Algorithm Template Library

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1 图论

1.1 线段树维护树直径

```
#include <bits/stdc++.h>
using namespace std;
const int N = 2e5 + 5;
int T, n, m;
int len, head[N], ST[20][N];
struct edge{int u, v, w;}ee[N];
int cnt, fa[N], log_2[N], st[N], en[N], dfn[N], dis[N],
\rightarrow dep[N], pos[N];
struct edges{int to, next, cost;}e[N];
void add(int u, int v, int w) {
    e[++ len] = (edges){v, head[u], w}, head[u] = len;
    e[++ len] = (edges)\{u, head[v], w\}, head[v] = len;
}
void dfs1(int u) {
    st[u] = ++ cnt, dfn[cnt] = u;
    for (int v, i = head[u]; i; i = e[i].next) {
        v = e[i].to;
        if (v == fa[u]) continue;
        fa[v] = u, dep[v] = dep[u] + 1;
        dis[v] = dis[u] + e[i].cost, dfs1(v);
    en[u] = cnt;
}
void dfs2(int u) {
    dfn[++ cnt] = u, pos[u] = cnt;
    for (int v, i = head[u]; i; i = e[i].next) {
        v = e[i].to;
        if (v == fa[u]) continue;
        dfs2(v), dfn[++ cnt] = u;
    }
```

```
}
int mmin(int x, int y) {
    if (dep[x] < dep[y]) return x;</pre>
    return y;
}
int lca(int u, int v) {
    static int w;
    if (pos[u] > pos[v]) swap(u, v);
    w = log_2[pos[v] - pos[u] + 1];
    return mmin(ST[w][pos[u]], ST[w][pos[v] - (1 << w) + 1]);
}
int dist(int u, int v) {
    int Lca = lca(u, v);
    return dis[u] + dis[v] - dis[Lca] * 2;
}
void build() {
    for (int i = 1; i <= cnt; i ++)
        ST[0][i] = dfn[i];
    for (int i = 1; i < 20; i ++)
        for (int j = 1; j <= cnt; j ++)
            if (j + (1 << (i - 1)) > cnt) ST[i][j] = ST[i -
             → 1][j];
            else ST[i][j] = mmin(ST[i - 1][j], ST[i - 1][j +
             \hookrightarrow (1 << (i - 1))]);
}
int M;
struct node {
    int 1, r, dis;
}tr[N << 1];</pre>
void update(int o, int o1, int o2) {
    static int d; static node tmp;
    if (tr[o1].dis == -1) {tr[o] = tr[o2]; return;}
    if (tr[o2].dis == -1) {tr[o] = tr[o1]; return;}
    if (tr[o1].dis > tr[o2].dis) tmp = tr[o1];
    else tmp = tr[o2];
    d = dist(tr[o1].1, tr[o2].1);
```

```
if (d > tmp.dis) tmp.l = tr[o1].l, tmp.r = tr[o2].l,
    \rightarrow tmp.dis = d;
    d = dist(tr[o1].1, tr[o2].r);
    if (d > tmp.dis) tmp.l = tr[o1].l, tmp.r = tr[o2].r,
    \rightarrow tmp.dis = d;
    d = dist(tr[o1].r, tr[o2].1);
    if (d > tmp.dis) tmp.l = tr[o1].r, tmp.r = tr[o2].l,
    \rightarrow tmp.dis = d;
    d = dist(tr[o1].r, tr[o2].r);
    if (d > tmp.dis) tmp.l = tr[o1].r, tmp.r = tr[o2].r,

    tmp.dis = d;

    tr[o] = tmp;
}
void ask(int s, int t) {
    if (s > t) return;
    for (s += M - 1, t += M + 1; s ^ t ^ 1; s >>= 1, t >>= 1)
        if (~s&1) update(0, 0, s ^ 1);
        if ( t&1) update(0, 0, t ^ 1);
    }
}
int main() {
    ios::sync_with_stdio(false);
    int u, v, w, ans; log_2[1] = 0;
    for (int i = 2; i <= 200000; i ++)
        if (i == 1 << (log_2[i - 1] + 1))
            log_2[i] = log_2[i - 1] + 1;
        else log_{2}[i] = log_{2}[i - 1];
    for (cin >> T; T --; ) {
        cin >> n >> m, cnt = len = 0;
        for (int i = 1; i <= n; i ++)
            head[i] = 0;
        for (int i = 1; i < n; i ++) {
            cin >> ee[i].u >> ee[i].v >> ee[i].w;
            add(ee[i].u, ee[i].v, ee[i].w);
        }
        dfs1(1);
        for (M = 1; M < n + 2; M <<= 1);
        for (int i = 1; i <= n; i ++)
            tr[i + M].l = tr[i + M].r = dfn[i], tr[i + M].dis
        for (int i = n + M + 1; i \le (M \le 1) + 1; i ++)
```

tr[i].dis = -1;

```
cnt = 0, dfs2(1), build();
        for (int i = M; i; i --)
            update(i, i << 1, i << 1 | 1);
        for (int i = 1; i < n; i ++)
             if (dep[ee[i].u] > dep[ee[i].v])
                 swap(ee[i].u, ee[i].v);
        for (int u, v, i = 1; i <= m; i ++) {
             cin >> u >> v, ans = 0;
            u = ee[u].v, v = ee[v].v, w = lca(u, v);
            if (w == u \mid \mid w == v) {
                 if (w != u) swap(u, v);
                 tr[0].dis = -1, ask(1, st[u] - 1), ask(en[u] +
                 \rightarrow 1, n), ans = max(ans, tr[0].dis);
                 tr[0].dis = -1, ask(st[u], st[v] - 1),
                 \rightarrow ask(en[v] + 1, en[u]), ans = max(ans,
                 \rightarrow tr[0].dis);
                 tr[0].dis = -1, ask(st[v], en[v]), ans =

→ max(ans, tr[0].dis);
            }
            else {
                 if (st[u] > st[v]) swap(u, v);
                 tr[0].dis = -1, ask(1, st[u] - 1), ask(en[u] +
                 \rightarrow 1, st[v] - 1), ask(en[v] + 1, n), ans =
                 \rightarrow max(ans, tr[0].dis);
                 tr[0].dis = -1, ask(st[u], en[u]), ans =
                 \rightarrow max(ans, tr[0].dis);
                 tr[0].dis = -1, ask(st[v], en[v]), ans =
                 \rightarrow max(ans, tr[0].dis);
            printf("%d\n", ans);
        }
    }
    return 0;
}
     有向图判断两个点能否到达
/* 时间 O(n*m/32) 空间 O(n*n/blk) blk 随便取 */
#include <bits/stdc++.h>
using namespace std;
const int N = 1e5 + 2;
```

```
const int BLK = 5000;
int n, k, p;
int d[N];
vector <int> e[N], f[N], ck[N];
int top, sta[N], in[N];
int cnt, dfn[N], low[N], vis[N];
int sum, bel[N];
bitset <BLK> a[N];
queue <int> q;
int topo[N];
void tarjan(int u) {
        vis[u] = in[u] = 1;
        sta[++ top] = u, dfn[u] = low[u] = ++ cnt;
        for (int v : e[u])
                if (!vis[v]) {
                        tarjan(v);
                        low[u] = min(low[v], low[u]);
                else if (in[v])
                        low[u] = min(low[v], low[u]);
        if (low[u] == dfn[u]) {
                sum ++; int i;
                while (1) {
                        i = sta[top --];
                        in[i] = 0, bel[i] = sum;
                        if (i == u) break;
                }
        }
}
int main() {
        freopen("input.txt", "r", stdin);
        freopen("output.txt", "w", stdout);
```

```
ios::sync_with_stdio(false);
cin >> n >> k;
for (int u, v, i = 1; i <= k; i ++) {
        cin >> u >> v;
        e[u].push_back(v);
}
cin >> p;
for (int u, v, i = 1; i <= p; i ++) {
        cin >> u >> v;
        f[u].push_back(v);
}
for (int i = 1; i <= n; i ++)
        if (!vis[i])
                tarjan(i);
for (int i = 1; i <= n; i ++) {
        for (int j : f[i]) {
                if (bel[i] == bel[j]) return

→ puts("NO"), 0;

                ck[bel[i]].push_back(bel[j]);
        }
        f[i].clear();
for (int i = 1; i <= n; i ++)
        for (int j : e[i]) {
                if (bel[i] == bel[j]) continue;
                f[bel[i]].push_back(bel[j]);
                d[bel[j]] ++;
        }
cnt = 0;
for (int i = 1; i <= sum; i ++)
        if (!d[i])
                q.push(i);
while (!q.empty()) {
        int now = q.front(); q.pop();
        topo[++ cnt] = now;
        for (int j : f[now]) {
                d[j] --;
                if (d[j] == 0) q.push(j);
        }
}
```

```
for (int i = 1, t = (sum + BLK - 1) / BLK; i <= t; i
         → ++) {
                for (int j = sum; j; j --) {
                         int u = topo[j];
                         a[u].reset();
                         if (BLK * (i - 1) < u \&\& u \le BLK * i)
                                 a[u][u - BLK * (i - 1) - 1] =
                                 for (int v : f[u])
                                 a[u] |= a[v];
                }
                for (int j = 1; j <= sum; j ++)
                         for (int v : ck[j])
                                 if (BLK * (i - 1) < v && v <=
                                 \rightarrow BLK * i && a[j][v - BLK *
                                  \hookrightarrow (i - 1) - 1] == 1) {
                                         puts("NO");
                                         return 0;
                                 }
        }
        printf("YES\n%d\n", k);
        for (int i = 1; i <= n; i ++)
                for (int j : e[i])
                        printf("%d %d\n", i, j);
        return 0;
}
1.3 spfa 费用流
const int N = 600, M = 800000, inf = 0x3f3f3f3f;
int s, t, ans, len, maxflow;
int T, n, m, K, W;
int head[N], incf[N], path[N], pre[N], vis[N], d[N];
struct edge{
        int to, next, cap, cost;
}e[M];
struct video {z
        int s, t, w, op;
```

```
}a[N];
void add(int u, int v, int w, int c) {
        e[++ len] = (edge)\{v, head[u], w, c\}, head[u] = len;
        e[++ len] = (edge)\{u, head[v], 0, -c\}, head[v] = len;
}
bool spfa() {
        deque <int> q;
        q.push_back(s), incf[s] = inf;
        for (int i = 1; i <= t; i ++) d[i] = inf;</pre>
        d[s] = 0;
        while (!q.empty()) {
                 int x = q.front();
                 q.pop_front(), vis[x] = 0;
                 for (int i = head[x]; i; i = e[i].next) {
                         if (e[i].cap \&\& d[e[i].to] > d[x] +
                          → e[i].cost) {
                                  d[e[i].to] = d[x] + e[i].cost;
                                  pre[e[i].to] = x,
                                  → path[e[i].to] = i;
                                  incf[e[i].to] = min(incf[x],
                                  \rightarrow e[i].cap);
                                  if (!vis[e[i].to]) {
                                          vis[e[i].to] = 1;
                                          if (q.empty() ||
                                           \rightarrow d[e[i].to] <

    d[q.front()])

¬ q.push_front(e[i].to);

                                          else

¬ q.push_back(e[i].to);

                                 }
                         }
                 }
        maxflow += incf[t];
        if (d[t] == inf) return 0;
        for (int i = t; i != s; i = pre[i]) {
                 e[path[i]].cap -= incf[t];
                 e[path[i] ^ 1].cap += incf[t];
        return ans += incf[t] * d[t], 1;
}
```

```
int main() {
        /*build graph*/
        while(spfa());
}
1.4 dinic 最大流
#include <bits/stdc++.h>
using namespace std;
const int N = 20000;
const int M = 500000;
const int inf = 0x3f3f3f3f;
int n, m;
int s, t, len = 1;
int to[M], cap[M], nex[M];
int g[N], p[N], q[N], d[N];
void add(int x, int y, int v) {
        to[++ len] = y, cap[len] = v, nex[len] = g[x], g[x] =
        \hookrightarrow len;
        to[++ len] = x, cap[len] = 0, nex[len] = g[y], g[y] =
        → len;
}
bool bfs() {
        int 1 = 1, r = 1, x, i;
        memset (d, 0, sizeof d);
        d[s] = 1, q[1] = s;
        while (1 <= r) {
                x = q[1 ++];
                for (i = g[x]; i; i = nex[i])
                         if (cap[i] && !d[to[i]])
                                 d[to[i]] = d[x] + 1, q[++ r] =
                                  \rightarrow to[i];
        }
        return d[t];
```

```
}
int dfs(int x, int y) {
       if (x == t \mid \mid y == 0) return y;
       int flow = 0;
       for (int &i = p[x]; i; i = nex[i]) {
               if (!cap[i] || d[to[i]] != d[x] + 1) continue;
               int f = dfs(to[i], min(y, cap[i]));
               flow += f, y -= f;
               cap[i] -= f, cap[i ^ 1] += f;
               if (!y) break;
       return flow;
}
int dinic() {
       int maxflow = 0;
       while (bfs()) {
               memcpy(p, g, sizeof g);
               maxflow += dfs(s, inf);
       return maxflow;
}
int main() {
       ios::sync_with_stdio(false);
       return 0;
}
1.5 DAG 删去无用边
无用边定义:对于边 (u,v) 如果存在从 u 到 v 不经过该边的另一条路
→ 径,则称该边无用
时间复杂度:0(n~3)
bool f[N][N];//i 是否可达 j
vector <int> e[N];
int main() {
       rep(i, 1, n)
               for (int j : e[i]) {
```

```
for (k, 1, n)
                                 if (i != k && j != k &&
                                  \rightarrow f[i][k] && f[k][j])
                                         no_use_edge;
                }
}
1.6 树分治
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
const int N = 1e5 + 5;
vector <int> e[N];
int n, a[N];
int root, _left, vis[N];
int siz[N], maxv[N];
void find_root(int u, int fr) {
        siz[u] = 1, maxv[u] = 0;
        for (int v : e[u]) {
                if (v == fr || vis[v]) continue;
                find_root(v, u);
                siz[u] += siz[v];
                maxv[u] = max(maxv[u], siz[v]);
        maxv[u] = max(maxv[u], _left - siz[u]);
        if (!root || maxv[u] < maxv[root])</pre>
                root = u;
}
void dfs(int u, int fr) {
        siz[u] = 1;
        for (int v : e[u]) {
                if (v == fr || vis[v]) continue;
                find_root(v, u);
                siz[u] += siz[v];
```

```
}
}
void solve(int u, int w) {
        dfs(u, u);//update siz[]
        a[u] = w, vis[u] = 1;
        for (int v : e[u]) {
                if (vis[v]) continue;
                _left = siz[v];
                root = 0;
                find_root(v, v);
                solve(root, w + 1);
        }
}
int main() {
    ios::sync_with_stdio(false);
    cin >> n;
    for (int u, v, i = 1; i < n; i ++) {
            cin >> u >> v;
            e[u].push_back(v);
            e[v].push_back(u);
    }
    _left = n, root = 0, find_root(1, 1);
    solve(root, 0);
    for (int i = 1; i <= n; i ++)
            printf("%c ", 'A' + a[i]);
    return 0;
}
```

res -> init(x);

2 数据结构

```
return res;
2.1 treap
                                    }
/* 容易实现的预开内存池 treap,
                                    void rot(node *&o, int d) {
→ 每次 head 清空即可
                                              node *tmp = o ->
如果初始要插入 n \land 1, 可改为类
                                              \rightarrow c[!d];
→ 似 splay 的 O(n)build 写法
                                              o \rightarrow c[!d] = tmp \rightarrow
                                              \hookrightarrow c[d], tmp -> c[d]
#include <bits/stdc++.h>
                                              \rightarrow = o;
                                              o -> update(), tmp ->
using namespace std;

    update(), o = tmp;

const int poolSize = 5e5 + 10;
                                    void insert(node *&o, int x) {
struct node {
                                              if (o == null) {
         node *c[2];
                                                       0 =
         int v, r, siz;
                                                       \rightarrow newnode(x);
         void update();
                                                       return;
         void init(int x);
                                              }
};
                                              int d = x > o \rightarrow v?
                                              node *null = new node(), *root
                                              insert(o -> c[d], x);
\rightarrow = null;
                                              if (o \rightarrow c[d] \rightarrow r < o
                                              \rightarrow -> r) rot(o, !d);
void node::update() {
                                              o -> update();
         siz = c[0] \rightarrow siz +
                                    }
         \rightarrow c[1] -> siz + 1;
}
                                    void del(node *&o, int x) {
                                              if (x == o -> v) {
void node::init(int x) {
                                                       if (o -> c[0]
         v = x, r = rand(), siz
                                                        \rightarrow == null)
         \hookrightarrow = 1;
                                                        \hookrightarrow {o = o ->
         c[0] = c[1] = null;
                                                        \hookrightarrow c[1];
}
                                                        → return;}
                                                       if (o -> c[1]
node nodesPool[poolSize];
                                                        \rightarrow == null)
                                                        \rightarrow {o = o ->
int head;//每次 head=0 清空
                                                       \rightarrow c[0];
                                                       → return;}
node *newnode(int x) {
         node *res =

→ ++];
```

```
int d = o ->
                                          */
                     \rightarrow c[0] \rightarrow r
                                        #include <bits/stdc++.h>
                    \rightarrow c[1] \rightarrow r
                                        using namespace std;
                    \rightarrow ? 1 : 0;
                    rot(o, d),
                                        const int N = 1e5 + 5;
                    \rightarrow del(o ->
                    \hookrightarrow c[d], x);
                                        const int Mod = 1e9 + 7;
          else del(o \rightarrow c[x <= o]
                                        int nowD, ans, x[3], y[3];
          \rightarrow -> v], x);
          o -> update();
                                        int n, m, a[N], b[N], c[N],
}
                                         \rightarrow d[N];
void build(node *&o, int 1,
                                        struct node {
\rightarrow int r) {
                                             int Max[3], Min[3], d[3];
          o = newnode(1);
                                             int val, maxv;
          if (1 == r) return;
                                             node *c[2];
          int mid = 1 + r >> 1;
          if (l < mid) build(o</pre>
                                             node() {
          \rightarrow -> c[0], 1, mid -
                                                  c[0] = c[1] = NULL;
                                                  val = maxv = 0;
          \rightarrow 1);
          if (o -> c[0] != null
                                             }
          \rightarrow && o -> c[0] -> r
          \rightarrow < o -> r) swap(o
                                             void pushup();
          \rightarrow -> c[0] -> r, o ->
                                             bool operator < (const</pre>
          \rightarrow r);
          if (mid < r) build(o</pre>
                                              → node &a) const {
          \rightarrow -> c[1], mid + 1,
                                                  return d[nowD] <
          \rightarrow r);
                                                   → a.d[nowD];
          if (o -> c[1] != null
          \hookrightarrow && o -> c[1] -> r
                                        }Null, nodes[N];
          \rightarrow < o \rightarrow r) swap(o
          \rightarrow -> c[1] -> r, o ->
                                        node *root = &Null;
          \hookrightarrow r);
                                        inline void node::pushup() {
          o -> update();
}
                                             if (c[0] != &Null) {
                                                  if (c[0] \rightarrow Max[1] >
                                                   \rightarrow Max[1]) Max[1] =
2.2 KDTree
                                                   \rightarrow c[0] -> Max[1];
2.2.1 3 维 KDtree
                                                  if (c[0] \rightarrow Max[2] >
                                                   \rightarrow Max[2]) Max[2] =
                                                   \rightarrow c[0] -> Max[2];
 *O(n*n^(1-1/k)),k 为维度
```

```
if (c[0] \rightarrow Min[0] <
                                               return res;
           \rightarrow Min[0]) Min[0] =
           \hookrightarrow c[0] -> Min[0];
          if (c[0] \rightarrow Min[2] <
                                          inline int calc(node *o) {
           \hookrightarrow Min[2]) Min[2] =
                                               if (y[0] < o -> Min[0] ||
           \hookrightarrow c[0] -> Min[2];
                                               \rightarrow x[1] > o -> Max[1] ||
          if (c[0] -> maxv >
                                                \rightarrow x[2] > o -> Max[2] ||
           \rightarrow maxv) maxv = c[0]
                                               \rightarrow y[2] < o -> Min[2])
           \hookrightarrow -> maxv;
                                                \hookrightarrow return -1;
     }
                                               return o -> maxv;
     if (c[1] != &Null) {
                                          }
          if (c[1] -> Max[1] >
           \rightarrow Max[1]) Max[1] =
                                          inline void query(node *o) {
           \hookrightarrow c[1] -> Max[1];
                                               if (o -> val > ans && y[0]
          if (c[1] -> Max[2] >
                                                \rightarrow >= o -> d[0] && x[1]
           \rightarrow Max[2]) Max[2] =
                                                \rightarrow <= o -> d[1] && x[2]
           \rightarrow c[1] -> Max[2];
                                               \rightarrow <= o -> d[2] && y[2]
          if (c[1] \rightarrow Min[0] <
                                               \rightarrow >= o -> d[2]) ans = o
           \rightarrow Min[0]) Min[0] =
                                               → -> val;
           \hookrightarrow c[1] -> Min[0];
                                               int dl, dr;
          if (c[1] \rightarrow Min[2] <
                                              if (o -> c[0] != &Null) dl
           \rightarrow Min[2]) Min[2] =
                                               \Rightarrow = calc(o \rightarrow c[0]);
           \hookrightarrow c[1] -> Min[2];
                                               else dl = -1;
          if (c[1] \rightarrow maxv >
                                              if (o -> c[1] != &Null) dr
           \rightarrow maxv) maxv = c[1]
                                               \Rightarrow = \operatorname{calc}(o \rightarrow c[1]);
           \hookrightarrow -> maxv;
                                               else dr = -1;
                                               if (dl > dr) {
     }
}
                                                    if (dl > ans) query(o
                                                     \rightarrow -> c[0]);
inline node *build(int 1, int
                                                    if (dr > ans) query(o
\rightarrow r) {
                                                     \rightarrow -> c[1]);
     int mid = 1 + r >> 1; nowD
                                               } else {
     \rightarrow = rand() % 3;
                                                    if (dr > ans) query(o
     nth_element(nodes + 1,
                                                     \rightarrow -> c[1]);
     \rightarrow nodes + mid, nodes + r
                                                    if (dl > ans) query(o
     \rightarrow + 1);
                                                     \rightarrow -> c[0]);
     node *res = &nodes[mid];
                                               }
     if (1 != mid) res -> c[0]
                                          }
     \rightarrow = build(1, mid - 1);
     else res \rightarrow c[0] = &Null;
                                          int main() {
     if (r != mid) res -> c[1]
     \rightarrow = build(mid + 1, r);

→ ios::sync_with_stdio(false);
     else res \rightarrow c[1] = &Null;
                                             cin >> n >> m;
     res -> pushup();
```

```
2.2.2 KDtree 二维空间区间覆盖
    for (int i = 1; i <= n; i
                                            单点查询
     → ++) {
         cin >> a[i];
                                    /* 类似线段树 */
        b[i] = d[a[i]];
                                    #include <bits/stdc++.h>
         d[a[i]] = i;
    }
                                    using namespace std;
    for (int i = 1; i <= n; i
     \rightarrow ++) d[i] = n + 1;
                                    const int N = 1e5 + 5;
    for (int i = n; i; i --) {
         c[i] = d[a[i]];
                                    const int Mod = 1e9 + 7;
         d[a[i]] = i;
                                    int nowD, x[2], y[2], z;
    for (int i = 1; i <= n; i
     → ++) {
                                    struct node {
        nodes[i].Min[0] =
                                         int Max[2], Min[2], d[2];
         \rightarrow nodes[i].d[0] =
                                         int val, lazy;
         \hookrightarrow b[i];
                                         node *c[2];
        nodes[i].Max[1] =
         \rightarrow nodes[i].d[1] =
                                         node() {
         \hookrightarrow c[i];
                                             c[0] = c[1] = NULL;
         nodes[i].Max[2] =
         \rightarrow nodes[i].Min[2] =
         \rightarrow nodes[i].d[2] = i;
                                         void pushup();
         nodes[i].val =
         → nodes[i].maxv =
                                         void pushdown();
         \rightarrow a[i];
                                         bool operator < (const</pre>
    root = build(1, n);
                                         → node &a) const {
    for (int 1, r; m --; ) {
                                             return d[nowD] <
         cin >> 1 >> r;

→ a.d[nowD];
         1 = (1 + ans) \% n + 1;
                                         }
         r = (r + ans) \% n + 1;
                                    }Null, nodes[N];
         if (l > r) swap(l, r);
         y[0] = 1 - 1;
                                    node *root = &Null;
         x[1] = r + 1;
         x[2] = 1, y[2] = r;
                                    inline void node::pushup() {
         ans = 0, query(root);
                                         if (c[0] != &Null) {
         cout << ans << endl;</pre>
                                             if (c[0] \rightarrow Max[0] >
    }
                                              \rightarrow Max[0]) Max[0] =
    cout << endl;</pre>
                                              \rightarrow c[0] -> Max[0];
    return 0;
                                             if (c[0] \rightarrow Max[1] >
}
                                              \rightarrow Max[1]) Max[1] =
                                              \rightarrow c[0] -> Max[1];
```

```
if (c[0] \rightarrow Min[0] <
                                               else res -> c[0] = &Null;
           \rightarrow Min[0]) Min[0] =
                                               if (r != mid) res -> c[1]
           \hookrightarrow c[0] -> Min[0];
                                               \rightarrow = build(mid + 1, r,
                                               \hookrightarrow !D);
          if (c[0] -> Min[1] <</pre>
           \rightarrow Min[1]) Min[1] =
                                               else res -> c[1] = &Null;
           \hookrightarrow c[0] -> Min[1];
                                               res -> pushup();
                                               return res;
     if (c[1] != &Null) {
                                         }
          if (c[1] \rightarrow Max[0] >
           \rightarrow Max[0]) Max[0] =
                                         inline int query(node *o) {
           \hookrightarrow c[1] -> Max[0];
                                               if (o == &Null) return -1;
          if (c[1] -> Max[1] >
                                               if (o \rightarrow lazy != -1) o \rightarrow
           \rightarrow Max[1]) Max[1] =
                                               → pushdown();
           \rightarrow c[1] -> Max[1];
                                               if (x[0] > o \rightarrow Max[0] \mid \mid
          if (c[1] \rightarrow Min[0] <
                                               \rightarrow y[0] > o -> Max[1] ||
           \rightarrow Min[0]) Min[0] =
                                               \rightarrow x[0] < o -> Min[0] ||
           \hookrightarrow c[1] -> Min[0];
                                               \rightarrow y[0] < o -> Min[1])
          if (c[1] -> Min[1] <</pre>
                                               \rightarrow return -1;
           \rightarrow Min[1]) Min[1] =
                                               if (x[0] == o -> d[0])
           \hookrightarrow c[1] -> Min[1];

    return o → val;

     }
                                               return max(query(o ->
}
                                                \rightarrow c[0]), query(o ->
                                                   c[1]));
inline void node::pushdown() {
     if (c[0] != &Null) c[0] ->
     \rightarrow val = c[0] -> lazy =
                                          inline void modify(node *o) {
                                               if (o == &Null) return;
     → lazy;
     if (c[1] != &Null) c[1] ->
                                               if (o \rightarrow lazy != -1) o \rightarrow
     \rightarrow val = c[1] -> lazy =
                                               → pushdown();
                                               if (x[0] > o -> Max[0] | |
     → lazy;
     lazy = -1;
                                                    y[0] > o -> Max[1] | |
}
                                                   x[1] < o -> Min[0] | |
                                                \rightarrow y[1] < o -> Min[1])
inline node *build(int 1, int

    return;

                                               if (x[0] \le o \rightarrow Min[0] &&
\rightarrow r, int D) {
     int mid = 1 + r >> 1; nowD
                                                \rightarrow y[0] <= o -> Min[1] &&
     \hookrightarrow = D;
                                                \rightarrow x[1] >= o -> Max[0] &&
     nth_element(nodes + 1,
                                                \rightarrow y[1] >= o -> Max[1]) {
     \rightarrow nodes + mid, nodes + r
                                                    o \rightarrow val = o \rightarrow lazy =
     \rightarrow + 1);
                                                     \hookrightarrow Z;
     node *res = &nodes[mid];
                                                    return;
     if (1 != mid) res -> c[0]
                                               }
     \rightarrow = build(1, mid - 1,
     → !D);
```

```
if (x[0] \le o -> d[0] \&\&
                                                      nodes[i].Min[0] =
     \rightarrow y[0] <= o -> d[1] &&
                                                       → nodes[i].Max[0]
     \rightarrow x[1] >= o -> d[0] &&
     \rightarrow y[1] >= o -> d[1]) o
                                                       \rightarrow nodes[i].d[0]
     \rightarrow -> val = z;
                                                       \hookrightarrow = i;
    modify(o \rightarrow c[0]),
                                                      nodes[i].Min[1] =
     \rightarrow modify(o -> c[1]);
                                                           nodes[i].Max[1]
}
                                                       \rightarrow nodes[i].d[1]
                                                       int n, m, k, a[N], c[N], d[N];
                                                      nodes[i].val = 1,
int cnt, st[N], en[N], dfn[N],
                                                       → nodes[i].lazy
→ dep[N];
                                                       \rightarrow = -1;
                                                 }
vector <int> e[N];
                                                 root = build(1, n, 0);
                                                 for (int u, v, w, i =
void dfs(int u) {
                                                  \rightarrow 1; i <= k; i ++) {
     st[u] = ++ cnt, dfn[cnt] =
                                                      cin >> u >> v >>
     \hookrightarrow u;
                                                       \hookrightarrow W;
     for (int v : e[u])
                                                      if (w == 0) {
         dep[v] = dep[u] + 1,
                                                           x[0] = st[u],
                                                            → y[0] =
          \rightarrow dfs(v);
     en[u] = cnt;
                                                            \rightarrow dep[u];
}
                                                           ans = (ans +
                                                            → 111 * i *
int main() {
                                                            → query(root)
                                                            \rightarrow % Mod) %
     → ios::sync_with_stdio(false);
                                                            \hookrightarrow Mod;
     int T, ans;
                                                      } else {
     for (cin >> T; T --; ) {
                                                           x[0] = st[u],
         cin >> n >> m >> k,
                                                            \rightarrow x[1] =
                                                            \rightarrow en[u];
          \rightarrow ans = cnt = 0;
         for (int i = 1; i <=</pre>
                                                           y[0] = dep[u],
          \hookrightarrow n; i ++)
                                                            \rightarrow y[1] =
              e[i].clear();
                                                            \rightarrow dep[u] +
         for (int u, i = 2; i
                                                            \hookrightarrow V;
          \rightarrow <= n; i ++) {
                                                           z = w,
              cin >> u;
                                                            → modify(root);
              e[u].push_back(i);
                                                      }
         }
                                                 }
         dfs(1);
                                                 cout << ans << endl;</pre>
         for (int i = 1; i <=
                                            }
          \rightarrow n; i ++) {
                                            return 0;
                                       }
```

```
2.2.3 KDtree 二维空间单点修改
                                              if (c[0] \rightarrow Max[0] >
       区间查询
                                               \rightarrow Max[0]) Max[0] =
                                               \rightarrow c[0] -> Max[0];
                                              if (c[0] \rightarrow Max[1] >
 * 调整重构系数可以影响常数
                                               \rightarrow Max[1]) Max[1] =
 * 询问多就让系数接近
                                               \rightarrow c[0] -> Max[1];
→ 0.70-0.75, 询问少就让系数在
                                              if (c[0] -> Min[0] <</pre>
→ 0.8-0.90
                                               \rightarrow Min[0]) Min[0] =
 */
                                               \rightarrow c[0] -> Min[0];
#include <bits/stdc++.h>
                                              if (c[0] -> Min[1] <</pre>
                                               \rightarrow Min[1]) Min[1] =
using namespace std;
                                               \rightarrow c[0] -> Min[1];
                                         }
const int inf = 1e9;
                                         if (c[1] != &Null) {
                                              if (c[1] \rightarrow Max[0] >
int n, m, tot, nowD;
                                               \rightarrow Max[0]) Max[0] =
                                               \hookrightarrow c[1] -> Max[0];
struct node {
                                              if (c[1] -> Max[1] >
    int Max[2], Min[2], d[2];
                                               \rightarrow Max[1]) Max[1] =
    int sum, siz, val;
                                               \rightarrow c[1] -> Max[1];
    node *c[2];
                                              if (c[1] -> Min[0] <</pre>
                                               \rightarrow Min[0]) Min[0] =
    node() {
                                               \rightarrow c[1] -> Min[0];
         Max[0] = Max[1] =
                                              if (c[1] \rightarrow Min[1] <
         → -inf;
                                               \rightarrow Min[1]) Min[1] =
         Min[0] = Min[1] = inf;
                                               \rightarrow c[1] -> Min[1];
         sum = val = siz = 0;
                                         }
         c[0] = c[1] = NULL;
                                     }
         d[0] = d[1] = 0;
    }
                                     inline bool cmp(const node *a,

    const node *b) {
    void update();
                                          return a -> d[nowD] < b ->
}Null, nodes[200010],
                                          \rightarrow d[nowD];
→ *temp[200010];
                                     }
node *root = &Null;
                                     inline void traverse(node *o)
                                      ← {
inline void node::update() {
                                          if (o == &Null) return;
    siz = c[0] \rightarrow siz + c[1]
                                         temp[++ tot] = o;
     → -> siz + 1;
                                         traverse(o -> c[0]);
    sum = c[0] \rightarrow sum + c[1]
                                         traverse(o -> c[1]);
     → -> sum + val;
                                     }
    if (c[0] != &Null) {
```

```
inline node *build(int 1, int
                                              if (p \rightarrow Max[1] > o \rightarrow
\rightarrow r, int D) {
                                               \rightarrow Max[1]) o \rightarrow Max[1] =
     int mid = l + r \gg 1; nowD
                                               \rightarrow p \rightarrow Max[1];
     \rightarrow = D;
                                              if (p -> Min[0] < o ->
     nth_element(temp + 1, temp
                                               \hookrightarrow Min[0]) o -> Min[0] =
                                               \rightarrow p \rightarrow Min[0];
     \rightarrow + mid, temp + r + 1,
     \rightarrow cmp);
                                              if (p -> Min[1] < o ->
     node *res = temp[mid];
                                               \rightarrow Min[1]) o -> Min[1] =
     res -> Max[0] = res ->
                                               \rightarrow p \rightarrow Min[1];
     \rightarrow Min[0] = res -> d[0];
                                              o -> siz ++, o -> sum += p
     res -> Max[1] = res ->

→ ¬> sum;

     \rightarrow Min[1] = res -> d[1];
                                              insert(o \rightarrow c[p \rightarrow c[D] >=
     if (1 != mid) res -> c[0]
                                              \rightarrow o \rightarrow c[D]], p, !D);
     \rightarrow = build(1, mid - 1,
                                              if (max(o \rightarrow c[0] \rightarrow siz,
                                               \rightarrow o -> c[1] -> siz) >
     → !D);
                                               \rightarrow int(o -> siz * 0.75 +
     else res -> c[0] = &Null;
     if (r != mid) res -> c[1]
                                               \rightarrow 0.5)) tmpD = D, tmp =
     \rightarrow = build(mid + 1, r,
                                               \hookrightarrow !D);
                                         }
     else res -> c[1] = &Null;
     res -> update();
                                         inline int query(node *o) {
                                              if (o == &Null) return 0;
     return res;
}
                                              if (x > o -> Max[0] || y >
                                               \rightarrow o \rightarrow Max[1] || a < o
                                               \rightarrow -> Min[0] || b < o ->
int x, y, a, b, tmpD;

→ Min[1]) return 0;

                                              if (x <= o -> Min[0] && y
node **tmp;
                                               \rightarrow <= o -> Min[1] && a >=
inline void rebuild(node *&o,
                                               \rightarrow o -> Max[0] && b >= o
\hookrightarrow int D) {
                                               → -> Max[1]) return o ->
     tot = 0;
                                                   sum;
                                              return (x \leq o \rightarrow d[0] &&
     traverse(o);
     o = build(1, tot, D);
                                               \rightarrow y <= o -> d[1] && a >=
                                               \rightarrow o -> d[0] && b >= o ->
}
                                               \rightarrow d[1] ? o \rightarrow val : 0)
                                                   + query(o -> c[1]) +
inline void insert (node *&o,
                                                    \rightarrow query(o -> c[0]);
\rightarrow node *p, int D) {
     if (o == &Null) {o = p;
                                         }
     → return;}
     if (p -> Max[0] > o ->
                                         int main() {
     \rightarrow Max[0]) o \rightarrow Max[0] =
     \rightarrow p \rightarrow Max[0];
                                               → ios::sync_with_stdio(false);
                                              cin >> m;
                                              node *ttt = &Null;
```

```
for (int t, ans = 0; ; ) {
                                              printf("%d\n",
    cin >> t;
                                               \rightarrow ans);
    if (t == 3) break;
                                         }
    if (t == 1) {
                                     }
         cin >> x >> y >>
                                    return 0;
                                }
         → a;
         x = ans, y =
         \rightarrow ans, n ++;
                                2.2.4 KDtree 找最近点
         nodes[n].sum =
         → nodes[n].val =
                                  * 为了维持树的平衡, 可以一开始
         \hookrightarrow a \hat{} ans,
                                 → 把所有点都读进来 build
             nodes[n].siz =
                                 * 然后打 flag 标记该点是否被激
         \hookrightarrow 1;
                                 → 活
         nodes[n].Max[0] =
                                 */
             nodes[n].Min[0]
                                #include <bits/stdc++.h>
         \rightarrow nodes[n].d[0]
                                using namespace std;
         \rightarrow = x:
         nodes[n].Max[1] =
                                const int N = 5e5 + 5;
         \rightarrow nodes[n].Min[1]
                                const int inf = 1 << 30;</pre>
         \rightarrow nodes[n].d[1]
         \rightarrow = y;
                                int n, m;
         nodes[n].c[0] =
         \rightarrow nodes[n].c[1]
                                int ql, qr, ans, tot, nowD;
         \rightarrow = &Null;
                                //nowD = rand() & 1 ?
         tmp = &(ttt),
                                struct Node {

→ insert(root,
                                     int d[2];
         \rightarrow &nodes[n], 0);
         if (*tmp != &Null)
                                    bool operator < (const</pre>
         → rebuild(*tmp,
                                     → Node &a) const {
         → tmpD);
                                         if (d[nowD] ==
    } else {
                                          → a.d[nowD]) return
         cin >> x >> y >> a
                                          \rightarrow d[!nowD] <
         \rightarrow >> b;
                                          → a.d[!nowD];
         x = ans, y =
                                         return d[nowD] <
         \rightarrow ans, a \hat{} ans,
                                          → a.d[nowD];
         \hookrightarrow b \hat{} ans;
                                    }
         if (x > a) swap(x,
                                }pot[N];
         → a);
         if (y > b) swap(y,
                                struct node {
         \rightarrow b);
                                     int min[2], max[2], d[2];
         ans = query(root);
                                    node *c[2];
```

```
if (c[1] -> max[1] >
     node() {
                                                     \rightarrow max[1]) max[1] =
          min[0] = min[1] =
                                                     \rightarrow c[1] -> max[1];
          \rightarrow max[0] = max[1] =
                                                    if (c[1] -> min[0] <
          \rightarrow d[0] = d[1] = 0;
                                                     \rightarrow min[0]) min[0] =
          c[0] = c[1] = NULL;
                                                     \hookrightarrow c[1] -> min[0];
     }
                                                    if (c[1] -> min[1] <</pre>
                                                     \rightarrow min[1]) min[1] =
                                                     \rightarrow c[1] \rightarrow min[1];
     node(int x, int y);
                                               }
     void update();
                                          }
                                          inline void build(node *&o,
}t[N], Null, *root;
                                          → int 1, int r, int D) {
                                               int mid = 1 + r >> 1;
node::node(int x, int y) {
                                               nowD = D;
     min[0] = max[0] = d[0] =
                                               nth_element(pot + 1, pot +
                                               \rightarrow mid, pot + r + 1);
     \hookrightarrow x:
     min[1] = max[1] = d[1] =
                                               o = new
                                               \rightarrow node(pot[mid].d[0],
     \rightarrow y;
     c[0] = c[1] = &Null;
                                               \rightarrow pot[mid].d[1]);
}
                                               if (1 != mid) build(o ->
inline void node::update() {
                                               \rightarrow c[0], 1, mid - 1, !D);
     if (c[0] != &Null) {
                                               if (r != mid) build(o ->
          if (c[0] -> max[0] >
                                               \rightarrow c[1], mid + 1, r, !D);
           \rightarrow max[0]) max[0] =
                                               o -> update();
           \rightarrow c[0] -> max[0];
                                         }
          if (c[0] -> max[1] >
           \rightarrow max[1]) max[1] =
                                          inline void insert(node *o) {
           \rightarrow c[0] -> max[1];
                                               node *p = root;
          if (c[0] -> min[0] <</pre>
                                               int D = 0;
           \rightarrow min[0]) min[0] =
                                               while (1) {
           \rightarrow c[0] -> min[0];
                                                    if (o \rightarrow max[0] > p \rightarrow
          if (c[0] -> min[1] <
                                                    \rightarrow max[0]) p \rightarrow
           \rightarrow min[1]) min[1] =
                                                    \rightarrow max[0] = o ->
           \rightarrow c[0] -> min[1];
                                                    \rightarrow max[0];
     }
                                                    if (o \rightarrow max[1] > p \rightarrow
     if (c[1] != &Null) {
                                                    \rightarrow max[1]) p \rightarrow
          if (c[1] \rightarrow max[0] >
                                                    \rightarrow max[1] = o ->
           \rightarrow max[0]) max[0] =
                                                     \rightarrow max[1];
           \rightarrow c[1] \rightarrow max[0];
```

```
if (o -> min[0] 
                                             d0 = abs(o \rightarrow d[0] - q1) +
                                             \rightarrow abs(o -> d[1] - qr);
          \rightarrow min[0]) p ->
          \rightarrow min[0] = o ->
                                             if (d0 < ans) ans = d0;
          \rightarrow min[0];
                                             if (o -> c[0] != &Null) dl
          if (o -> min[1] 
                                             \rightarrow = dist(o -> c[0]);
          \rightarrow min[1]) p ->
                                             else dl = inf;
          \rightarrow min[1] = o ->
                                             if (o -> c[1] != &Null) dr
          \rightarrow min[1];
                                             \rightarrow = dist(o -> c[1]);
                                             else dr = inf;
         if (o \rightarrow d[D] \rightarrow p \rightarrow
                                             if (dl < dr) {
          \rightarrow d[D]) {
               if (p \rightarrow c[1] ==
                                                  if (dl < ans) query(o

→ &Null) {
                                                  \rightarrow -> c[0]);
                                                  if (dr < ans) query(o</pre>
                   p \rightarrow c[1] = o;
                   return;
                                                  \rightarrow -> c[1]);
                                             } else {
               } else p = p \rightarrow
               \hookrightarrow c[1];
                                                  if (dr < ans) query(o
         } else {
                                                  \rightarrow -> c[1]);
               if (p -> c[0] ==
                                                  if (dl < ans) query(o</pre>
               \rightarrow -> c[0]);
                   p -> c[0] = o;
                                             }
                                        }
                   return;
               } else p = p \rightarrow
               \rightarrow c[0];
                                        int main() {
         }
         D = 1;
                                              → ios::sync_with_stdio(false);
                                             cin >> n >> m;
    }
}
                                             for (int i = 1; i <= n; i
                                             → ++)
inline int dist(node *o) {
                                                  cin >> pot[i].d[0] >>
     int dis = 0;
                                                  \rightarrow pot[i].d[1];
     if (q1 < o -> min[0]) dis
                                             build(root, 1, n, 0);
     \rightarrow += o -> min[0] - ql;
     if (ql > o \rightarrow max[0]) dis
                                             for (int x, y, z; m --; )
     \rightarrow += ql - o -> max[0];
     if (qr < o \rightarrow min[1]) dis
                                                  cin >> x >> y >> z;
     \rightarrow += o -> min[1] - qr;
                                                  if (x == 1) {
     if (qr > o \rightarrow max[1]) dis
                                                       t[tot].max[0] =
     \rightarrow += qr - o -> max[1];

    t[tot].min[0]

    return dis;
                                                       \rightarrow = t[tot].d[0]
}
                                                       \rightarrow = y;
inline void query(node *o) {
     int dl, dr, d0;
```

```
t[tot].max[1] =
                \rightarrow t[tot].min[1]
                \rightarrow = t[tot].d[1]
                \hookrightarrow = z;
               t[tot].c[0] =
                \rightarrow t[tot].c[1] =

→ &Null;

                insert(&t[tot
                → ++]);
          } else {
               ans = inf, ql = y,
                \rightarrow qr = z;
               query(root),
                \rightarrow printf("%d\n",
                \hookrightarrow ans);
          }
     }
     return 0;
}
```

构造 3 int t = a[n -→ 1]; 3.1 若干排列使所有数对都出现 for (int j = n)一次 \hookrightarrow -1; j > → 0; j --) /*n/2 个排列使得所有数对 (i,j) a[j] =→ 且 *i<j* 都出现一次 → a[j *n 为奇数则首尾相连, 偶数不连 \hookrightarrow 1]; #include <bits/stdc++.h> a[0] = t;} using namespace std; } typedef vector<int> vi; void get_odd(int n, vi ans[]) void get_even(int n, vi ans[]) get_even(n - 1, ans); ← { for (int i = 1; i <= n</pre> vi a(n); → / 2; i ++) { for (int i = 0; i < n;</pre> for (int j = n)→ i ++) \rightarrow -1; j > a[i] = i + 1;→ 0; j --) for (int i = 1; i <= n</pre> ans[i][j] → / 2; i ++) { ans[i].resize(n ans[i][j \rightarrow + 1); for (int j = 1] \rightarrow 0; j < n / \rightarrow 2; j ++) \hookrightarrow 1; ans[i][j ans[i][0] = 1;} 2] } a[j] ans[i]t main() { ans[i]tj vi ans[2019]; int n; cin >> n; if (n & 1) get_odd(n, → ans); 1] else get_even(n, → ans); a[n return 0; 1 j];

```
3.2 rec-free
                                                           c[k * M +
                                                            \rightarrow i][a[i][j]]
/* 不存在四个 1 构成一个矩形,
                                                            \hookrightarrow = 1;
→ 并使得 1 尽量多,输出 01 矩
                                            }
→ 阵 */
                                       }
#include <bits/stdc++.h>
                                        int main() {
using namespace std;
                                                 make();
                                                 return 0;
const int N = 200, n = 150, M
\rightarrow = 13;
//M 为质数, N>M*M>n
int a[N][N], b[N][N], c[N][N];
void make() {
         for (int i = 1; i <=
          \hookrightarrow M; i ++)
          for (int j = 1; j <=</pre>
          \hookrightarrow M; j ++)
              a[i][j + 1] = M *
              \hookrightarrow (j - 1) + i;
     for (int i = 1; i <= M; i</pre>
     → ++)
          for (int j = 1; j <=</pre>
          \rightarrow M; j ++)
              c[i][a[i][j]] = 1;
     for (int k = 1; k < M; k

→ ++) {
         memcpy(b, a, sizeof
          → b);
         for (int i = 1; i <=</pre>
          \hookrightarrow M; i ++)
              for (int j = 1; j
               \hookrightarrow <= M; j ++)
                   a[i][j] = (b[i
                    \hookrightarrow + j - 1 -
                    \rightarrow ((i + j -
                    \hookrightarrow 1) > M ? M
                    → : 0)][j]);
          for (int i = 1; i <=
          \hookrightarrow M; i ++)
              for (int j = 1; j
               \hookrightarrow <= M; j ++)
```

计算几何 return → fabs(y 4.1 最小矩形覆盖含凸包和旋转 卡壳 \rightarrow a.y) < \hookrightarrow eps * 最小矩形覆盖,保留六位小数, ? → 逆时针输出四个顶点坐标 #include <bits/stdc++.h> \hookrightarrow a.x using namespace std; $\hookrightarrow \quad \lambda$ \hookrightarrow < namespace minRectCover { \hookrightarrow a.y; } const int N = 1e5 + 5; const double eps = point operator → 1e-8; → - (const → point &a) struct point{ \hookrightarrow const { double x, y; return point(){} → point(x point(double \hookrightarrow x, double \hookrightarrow a.x, \rightarrow y):x(x), \hookrightarrow y \rightarrow $y(y){}$ \rightarrow a.y); bool operator } $_{\hookrightarrow}$ < (const → point &a) point operator $\hookrightarrow \quad \texttt{const} \ \{$ → + (const → point &a) \hookrightarrow const { return \rightarrow point(x → + \hookrightarrow a.x, $\hookrightarrow \quad \text{y}$ \hookrightarrow + \rightarrow a.y);

}

```
point operator
                                 double
\rightarrow / (const
                                  → operator *
   double &a)
                                     (const
                                     point &a)
   const {
        return
                                     const { //
            point(x
                                     X
            /
                                          return
                                             X
           a,
            У
            /
                                             a.y
           a);
}
                                              У
point operator
                                             a.x;
\rightarrow * (const
                                 }

→ double &a)

    const {
                         }p[N], q[N], rc[4];
        return
                         double sqr(double x)
            point(x
                         \rightarrow {return x * x;}
                         double abs(point a)
                         У
                         \rightarrow a);}
           a);
                         int sgn(double x)
}
                         \rightarrow eps ? 0 : (x < 0 ?
double
                         \rightarrow -1 : 1);}
                         point vertical(point
\hookrightarrow operator /
    (const
                         \rightarrow a, point b)
   point &a)
                         const { //
                         \hookrightarrow + a.y - b.y, a.y -
                         \rightarrow a.x + b.x) -
                         → a;}//与 ab 向量垂
        return
                         → 直的向量
                         point vec(point
                         → a){return a /
            a.x
                         \rightarrow abs(a);}
            У
                         void convexhull(int n,
                         → point *hull, int
            a.y;
}
                         → &top) {//如果要计
                            算周长需要特判 n==2
```

```
for (int i =
                                            int cnt =
\hookrightarrow 1; i < n;
                                            → 0;//去重
\hookrightarrow i ++)
                                            for (int i =
                                            \rightarrow 1; i < n;
           if
              (p[i]
                                            \hookrightarrow i ++)
           p[0])
                                                            (sgn(p[i].x
                     swap(p[i],
                      \hookrightarrow p[0]);
                                                            p[cnt].x)
sort (p + 1, p
\hookrightarrow + n,
                                                            0
\hookrightarrow [&](point
                                                            \hookrightarrow a, point
                                                           sgn(p[i].y
→ b){
           double
                                                           p[cnt].y)
           \hookrightarrow \quad \text{t}
                                                           !=
                                                           0)
                (a
                                                                 p[++

    cnt]

                p[0])
                                                                     p[i];
                (b
                                           n = cnt + 1;
                                           hull[top = 1]
                p[0]);
                                            \rightarrow = p[0];
           if
                                            for (int i =
                (fabs(t)
                <
                                            \rightarrow 1; i < n;
                eps)
                                            \hookrightarrow i ++) {
                return
                sgn(abs(p[0]
                a)
                abs(p[0]
                b))
                <
                0;
           return
                t
               >
              0;
});
```

```
while
                                                     int 1 = 1, r =
                                                     \rightarrow 1, t = 1;
                          (top
                         >
                                                    double L, R,
                         1
                                                     \hookrightarrow D, H;
                         &&
                                                    for (int i =
                          (hull[top]
                                                     \rightarrow 0; i < n;
                                                     → i ++) {
                          hull[top
                                                               D =
                                                                \rightarrow abs(q[i]
                         1])
                                                                \,\hookrightarrow\, q[\text{i}
                          (p[i]
                                                                \hookrightarrow
                                                                   1]);
                         hull[top])
                                                               while
                                                                    (sgn((q[i
                          eps)
                         top
                                                                    1]
                         --;
                     hull[++
                                                                    q[i])
                     → top]
                                                                   (q[t
                        p[i];
                                                                    +
          }
                                                                    1]
          hull[0] =
           → hull[top];
                                                                    q[i])
}
                                                                    (q[i
void main() {
          int n;
                                                                    1]
          scanf("%d",
           \rightarrow &n);
                                                                    q[i])
          for (int i =
                                                                    *
           \hookrightarrow 0; i < n;
                                                                    (q[t]
           \hookrightarrow i ++)
                     scanf("%lf
                                                                    q[i]))
                     \rightarrow &p[i].x,
                                                                    -1)
                     \rightarrow &p[i].y);
          convexhull(n,
                                                                    (t
           \rightarrow q, n);
                                                                    1)
          double ans =
           → 1e20;
                                                                \hookrightarrow n;
```

whi	while		while	
\hookrightarrow	(sgn((q[i	\hookrightarrow	(sgn((q[i	
\hookrightarrow	+	\hookrightarrow	+	
	1]	\hookrightarrow	1]	
\hookrightarrow	- q[i])	\hookrightarrow	- q[i])	
	/	$\overset{\hookrightarrow}{\hookrightarrow}$		
	(q[r		(q[l	
	+	\hookrightarrow		
	1]	\hookrightarrow		
\hookrightarrow \hookrightarrow	- q[i])	\hookrightarrow \hookrightarrow	- q[i])	
\hookrightarrow	_	\hookrightarrow	_	
	(q[i	\hookrightarrow	(q[i	
	+	\hookrightarrow		
	1]	\hookrightarrow	1]	
\hookrightarrow	- q[i])	$\overset{\hookrightarrow}{\hookrightarrow}$	- q[i])	
	/	\rightarrow \rightarrow		
	(q[r]		(q[1]	
\hookrightarrow		\hookrightarrow	-	
	q[i]))	\hookrightarrow	_	
	> -1)	$\overset{\hookrightarrow}{\hookrightarrow}$		
\rightarrow \rightarrow		\rightarrow \rightarrow	_	
\hookrightarrow		\hookrightarrow		
	(r	\hookrightarrow	(1	
	+	\hookrightarrow		
	1)	\hookrightarrow		
	%	\hookrightarrow		
\hookrightarrow if	n; (i		n;	
11	==	ь – ⇔	fabs((q[i	
\hookrightarrow	0)	\hookrightarrow	+	
\hookrightarrow	1	\hookrightarrow	1]	
\hookrightarrow	=	\hookrightarrow	-	
\hookrightarrow	r;	\hookrightarrow	q[i])	
		\hookrightarrow	/ (a[]]	
		\hookrightarrow \hookrightarrow	(q[1] -	
		\rightarrow	q[i])	
		\hookrightarrow	/	
		\hookrightarrow	D);	

```
rc[0]
R =
    fabs((q[i
                                               q[i]
    1]
                                               (q[i
    q[i])
                                               1]
    (q[r]
                                               q[i])
    q[i])
                                               (R
    D);
                                               /
H =
                                               D);//右
    fabs((q[i
                                               下
                                           rc[1]
    1]
                                               rc[0]
   q[i])
                                               vec(vertical(q[
    (q[t]
                                               q[i
    q[i])
                                               1]))
    D);
                                               H;//右
                                               上
double
                                           rc[2]
    tmp
    (R
                                               rc[1]
    +
    L)
                                               (rc[0]
    Н;
                                               q[i])
if
    (tmp
                                               ((R
    <
    ans)
                                               L)
    {
                                               abs(q[i]
        ans
                                               rc[0]));//左
            tmp;
                                               上
```

rc[3]

```
rc[2]
                                               (rc[1]
                                              rc[0]);
                                }
                     }
                     printf("%.6f\n",
                      \hookrightarrow ans);
                     int fir = 0;
                     for (int i =
                      \rightarrow 1; i < 4;
                      \hookrightarrow i ++)
                                if
                                    (rc[i]
                                     rc[fir])
                                           fir
                                           \hookrightarrow i;
                     for (int i =
                      \hookrightarrow 0; i < 4;
                      \hookrightarrow i ++)
                                printf("%.6f
                                 \rightarrow %.6f\n",
                                     rc[(fir
                                    i)
                                    %
                                    4].x,
                                    rc[(fir
                                    i)
                                 \rightarrow 4].y);
          }
}
int main() {
          minRectCover::main();
          return 0;
```

其他 5 for (p = g[u];p; p = p5.1 数字哈希 -> nxt) if (p -> v == namespace my_hash { v) return const int $N = (1 \ll$ → 19) - 1;//散列大小, return 0; → 一定要取 2~k-1, 不 } → 超内存的情况下, N → 越大碰撞越少 void init() {T ++, cur → = pool;}//应对多组 struct E { → 数据使用 int v; } E *nxt; *g[N + 1], pool[N],海岛分金币 5.2 *cur = pool, *p; 海岛分金币 1 5.2.1int vis[N + 1], T; /* 非朴素模型,有额外条件: void ins(int v) { 每个人做决定时如果有多种方案可 int u = v & N;→ 以使自己获得最大收益 if (vis[u] <</pre> 那么他会让决策顺序靠前的人获得 \rightarrow T) vis[u] → 的收益尽可能的大! \hookrightarrow = T, g[u] solution: \rightarrow = NULL; 贪心模拟 for (p = g[u];*/ \rightarrow p; p = p #include <bits/stdc++.h> \rightarrow -> nxt) if \hookrightarrow (p -> v ==#define v first → v) return; #define id second p = cur ++; p-> v = v;using namespace std; \hookrightarrow p -> nxt = g[u]; g[u] typedef pair<int, int> pr; \rightarrow = p; } const int N = 1010; int ask(int v) { int a[N][N]; int u = v & N;if (vis[u] <</pre> → T) return pr b[N]; \rightarrow 0; int n, m;

```
int main() {
                                    printf("%d ",
   cin >> n >> m;
                                     \rightarrow a[n][i]);
   a[1][1] = m;
                                 return 0;
   for (int i = 2; i <= n; i
                             }

→ ++) {
       for (int j = 1; j < i;
                             5.2.2 海岛分金币 2

→ j ++)

           b[j] = pr(a[i -
                             海盗分金币朴素模型:
           → 1][j], j);
                             n 个海盗分 m 个金币, 依次做决策,
       sort (b + 1, b + i,
                             → 如果不少于半数的人同意则方
       → 案通过, 否则当前做决策的人会

    y) {return x.v !=

                             → 被淘汰 (收益视为-1), 由下一
       \rightarrow y.v ? (x.v < y.v)
                             → 人做出决策
          : (x.id >
                             如果一个海盗有多种方案均为最大
       → y.id);});
                             → 收益, 那么他会希望淘汰的人越
       //按照是否容易满足来排
                             → 多越好
       → 序,因为容易满足的
                             求出第 x 个做决策的海盗的最大可
       → 人消耗掉的金币比较
                             → 能受益和最小可能收益
       → 少,也就使得当前的
                             */
       → 人获利最大
                             #include <bits/stdc++.h>
       int s = m, nd = (i -
       \rightarrow 1) / 2;
                             using namespace std;
       for (int j = 1; j < i
       struct node {
          nd --;
                                 int min_v, max_v;
           s -=
                                 node():min_v(0), max_v(0)
           \rightarrow (a[i][b[j].id]
                                 → {}
           \rightarrow = a[i -
                                 node(int min_v, int
           \rightarrow 1][b[j].id] +

    max_v):min_v(min_v),

→ 1);
                                 \rightarrow max_v(max_v) {}
                             };
       if (s < 0) {
          for (int j = 1; j
                             node ask(int n, int m, int x)
           → < i; j ++)</pre>
                             → {//n 个人分 m 个金币, 第 x
              a[i][j] = a[i]
                              → 个做决策的人最少/最多分到多
              → - 1][j];
                             → 少个金币
           a[i][i] = -1;
                                 int y = n + 1 - x;
       }
                                 if (n >= (m + 2) * 2) {
       else {
                                    int a = (m + 1) * 2, b
           a[i][i] = s;
                                     \rightarrow = 2, c = 4;
                                    //前 a 个为 [0,1], 后 b
                                     → 个为 [0,0],将持续
   for (int i = n; i; i --)

→ c ↑
```

```
while (a + b + c \le n)
                                             else return node(m
                                             \rightarrow - (n - 1) / 2,
     → {
        a += b;
                                             \hookrightarrow m - (n - 1) /
        b *= 2;

→ 2);
        c *= 2;
                                        }
                                        else return node(0,
    if (y <= a) return
                                         \rightarrow 0);
    \rightarrow node(0, 1);
                                   }
    else if (y \le a + b)
                               }
    \rightarrow return node(0, 0);
    else return node(-1,
                               int main() {
     \rightarrow -1);
}
                                    → ios::sync_with_stdio(false);
else if (n == m * 2 + 3) {
                                    int x, n, m, k; node y;
    if (x == 1) return
                                    cin >> n >> m >> k;
    \rightarrow node(-1, -1);
                                    while (k --) {
    else if (y \le m * 2 \&\&
                                        cin >> x;
     \rightarrow y % 2 == 1 | | x ==
                                        y = ask(n, m, x);
     \rightarrow 2) return node(0,
                                        printf("%d %d\n",
     \rightarrow 0);

    y.min_v, y.max_v);

    else return node(0,
                                    }
     \rightarrow 1);
                                   return 0;
}
                               }
else if (n == m * 2 + 2) {
    if (y <= m * 2 && y %
                               /*
    \rightarrow 2 == 1 | | x == 1)
                               m = 5
    \rightarrow return node(0, 0);
    else return node(0,
                               1 5
                               2 0 5
     \rightarrow 1);
}
                               3 1 0 4
else if (n == m * 2 + 1) {
                               4 0 1 0 4
    if (y <= m * 2 && y %
                               5 1 0 1 0 3

→ 2 == 1) return

                               6 0 1 0 1 0 3
    \rightarrow node(1, 1);
                               7 1 0 1 0 1 0 2
    else return node(0,
                               8 0 1 0 1 0 1 0 2
     → 0);
                               9 1 0 1 0 1 0 1 0 1
}
                               10 0 1 0 1 0 1 0 1 0 1
                               11 1 0 1 0 1 0 1 0 1 0 0
else {
    if (x & 1) {
                               12 0 _ 0 _ 0 _ 0 _ 0 _ 0
         if (x != 1) return 13 0 _ 0 _ 0 _ 0 _ 0 _ 0 -1
         \rightarrow node(1, 1);
                               14 _ _ _ _ 0 0
                               15 _ _ _ _ 0 0
```

```
16 _ _ _ _ 0 0
                                           if (S == T) {
T = fread(buf, 1,

→ SIZE, stdin) +

17 _ _ _ _ 0 0
\hookrightarrow (S = buf);
18 _ _ _ _ _
                                               if (S == T) return
→ 0 0 0 0
                                                \rightarrow 0 0 0 0 -1
                                           return *S ++;
\rightarrow 0 0 0 0 -1 -1
                                       inline void in(int &x)
\rightarrow 0 0 0 0 -1 -1 -1
                                       \hookrightarrow {//for int
                                           static int ch;

→ 0 0 0 0 −1 −1 −1 −1
                                           while (ch =

    get_char(), ch <</pre>
\rightarrow 0 0 0 0 -1 -1 -1 -1 -1
                                           \rightarrow 48);x = ch ^{\circ} 48;
                                           while (ch =
\rightarrow 0 0 0 0 -1 -1 -1 -1 -1 -1

    get_char(), ch >

                                           \rightarrow 47) x = x * 10 +
\rightarrow 0 0 0 0 -1 -1 -1 -1 -1 -1
                                           \rightarrow (ch ^{^{\circ}} 48);
}
   _ _ _ 0 0 0 0 0 0 0 0
                                       char buffer[SIZE];
                                       char *s = buffer;
                                       void flush() {//最后需要
     根号枚举
5.3

    flush!!

for (int i = 1, last; i <= n;</pre>
                                           fwrite(buffer, 1, s -
\rightarrow i = last + 1) {
                                           \hookrightarrow buffer, stdout);
        last = n / (n / i);
                                           s = buffer;
        //当前枚举区间为 [i,
                                           fflush(stdout);
         \rightarrow last]
                                       }
}
                                       inline void print(const
5.4 读入输出外挂
                                       → char ch) {
                                           if(s - buffer > SIZE -
namespace IO {//only for

→ 2) flush();
\hookrightarrow int!!!
                                           *s++ = ch;
    static const int SIZE = 1
                                       }
    inline void print(char
    inline int get_char() {
                                       → *str) {//for string
        static char *S, *T =
                                           while(*str != 0)

→ S, buf[SIZE];
```

```
print(char(*str
                                         lp, lq = lp + rp * cnt, lq
              → ++));

→ + rq * cnt

    }
                                        else:
                                         1, r, mid, cnt = 1, (inf -
    inline void print(int x) {
                                          \rightarrow rq) // lq + 1, -1, -1
         static char buf[25];
                                         while 1 <= r:
                                          mid = 1 + r >> 1
         static char *p = buf;
         if (x < 0)
                                          if (rp + lp * mid) * inff
                                           \rightarrow > (rq + lq * mid) *
         \rightarrow print('-'), x =
                                           \hookrightarrow n:
         \hookrightarrow -x;
         if (x == 0)
                                           cnt, l = mid, mid + 1
         → print('0');
                                           else:
         while(x) *(++ p) = x \%
                                           r = mid - 1
         \rightarrow 10, x /= 10;
                                         rp, rq = rp + lp * cnt, rq
         while(p != buf)
                                          \rightarrow + lq * cnt
                                       if lq <= inf: print(lp, lq)</pre>

→ print(char(*(p --)))

         → ^ 48));
                                       else: print(rp, rq)
    }
};
```

5.5 给定小数化成分数

```
# 本题答案的分母不超过 1e9, 给
→ 定小数的小数点位为 18 位
inf, inff = 10 ** 9, 10 ** 18
for i in range(int(input())):
n = int(input()[2:])
 if n == 0: print('0 1')
  lp, lq, rp, rq = 0, 1, 1, 1
  while max(lq, rq) <= inf:</pre>
   mp, mq = lp + rp, lq + rq
   if mp * inff <= mq * n:</pre>
    1, r, mid, cnt = 1, (inf -
    \rightarrow lq) // rq + 1, -1, -1
    while 1 <= r:
     mid = 1 + r >> 1
     if (lp + rp * mid) * inff
     \rightarrow <= (lq + rq * mid) *
     \hookrightarrow n:
      cnt, 1 = mid, mid + 1
     else:
      r = mid - 1
```