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CSC332-Activity 4

Let  $N = 5$  because there are 5 processes and  $M = 4$  because there are 4 resources. As well as  $Work = Availability$ .

From using the banker's algorithm, we can solve if the state is safe or not by using  $Need = Max - Allocation$ . Thus,

$$T_0 \Rightarrow 2\ 4\ 3\ 1 - 0\ 1\ 2\ 0 = 2\ 3\ 1\ 1$$

$$T_1 \Rightarrow 2\ 4\ 2\ 4 - 0\ 1\ 1\ 2 = 2\ 3\ 1\ 2$$

$$T_2 \Rightarrow 3\ 6\ 5\ 1 - 1\ 2\ 4\ 0 = 2\ 4\ 1\ 1$$

$$T_3 \Rightarrow 1\ 2\ 6\ 2 - 3\ 1\ 2\ 0 = 1\ 8\ 5\ 8$$

$$T_4 \Rightarrow 3\ 1\ 1\ 2 - 1\ 0\ 0\ 1 = 2\ 1\ 1\ 1$$

	Allocation Matrix	Max	Need
	A B C D	A B C D	A B C D
$T_0$	0 1 2 0	2 4 3 1	2 3 1 1
$T_1$	0 1 1 2	2 4 2 4	2 3 1 2
$T_2$	1 2 4 0	3 6 5 1	2 4 1 1
$T_3$	3 1 2 0	1 2 6 2	1 8 5 8
$T_4$	1 0 0 1	3 1 1 2	2 1 1 1

A) Available = 2223

Now we check if  $Need \leq Availability$ , then available = available + allocation.

We will start with  $T_4 \rightarrow 2223 - 2111 = 0112$  remaining.

$T_4$  will release all resources to  $T_0 \rightarrow 3112 + 0112 = 3224$

Then, we will subtract  $T_0$  to available  $\rightarrow 3224 - 2311 = 0913$

Then, we add  $T_0$  will release all resources  $\rightarrow 0913 + 2431 = 3344$

Then, we subtract  $T_1$  to the current available resources  $\rightarrow 3344 - 2312 = 1032$

Then,  $T_1$  will release all its resources  $\rightarrow 1032 + 2424 = 3456$

Then, its  $T_2$  time  $\rightarrow 3456 - 2411 = 1045$

Now, we release  $T_2 \rightarrow 1045 + 3651 = 4696$

Finally, its  $T_3$  time  $\rightarrow 4696 - 1858 = 2838$

And now we release  $T_3 \rightarrow 2838 + 1262 = 4100$

So, we get a safe state and the order of execution is  $T_4, T_0, T_1, T_2, T_3$

**It is taking way too long to type everything, so I have decided to do the rest by handwriting the rest.**

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B) Available = (4, 4, 1, 1)

add allocation of  $T_4$  to current available

$$\text{Available} = (4, 4, 1, 1) + (1, 0, 1, 1) = (5, 4, 1, 2)$$

This satisfy the need of  $T_1$  so we can add the allocation of  $T_1$  to current Available

$$(5, 4, 1, 2) + (0, 1, 1, 2) = (5, 5, 2, 4)$$

This satisfy the need of  $T_0$ , so lets add the allocation of  $T_0$  to current available

$$(5, 5, 2, 4) + (0, 1, 2, 0) = (5, 6, 6, 4)$$

because the above available, we know it can satisfy  $T_2$  and  $T_3$   
so we can say we have a safe sequence,  
and the safe sequence is

$$T_4, T_1, T_0, T_2, T_3$$

C) Available = (3, 0, 1, 4)

we cannot satisfy the need of any thread so the state is unsafe.

d) Available = (1, 5, 2, 2)

add allocation ~~of  $T_3$~~  of  $T_3$  to available resources

$$(1, 5, 2, 2) + (3, 1, 2, 0) = (4, 6, 4, 2)$$

This satisfy the need of  $T_4$  to available resources, now we can add the allocation of  $T_4$

$$(4, 6, 4, 2) + (1, 0, 0, 1) = (5, 6, 4, 3)$$

This can satisfy the need of  $T_0$  and  $T_1$  and  $T_2$  which means it is in a safe state. and the sequence can be  $T_3, T_4, T_0, T_1, T_2$