a[N];
queue <int> q;
int topo[N];
void tarjan(int u) {
        vis[u] = in[u] = 1;
        sta[++ top] = u, dfn[u] = low[u] = ++ cnt;
        for (int v : e[u])
                if (!vis[v]) {
                        tarjan(v);
                        low[u] = min(low[v], low[u]);
                }
                else if (in[v])
                        low[u] = min(low[v], low[u]);
        if (low[u] == dfn[u]) {
                sum ++; int i;
                while (1) {
                        i = sta[top --];
```

```
in[i] = 0, bel[i] = sum;
                         if (i == u) break;
                }
        }
}
int main() {
        cin >> n >> k;
        for (int u, v, i = 1; i <= k; i ++) {
                cin >> u >> v;
                e[u].push_back(v);
        }
        cin >> p;
        for (int u, v, i = 1; i <= p; i ++) {
                cin >> u >> v;
                f[u].push_back(v);
        for (int i = 1; i <= n; i ++)
                if (!vis[i])
                         tarjan(i);
        for (int i = 1; i <= n; i ++) {
                for (int j : f[i]) {
                         if (bel[i] == bel[j]) return puts("NO"), 0;
                         ck[bel[i]].push_back(bel[j]);
                f[i].clear();
        for (int i = 1; i <= n; i ++)
                for (int j : e[i]) {
                         if (bel[i] == bel[j]) continue;
                         f[bel[i]].push_back(bel[j]);
                         d[bel[j]] ++;
                }
        cnt = 0;
        for (int i = 1; i <= sum; i ++)</pre>
                if (!d[i])
                         q.push(i);
        while (!q.empty()) {
                int now = q.front(); q.pop();
                topo[++ cnt] = now;
                for (int j : f[now]) {
                         d[j] --;
                         if (d[j] == 0) q.push(j);
                }
        }
        for (int i = 1, t = (sum + BLK - 1) / BLK; <math>i \le t; i \leftrightarrow t) {
                for (int j = sum; j; j --) {
                         int u = topo[j];
                         a[u].reset();
                         if (BLK * (i - 1) < u \&\& u \le BLK * i)
                                 a[u][u - BLK * (i - 1) - 1] = 1;
                         for (int v : f[u])
                                 a[u] |= a[v];
                }
                for (int j = 1; j <= sum; j ++)
```

```
for (int v : ck[i])
                                 if (BLK * (i - 1) < v && v \le BLK * i && a[j][v - BLK *
                                 \rightarrow (i - 1) - 1] == 1) {
                                         puts("NO");
                                         return 0;
                                 }
       printf("YES\n%d\n", k);
        for (int i = 1; i <= n; i ++)
                for (int j : e[i])
                        printf("%d %d\n", i, j);
        return 0;
}
1.3 spfa 费用流
const int N = 600, M = 800000, inf = 0x3f3f3f3f;
int s, t, ans, len, maxflow;
int T, n, m, K, W;
int head[N], incf[N], path[N], pre[N], vis[N], d[N];
struct edge{int to, next, cap, cost;}e[M];
struct video {int s, t, w, op; }a[N];
void add(int u, int v, int w, int c) {
        e[++ len] = (edge)\{v, head[u], w, c\}, head[u] = len;
        e[++ len] = (edge)\{u, head[v], 0, -c\}, head[v] = len;
}
bool spfa() {
        deque <int> q;
        q.push_back(s), incf[s] = inf;
        for (int i = 1; i <= t; i ++) d[i] = inf;
        d[s] = 0;
        while (!q.empty()) {
                int x = q.front();
                q.pop_front(), vis[x] = 0;
                for (int i = head[x]; i; i = e[i].next) {
                        if (e[i].cap \&\& d[e[i].to] > d[x] + e[i].cost) {
                                 d[e[i].to] = d[x] + e[i].cost;
                                 pre[e[i].to] = x, path[e[i].to] = i;
                                 incf[e[i].to] = min(incf[x], e[i].cap);
                                 if (!vis[e[i].to]) {
                                         vis[e[i].to] = 1;
                                         if (q.empty() || d[e[i].to] < d[q.front()])</pre>

¬ q.push_front(e[i].to);

                                         else q.push_back(e[i].to);
                                }
                        }
                }
        maxflow += incf[t];
        if (d[t] == inf) return 0;
        for (int i = t; i != s; i = pre[i]) {
                e[path[i]].cap -= incf[t];
                e[path[i] ^ 1].cap += incf[t];
        }
```

```
return ans += incf[t] * d[t], 1;
}
int main() {
       /*build graph*/
       while(spfa());
}
1.4 dsu
/*DSU
 * 用途:O(nlogn) 解决无修改的子树询问问题, 需要保证操作支持删除
 *解决方法:对于每个节点,先对所有轻儿子, dfs 下去求一遍,再消除影响
         然后再 dfs 自己的重儿子, 然后不消除影响, 再加上所有轻儿子
         就得到当前节点为根的子树的答案了
 */
int n, c[N];
int cnt[N], maxCnt;
int siz[N], son[N];
vector <int> e[N];
11 ans[N], sum[N];
void dfs1(int u, int fr) {
       siz[u] = 1;
       for (int v : e[u]) {
              if (v == fr) continue;
              dfs1(v, u);
              siz[u] += siz[v];
               if (siz[v] > siz[son[u]]) son[u] = v;
}
void update(int x, int y) {
       sum[cnt[x]] -= x;
       cnt[x] += y;
       sum[cnt[x]] += x;
       if (cnt[x] > maxCnt) maxCnt = cnt[x];
       if (sum[maxCnt] == 0) maxCnt --;
void dfs3(int u, int fr, int val) {
       update(c[u], val);
       for (int v : e[u]) {
              if (v == fr) continue;
               dfs3(v, u, val);
       }
}
void dfs2(int u, int fr) {
       for (int v : e[u]) {
               if (v == fr || v == son[u]) continue;
              dfs2(v, u), dfs3(v, u, -1);
       }
       if (son[u]) dfs2(son[u], u);
       for (int v : e[u]) {
               if (v == fr || v == son[u]) continue;
               dfs3(v, u, 1);
       update(c[u], 1);
```

```
ans[u] = sum[maxCnt];
int main() {
       ios::sync_with_stdio(false);
       cin >> n;
       for (int i = 1; i <= n; i ++)
              cin >> c[i]:
       for (int u, v, i = 1; i < n; i ++) {
              cin >> u >> v;
              e[u].push_back(v);
              e[v].push_back(u);
       }
       dfs1(1, 1), dfs2(1, 1);
       for (int i = 1; i <= n; i ++)
              cout << ans[i] << ' ';
       return 0;
}
1.5 长链剖分
/* 长链剖分, 选择深度最大的儿子作为重儿子, 用于合并以深度为下标的信息
 *像 dsu 一样,直接继承重儿子信息,然后按深度暴力合并其他儿子信息
 * 时间复杂度考虑每个节点作为轻儿子里的节点被合并只会有一次, 所以 O(n)
 * 另一种用法,可以 O(nlogn) 预处理后,O(1) 找到 k 级祖先
 * example problem: 给个树, 第 i 个点有两个权值 ai 和 bi
 * 现在求一条长度为 m 的路径, 使得 \Sigma ai/\Sigma bi 最小
 * 防爆栈 trick: 像重链剖分一样改成 bfs
 */
int n, m;
double k, a[N], b[N];
int len[N], son[N];
vector <int> e[N];
double tmp[N], *ptr, *f[N], temp[N];
void dfs(int u, int fr) {
   for (int v : e[u]) {
       if (v == fr) continue;
       dfs(v, u);
       if (len[v] > len[son[u]]) son[u] = v;
   len[u] = len[son[u]] + 1;
inline double F(int x, int y) {return y >= len[x] ? 0 : f[x][y];}
bool solve(int u, int fr) {
   /* 为实现 O(1) 继承, 采用 f[u]-f[v] 来保存 u-fa[v] 路径上的最小权值和 (dep[v]>dep[u])
    * 即自底向上累加
    */
   if (son[u]) {
       f[son[u]] = f[u] + 1;
       if (solve(son[u], u)) return 1;
       f[u][0] = val[u] + f[u][1];
       if (len[u] >= m \&\& f[u][0] - F(u, m) <= 0) return 1;
       for (int v : e[u]) {
          if (v == son[u] || v == fr) continue;
          f[v] = ptr, ptr += len[v];
```

```
if (solve(v, u)) return 1;
            for (int j = 1; j \le len[v] && j \le m; j ++) {
                if (len[u] + j < m) continue;</pre>
                if (f[v][0] - F(v, j) + f[u][0] - F(u, m - j) \le 0) return 1;
            temp[0] = val[u];
            for (int j = 1; j <= len[v]; j ++)
                temp[j] = val[u] + min(f[u][1] - F(u, j + 1), f[v][0] - F(v, j));
            if (len[v] + 1 == len[u]) f[u][0] = temp[len[v]];
            for (int j = 1; j <= len[v]; j ++)
                f[u][j] = f[u][0] - temp[j - 1];
            if (len[v] + 1 != len[u]) f[u][len[v] + 1] = f[u][0] - temp[len[v]];
       }
   }
   else {
        f[u][0] = val[u];
        if (m == 1 && f[u][0] <= 0) return 1;
   return 0;
}
bool judge(double mid) {
   f[1] = ptr = tmp, ptr += len[1], k = mid;
   return solve(1, 1);
int main() {
   cin >> n >> m;
   for (int i = 1; i <= n; i ++) cin >> a[i];
   for (int i = 1; i <= n; i ++) cin >> b[i];
    if (m == -1) {
       double ans = 1e9;
        for (int i = 1; i <= n; i ++)
            ans = min(ans, a[i] / b[i]);
       printf("%.2f\n", ans);
       return 0;
   }
   for (int u, v, i = 1; i < n; i ++) {
       cin >> u >> v;
        e[u].push_back(v);
        e[v].push_back(u);
   }
   dfs(1, 1);
   int flag = 0;
   double l = 0, r = 2e5, mid, ans;
   for (int i = 0; i < 50; i ++) {
       mid = (1 + r) / 2;
        if (judge(mid)) r = mid - eps, flag = 1, ans = mid;
       else l = mid + eps;
   }
   if (flag) printf("%.2f\n", ans);
    else puts("-1");
   return 0;
}
```

1.6 LCT

```
const int N = 3e5 + 5;
int n, m;
int fa[N], ch[N][2], rev[N], val[N], sum[N];
int sta[N], top;
bool isroot(int x) {
   return ch[fa[x]][0] != x && ch[fa[x]][1] != x;
void reverse(int x) {
   if (!x) return;
   swap(ch[x][0], ch[x][1]);
   rev[x] ^= 1;
}
void pushdown(int x) {
   if (!rev[x]) return;
   reverse(ch[x][0]);
   reverse(ch[x][1]);
   rev[x] = 0;
void pushup(int x) {
    sum[x] = sum[ch[x][0]] ^ sum[ch[x][1]] ^ val[x];
void rot(int x) {
    int y = fa[x], z = fa[y], d = ch[y][1] == x, c = ch[x][!d];
   fa[x] = z; if (!isroot(y)) ch[z][ch[z][1] == y] = x;
    ch[y][d] = c; if (c) fa[c] = y;
   fa[y] = x, ch[x][!d] = y;
   pushup(y), pushup(x);
void splay(int x) {
   int u = x, top = 0, y, z;
   while (!isroot(u)) sta[++ top] = u, u = fa[u];
   sta[++ top] = u;
   while (top) pushdown(sta[top --]);
   while (!isroot(x)) {
       y = fa[x], z = fa[y];
        if (!isroot(y)) {
            if ((ch[z][0] == y) ^ (ch[y][0] == x)) rot(x);
            else rot(y);
       }
       rot(x);
   }
void access(int x) {//把 x 到根的路径拎出来
   for (int y = 0; x != 0; y = x, x = fa[x]) {
       splay(x), ch[x][1] = y, pushup(x);
   }
void makeroot(int x) {//令 x 成为这棵树的根
   access(x), splay(x), reverse(x);
int findroot(int x) {//找根
   access(x), splay(x);
```

```
while (ch[x][0]) pushdown(x), x = ch[x][0];
    splay(x);//把根转到顶保证复杂度
   return x;
}
void split(int x, int y) {//拉出 x-y 的路径
   makeroot(x);
   access(y), splay(y);
    //y 存了这条路径的信息
}
void link(int x, int y) {
   makeroot(x);
    if (findroot(y) != x) fa[x] = y;
void cut(int x, int y) {
   makeroot(x);
    if (findroot(y) == x \&\& fa[y] == x \&\& ch[x][1] == y) {
       fa[y] = ch[x][1] = 0;
       pushup(x);
   }
}
int main() {
   ios::sync_with_stdio(false);
   cin >> n >> m;
   for (int i = 1; i <= n; i ++)
       cin >> val[i];
   for (int op, x, y; m --; ) {
       cin >> op >> x >> y;
        switch(op) {
            case 0:split(x, y);printf("%d\n", sum[y]);break;
            case 1:link(x, y);break;
            case 2:cut(x, y);break;
            case 3:splay(x);val[x] = y;pushup(x);break;
            //单点更新完记得 pushup
       }
   }
}
1.7 dinic 最大流
const int N = 20000;
const int M = 500000;
const int inf = 0x3f3f3f3f;
int n, m;
int s, t, len = 1;
int to[M], cap[M], nex[M];
int g[N], p[N], q[N], d[N];
void add(int x, int y, int v) {
       to[++ len] = y, cap[len] = v, nex[len] = g[x], g[x] = len;
       to[++ len] = x, cap[len] = 0, nex[len] = g[y], g[y] = len;
}
bool bfs() {
        int 1 = 1, r = 1, x, i;
       memset (d, 0, sizeof d);
       d[s] = 1, q[1] = s;
```

```
while (1 \leftarrow r) {
               x = q[1 ++];
               for (i = g[x]; i; i = nex[i])
                       if (cap[i] && !d[to[i]])
                               d[to[i]] = d[x] + 1, q[++ r] = to[i];
       }
       return d[t];
}
int dfs(int x, int y) {
        if (x == t || y == 0) return y;
        int flow = 0;
       for (int &i = p[x]; i; i = nex[i]) {
               if (!cap[i] || d[to[i]] != d[x] + 1) continue;
               int f = dfs(to[i], min(y, cap[i]));
               flow += f, y -= f;
               cap[i] -= f, cap[i ^ 1] += f;
               if (!y) break;
       }
       return flow;
}
int dinic() {
       int maxflow = 0;
       while (bfs()) {
               memcpy(p, g, sizeof g);
               maxflow += dfs(s, inf);
       return maxflow;
}
1.8 DAG 删去无用边
/* 无用边定义: 对于边 (u,v) 如果存在从 u 到 v 不经过该边的另一条路径,则称该边无用
* 时间复杂度:0(n~3) */
bool f[N][N];//i 是否可达 j
vector <int> e[N];
int main() {
       rep(i, 1, n)
               for (int j : e[i]) {
                       rep (k, 1, n)
                               if (i != k && j != k && f[i][k] && f[k][j])
                                       no_use_edge;
               }
}
     树分治
1.9
const int N = 1e5 + 5;
vector <int> e[N];
int n, a[N];
int root, _left, vis[N];
int siz[N], maxv[N];
void find_root(int u, int fr) {
        siz[u] = 1, maxv[u] = 0;
       for (int v : e[u]) {
```

```
if (v == fr || vis[v]) continue;
                find_root(v, u);
                siz[u] += siz[v];
                maxv[u] = max(maxv[u], siz[v]);
        maxv[u] = max(maxv[u], _left - siz[u]);
        if (!root || maxv[u] < maxv[root])</pre>
                root = u;
}
void dfs(int u, int fr) {
        siz[u] = 1;
        for (int v : e[u]) {
                if (v == fr || vis[v]) continue;
                find_root(v, u);
                siz[u] += siz[v];
        }
void solve(int u, int w) {
        dfs(u, u);//update siz[]
        a[u] = w, vis[u] = 1;
        for (int v : e[u]) {
                if (vis[v]) continue;
                _left = siz[v];
                root = 0;
                find_root(v, v);
                solve(root, w + 1);
        }
int main() {
    ios::sync_with_stdio(false);
    cin >> n;
    for (int u, v, i = 1; i < n; i ++) {
            cin >> u >> v;
            e[u].push_back(v);
            e[v].push_back(u);
    }
    _left = n, root = 0, find_root(1, 1);
    solve(root, 0);
    for (int i = 1; i <= n; i ++)
            printf("%c ", 'A' + a[i]);
    return 0;
}
```

2 数据结构

2.1 splay

```
#define mid (l + r >> 1)
int n, m, a, b, len, tot;
struct node {
        bool rev; //翻转标记
        int v, siz;
        node *c[2];
        node():rev(0),v(0),siz(0),c{NULL, NULL}{}
        node *init(int x);
        void pushdown();
        void mata() {siz = c[0] -> siz + c[1] -> siz + 1;}
        int cmp(int k) {return k<=c[0]->siz?0:(k==c[0]->siz+1?-1:1);}
        void print();
}pool[N], *null = new node();
node *node::init(int x){rev=0, v=x, siz=1, c[0]=c[1]=null;return this;}
void node::pushdown() {
        if (!rev) return;
        if (c[0] != null) c[0] -> rev ^= 1;
        if (c[1] != null) c[1] -> rev ^= 1;
        swap(c[0], c[1]), rev = 0;
}
void node::print() {
        pushdown();
        if (c[0] != null) c[0] -> print();
        if (1 \le v \&\& v \le n) printf("%d", v);
        if (c[1] != null) c[1] -> print();
node *build(int 1, int r) {//初始序列为 1-n
        if (l == r) return pool[tot ++].init(r);
        node *tmp = pool[tot ++].init(mid);
        if (1 < mid) tmp \rightarrow c[0] = build(1, mid - 1);
        if (mid < r) tmp \rightarrow c[1] = build(mid + 1, r);
        tmp -> mata(); return tmp;
void rot(node *&o, int k) {//把 k 儿子提上来
        o -> pushdown(); node *tmp = o -> c[k];
        tmp -> pushdown(); o -> c[k] = tmp -> c[!k];
        tmp \rightarrow c[!k] \rightarrow pushdown(); tmp \rightarrow c[!k] = o;
        o -> mata(), tmp -> mata(), o = tmp;
void splay(node *&o, int k) {//把以 o 为根的 splayTree 中 rk 为 k 的点提到根
        int k1 = o -> cmp(k); o -> pushdown();
        if (k1 == -1) return; o \rightarrow c[k1] \rightarrow pushdown();
        if (k1) k = 0 -> c[0] -> siz + 1;
        int k2 = o \rightarrow c[k1] \rightarrow cmp(k);
        if (\simk2) {//k2 != -1
                 if (k2) k = 0 \rightarrow c[k1] \rightarrow c[0] \rightarrow siz + 1;
                 o \rightarrow c[k1] \rightarrow c[k2] \rightarrow pushdown();
                 splay(o \rightarrow c[k1] \rightarrow c[k2], k);
                 if (k2 == k1) rot(o, k1);
                 else rot(o -> c[k1], k2);
```

```
}
        rot(o, k1);
int main() {
        scanf("%d %d", &n, &m);
        node *root = build(0, n + 1); //方便边界左右处理各多开一个
        for (; m --; ) {
                scanf("%d %d", &a, &b), a ++, b ++, len = b - a + 1;
                 splay(root, a - 1), splay(root -> c[1], len + 1);
                root -> c[1] -> c[0] -> rev ^= 1, root -> c[1] -> c[0] -> pushdown();
        }
        root -> print(); return 0;
}
2.2 treap
/* 容易实现的预开内存池 treap, 每次 head 清空即可
如果初始要插入 n 个 1, 可改为类似 splay 的 O(n)build 写法
const int poolSize = 5e5 + 10;
struct node {
        node *c[2];
        int v, r, siz;
        void update();
        void init(int x);
};
node *null = new node(), *root = null;
void node::update() {
        siz = c[0] \rightarrow siz + c[1] \rightarrow siz + 1;
}
void node::init(int x) {
        v = x, r = rand(), siz = 1;
        c[0] = c[1] = null;
}
node nodesPool[poolSize];
int head;//每次 head=0 清空
node *newnode(int x) {
        node *res = &nodesPool[head ++];
        res -> init(x);
        return res;
}
void rot(node *&o, int d) {
        node *tmp = o \rightarrow c[!d];
        o \rightarrow c[!d] = tmp \rightarrow c[d], tmp \rightarrow c[d] = o;
        o -> update(), tmp -> update(), o = tmp;
}
void insert(node *&o, int x) {
        if (o == null) {
                o = newnode(x);
                return;
        }
        int d = x > o -> v ? 0 : 1;
        insert(o -> c[d], x);
        if (o \rightarrow c[d] \rightarrow r < o \rightarrow r) rot(o, !d);
```

```
o -> update();
void del(node *&o, int x) {
         if (x == o -> v) {
                   if (o \rightarrow c[0] == null) \{o = o \rightarrow c[1]; return;\}
                   if (o \rightarrow c[1] == null) \{o = o \rightarrow c[0]; return;\}
                   int d = o \rightarrow c[0] \rightarrow r < o \rightarrow c[1] \rightarrow r ? 1 : 0;
                   rot(o, d), del(o -> c[d], x);
         else del(o \rightarrow c[x <= o \rightarrow v], x);
         o -> update();
}
void build(node *&o, int 1, int r) {
         o = newnode(1);
         if (1 == r) return;
         int mid = 1 + r >> 1;
         if (1 < mid) build(o -> c[0], 1, mid - 1);
         if (o \rightarrow c[0] = null \&\& o \rightarrow c[0] \rightarrow r < o \rightarrow r) swap(o \rightarrow c[0] \rightarrow r, o \rightarrow r);
         if (mid < r) build(o -> c[1], mid + 1, r);
         if (o \rightarrow c[1] != null && o \rightarrow c[1] \rightarrow r < o \rightarrow r) swap(o \rightarrow c[1] \rightarrow r, o \rightarrow r);
         o -> update();
}
2.3 吉老师线段树
/* 区间每个数变为 min(a[i],t) + 区间最大 + 区间和 O(nlog)*/
const int N = (1 << 20) + 5;
#define lc (o << 1)
#define rc (lc / 1)
struct node {
         int max1, max2, cnt;
         ll sum;
}tr[N << 1];
int T, n, m;
int op, x, y, t;
int a[N];
void pushup(int o) {
         if (tr[lc].max1 == tr[rc].max1) {
                   tr[o].max1 = tr[lc].max1;
                   tr[o].cnt = tr[lc].cnt + tr[rc].cnt;
                   tr[o].max2 = max(tr[lc].max2, tr[rc].max2);
         }
         else {
                   if (tr[lc].max1 > tr[rc].max1) {
                            tr[o] = tr[lc];
                            tr[o].max2 = max(tr[o].max2, tr[rc].max1);
                   }
                   else {
                            tr[o] = tr[rc];
                            tr[o].max2 = max(tr[o].max2, tr[lc].max1);
                   }
         tr[o].sum = tr[lc].sum + tr[rc].sum;
}
```

```
void pushdown(int o) {
        if (tr[o].max1 < tr[lc].max1) {</pre>
                tr[lc].sum += 111 * (tr[o].max1 - tr[lc].max1) * tr[lc].cnt;
                tr[lc].max1 = tr[o].max1;
        if (tr[o].max1 < tr[rc].max1) {</pre>
                tr[rc].sum += 111 * (tr[o].max1 - tr[rc].max1) * tr[rc].cnt;
                tr[rc].max1 = tr[o].max1;
        }
void build(int o, int 1, int r) {
        if (1 == r) {
                tr[o].max1 = tr[o].sum = a[r];
                tr[o].cnt = 1, tr[o].max2 = 0;
                return;
        }
        int mid = 1 + r >> 1;
        build(lc, l, mid);
        build(rc, mid + 1, r);
        pushup(o);
void update(int o, int l, int r, int s, int t, int v) {
        if (v >= tr[o].max1) return;
        pushdown(o);
        if (s <= 1 && r <= t) {
                if (v > tr[o].max2) {
                         tr[o].sum += 111 * (v - tr[o].max1) * tr[o].cnt;
                         tr[o].max1 = v;
                         return;
                }
        }
        int mid = 1 + r >> 1;
        if (s <= mid) update(lc, l, mid, s, t, v);</pre>
        if (t > mid) update(rc, mid + 1, r, s, t, v);
        pushup(o);
}
11 ask_max(int o, int 1, int r, int s, int t) {
        if (s <= 1 && r <= t) return tr[o].max1;</pre>
        pushdown(o);
        int mid = 1 + r >> 1;
        11 \text{ res} = 0;
        if (s <= mid) res = max(res, ask_max(lc, l, mid, s, t));</pre>
        if (t > mid) res = max(res, ask_max(rc, mid + 1, r, s, t));
        return res;
11 ask_sum(int o, int 1, int r, int s, int t) {
        if (s <= 1 && r <= t) return tr[o].sum;</pre>
        pushdown(o);
        int mid = 1 + r >> 1;
        11 \text{ res} = 0;
        if (s \le mid) res += ask_sum(lc, l, mid, s, t);
        if (t > mid) res += ask_sum(rc, mid + 1, r, s, t);
        return res;
}
```

```
int main() {
        ios::sync_with_stdio(false);
        for (cin >> T; T --; ) {
                 cin >> n >> m;
                 for (int i = 1; i <= n; i ++)
                          cin >> a[i];
                 build(1, 1, n);
                 while (m --) {
                          cin >> op;
                          if (op == 0) {
                                   cin >> x >> y >> t;
                                   update(1, 1, n, x, y, t);
                          else if (op == 1) {
                                   cin >> x >> y;
                                   cout << ask_max(1, 1, n, x, y) << '\n';
                          }
                          else {
                                   cin >> x >> y;
                                   cout << ask_sum(1, 1, n, x, y) << '\n';
                          }
                 }
        }
        return 0;
}
2.4 KDTree
2.4.1 3 维 KDtree
/*O(n*n^(1-1/k)),k 为维度 */
const int N = 1e5 + 5;
const int Mod = 1e9 + 7;
int nowD, ans, x[3], y[3];
int n, m, a[N], b[N], c[N], d[N];
struct node {
    int Max[3], Min[3], d[3];
    int val, maxv;
    node *c[2];
    node() {
        c[0] = c[1] = NULL;
        val = maxv = 0;
    void pushup();
    bool operator < (const node &a) const {</pre>
        return d[nowD] < a.d[nowD];</pre>
    }
}Null, nodes[N];
node *root = &Null;
inline void node::pushup() {
    if (c[0] != &Null) {
        if (c[0] \rightarrow Max[1] > Max[1]) Max[1] = c[0] \rightarrow Max[1];
        if (c[0] \rightarrow Max[2] > Max[2]) Max[2] = c[0] \rightarrow Max[2];
        if (c[0] \rightarrow Min[0] < Min[0]) Min[0] = c[0] \rightarrow Min[0];
        if (c[0] \rightarrow Min[2] < Min[2]) Min[2] = c[0] \rightarrow Min[2];
```

```
if (c[0] \rightarrow maxv > maxv) maxv = c[0] \rightarrow maxv;
    }
    if (c[1] != &Null) {
         if (c[1] \rightarrow Max[1] > Max[1]) Max[1] = c[1] \rightarrow Max[1];
         if (c[1] \rightarrow Max[2] > Max[2]) Max[2] = c[1] \rightarrow Max[2];
         if (c[1] \rightarrow Min[0] < Min[0]) Min[0] = c[1] \rightarrow Min[0];
         if (c[1] \rightarrow Min[2] < Min[2]) Min[2] = c[1] \rightarrow Min[2];
         if (c[1] \rightarrow maxv > maxv) maxv = c[1] \rightarrow maxv;
    }
inline node *build(int 1, int r) {
    int mid = 1 + r >> 1; nowD = rand() % 3;
    nth_element(nodes + 1, nodes + mid, nodes + r + 1);
    node *res = &nodes[mid];
    if (1 != mid) res -> c[0] = build(1, mid - 1);
    else res \rightarrow c[0] = &Null;
    if (r != mid) res \rightarrow c[1] = build(mid + 1, r);
    else res -> c[1] = &Null;
    res -> pushup();
    return res;
}
inline int calc(node *o) {
    if (y[0] < o \rightarrow Min[0] \mid | x[1] > o \rightarrow Max[1] \mid | x[2] > o \rightarrow Max[2] \mid | y[2] < o \rightarrow
     \rightarrow Min[2]) return -1;
    return o -> maxv;
inline void query(node *o) {
    if (o -> val > ans && y[0] >= o -> d[0] && x[1] <= o -> d[1] && x[2] <= o -> d[2] &&
     \rightarrow y[2] >= o -> d[2]) ans = o -> val;
    int dl, dr;
    if (o \rightarrow c[0] != \&Null) dl = calc(o \rightarrow c[0]);
    else dl = -1;
    if (o \rightarrow c[1] != \&Null) dr = calc(o \rightarrow c[1]);
    else dr = -1;
    if (dl > dr) {
         if (dl > ans) query(o \rightarrow c[0]);
         if (dr > ans) query(o \rightarrow c[1]);
    } else {
         if (dr > ans) query(o -> c[1]);
         if (dl > ans) query(o -> c[0]);
    }
}
int main() {
    ios::sync_with_stdio(false);
    cin >> n >> m;
    for (int i = 1; i <= n; i ++) {
         cin >> a[i];
         b[i] = d[a[i]];
         d[a[i]] = i;
    }
    for (int i = 1; i <= n; i ++) d[i] = n + 1;
    for (int i = n; i; i --) {
         c[i] = d[a[i]];
         d[a[i]] = i;
```

```
}
    for (int i = 1; i <= n; i ++) {
        nodes[i].Min[0] = nodes[i].d[0] = b[i];
        nodes[i].Max[1] = nodes[i].d[1] = c[i];
        nodes[i].Max[2] = nodes[i].Min[2] = nodes[i].d[2] = i;
        nodes[i].val = nodes[i].maxv = a[i];
    root = build(1, n);
    for (int 1, r; m --; ) {
        cin >> 1 >> r;
        1 = (1 + ans) \% n + 1;
        r = (r + ans) \% n + 1;
         if (1 > r) swap(1, r);
        y[0] = 1 - 1;
        x[1] = r + 1;
        x[2] = 1, y[2] = r;
        ans = 0, query(root);
        cout << ans << endl;</pre>
    }
    cout << endl;</pre>
    return 0;
}
2.4.2 KDtree 二维空间区间覆盖单点查询
/* 类似线段树 */
const int N = 1e5 + 5;
const int Mod = 1e9 + 7;
int nowD, x[2], y[2], z;
struct node {
    int Max[2], Min[2], d[2];
    int val, lazy;
    node *c[2];
    node() {
         c[0] = c[1] = NULL;
    }
    void pushup();
    void pushdown();
    bool operator < (const node &a) const {</pre>
        return d[nowD] < a.d[nowD];</pre>
    }
}Null, nodes[N];
node *root = &Null;
inline void node::pushup() {
    if (c[0] != &Null) {
         if (c[0] \rightarrow Max[0] > Max[0]) Max[0] = c[0] \rightarrow Max[0];
         if (c[0] \rightarrow Max[1] > Max[1]) Max[1] = c[0] \rightarrow Max[1];
         if (c[0] \rightarrow Min[0] < Min[0]) Min[0] = c[0] \rightarrow Min[0];
         if (c[0] \rightarrow Min[1] < Min[1]) Min[1] = c[0] \rightarrow Min[1];
    }
    if (c[1] != &Null) {
         if (c[1] \rightarrow Max[0] > Max[0]) Max[0] = c[1] \rightarrow Max[0];
         if (c[1] \rightarrow Max[1] > Max[1]) Max[1] = c[1] \rightarrow Max[1];
         if (c[1] \rightarrow Min[0] < Min[0]) Min[0] = c[1] \rightarrow Min[0];
```

```
if (c[1] \rightarrow Min[1] < Min[1]) Min[1] = c[1] \rightarrow Min[1];
    }
}
inline void node::pushdown() {
    if (c[0] != &Null) c[0] -> val = c[0] -> lazy = lazy;
    if (c[1] != &Null) c[1] -> val = c[1] -> lazy = lazy;
    lazy = -1;
inline node *build(int 1, int r, int D) {
    int mid = 1 + r >> 1; nowD = D;
    nth_element(nodes + 1, nodes + mid, nodes + r + 1);
    node *res = &nodes[mid];
    if (1 != mid) res -> c[0] = build(1, mid - 1, !D);
    else res -> c[0] = &Null;
    if (r != mid) res -> c[1] = build(mid + 1, r, !D);
    else res \rightarrow c[1] = &Null;
    res -> pushup();
    return res;
}
inline int query(node *o) {
    if (o == &Null) return -1;
    if (o \rightarrow lazy != -1) o \rightarrow pushdown();
    if (x[0] > o -> Max[0] \mid | y[0] > o -> Max[1] \mid | x[0] < o -> Min[0] \mid | y[0] < o ->
     \rightarrow Min[1]) return -1;
    if (x[0] == o \rightarrow d[0]) return o \rightarrow val;
    return max(query(o -> c[0]), query(o -> c[1]));
}
inline void modify(node *o) {
    if (o == &Null) return;
    if (o -> lazy != -1) o -> pushdown();
    if (x[0] > o -> Max[0] \mid | y[0] > o -> Max[1] \mid | x[1] < o -> Min[0] \mid | y[1] < o ->

→ Min[1]) return;

     \text{if } (x[0] \mathrel{<=} o \mathrel{->} \texttt{Min}[0] \; \&\& \; y[0] \mathrel{<=} o \mathrel{->} \texttt{Min}[1] \; \&\& \; x[1] \mathrel{>=} o \; \mathrel{->} \texttt{Max}[0] \; \&\& \; y[1] \mathrel{>=} o \; \mathrel{->} \\ \\
     \rightarrow Max[1]) {
         o \rightarrow val = o \rightarrow lazy = z;
         return;
    }
    if (x[0] \le o -> d[0] \&\& y[0] \le o -> d[1] \&\& x[1] >= o -> d[0] \&\& y[1] >= o -> d[1])
     \rightarrow o -> val = z;
    modify(o \rightarrow c[0]), modify(o \rightarrow c[1]);
}
int n, m, k, a[N], c[N], d[N];
int cnt, st[N], en[N], dfn[N], dep[N];
vector <int> e[N];
void dfs(int u) {
    st[u] = ++ cnt, dfn[cnt] = u;
    for (int v : e[u])
         dep[v] = dep[u] + 1, dfs(v);
    en[u] = cnt;
}
int main() {
    ios::sync_with_stdio(false);
    int T, ans;
    for (cin >> T; T --; ) {
```

```
cin >> n >> m >> k, ans = cnt = 0;
       for (int i = 1; i <= n; i ++)
           e[i].clear();
       for (int u, i = 2; i <= n; i ++) {
           cin >> u;
           e[u].push_back(i);
       }
       dfs(1);
       for (int i = 1; i <= n; i ++) {
           nodes[i].Min[0] = nodes[i].Max[0] = nodes[i].d[0] = i;
           nodes[i].Min[1] = nodes[i].Max[1] = nodes[i].d[1] = dep[dfn[i]];
           nodes[i].val = 1, nodes[i].lazy = -1;
       }
       root = build(1, n, 0);
       for (int u, v, w, i = 1; i <= k; i ++) {
           cin >> u >> v >> w;
           if (w == 0) {
               x[0] = st[u], y[0] = dep[u];
               ans = (ans + 111 * i * query(root) % Mod) % Mod;
           } else {
               x[0] = st[u], x[1] = en[u];
               y[0] = dep[u], y[1] = dep[u] + v;
               z = w, modify(root);
       }
       cout << ans << endl;</pre>
   }
   return 0;
}
2.4.3 KDtree 二维空间单点修改区间查询
* 调整重构系数可以影响常数
 * 询问多就让系数接近 0.70-0.75, 询问少就让系数在 0.8-0.90
const int inf = 1e9;
int n, m, tot, nowD;
struct node {
   int Max[2], Min[2], d[2];
   int sum, siz, val;
   node *c[2];
   node() {
       Max[0] = Max[1] = -inf;
       Min[0] = Min[1] = inf;
       sum = val = siz = 0;
       c[0] = c[1] = NULL;
       d[0] = d[1] = 0;
   }
   void update();
Null, nodes[200010], *temp[200010];
node *root = &Null;
inline void node::update() {
   siz = c[0] -> siz + c[1] -> siz + 1;
```

```
sum = c[0] -> sum + c[1] -> sum + val;
    if (c[0] != &Null) {
         if (c[0] \rightarrow Max[0] > Max[0]) Max[0] = c[0] \rightarrow Max[0];
         if (c[0] \rightarrow Max[1] > Max[1]) Max[1] = c[0] \rightarrow Max[1];
         if (c[0] -> Min[0] < Min[0]) Min[0] = c[0] -> Min[0];
         if (c[0] -> Min[1] < Min[1]) Min[1] = c[0] -> Min[1];
    if (c[1] != &Null) {
         if (c[1] \rightarrow Max[0] > Max[0]) Max[0] = c[1] \rightarrow Max[0];
         if (c[1] \rightarrow Max[1] > Max[1]) Max[1] = c[1] \rightarrow Max[1];
         if (c[1] -> Min[0] < Min[0]) Min[0] = c[1] -> Min[0];
         if (c[1] \rightarrow Min[1] < Min[1]) Min[1] = c[1] \rightarrow Min[1];
    }
}
inline bool cmp(const node *a, const node *b) {
    return a -> d[nowD] < b -> d[nowD];
inline void traverse(node *o) {
    if (o == &Null) return;
    temp[++ tot] = o;
    traverse(o -> c[0]);
    traverse(o -> c[1]);
}
inline node *build(int 1, int r, int D) {
    int mid = 1 + r >> 1; nowD = D;
    nth_element(temp + 1, temp + mid, temp + r + 1, cmp);
    node *res = temp[mid];
    res -> Max[0] = res -> Min[0] = res -> d[0];
    res -> Max[1] = res -> Min[1] = res -> d[1];
    if (1 != mid) res -> c[0] = build(1, mid - 1, !D);
    else res -> c[0] = &Null;
    if (r != mid) res \rightarrow c[1] = build(mid + 1, r, !D);
    else res -> c[1] = &Null;
    res -> update();
    return res;
}
int x, y, a, b, tmpD;
node **tmp;
inline void rebuild(node *&o, int D) {
    tot = 0;
    traverse(o);
    o = build(1, tot, D);
inline void insert(node *&o, node *p, int D) {
    if (o == &Null) {o = p; return;}
    if (p -> Max[0] > o -> Max[0]) o -> Max[0] = p -> Max[0];
    if (p \rightarrow Max[1] > o \rightarrow Max[1]) o \rightarrow Max[1] = p \rightarrow Max[1];
    if (p \rightarrow Min[0] < o \rightarrow Min[0]) o \rightarrow Min[0] = p \rightarrow Min[0];
    if (p \rightarrow Min[1] < o \rightarrow Min[1]) o \rightarrow Min[1] = p \rightarrow Min[1];
    o -> siz ++, o -> sum += p -> sum;
    insert(o \rightarrow c[p \rightarrow c[D] >= o \rightarrow c[D]], p, !D);
    if (\max(o \rightarrow c[0] \rightarrow siz, o \rightarrow c[1] \rightarrow siz) > int(o \rightarrow siz * 0.75 + 0.5)) tmpD = D,
     \rightarrow tmp = &o;
```

```
inline int query(node *o) {
    if (o == &Null) return 0;
    if (x > o -> Max[0] \mid | y > o -> Max[1] \mid | a < o -> Min[0] \mid | b < o -> Min[1]) return
   if (x \le o -> Min[0] \&\& y \le o -> Min[1] \&\& a >= o -> Max[0] \&\& b >= o -> Max[1])

→ return o -> sum;

   return (x <= o -> d[0] && y <= o -> d[1] && a >= o -> d[0] && b >= o -> d[1] ? o ->
    \rightarrow val : 0)
        + query(o -> c[1]) + query(o -> c[0]);
}
int main() {
   ios::sync_with_stdio(false);
   cin >> m;
   node *ttt = &Null;
   for (int t, ans = 0; ; ) {
        cin >> t;
        if (t == 3) break;
        if (t == 1) {
            cin >> x >> y >> a;
            x = ans, y = ans, n ++;
            nodes[n].sum = nodes[n].val = a ^ ans, nodes[n].siz = 1;
            nodes[n].Max[0] = nodes[n].Min[0] = nodes[n].d[0] = x;
            nodes[n].Max[1] = nodes[n].Min[1] = nodes[n].d[1] = y;
            nodes[n].c[0] = nodes[n].c[1] = &Null;
            tmp = &(ttt), insert(root, &nodes[n], 0);
            if (*tmp != &Null) rebuild(*tmp, tmpD);
        } else {
            cin >> x >> y >> a >> b;
            x = ans, y = ans, a = ans, b = ans;
            if (x > a) swap(x, a);
            if (y > b) swap(y, b);
            ans = query(root);
            printf("%d\n", ans);
       }
   }
   return 0;
}
2.4.4 KDtree 找最近点
 * 为了维持树的平衡,可以一开始把所有点都读进来 build
 * 然后打 flag 标记该点是否被激活
const int N = 5e5 + 5;
const int inf = 1 << 30;</pre>
int n, m;
int ql, qr, ans, tot, nowD;
//nowD = rand() \& 1 ?
struct Node {
    int d[2];
   bool operator < (const Node &a) const {</pre>
        if (d[nowD] == a.d[nowD]) return d[!nowD] < a.d[!nowD];</pre>
        return d[nowD] < a.d[nowD];</pre>
```

```
}
}pot[N];
struct node {
    int min[2], max[2], d[2];
    node *c[2];
    node() {
         min[0] = min[1] = max[0] = max[1] = d[0] = d[1] = 0;
         c[0] = c[1] = NULL;
    node(int x, int y);
    void update();
}t[N], Null, *root;
node::node(int x, int y) {
    min[0] = max[0] = d[0] = x;
    min[1] = max[1] = d[1] = y;
    c[0] = c[1] = &Null;
}
inline void node::update() {
    if (c[0] != &Null) {
         if (c[0] \rightarrow max[0] > max[0]) max[0] = c[0] \rightarrow max[0];
         if (c[0] \rightarrow max[1] > max[1]) max[1] = c[0] \rightarrow max[1];
         if (c[0] \rightarrow min[0] < min[0]) min[0] = c[0] \rightarrow min[0];
         if (c[0] \rightarrow min[1] < min[1]) min[1] = c[0] \rightarrow min[1];
    }
    if (c[1] != &Null) {
         if (c[1] \rightarrow max[0] > max[0]) max[0] = c[1] \rightarrow max[0];
         if (c[1] \rightarrow max[1] > max[1]) max[1] = c[1] \rightarrow max[1];
         if (c[1] \rightarrow min[0] < min[0]) min[0] = c[1] \rightarrow min[0];
         if (c[1] \rightarrow min[1] < min[1]) min[1] = c[1] \rightarrow min[1];
    }
}
inline void build(node *&o, int 1, int r, int D) {
    int mid = 1 + r >> 1;
    nowD = D;
    nth_element(pot + 1, pot + mid, pot + r + 1);
    o = new node(pot[mid].d[0], pot[mid].d[1]);
    if (1 != mid) build(o -> c[0], 1, mid - 1, !D);
    if (r != mid) build(o \rightarrow c[1], mid + 1, r, !D);
    o -> update();
inline void insert(node *o) {
    node *p = root;
    int D = 0;
    while (1) {
         if (o \rightarrow max[0] > p \rightarrow max[0]) p \rightarrow max[0] = o \rightarrow max[0];
         if (o \rightarrow max[1] > p \rightarrow max[1]) p \rightarrow max[1] = o \rightarrow max[1];
         if (o \rightarrow min[0] 
         if (o \rightarrow min[1] 
         if (o -> d[D] >= p -> d[D]) {
             if (p \rightarrow c[1] == &Null) {
                  p \rightarrow c[1] = o;
                  return;
             } else p = p -> c[1];
```

```
} else {
             if (p -> c[0] == &Null) {
                 p -> c[0] = o;
                 return;
             } else p = p -> c[0];
        D = 1;
    }
}
inline int dist(node *o) {
    int dis = 0;
    if (ql < o \rightarrow min[0]) dis += o \rightarrow min[0] - ql;
    if (ql > o \rightarrow max[0]) dis += ql - o \rightarrow max[0];
    if (qr < o \rightarrow min[1]) dis += o \rightarrow min[1] - qr;
    if (qr > o \rightarrow max[1]) dis += qr - o \rightarrow max[1];
    return dis;
inline void query(node *o) {
    int dl, dr, d0;
    d0 = abs(o \rightarrow d[0] - q1) + abs(o \rightarrow d[1] - qr);
    if (d0 < ans) ans = d0;
    if (o -> c[0] != &Null) dl = dist(o -> c[0]);
    else dl = inf;
    if (o -> c[1] != &Null) dr = dist(o -> c[1]);
    else dr = inf;
    if (dl < dr) {
        if (dl < ans) query(o -> c[0]);
        if (dr < ans) query(o -> c[1]);
        if (dr < ans) query(o -> c[1]);
        if (dl < ans) query(o -> c[0]);
    }
}
int main() {
    ios::sync_with_stdio(false);
    cin >> n >> m;
    for (int i = 1; i <= n; i ++)
        cin >> pot[i].d[0] >> pot[i].d[1];
    build(root, 1, n, 0);
    for (int x, y, z; m --; ) {
        cin >> x >> y >> z;
        if (x == 1) {
             t[tot].max[0] = t[tot].min[0] = t[tot].d[0] = y;
             t[tot].max[1] = t[tot].min[1] = t[tot].d[1] = z;
             t[tot].c[0] = t[tot].c[1] = &Null;
             insert(&t[tot ++]);
        } else {
             ans = inf, ql = y, qr = z;
             query(root), printf("%d\n", ans);
        }
    }
    return 0;
}
```

3 构造

3.1 若干排列使所有数对都出现一次

```
/*n/2 个排列使得所有数对 (i,j) 且 i<j 都出现一次
*n 为奇数则首尾相连, 偶数不连
typedef vector<int> vi;
void get_even(int n, vi ans[]) {
       vi a(n);
       for (int i = 0; i < n; i ++)
               a[i] = i + 1;
       for (int i = 1; i <= n / 2; i ++) {
               ans[i].resize(n + 1);
               for (int j = 0; j < n / 2; j ++)
                       ans[i][j * 2] = a[j], ans[i][j * 2 + 1] = a[n - 1 - j];
               int t = a[n - 1];
               for (int j = n - 1; j > 0; j --)
                       a[j] = a[j - 1];
               a[0] = t;
       }
}
void get_odd(int n, vi ans[]) {
       get_even(n - 1, ans);
       for (int i = 1; i <= n / 2; i ++) {
               for (int j = n - 1; j > 0; j --)
                       ans[i][j] = ans[i][j - 1] + 1;
               ans[i][0] = 1;
       }
}
int main() {
       vi ans[2019];
       int n; cin >> n;
       if (n & 1) get_odd(n, ans);
       else get_even(n, ans);
       return 0;
}
3.2 rec-free
/* 不存在四个 1 构成一个矩形, 并使得 1 尽量多, 输出 01 矩阵 */
const int N = 200, n = 150, M = 13; //M 为质数, N>M*M>n
int a[N][N], b[N][N], c[N][N];
void make() {
       for (int i = 1; i <= M; i ++)
       for (int j = 1; j <= M; j ++)
           a[i][j + 1] = M * (j - 1) + i;
   for (int i = 1; i <= M; i ++)
       for (int j = 1; j <= M; j ++)
           c[i][a[i][j]] = 1;
   for (int k = 1; k < M; k ++) {
       memcpy(b, a, sizeof b);
       for (int i = 1; i <= M; i ++)
           for (int j = 1; j <= M; j ++)
```

4 计算几何

4.1 最小矩形覆盖含凸包和旋转卡壳

```
/* 最小矩形覆盖, 保留六位小数, 逆时针输出四个顶点坐标 */
namespace minRectCover {
       const int N = 1e5 + 5;
       const double eps = 1e-8;
       struct point{
               double x, y;
               point(){}
               point(double x, double y):x(x), y(y){}
               bool operator < (const point &a) const {
                       return fabs(y - a.y) < eps ? x < a.x : y < a.y;}
               point operator - (const point &a) const {
                       return point(x - a.x, y - a.y);}
               point operator + (const point &a) const {
                       return point(x + a.x, y + a.y);}
               point operator / (const double &a) const {
                       return point(x / a, y / a);}
               point operator * (const double &a) const {
                       return point(x * a, y * a);}
               double operator / (const point &a) const { // .
                       return x * a.x + y * a.y;}
               double operator * (const point &a) const { // X
                       return x * a.y - y * a.x;}
       }p[N], q[N], rc[4];
       double sqr(double x) {return x * x;}
       double abs(point a) {return sqrt(a / a);}
       int sgn(double x) {return fabs(x) < eps ? 0 : (x < 0 ? -1 : 1);}
       point vertical(point a, point b) {//与 ab 向量垂直的向量
               return point(a.x + a.y - b.y, a.y - a.x + b.x) - a;}
       point vec(point a){return a / abs(a);}
       void convexhull(int n, point *hull, int &top) {//如果要计算周长需要特判 n==2
               for (int i = 1; i < n; i ++)
                       if (p[i] < p[0])
                               swap(p[i], p[0]);
               sort (p + 1, p + n, [\&] (point a, point b) {
                       double t = (a - p[0]) * (b - p[0]);
                       if (fabs(t) < eps) return sgn(abs(p[0] - a) - abs(p[0] - b)) < 0;
                       return t > 0;
               });
               int cnt = 0;//去重
               for (int i = 1; i < n; i ++)
                       if (sgn(p[i].x - p[cnt].x) != 0 || sgn(p[i].y - p[cnt].y) != 0)
                               p[++ cnt] = p[i];
               n = cnt + 1;
               hull[top = 1] = p[0];
               for (int i = 1; i < n; i ++) {
                       while (top > 1 &&
                                [hull[top] - hull[top - 1]) * (p[i] - hull[top]) < eps)

    top --;
```

```
hull[++ top] = p[i];
                hull[0] = hull[top];
        }
        void main() {
                int n;
                scanf("%d", &n);
                for (int i = 0; i < n; i ++)</pre>
                         scanf("%lf %lf", &p[i].x, &p[i].y);
                convexhull(n, q, n);
                double ans = 1e20;
                int 1 = 1, r = 1, t = 1;
                double L, R, D, H;
                for (int i = 0; i < n; i ++) {
                         D = abs(q[i] - q[i + 1]);
                         while (sgn((q[i + 1] - q[i]) * (q[t + 1] - q[i]) -
                                 (q[i + 1] - q[i]) * (q[t] - q[i])) > -1) t = (t + 1) % n;
                         while (sgn((q[i + 1] - q[i]) / (q[r + 1] - q[i]) - (q[i + 1] -
                         \rightarrow q[i]) / (q[r] - q[i])) > -1) r = (r + 1) % n;
                         if (i == 0) 1 = r;
                         while (sgn((q[i + 1] - q[i]) / (q[1 + 1] - q[i]) - (q[i + 1] -
                         \rightarrow q[i]) / (q[1] - q[i])) < 1) 1 = (1 + 1) % n;
                         L = fabs((q[i + 1] - q[i]) / (q[1] - q[i]) / D);
                         R = fabs((q[i + 1] - q[i]) / (q[r] - q[i]) / D);
                         H = fabs((q[i + 1] - q[i]) * (q[t] - q[i]) / D);
                         double tmp = (R + L) * H;
                         if (tmp < ans) {</pre>
                                 ans = tmp;
                                 rc[0] = q[i] + (q[i + 1] - q[i]) * (R / D); //右下
                                 rc[1] = rc[0] + vec(vertical(q[i], q[i + 1])) * H;//右上
                                 rc[2] = rc[1] - (rc[0] - q[i]) * ((R + L) / abs(q[i] -
                                 → rc[0]));//左上
                                 rc[3] = rc[2] - (rc[1] - rc[0]);
                         }
                }
                printf("%.6f\n", ans);
                int fir = 0;
                for (int i = 1; i < 4; i ++)
                         if (rc[i] < rc[fir])</pre>
                                 fir = i;
                for (int i = 0; i < 4; i ++)
                         printf("\frac{.6f}{.6f}", rc[(fir + i) \frac{.4}{.x}, rc[(fir + i) \frac{.4}{.y});
        }
}
```

5 数论

5.1 CRT

```
void exgcd(ll a, ll b, ll &d, ll &x, ll &y) {
       if (!b) {
               d = a, x = 1, y = 0;
               return;
       }
       exgcd(b, a % b, d, y, x);
       y = a / b * x;
ll crt(ll *m, ll *a, int n) {//n 个式子: y = mx + a, 下标从 1 开始
       11 A = a[1], M = m[1], d, x, y, m2;
       for (int i = 2; i <= n; i ++) \{// k1 * m1 - k2 * m2 = a2 - a1\}
               exgcd(M, m[i], d, x, y);
                if ((a[i] - A) \% d) return -1;
               m2 = M / d * m[i];
               x = (a[i] - A) / d * x % m[i];
                A = (A + x * M \% m2) \% m2;
                if (A < 0) A += m2;//保证 A>=0
               M = m2;
       }
       return A; //y = Mx + A
}
```

6 字符串

6.1 KMP

```
int nxt[N];
void kmp(int n, char *a, int m, char *b) {
       //长度为 m 的 b 中找 a,下标从 o 开始,得到的是匹配成功的末尾位置
       static int i, j, cnt, tmp[2]; cnt = 0;
       for (nxt[0] = j = -1, i = 1; i < n; nxt[i ++] = j) {
               while (~j \&\& a[j + 1] != a[i]) j = nxt[j];
               if (a[j + 1] == a[i]) j ++;
       }
       for (j = -1, i = 0; i < m; i ++) {
               while (~j \&\& a[j + 1] != b[i]) j = nxt[j];
               if (a[j + 1] == b[i]) j ++;
               if (j == n - 1) {
                      printf("%d ", i);
                       j = nxt[j];
               }
       }
}
```

其他

7.1 数字哈希

a[1][1] = m;

for (int i = 2; i <= n; i ++) { for (int j = 1; j < i; j ++) b[j] = pr(a[i - 1][j], j);

int s = m, nd = (i - 1) / 2;

→ y.id);});

```
namespace my_hash {
       const int N = (1 << 19) - 1; //散列大小, 一定要取 2^{-n}-1, 不超内存的情况下, N越大碰撞
        → 越少
       struct E {
               int v;
               E *nxt;
       *g[N + 1], pool[N], *cur = pool, *p;
       int vis[N + 1], T;
       void ins(int v) {
               int u = v \& N;
               if (vis[u] < T) vis[u] = T, g[u] = NULL;</pre>
               for (p = g[u]; p; p = p \rightarrow nxt) if (p \rightarrow v == v) return;
               p = cur ++; p -> v = v; p -> nxt = g[u]; g[u] = p;
       }
       int ask(int v) {
               int u = v \& N;
               if (vis[u] < T) return 0;</pre>
               for (p = g[u]; p; p = p \rightarrow nxt) if (p \rightarrow v == v) return 1;
               return 0;
       }
       void init() {T ++, cur = pool;}//应对多组数据使用
}
7.2 海岛分金币
7.2.1 海岛分金币 1
非朴素模型,有额外条件:
每个人做决定时如果有多种方案可以使自己获得最大收益
那么他会让决策顺序靠前的人获得的收益尽可能的大!
solution:
贪心模拟
*/
#define v first
#define id second
typedef pair<int, int> pr;
const int N = 1010;
int a[N][N];
pr b[N];
int n, m;
int main() {
   cin >> n >> m;
```

sort $(b + 1, b + i, [\&] (pr x, pr y) \{return x.v != y.v ? (x.v < y.v) : (x.id > y.v) \}$

//按照是否容易满足来排序,因为容易满足的人消耗掉的金币比较少,也就使得当前的人获利最大

```
for (int j = 1; j < i && nd; j ++) {
          nd --;
          s = (a[i][b[j].id] = a[i - 1][b[j].id] + 1);
      }
      if (s < 0) {
          for (int j = 1; j < i; j ++)
             a[i][j] = a[i - 1][j];
          a[i][i] = -1;
      }
      else {
          a[i][i] = s;
   }
   for (int i = n; i; i --)
      printf("%d ", a[n][i]);
   return 0;
}
7.2.2 海岛分金币 2
海盗分金币朴素模型:
n 个海盗分 m 个金币, 依次做决策, 如果不少于半数的人同意则方案通过, 否则当前做决策的人会被淘汰
→ (收益视为-1), 由下一人做出决策
如果一个海盗有多种方案均为最大收益, 那么他会希望淘汰的人越多越好
求出第 x 个做决策的海盗的最大可能受益和最小可能收益
struct node {
   int min_v, max_v;
   node():min_v(0), max_v(0)  {}
   node(int min_v, int max_v):min_v(min_v), max_v(max_v) {}
};
int y = n + 1 - x;
   if (n >= (m + 2) * 2) {
      int a = (m + 1) * 2, b = 2, c = 4;
      //前 a 个为 [0,1], 后 b 个为 [0,0], 将持续 c 个
      while (a + b + c \le n) {
          a += b;
          b *= 2;
          c *= 2;
      if (y <= a) return node(0, 1);</pre>
      else if (y \le a + b) return node(0, 0);
      else return node(-1, -1);
   else if (n == m * 2 + 3) {
      if (x == 1) return node(-1, -1);
      else if (y \le m * 2 \&\& y \% 2 == 1 || x == 2) return node(0, 0);
      else return node(0, 1);
   else if (n == m * 2 + 2) {
      if (y \le m * 2 \&\& y \% 2 == 1 || x == 1) return node(0, 0);
      else return node(0, 1);
```

```
}
   else if (n == m * 2 + 1) {
      if (y \le m * 2 \&\& y \% 2 == 1) return node(1, 1);
      else return node(0, 0);
   }
   else {
      if (x & 1) {
          if (x != 1) return node(1, 1);
          else return node(m - (n - 1) / 2, m - (n - 1) / 2);
      }
      else return node(0, 0);
   }
}
int main() {
   ios::sync_with_stdio(false);
   int x, n, m, k; node y;
   cin >> n >> m >> k;
   while (k --) {
      cin >> x;
      y = ask(n, m, x);
      printf("%d %d\n", y.min_v, y.max_v);
   }
   return 0;
}
/*
m = 5
1 5
2 0 5
3 1 0 4
4 0 1 0 4
5 1 0 1 0 3
6 0 1 0 1 0 3
7 1010102
8 0 1 0 1 0 1 0 2
9 1 0 1 0 1 0 1 0 1
10 0 1 0 1 0 1 0 1 0 1
11 1 0 1 0 1 0 1 0 1 0 0
12 0 _ 0 _ 0 _ 0 _ 0 _ 0
13 0 _ 0 _ 0 _ 0 _ 0 _ 0 _ 1
14 _ _ _ _ 0 0
15 _ _ _ _ 0 0 -1
16 _ _ _ _ 0 0 -1 -1
17 _ _ _ _ 0 0 -1 -1 -1
18 _ _ _ _ 0 0 0 0
19 _ _ _ _ 0 0 0 0 -1
20 _ _ _ _ 0 0 0 0 -1 -1
21 _ _ _ _ _ 0 0 0 0 -1 -1 -1
22 _ _ _ _ 0 0 0 0 -1 -1 -1 -1
23 _ _ _ _ _ _ _ _ _ _ _ _ 0 0 0 0 -1 -1 -1 -1 -1
24 _ _ _ _ _ _ 0 0 0 0 -1 -1 -1 -1 -1 -1
25 _ _ _ _ _ 0 0 0 0 -1 -1 -1 -1 -1 -1 -1
```

```
7.3 根号枚举
for (int i = 1, last; i <= n; i = last + 1) {</pre>
       last = n / (n / i);
       //当前枚举区间为 [i, last]
}
7.4 读入输出外挂
namespace IO {//only for int!!!
    static const int SIZE = 1 << 20;</pre>
    inline int get_char() {
       static char *S, *T = S, buf[SIZE];
       if (S == T) {
           T = fread(buf, 1, SIZE, stdin) + (S = buf);
           if (S == T) return -1;
       }
       return *S ++;
   }
    inline void in(int &x) {//for int
       static int ch;
       while (ch = get_char(), ch < 48);x = ch ^ 48;
       while (ch = get char(), ch > 47) x = x * 10 + (ch^48);
   }
   char buffer[SIZE];
    char *s = buffer;
   void flush() {//最后需要 flush!!
       fwrite(buffer, 1, s - buffer, stdout);
       s = buffer;
       fflush(stdout);
   }
    inline void print(const char ch) {
       if(s - buffer > SIZE - 2) flush();
       *s++ = ch;
   }
    inline void print(char *str) {//for string
       while(*str != 0)
           print(char(*str ++));
   }
    inline void print(int x) {
       static char buf [25];
       static char *p = buf;
       if (x < 0) print('-'), x = -x;
       if (x == 0) print('0');
       while(x) *(++ p) = x \% 10, x /= 10;
       while(p != buf) print(char(*(p --) ^ 48));
   }
};
```

7.5 给定小数化成分数

```
# 本题答案的分母不超过 1e9, 给定小数的小数点位为 18 位
# 单次 O(log2n)
inf, inff = 10 ** 9, 10 ** 18
for i in range(int(input())):
   n = int(input()[2:])
    if n == 0: print('0 1')
    else:
        lp, lq, rp, rq = 0, 1, 1, 1
        while max(lq, rq) <= inf:</pre>
            mp, mq = lp + rp, lq + rq
            if mp * inff <= mq * n:</pre>
                1, r, mid, cnt = 1, (inf - lq) // rq + 1, -1, -1
                while 1 <= r:
                    mid = 1 + r >> 1
                    if (lp + rp * mid) * inff <= (lq + rq * mid) * n:
                        cnt, 1 = mid, mid + 1
                    else:
                        r = mid - 1
                lp, lq = lp + rp * cnt, lq + rq * cnt
            else:
                1, r, mid, cnt = 1, (inf - rq) // lq + 1, -1, -1
                    while 1 <= r:
                              mid = 1 + r >> 1
                        if (rp + lp * mid) * inff > (rq + lq * mid) * n:
                            cnt, 1 = mid, mid + 1
                            r = mid - 1
           rp, rq = rp + lp * cnt, rq + lq * cnt
        if lq <= inf: print(lp, lq)</pre>
        else: print(rp, rq)
```