Chapter 21 Custom Generic Data Structures

Java How to Program, 11/e Questions? E-mail paul.deitel@deitel.com

OBJECTIVES

In this chapter you'll:

- Form linked data structures using references, selfreferential classes, recursion and generics.
- Create and manipulate dynamic data structures, such as linked lists, queues, stacks and binary trees.
- Learn various important applications of linked data structures.
- Create reusable data structures with composition.
- Organize classes in packages to promote reuse.

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21.1 Introduction

- Dynamic data structures grow and shrink at execution time.
- Linked lists are collections of data items "linked up in a chain" insertions and deletions can be made anywhere in a linked list.
- Stacks are important in compilers and operating systems; insertions and deletions are made only at one end of a stack—its top.
- Queues represent waiting lines; insertions are made at the back (also referred to as the tail) of a queue and deletions are made from the front (also referred to as the head).
- Binary trees facilitate high-speed searching and sorting of data, eliminating duplicate data items efficiently, representing file-system directories, compiling expressions into machine language and many other interesting applications.



Software Engineering Observation 21.1

The vast majority of software developers should use the predefined generic collection classes that we discussed in Chapter 16, rather than developing customized linked data structures.

21.2 Self-Referential Classes

- A self-referential class contains an instance variable that refers to another object of the same class type.
- For example, the generic class declaration

```
class Node<T>
{
   private T data;
   private Node<T> nextNode; // reference to next node

   public Node(T data) { /* constructor body */ }
   public void setData(T data) { /* method body */ }
   public T getData() { /* method body */ }
   public void setNext(Node<T> next) { /* method body */ }
   public Node<T> getNext() { /* method body */ }
} // end class Node<T>
```

declares class Node, which has two private instance variables—data (of the generic type T) and Node<T> variable nextNode.

21.2 Self-Referential Classes (cont.)

- Variable nextNode references a Node<T> object, an object of the same class being declared here—hence, the term "self-referential class."
- ▶ Field nextNode is a link—it "links" an object of type Node<T> to another object of the same type.
- Programs can link self-referential objects together to form such useful data structures as lists, queues, stacks and trees.
- Figure 21.1 illustrates two self-referential objects linked together to form a list.



Fig. 21.1 | Self-referential-class objects linked together.

21.3 Dynamic Memory Allocation

- Creating and maintaining dynamic data structures requires dynamic memory allocation—allowing a program to obtain more memory space at execution time to hold new nodes and to release space no longer needed.
- The limit for dynamic memory allocation can be as large as the amount of available physical memory in the computer or the amount of available disk space in a virtual-memory system.
- Often the limits are much smaller, because the computer's available memory must be shared among many applications.

21.4 Linked Lists

- A linked list is a linear collection (i.e., a sequence) of self-referential-class objects, called nodes, connected by reference links—hence, the term "linked" list.
- Typically, a program accesses a linked list via a reference to its first node.
- The program accesses each subsequent node via the link reference stored in the previous node.
- By convention, the link reference in the last node of the list is set to null to indicate "end of list."
- A linked list is appropriate when the number of data elements to be represented in the data structure is unpredictable.
- Linked lists become full only when the system has *insufficient memory* to satisfy dynamic storage allocation requests.



Performance Tip 21.1

After locating the insertion point, insertion into a linked list is fast—only two references have to be modified. All existing node objects remain at their current locations in memory.



Performance Tip 21.2

Insertion and deletion in a sorted array can be time consuming—all the elements following the inserted or deleted element must be shifted appropriately.

21.4.1 Singly Linked Lists

- Linked list nodes normally are not stored contiguously in memory.
- Rather, they are logically contiguous.
- illustrates a linked list with several nodes.
- ▶ This diagram presents a singly linked list—each node contains one reference to the next node in the list.
- Often, linked lists are implemented as doubly linked lists—each node contains a reference to the next node in the list and a reference to the preceding one.



Performance Tip 21.3

The elements of an array are contiguous in memory. This allows immediate access to any array element, because its address can be calculated directly as its offset from the beginning of the array. Linked lists do not afford such immediate access—an element can be accessed only by traversing the list from the front (or the back in a doubly linked list).

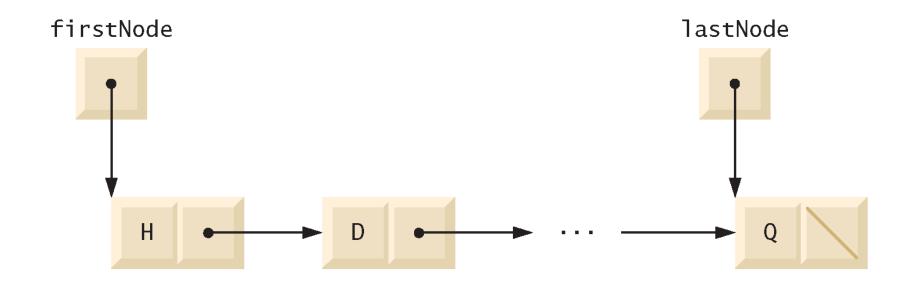


Fig. 21.2 | Linked-list graphical representation.

21.4.2 Implementing a Generic List Class

- The program of Figs. 21.3–21.5 uses an object of our generic List class to manipulate a list of miscellaneous objects.
- ▶ The program consists of four classes:
- ▶ ListNode (Fig. 21.3, lines 6-37),
- ▶ List (Fig. 21.3, lines 40–147),
- ▶ EmptyListException (Fig. 21.4)
- ListTest (Fig. 21.5).

```
// Fig. 21.3: List.java
  // ListNode and List class declarations.
    package com.deitel.datastructures;
    import java.util.NoSuchElementException;
    // class to represent one node in a list
    class ListNode<E> {
       // package access members; List can access these directly
       E data; // data for this node
       ListNode<E> nextNode; // reference to the next node in the list
11
12
13
    // constructor creates a ListNode that refers to object
       ListNode(E object) {this(object, null);}
14
15
```

Fig. 21.3 | ListNode and List class declarations. (Part 1 of 8.)

```
16
       // constr\uctor creates ListNode that refers to the specified
       // object and to the next ListNode
17
18
       ListNode(E object, ListNode<E> node) {
          data = object;
19
20
          nextNode = node;
21
22
23
       // return reference to data in node
24
       E getData() {return data;}
25
       // return reference to next node in list
26
27
       ListNode<E> getNext() {return nextNode;}
28
29
```

Fig. 21.3 | ListNode and List class declarations. (Part 2 of 8.)

```
30
    // class List definition
31
    public class List<E> {
32
       private ListNode<E> firstNode;
33
       private ListNode<E> lastNode;
       private String name; // string like "list" used in printing
34
35
       // constructor creates empty List with "list" as the name
36
37
       public List() {this("list");}
38
39
       // constructor creates an empty List with a name
       public List(String listName) {
40
41
          name = listName;
42
          firstNode = lastNode = null:
43
44
```

Fig. 21.3 | ListNode and List class declarations. (Part 3 of 8.)

```
// insert item at front of List
45
46
       public void insertAtFront(E insertItem) {
          if (isEmpty()) { // firstNode and lastNode refer to same object
47
             firstNode = lastNode = new ListNode<E>(insertItem);
48
49
          else { // firstNode refers to new node
50
51
             firstNode = new ListNode<E>(insertItem, firstNode);
52
53
54
55
       // insert item at end of List
56
       public void insertAtBack(E insertItem) {
          if (isEmpty()) { // firstNode and lastNode refer to same object
57
             firstNode = lastNode = new ListNode<E>(insertItem);
58
59
          else { // lastNode's nextNode refers to new node
60
              lastNode = lastNode.nextNode = new ListNode<E>(insertItem);
61
62
63
```

Fig. 21.3 ListNode and List class declarations. (Part 4 of 8.)

```
64
65
       // remove first node from List
66
       public E removeFromFront() throws NoSuchElementException {
          if (isEmpty()) { // throw exception if List is empty
67
             throw new NoSuchElementException(name + " is empty");
68
69
70
71
          E removedItem = firstNode.data; // retrieve data being removed
72
73
          // update references firstNode and lastNode
          if (firstNode == lastNode) {
74
             firstNode = lastNode = null;
75
76
77
          else {
             firstNode = firstNode.nextNode;
78
79
80
81
          return removedItem; // return removed node data
82
```

Fig. 21.3 | ListNode and List class declarations. (Part 5 of 8.)

```
83
84
       // remove last node from List
       public E removeFromBack() throws NoSuchElementException {
85
86
          if (isEmpty()) { // throw exception if List is empty
87
             throw new NoSuchElementException(name + " is empty");
88
89
90
          E removedItem = lastNode.data; // retrieve data being removed
91
92
          // update references firstNode and lastNode
93
          if (firstNode == lastNode) {
             firstNode = lastNode = null;
94
95
```

Fig. 21.3 | ListNode and List class declarations. (Part 6 of 8.)

```
96
           else { // locate new last node
              ListNode<E> current = firstNode;
97
98
              // loop while current node does not refer to lastNode
99
100
              while (current.nextNode != lastNode) {
101
                 current = current.nextNode;
102
103
104
              lastNode = current; // current is new lastNode
              current.nextNode = null;
105
106
107
108
           return removedItem; // return removed node data
109
110
```

Fig. 21.3 | ListNode and List class declarations. (Part 7 of 8.)

```
// determine whether list is empty; returns true if so
public boolean isEmpty() {return firstNode == null;}
112
113
114
       // output list contents
115
       public void print() {
116
          if (isEmpty()) {
              System.out.printf("Empty %s%n", name);
117
118
              return;
119
120
121
           System.out.printf("The %s is: ", name);
          ListNode<E> current = firstNode;
122
123
          // while not at end of list, output current node's data
124
          while (current != null) {
125
              System.out.printf("%s ", current.data);
126
127
              current = current.nextNode;
128
129
130
           System.out.println();
131
132 }
```

Fig. 21.3 | ListNode and List class declarations. (Part 8 of 8.)

```
// Fig. 21.4: ListTest.java
   // ListTest class to demonstrate List capabilities.
    import com.deitel.datastructures.List;
    import java.util.NoSuchElementException;
    public class ListTest {
       public static void main(String[] args) {
          List<Integer> list = new List<>();
          // insert integers in list
10
          list.insertAtFront(-1);
\mathbf{I}
          list.print();
12
          list.insertAtFront(0);
13
          list.print();
14
          list.insertAtBack(1);
15
16
          list.print();
          list.insertAtBack(5);
          list.print();
18
19
```

Fig. 21.4 ListTest class to demonstrate List capabilities. (Part 1 of 3.)

```
20
          // remove objects from list; print after each removal
21
          try {
              int removedItem = list.removeFromFront();
22
              System.out.printf("%n%d removed%n", removedItem);
23
              list.print();
24
25
              removedItem = list.removeFromFront();
26
              System.out.printf("%n%d removed%n", removedItem);
27
              list.print();
28
29
              removedItem = list.removeFromBack();
30
              System.out.printf("%n%d removed%n", removedItem);
31
              list.print();
32
33
              removedItem = list.removeFromBack();
34
              System.out.printf("%n%d removed%n", removedItem);
35
              list.print();
36
37
38
           catch (NoSuchElementException noSuchElementException) {
              noSuchElementException.printStackTrace();
39
40
41
42
```

Fig. 21.4 | ListTest class to demonstrate List capabilities. (Part 2 of 3.)

```
The list is: -1
The list is: 0-1
The list is: 0 - 11
The list is: 0 -1 1 5
0 removed
The list is: -1 1 5
-1 removed
The list is: 1 5
5 removed
The list is: 1
1 removed
Empty list
```

Fig. 21.4 ListTest class to demonstrate List capabilities. (Part 3 of 3.)

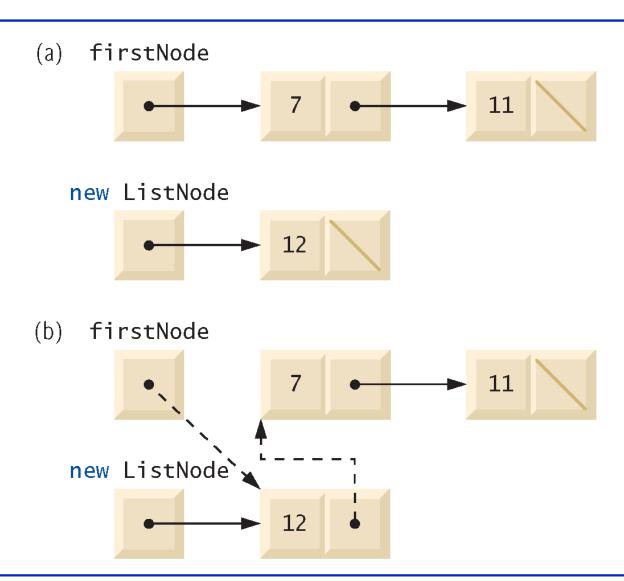


Fig. 21.5 Graphical representation of operation insertAtFront.

21.4.6 List Method insertAtBack

- ▶ Method insertAtBack (lines 69–75 of Fig. 21.3) places a new node at the back of the list.
- The steps are:
 - Call isEmpty to determine whether the list is empty.
 - If the list is empty, assign to firstNode and lastNode the new ListNode that was initialized with insertItem.
 - If the list is not empty, link the new node into the list by assigning to lastNode and lastNode.nextNode the reference to the new ListNode that was initialized with insertItem.

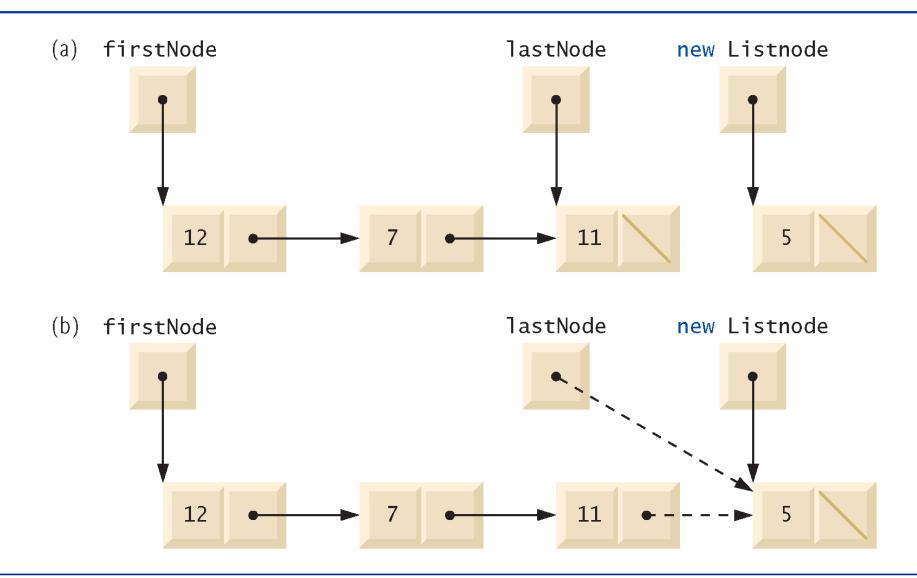


Fig. 21.6 Graphical representation of operation **insertAtBack**.

21.4.7 List Method removeFromFront

- ▶ Method removeFromFront (lines 78–92 of) removes the first node of the list and returns a reference to the removed data.
- The steps are:
 - Assign firstNode.data to removedItem.
 - If firstNode and lastNode refer to the same object, the list has only one element at this time. So, the method sets firstNode and lastNode to null to remove the node from the list (leaving the list empty).
 - If the list has more than one node, then the method leaves reference lastNode as is and assigns the value of firstNode.nextNode to firstNode. Thus, firstNode references the node that was previously the second node in the list.
 - Return the removedItem reference (line 91).

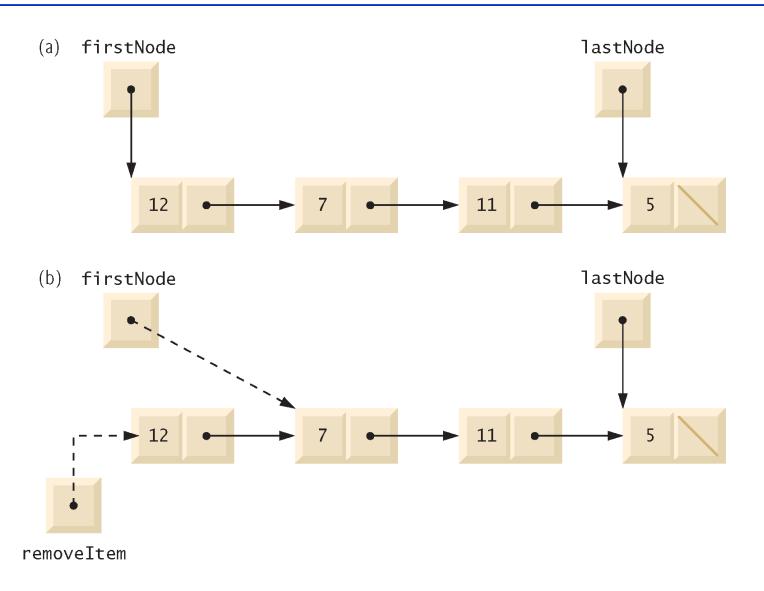


Fig. 21.7 | Graphical representation of operation removeFromFront.

21.4.8 List Method removeFromBack

- ▶ Method removeFromBack (lines 95–118 of Fig. 21.3) removes the last node of a list and returns a reference to the removed data.
- The steps are:
 - Assign lastNode.data to removedItem.
 - If the firstNode and lastNode refer to the same object, the list has only one element at this time. So, set firstNode and lastNode to null to remove that node from the list (leaving the list empty).
 - If the list has more than one node, create the ListNode reference current and assign it firstNode.
 - Now "walk the list" with current until it references the node before the last node. The while loop assigns current.nextNode to current as long as current.nextNode is not lastNode.

21.4.8 List Method removeFromBack

- After locating the second-to-last node, assign current to lastNode to update which node is last in the list.
- Set the current.nextNode to null to remove the last node from the list and terminate the list at the current node.
- Return the removedItem reference.

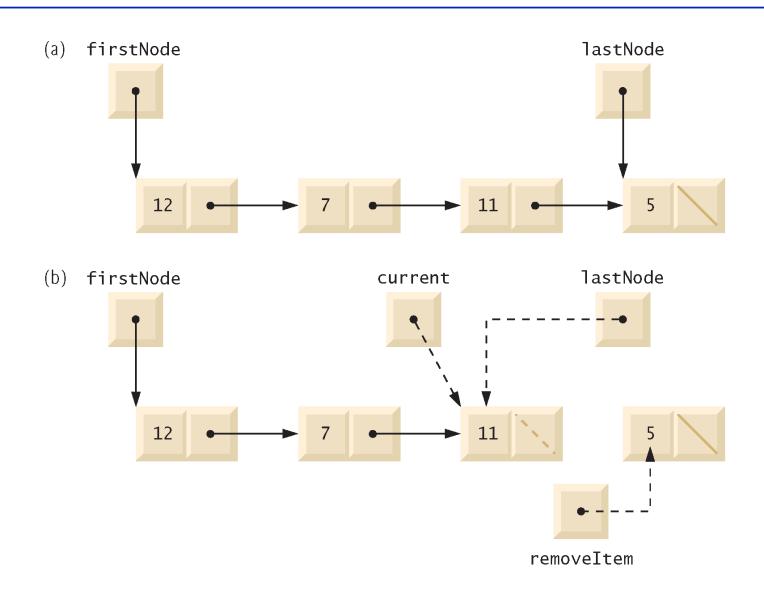


Fig. 21.8 | Graphical representation of operation removeFromBack.

21.4.9 List Method print

- ▶ Method print (lines 127–146 of Fig. 21.3) first determines whether the list is empty (lines 129–133).
- If so, print displays a message indicating that the list is empty and returns control to the calling method.
- Otherwise, print outputs the list's data. Line 136 creates ListNode current and initializes it with firstNode.

21.4.10 Creating Your Own Packages

- The Java API types (classes, interfaces and enums) are organized in *packages* that group related types.
- Packages facilitate software reuse by enabling programs to import existing classes, rather than *copying* them into the folders of each program that uses them.
- Programmers use packages to organize program components, especially in large programs.
- Packages help you specify unique names for every type you declare, which (as we'll discuss) helps prevent class-name conflicts.

Steps for Declaring a Reusable Class

- Before a class can be imported into multiple programs, it must be placed in a package to make it reusable.
 - Declare one or more public types (classes, interfaces and enums). Only public types can be reused outside the package in which they're declared.
 - Choose a unique package name and add a package declaration to the source-code file for each reusable type that should be part of the package.
 - Compile the types so that they're placed in the appropriate package directory.
 - Import the reusable types into a program and use them.

Step 1: Creating public Types for Reuse

- For *Step 1*, you declare the types that will be placed in the package, including both the reusable types and any supporting types.
- In Fig. 21.3, class List is public, so it's reusable outside its package. Class ListNode, however, is not public, so it can be used only by class List and any other types declared in the same package.
- ▶ In Fig. 21.4, class EmptyListException is public, so it, too, is reusable.
- If a source-code file contains more than one type, all types in the file are placed in the same package when the file is compiled.

Step 2: Adding the package Statements

- For Step 2, you provide a package declaration containing the package's name.
- All source-code files containing types that should be part of the same package must contain the same package declaration.
- Figures 21.3 and 21.4 each contain:
 - package com.deitel.datastructures;
- indicating that all the types declared in those files—ListNode and List in Fig. 21.3 and EmptyListException in Fig. 21.4—are part of the com.deitel.datastructures package.

Package Naming Conventions

- A package name's parts are separated by dots (.), and there typically are two or more parts.
- To ensure unique package names, you typically begin the name with your institution's or company's Internet domain name in reverse order—e.g., our domain name is deitel.com, so we begin our package names with com.deitel.
- After the reversed domain name, you can specify additional parts in a package name.
- We chose datastructures as the next part in our package name to indicate that classes ListNode, List and EmptyListException are from this data structures chapter.

Fully Qualified Names

- The package name is part of the fully qualified type name, so the name of class List is actually com.deitel.datastructures.List.
- You can use this fully qualified name in your programs, or you can import the class and use its simple name (the class name by itself—List) in the program.
- If another package also contains a List class, the fully qualified class names can be used to distinguish between the classes in the program and prevent a name conflict (also called a name collision).

Step 3: Compiling Packaged Types

- When a Java file containing a package declaration is compiled, the resulting class file is placed in a directory specified by the declaration.
- Classes in the package com.deitel.datastructures are placed in the directory
 - com
 - deitel
 - datastructures
- The names in the package declaration specify the exact location of the package's classes.

- The javac command-line option -d causes the compiler to create the directories based on the package declaration.
- The option also specifies where the top-level directory in the package name should be placed on your system—you may specify a relative or complete path to this location. For example, the command
 - javac -d . List.java EmptyListException.java
- specifies that the first directory in our package name (com) should be placed in the current directory.
 - The period (.) after -d in the preceding command represents the current directory on the Windows, UNIX, Linux and Mac OS X operating systems (and several others as well).
- Similarly, the command
 - javac -d .. List.java EmptyListException.java
- specifies that the first directory in our package name (com) should be placed in the parent directory.

Step 4: Importing Types from Your Package

- Once types are compiled into a package, they can be imported (Step 4).
- Class ListTest (Fig. 21.5) is in the default package because its .java file does not contain a package declaration.
- Because class ListTest is in a different package from List and EmptyListException, you must either import these classes so that class ListTest can use them (lines 3-4 of Fig. 21.5) or you must fully qualify the names List and EmptyListException everywhere they're used throughout class ListTest.
- For example, line 10 of Fig. 21.5 could have been written as:
 com.deitel.datastructures.List<Integer> list =
 new com.deitel.datastructures.List<>();

Single-Type-Import vs. Type-Import-On-Demand Declarations

- Lines 3–4 of Fig. 21.5 are single-type-import declarations—they each specify one class to import. When a source-code file uses multiple classes from a package, you can import those classes with a a type-import-on-demand declaration of the form
 - import packagename.*;
- which uses an asterisk (*) at its end to inform the compiler that all public classes from the packagename package can be used in the file containing the import.
- Only those classes that are used are loaded at execution time.



Common Programming Error 21.1

Using the import declaration import java.*; causes a compilation error. You must specify the full package name from which you want to import classes.



Error-Prevention Tip 21.1

Using single-type-import declarations helps avoid naming conflicts by importing only the types you actually use in your code.

Specifying the Classpath When Compiling a Program

- When compiling ListTest, javac must locate the .class files for classes List and EmptyListException to ensure that class ListTest uses them correctly.
- The compiler uses a special object called a class loader to locate the classes it needs.
 - The class loader begins by searching the standard Java classes that are bundled with the JDK.
 - Then it searches for optional packages. Java provides an extension mechanism that enables new (optional) packages to be added to Java for development and execution purposes.
 - If the class is not found in the standard Java classes or in the extension classes, the class loader searches the classpath—a list of directories or archive files containing reusable types.
 - Each directory or archive file is separated from the next by a directory separator—a semicolon (;) on Windows or a colon (:) on UNIX/Linux/Mac OS X.

- By default, the classpath consists only of the current directory. However, the classpath can be modified by
 - providing the -classpath listOfDirectories option to the javac compiler or
 - setting the CLASSPATH environment variable (a special variable that you define and the operating system maintains so that programs can search for classes in the specified locations).
- If you compile ListTest.java without specifying the -classpath option, as in
 - javac ListTest.java
- the class loader assumes that the additional package(s) used by the ListTest program are in the current directory.

- To compile ListTest.java, use the command
 - javac -classpath .;.. ListTest.java
- on Windows or the command
 - javac -classpath .:.. ListTest.java
- on UNIX/Linux/Mac OS X. The . in the classpath enables the class loader to locate ListTest in the current directory.
- The .. enables the class loader to locate the contents of package com.deitel.datastructures in the parent directory.



Common Programming Error 21.2

Specifying an explicit classpath eliminates the current directory from the classpath. This prevents classes in the current directory (including packages in the current directory) from loading properly. If classes must be loaded from the current directory, include a dot (.) in the classpath to specify the current directory.



Software Engineering Observation 21.2

In general, it's a better practice to use the - classpath option of the compiler, rather than the CLASSPATH environment variable, to specify the classpath for a program. This enables each program to have its own classpath.



Error-Prevention Tip 21.2

Specifying the classpath with the CLASSPATH environment variable can cause subtle and difficult-to-locate errors in programs that use different versions of the same package.

Specifying the Classpath When Executing a Program

- When you execute a program, the JVM must be able to locate the .class files for the program's classes.
- Like the compiler, the java command uses a *class loader* that searches the standard classes and extension classes first, then searches the classpath (the current directory by default).
- The classpath can be specified explicitly by using the same techniques discussed for the compiler.
- As with the compiler, it's better to specify an individual program's classpath via command-line JVM options.
- You can specify the classpath in the java command via the -classpath or -cp command-line options, followed by a list of directories or archive files.

- Again, if classes must be loaded from the current directory, be sure to include a dot (.) in the classpath to specify the current directory.
- ▶ To execute the ListTest program, use the following command:
 - ∘ java -classpath .:.. ListTest

21.5 Stacks

- A stack is a constrained version of a list—new nodes can be added to and removed from a stack only at the top.
- A stack is referred to as a last-in, first-out (LIFO) data structure.
- A stack is not required to be implemented as a linked list—it can also be implemented using an array.
- The primary methods for manipulating a stack are push and pop, which add a new node to the top of the stack and remove a node from the top of the stack, respectively.

```
// Fig. 21.9: Stack.java
2 // Stack uses a composed List object.
    package com.deitel.datastructures;
    import java.util.NoSuchElementException;
    public class Stack<E> {
       private List<E> stackList;
10
      // constructor
       public Stack() {stackList = new List<E>("stack");}
12
13
      // add object to stack
       public void push(E object) {stackList.insertAtFront(object);}
14
15
```

Fig. 21.9 | Stack uses a composed List object. (Part I of 2.)

```
16
       // remove object from stack
       public E pop() throws NoSuchElementException {
17
          return stackList.removeFromFront();
18
19
20
       // determine if stack is empty
21
       public boolean isEmpty() {return stackList.isEmpty();}
22
23
24
       // output stack contents
       public void print() {stackList.print();}
25
26
```

Fig. 21.9 | Stack uses a composed List object. (Part 2 of 2.)

```
// Fig. 21.10: StackTest.java
   // Stack manipulation program.
    import com.deitel.datastructures.Stack;
    import java.util.NoSuchElementException;
    public class StackTest {
       public static void main(String[] args) {
          Stack<Integer> stack = new Stack<>();
          // use push method
10
          stack.push(-1);
\mathbf{I}
           stack.print();
12
          stack.push(0);
13
           stack.print();
14
          stack.push(1);
15
          stack.print();
16
           stack.push(5);
17
           stack.print();
18
19
```

Fig. 21.10 | Stack manipulation program. (Part 1 of 3.)

```
// remove items from stack
20
21
          boolean continueLoop = true;
22
23
          while (continueLoop) {
24
              try {
                 int removedItem = stack.pop(); // remove top element
25
                 System.out.printf("%n%d popped%n", removedItem);
26
                 stack.print();
27
28
29
              catch (NoSuchElementException noSuchElementException) {
30
                 continueLoop = false;
31
                 noSuchElementException.printStackTrace();
32
33
34
35
```

Fig. 21.10 | Stack manipulation program. (Part 2 of 3.)

```
The stack is: -1
The stack is: 0 -1
The stack is: 1 \quad 0 \quad -1
The stack is: 5 1 0 -1
5 popped
The stack is: 1 0 -1
1 popped
The stack is: 0 -1
0 popped
The stack is: -1
-1 popped
Empty stack
java.util.NoSuchElementException: stack is empty
        at com.deitel.datastructures.List.removeFromFront(List.java:68)
        at com.deitel.datastructures.Stack.pop(Stack.java:18)
        at StackTest.main(StackTest.java:25)
```

Fig. 21.10 Stack manipulation program. (Part 3 of 3.)

21.6 Queues

- ▶ A queue is similar to a checkout line in a supermarket—the cashier services the person at the beginning of the line first.
- Other customers enter the line only at the end and wait for service.
- Queue nodes are removed only from the head (or front) of the queue and are inserted only at the tail (or end).
- ▶ For this reason, a queue is a first-in, first-out (FIFO) data structure.
- ▶ The insert and remove operations are known as enqueue and dequeue.
- Next program creates a Queue<T> class that contains a List<T> object and provides methods enqueue, dequeue, isEmpty and print.

```
// Fig. 21.11: Queue.java
2 // Queue uses class List.
    package com.deitel.datastructures;
    import java.util.NoSuchElementException;
    public class Queue<E> {
       private List<E> queueList;
10
      // constructor
       public Queue() {queueList = new List<E>("queue");}
12
13
      // add object to queue
       public void enqueue(E object) {queueList.insertAtBack(object);}
14
```

Fig. 21.11 Queue uses class List. (Part 1 of 2.)

```
15
16
       // remove object from queue
17
       public E dequeue() throws NoSuchElementException {
18
          return queueList.removeFromFront();
19
20
21
       // determine if queue is empty
       public boolean isEmpty() {return queueList.isEmpty();}
22
23
24
       // output queue contents
25
       public void print() {queueList.print();}
26
```

Fig. 21.11 Queue uses class List. (Part 2 of 2.)

```
// Fig. 21.12: QueueTest.java
  // Class QueueTest.
    import com.deitel.datastructures.Queue;
    import java.util.NoSuchElementException;
    public class QueueTest {
       public static void main(String[] args) {
          Queue<Integer> queue = new Queue<>();
          // use enqueue method
10
\mathbf{I}
          queue.enqueue(-1);
          queue.print();
12
          queue.enqueue(0);
13
14
          queue.print();
15
          queue.enqueue(1);
16
          queue.print();
17
          queue.enqueue(5);
18
          queue.print();
19
```

Fig. 21.12 | Class QueueTest. (Part 1 of 3.)

```
20
          // remove objects from queue
21
          boolean continueLoop = true;
22
23
          while (continueLoop) {
24
              try {
25
                 int removedItem = queue.dequeue(); // remove head element
                 System.out.printf("%n%d dequeued%n", removedItem);
26
                 queue.print();
27
28
29
              catch (NoSuchElementException noSuchElementException) {
30
                 continueLoop = false;
                 noSuchElementException.printStackTrace();
31
32
33
34
35
```

Fig. 21.12 Class QueueTest. (Part 2 of 3.)

```
The queue is: -1
The queue is: -1 0
The queue is: -1 0 1
The queue is: -1 0 1 5
-1 dequeued
The queue is: 0 1 5
0 dequeued
The queue is: 1 5
1 dequeued
The queue is: 5
5 dequeued
Empty queue
java.util.NoSuchElementException: queue is empty
        at com.deitel.datastructures.List.removeFromFront(List.java:68)
        at com.deitel.datastructures.Queue.dequeue(Queue.java:18)
        at QueueTest.main(QueueTest.java:25)
```

Fig. 21.12 | Class QueueTest. (Part 3 of 3.)

21.7 Trees

- A tree is a nonlinear, two-dimensional data structure with special properties.
- Tree nodes contain two or more links.
- Binary tree nodes each contain two links (one or both of which may be null).
- The root node is the first node in a tree.
- Each link in the root node refers to a child.
- ▶ The left child is the first node in the left subtree (also known as the root node of the left subtree), and the right child is the first node in the right subtree (also known as the root node of the right subtree).
- The children of a specific node are called siblings.
- A node with no children is called a leaf node.
- Computer scientists normally draw trees from the root node down—the opposite of the way most trees grow in nature.

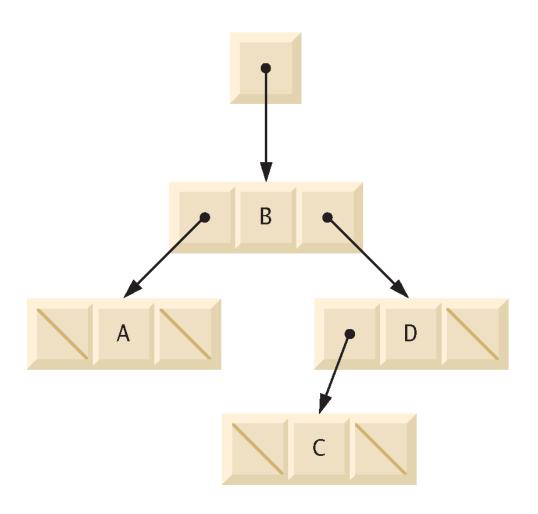


Fig. 21.13 | Binary tree graphical representation.

21.7 Trees (Cont.)

- A binary search tree (with no duplicate node values) has the characteristic that the values in any left subtree are less than the value in that subtree's parent node, and the values in any right subtree are greater than the value in that subtree's parent node.
- Fig. 21.14 illustrates a binary search tree with 12 integer values.
- ▶ Figs. 21.15 and 21.16 create a generic binary search tree class and use it to manipulate a tree of integers.
- ▶ The application in traverses the tree (i.e., walks through all its nodes) three ways—using recursive inorder, preorder and postorder traversals (trees are almost always processed recursively).

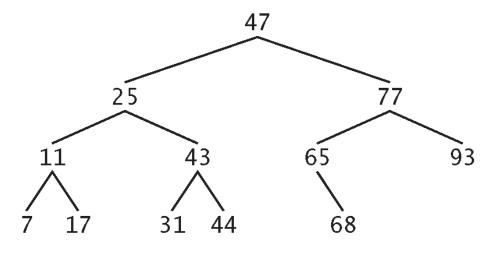


Fig. 21.14 | Binary search tree containing 12 values.

```
// Fig. 21.15: Tree.java
  // TreeNode and Tree class declarations for a binary search tree.
    package com.deitel.datastructures;
    // class TreeNode definition
    class TreeNode<E extends Comparable<E>> {
       // package access members
       TreeNode<E> leftNode;
       E data; // node value
       TreeNode<E> rightNode;
10
12
       // constructor initializes data and makes this a leaf node
13
       public TreeNode(E nodeData) {
14
          data = nodeData;
15
          leftNode = rightNode = null; // node has no children
16
17
```

Fig. 21.15 | TreeNode and Tree class declarations for a binary search tree. (Part 1 of 6.)

```
18
       // locate insertion point and insert new node; ignore duplicate values
       public void insert(E insertValue) {
19
          // insert in left subtree
20
21
           if (insertValue.compareTo(data) < 0) {</pre>
22
              // insert new TreeNode
              if (leftNode == null) {
23
                 leftNode = new TreeNode<E>(insertValue);
24
25
              else { // continue traversing left subtree recursively
26
                 leftNode.insert(insertValue);
27
28
29
           // insert in right subtree
30
           else if (insertValue.compareTo(data) > 0) {
3 I
              // insert new TreeNode
32
              if (rightNode == null) {
33
                 rightNode = new TreeNode<E>(insertValue);
34
35
              else { // continue traversing right subtree recursively
36
                 rightNode.insert(insertValue);
37
38
39
40
41
```

Fig. 21.15 | TreeNode and Tree class declarations for a binary search tree. (Part 2 of 6.)

```
42
    // class Tree definition
43
    public class Tree<E extends Comparable<E>>> {
44
45
       private TreeNode<E> root;
46
       // constructor initializes an empty Tree of integers
47
       public Tree() {root = null;}
48
49
50
       // insert a new node in the binary search tree
       public void insertNode(E insertValue) {
51
52
           if (root == null) {
              root = new TreeNode<E>(insertValue); // create root node
53
54
55
          else {
              root.insert(insertValue); // call the insert method
56
57
58
59
```

Fig. 21.15 | TreeNode and Tree class declarations for a binary search tree. (Part 3 of 6.)

```
60
       // begin preorder traversal
61
       public void preorderTraversal() {preorderHelper(root);}
62
63
       // recursive method to perform preorder traversal
       private void preorderHelper(TreeNode<E> node) {
64
65
           if (node == null) {
66
              return;
67
68
69
           System.out.printf("%s ", node.data); // output node data
          preorderHelper(node.leftNode); // traverse left subtree
70
          preorderHelper(node.rightNode); // traverse right subtree
71
72
73
```

Fig. 21.15 | TreeNode and Tree class declarations for a binary search tree. (Part 4 of 6.)

```
// begin inorder traversal
74
75
       public void inorderTraversal() {inorderHelper(root);}
76
       // recursive method to perform inorder traversal
77
       private void inorderHelper(TreeNode<E> node) {
78
           if (node == null) {
79
80
              return;
81
82
83
          inorderHelper(node.leftNode); // traverse left subtree
           System.out.printf("%s ", node.data); // output node data
84
          inorderHelper(node.rightNode); // traverse right subtree
85
86
87
```

Fig. 21.15 | TreeNode and Tree class declarations for a binary search tree. (Part 5 of 6.)

```
// begin postorder traversal
88
89
       public void postorderTraversal() {postorderHelper(root);}
90
91
       // recursive method to perform postorder traversal
       private void postorderHelper(TreeNode<E> node) {
92
93
           if (node == null) {
94
              return;
95
96
          postorderHelper(node.leftNode); // traverse left subtree
97
          postorderHelper(node.rightNode); // traverse right subtree
98
          System.out.printf("%s ", node.data); // output node data
99
100
101
```

Fig. 21.15 | TreeNode and Tree class declarations for a binary search tree. (Part 6 of 6.)

```
// Fig. 21.16: TreeTest.java
   // Binary tree test program.
    import java.security.SecureRandom;
    import com.deitel.datastructures.Tree;
    public class TreeTest {
       public static void main(String[] args) {
          Tree<Integer> tree = new Tree<Integer>();
          SecureRandom randomNumber = new SecureRandom();
10
          System.out.println("Inserting the following values: ");
12
          // insert 10 random integers from 0-99 in tree
13
          for (int i = 1; i <= 10; i++) {
14
             int value = randomNumber.nextInt(100);
15
             System.out.printf("%d ", value);
16
             tree.insertNode(value);
17
18
```

Fig. 21.16 | Binary tree test program. (Part 1 of 3.)

```
19
          System.out.printf("%n%nPreorder traversal%n");
20
          tree.preorderTraversal();
21
22
          System.out.printf("%n%nInorder traversal%n");
23
          tree.inorderTraversal();
24
25
          System.out.printf("%n%nPostorder traversal%n");
26
          tree.postorderTraversal();
27
          System.out.println();
28
29
30
```

Fig. 21.16 Binary tree test program. (Part 2 of 3.)

```
Inserting the following values: 49 64 14 34 85 64 46 14 37 55

Preorder traversal 49 14 34 46 37 64 55 85

Inorder traversal 14 34 37 46 49 55 64 85

Postorder traversal 37 46 34 14 55 85 64 49
```

Fig. 21.16 Binary tree test program. (Part 3 of 3.)

- Class Tree requires Comparable objects, so that each value inserted in the tree can be compared with the existing values to determine the insertion point.
- ▶ Class Tree's method insertNode first determines whether the tree is empty.
- If so, it allocates a new TreeNode, initializes the node with the value being inserted in the tree and assigns the new node to reference root.
- If the tree is not empty, the method calls TreeNode method insert
- ▶ This method uses recursion to determine the location for the new node in the tree and inserts the node at that location.
- A node can be inserted only as a leaf node in a binary search tree.

- TreeNode method insert compares the value to insert with the data value in the root node.
 - If the insert value is less than the root node data, the program determines if the left subtree is empty.
 - If so, allocates a new TreeNode, initializes it with the value being inserted and assigns the new node to reference leftNode.
 - Otherwise, recursively calls insert for the left subtree to insert the value into the left subtree.
 - If the insert value is greater than the root node data, the program determines if the right subtree is empty.
 - If so, allocates a new TreeNode, initializes it with the value being inserted and assigns the new node to reference rightNode.
 - Otherwise, recursively calls insert for the right subtree to insert the value in the right subtree.
 - If the insertValue is already in the tree, it's simply ignored.

- Methods inorderTraversal, preorderTraversal and postorderTraversal call Tree helper methods inorderHelper, preorderHelper and postorderHelper, respectively, to traverse the tree and print the node values.
 - Helper methods in class Tree enable you to start a traversal without having to pass the root node to the method.
 - Reference root is an implementation detail that a programmer should not be able to access.
- Methods inorderTraversal, preorderTraversal and postorderTraversal simply take the private root reference and pass it to the appropriate helper method to initiate a traversal of the tree.

- The base case for each helper method determines whether the reference it receives is null and, if so, returns immediately.
- Method inorderHelper defines the steps for an inorder traversal:
 - Traverse the left subtree with a call to inorderHelper
 - Process the value in the node
 - Traverse the right subtree with a call to inorderHelper
- The inorder traversal does not process the value in a node until the values in that node's left subtree are processed.
- Note that the inorder traversal of a binary search tree prints the node values in ascending order.
- The process of creating a binary search tree actually sorts the data; thus, it's called the binary tree sort.

- Method preorderHelper defines the steps for a preorder traversal:
 - Process the value in the node
 - Traverse the left subtree with a call to preorderHelper
 - Traverse the right subtree with a call to preorderHelper
- The preorder traversal processes the value in each node as the node is visited.
- After processing the value in a particular node, the preorder traversal processes the values in the left subtree, then processes the values in the right subtree.

- Method postorderHelper defines the steps for a postorder traversal:
 - Traverse the left subtree with a call to postorderHelper
 - Traverse the right subtree with a call to postorderHelper
 - Process the value in the node
- The postorder traversal processes the value in each node after the values of all that node's children are processed.

- A binary search tree facilitates duplicate elimination.
- Insert operation recognizes a duplicate, because a duplicate follows the same "go left" or "go right" decisions on each comparison as the original value did.
- Searching a binary tree for a value that matches a key value is fast, especially for tightly packed (or balanced) trees.
 - In a tightly packed tree, each level contains about twice as many elements as the previous level.
 - A tightly packed binary search tree (e.g.,) with n elements has $\log_2 n$ levels.
 - Thus, at most $\log_2 n$ comparisons are required either to find a match or to determine that no match exists.
 - Searching a (tightly packed) 1000-element binary search tree requires at most 10 comparisons, because $2^{10} > 1000$.
 - Searching a (tightly packed) 1,000,000-element binary search tree requires at most 20 comparisons, because $2^{20} > 1,000,000$.

- The chapter exercises present algorithms for several other binary tree operations, such as deleting an item from a binary tree, printing a binary tree in a two-dimensional tree format and performing a level-order traversal of a binary tree.
- The level-order traversal visits the nodes of the tree row by row, starting at the root node level.