

## **Algorithms for Scientific Computing**

Hierarchical Methods and Sparse Grids

Algorithms and Data Structures for Sparse Grids –

Michael Bader Technical University of Munich

Summer 2022





#### Part I

# Algorithms vs. Data Structures

- Consider typical sparse grid algorithms, such as: hierarchization/dehierarchization, integration, classification, data mining, solution of PDE, ...
- Important: adaptive representation
- Algorithms depend on data structure:
  - Efficient traversal of sparse grid necessary
  - Thus, we deal with data structures for sparse grids, too

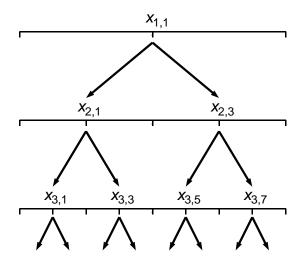


### Data Structures (d = 1)

- How to store function  $u:[0,1] \to \mathbb{R}$  in hierarchical representation (i.e. surplusses  $v_{77}$ )?
- Simplest choice: array → does not allow adaptivity
- Order and store grid points and associated values in binary tree
  - Root is node  $x_{1,1} = 1/2$
  - Children of node x<sub>l,i</sub> are if existent the grid points x<sub>l+1,2i-1</sub> and x<sub>l+1,2i+1</sub> of level l + 1
  - Alternative point of view if child does not exist:
     Complete subtree of binary tree starting from child with all surplusses set to 0



## Data Structures (d = 1) (2)





### Typical Algorithms (d = 1)

#### Hierarchization and Dehierarchization

- Prototype for typical algorithm (c.f. tutorials)
- Our data structure has to allow
  - 1. Iteration over all grid points, considering the hierarchical relations
    - E.g. for hierarchization: first handle all grid points in the support of  $\phi_{l,i}$ , then compute  $v_{l,i}$
  - **2.** Access to *hierarchical neighbors*: grid points at interval boundaries of support of  $\phi_{l,i}$  (if possible exception for points 0 and 1 as not in the tree), e.g. to compute

$$v_{l,i} = u_{l,i} - \frac{1}{2}(u_l + u_r).$$



## Typical Algorithms (d = 1) (2)

Hierarchical neighbors are easy to find geometrically

$$X_{l,i-1}, X_{l,i+1}$$

- But have even indices ⇒ really are on another level (< I)</li>
- Thus, in the binary tree structure:
  - Can be found on way from root to node
  - One of the two indices is the parent node
- For hierarchization/dehierarchization: pass hierarchical neighbors as additional parameters
- Developing algorithms:
  - Try to store all information to process one node at the node and its hierarchical neighbors
  - Access to other nodes may be expensive (esp. for trees)
  - Note: complexity of a tree traversal with "supply of hierarchical neighbors" is at most linear in number of nodes

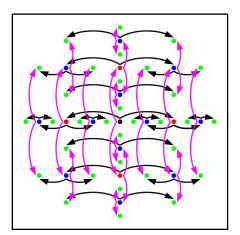


### Data Structures and Typical Algorithms (d > 1)

- What data structure to use in more than one dimension?
- Algorithmically: use construction of basis functions as product of 1D hat functions. Ideally:
  - Use a loop 1,..., d over the dimensions
  - Apply 1d algorithm on one-dimensional structures in each dimension (see also worksheet 7)
- ⇒ Need access to hierarchical neighbors in each spacial direction; implies to create binary tree structure in each dimension
  - Disadvantages:
    - Storage requirements (2*d* pointers)
    - High effort to keep structure consistent when inserting or deleting points



# Data Structures and Typical Algorithms (d > 1) (2)



If you watch closely, you recognize separate binary tree structures for rows (black) and columns (magenta)



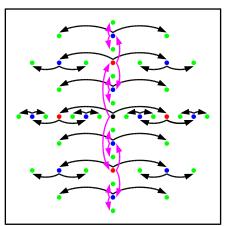
# Data Structures and Typical Algorithms (d > 1) (3)

#### Often better:

- Store in a node only two pointers for one direction (e.g.  $x_1$ )
- A binary tree of nodes is a row (a 1d structure parallel to the  $x_1$  axis)
- For next spacial direction  $x_2$ , only a binary tree in  $x_2$  direction required
- Stores one plane parallel to x<sub>1</sub>-x<sub>2</sub> coordinate plane; nodes are the binary trees with 1d structures
- For each additional spatial direction  $x_d$  build binary tree with (d-1)-dimensional structures as nodes
- Disadvantage: Access to hierarchical neighbors not that easy any more (except for x<sub>1</sub>-direction)
- But can be achieved without much more computational effort by suitable reordering of loops and tree traversals



# Data Structures and Typical Algorithms (d > 1) (4)



Already more clear: One plane (two-dimensional structure) consists of one binary tree (magenta) of which the nodes are binary trees (black) for each row



# Data Structures and Typical Algorithms (d > 1) (5)

#### Hash table

- Much more comfortable (and not too inefficient) alternative
- Store coefficients as target values, with, e.g.,  $(\vec{l}, \vec{i})$  as keys
- No need to care about tree structures
- Only requires computation of indices of accessed nodes (hierarchical neighbor, ...)
- ⇒ Best solution for your own sparse grid experiments

#### Further assumptions on data structures

- Algorithms will assume that all hierarchical neighbors exist for each grid point
- ⇒ If creating grid points adaptively, create them if necessary
  - No further assumptions



## **Data Structures for Regular Sparse Grids**

#### **Array-Based Data Structures**

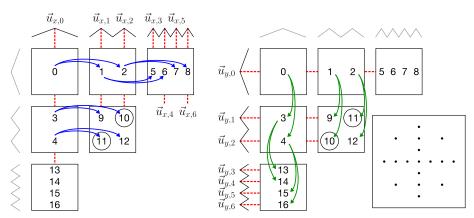
- Cartesian meshes with  $2^{l_1} \times 2^{l_2} \times ... \times 2^{l_d}$  grid points
- suggests classical array indexing similar to  $i \cdot n + j$ 
  - → question: what is the "fastest-running" index?
- number of grid points per subspace is constant along "diagonals", i.e., for constant |I|1
  - → sequentialized storage scheme for subspaces
  - → start of each subgrid can be easily computed
  - → index offset to hierarchical neighbours may be computed
- additional considerations: best layout for vectorization, parallelization, etc.

#### **Towards Dimensional Adaptivity**

- add or remove entire subspaces/subgrids in an adaptive fashion
- may introduce higher accuracy only in selected dimensions



## **Data Structures for Regular Sparse Grids (2)**



#### Example - Array-Based Data Structures (Buse et al., ISPDC 2012)

 note uniform vs. non-uniform index offset for access to hierarchical parents/neighbours in x- vs. y direction



#### **Summary**

#### **Data Structures**

- array-based for regular sparse grids and combination technique (see tutorials)
- hierarchical adaptivity reflected by tree-based data structures (but: more complicated in higher dimensions)
- hash-based data structures
- dimensional adaptivity allows to stick to array-based data structures

#### **Algorithms**

- hierarchisation and dehierarchisation: tree-based recursion plus "hierarchical neighbours"
- archimedes quadrature → recursion on dimensions
- much more complicated algorithms, if we want to use sparse grids for solution of partial differential equations



### Part II

# **Classification with Sparse Grids**



#### **Recall: Classification in 1D**

Given: training set (normalized)

$$S := \{(\vec{x}_i, y_i) \in [0, 1] \times \{+1, -1\}\}_{i=1}^m$$

Find approximation f<sub>N</sub>:

$$f_N(x) = \sum_{i=1}^N v_i \phi_i(x)$$

Classical approach: minimize quadratic error

$$\sum_{i=1}^{m} (f_N(x_i) - y_i)^2 \stackrel{!}{=} \min \quad \Leftrightarrow \quad \sum_{i=1}^{m} \left( \sum_{i=1}^{N} v_i \phi_i(x_i) - y_i \right)^2 \stackrel{!}{=} \min$$

• Solution obtained via "least squares":  $G^TGv = G^Ty$ , where  $G_{ij} = \phi_i(x_i)$ 



### **Recall: Least Squares Solution**

minimize quadratic error → find values v<sub>i</sub> that minimize

$$\sum_{i=1}^{m} \left( \sum_{i=1}^{N} v_{i} \phi_{j}(x_{i}) - y_{i} \right)^{2} \quad \text{or} \quad \sum_{i=1}^{m} \left( \sum_{i=1}^{N} G_{ij} v_{j} - y_{i} \right)^{2} \quad \text{with } G_{ij} := \phi_{j}(x_{i})$$

• approach: set all partial derivatives  $\frac{\partial}{\partial v_i}$  to zero

$$\frac{\partial}{\partial v_{k}} \left( \sum_{i=1}^{m} \left( \sum_{j=1}^{N} G_{ij} v_{j} - y_{i} \right)^{2} \right) = \sum_{i=1}^{m} \frac{\partial}{\partial v_{k}} \left( \sum_{j=1}^{N} G_{ij} v_{j} - y_{i} \right)^{2} = 0$$

$$\Leftrightarrow \sum_{i=1}^{m} 2 \left( \sum_{j=1}^{N} G_{ij} v_{j} - y_{i} \right) G_{ik} = 2 \sum_{i=1}^{m} \left( \sum_{j=1}^{N} G_{ik} G_{ij} v_{j} - G_{ik} y_{i} \right) = 0$$

$$\Leftrightarrow \sum_{i=1}^{m} \sum_{j=1}^{N} G_{ik} G_{ij} v_{j} - \sum_{i=1}^{m} G_{ik} y_{i} = \sum_{i=1}^{m} G_{ik} \sum_{j=1}^{N} G_{ij} v_{j} - \sum_{i=1}^{m} G_{ik} y_{i} = 0$$

$$\underbrace{ = (G^{T}G_{V})_{k}}_{=(G^{T}G_{V})_{k}}$$



## **Classification and Regularization**

- Possible problem: "overfitting" include penalty term to minimize gradient (or similar property) of f<sub>N</sub> to avoid oscillations due to noise in training data
- Solve regularized least squares problem

$$f_N \stackrel{!}{=} \underset{f_N \in V_N}{\arg \min} \left( \frac{1}{m} \sum_{i=1}^m (y_i - f_N(\vec{x}_i))^2 + \lambda \|\nabla f_N\|_{L_2}^2 \right)$$

with 
$$\|g\|_{L_2}^2 := \int g^2 \, d\vec{x}$$

- minimize guadratical error and prevent overfitting
  - $\rightarrow$  Parameter  $\lambda$  to control trade-off



### Classification and Regularization (2)

How to minimize  $\lambda \|\nabla f_N\|_{L_2}^2$ :

• Piecewise linear function  $f_N$ :

$$\nabla f_N(x) = f'_N(x) = \sum_{i=1}^N v_i \phi'_i(x) \quad \Rightarrow \quad \|\nabla f_N\|_{L_2}^2 = \int \left(\sum_{i=1}^N v_i \phi'_i(x)\right)^2 dx$$

Minimize → set partial derivatives w.r.t. all v<sub>i</sub> to 0:

$$ightsquigarrow rac{\partial}{\partial v_k} ig( \| 
abla f_N \|_{L_2}^2 ig) = \cdots = 2 \sum_j C_{jk} v_j, \quad \text{with} \quad C_{jk} := \int \phi_j'(x) \phi_k'(x), \, \mathrm{d}x$$

Thus: to solve regularized least squares problem

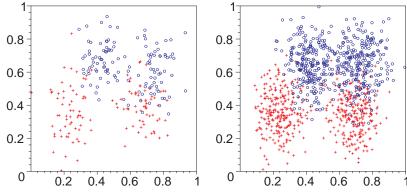
$$f_N \stackrel{!}{=} \underset{f_N \in V_N}{\operatorname{arg \, min}} \left( \frac{1}{m} \sum_{i=1}^m (y_i - f_N(\vec{x}_i))^2 + \lambda \|\nabla f_N\|_{L_2}^2 \right)$$

solve linear system of the following form:  $\frac{1}{m}G^TGv + \lambda Cv = \frac{1}{m}G^Ty$ 



### Example 1 – Ripley Data Set

- Artificial, 2d data set, frequently used as a benchmark (mixture of Gaussian distributions plus noise)
- 250 points for training, 1000 to test on

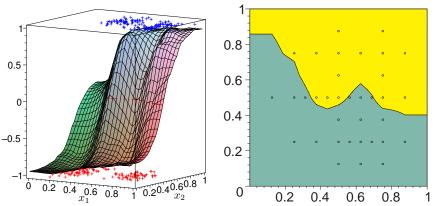


Constructed to contain 8% of noise



### **Ripley Data Set Using Sparse Grids**

Compute adaptive sparse grid classifier, e.g.:



- Best accuracy: 91.5% on test data (max. 92%)
- Suitable treatment of boundary needed



## From Minimization to System of Linear Equations

• d-dim. problem; find function  $f_N(\vec{x}) = \sum_{\vec{l}} \sum_{\vec{i}} v_{\vec{l},\vec{i}} \phi_{\vec{l},\vec{i}}(\vec{x})$  such that

$$f_N \stackrel{!}{=} \underset{f_N \in V_N}{\operatorname{arg \, min}} \left( \frac{1}{m} \sum_{i=1}^m (y_i - f_N(\vec{x}_i))^2 + \lambda \|\nabla f_N\|_{L_2}^2 \right)$$

Again leads to N linear equations for N unknowns (m data points)

$$\left(\frac{1}{m}G^{T}G + \lambda C\right)\vec{v} = \frac{1}{m}G^{T}\vec{y},$$

#### **Questions when using Sparse Grids:**

- How do the matrices G and C look like?
- Should we explicitly set up G and C or is there a better solution?
- In 1D: C is a diagonal matrix! (However: G is complicated)
- In general: level-wise hierarchical/recursive algorithm!



#### Recall: Structure of G for Hierarchical Basis

Structure of the matrix *G*, where  $G_{ij} = \phi_j(x_i)$ , where  $\{\phi_1, \phi_2, \phi_3\} = \{\phi_{2,1}, \phi_{1,1}, \phi_{2,3}\}$ :

- again assume  $N_j$  data points per interval  $[\xi_{j-1}, \xi_j]$
- Example structure then (again with 3 basis functions, 4 intervals, 10 data points):

G<sup>T</sup>G no longer tridiagonal → expect a denser matrix



### Towards a Hierachical/Recursive Algorithm

• Consider right-hand side  $\frac{1}{m}G^T\vec{y}$ :

$$(G^T \vec{y})_i = \sum_j G_{ij}^T y_j = \sum_j \phi_i(\vec{x}_j) y_j$$

(note that we switch to 1D numbering of basis functions  $\rightsquigarrow \psi_i$ )

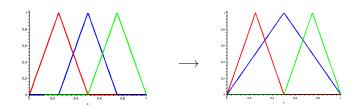
- Consider a nodal basis  $\phi_i(\vec{x}_{n,j}) = \delta_{ij}$ , then  $G^T$  easy to set up and  $(G^T \vec{y})_i$  easy to compute
- Hierarchical transform when using hierarchical basis  $\psi_i(x)$ ?

#### Approach:

- Consider vectors of basis functions  $\vec{\psi} = (\psi_i)_i$  and  $\vec{\phi} = (\phi)_i$
- Show that then  $\psi = H\phi$  (matrix-vector product)
- Then:  $\sum_i \psi_i(x_j) y_j = \sum_i (H\phi)_i(x_j) y_j = (HG^T \vec{y})_i$
- Do not set up matrix  $HG^T \rightarrow perform$  as two matrix-vector products
- Now: how does H look like and how do we compute  $H\vec{y}$ ?



#### **Recall: Hierarchical Basis Transformation**



• represent "wider" hat function  $\phi_{1,1}(x)$  via basis functions  $\phi_{2,j}(x)$ 

$$\phi_{1,1}(x) = \frac{1}{2}\phi_{2,1}(x) + \phi_{2,2}(x) + \frac{1}{2}\phi_{2,3}(x)$$

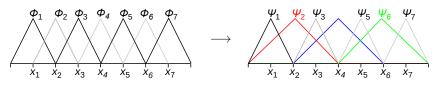
 consider vector of hierarchical/nodal basis functions and write transformation as matrix-vector product:

$$\begin{pmatrix} \psi_{2,1}(x) \\ \psi_{2,2}(x) \\ \psi_{2,3}(x) \end{pmatrix} := \begin{pmatrix} \phi_{2,1}(x) \\ \phi_{1,1}(x) \\ \phi_{2,3}(x) \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 \\ \frac{1}{2} & 1 & \frac{1}{2} \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} \phi_{2,1}(x) \\ \phi_{2,2}(x) \\ \phi_{2,3}(x) \end{pmatrix}$$



### Recall: Hierarchical Basis Transformation (2)

#### Consider "semi-hierarchical" transform:



Matrices for change of basis are then:  $(H_3^{(2)})$  to transform to hierarchical basis)

$$H_3^{(1)} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ \frac{1}{2} & 1 & \frac{1}{2} & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{2} & 1 & \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2} & 1 & \frac{1}{2} \\ 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \qquad H_3^{(2)} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{2} & 0 & 1 & 0 & \frac{1}{2} & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$$H_3^{(2)} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{2} & 0 & 1 & 0 & \frac{1}{2} & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$



## Recall: Hierarchical Basis Transformation (3)

#### Level-wise hierarchical transform:

- hierarchical basis transformation:  $\psi_{n,i}(x) = \sum_{i} H_{i,j} \phi_{n,j}(x)$
- written as matrix-vector product:  $\vec{\psi}_n = H_n \vec{\phi}_n$
- $H_n\vec{\phi}_n$  can be performed as a sequence of level-wise transforms:

For k from 1 to n-1 
$$\vec{\phi}_n := H_n^{(k)} \vec{\phi}_n$$

matrix H<sub>n</sub> for hierarchical basis transformation is thus:

$$H_n = H_n^{(n-1)} H_n^{(n-2)} \dots H_n^{(2)} H_n^{(1)}$$

• where each level-wise transform  $H_n^{(k)} \vec{\phi}_n$  has a simple loop implementation:

For j from 
$$2^k$$
 to  $2^n$  step  $2^k$ 

$$\phi_{n,j} := \frac{1}{2}\phi_{n,j-2^{k-1}} + \phi_{n,j} + \frac{1}{2}\phi_{n,j+2^{k-1}}$$



## **Classification with Sparse Grids**

#### **Notes on Implementation**

- in higher dimensions: nodal basis leads to simple matrix structures, but the systems of equations are difficult to solve
  - → most important: curse of dimensionality kicks in
- hierachical basis leads to system of equations that can be solved efficiently;
  - → complicated matrix structures,
  - → algorithms based on hierarchization/dehierarchization
- sparse grids: implementation is not just a hierarchization of node basis
  - → complicated hierarchical, recursive algorithms
  - → mitigates curse of dimensionality!