

ps5

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## Problem 1

According to the question,  $A$  is a symmetric matrix. Suppose the dimension of  $A$  is  $n \times n$ . Now, we have

$$A = \Gamma \Lambda \Gamma^T$$

Then, what we need to prove is that the determinant of  $A$  is equal to the product of the eigenvalues.

$$\begin{aligned} |A| &= |\Gamma \Lambda \Gamma^T| \\ &= |\Gamma| |\Lambda| |\Gamma^T| \\ &= |\Gamma| |\Gamma^T| |\Lambda| \\ &= |\Gamma \Gamma^T| |\Lambda| \end{aligned}$$

Because  $\Gamma$  is an orthogonal matrix of eigenvectors and  $\Lambda$  is a diagonal matrix of eigenvalues. Therefore, we have

$$\begin{aligned} \Gamma \Gamma^T &= I \\ |A| &= |I| |\Lambda| \\ &= |\Lambda| \\ &= \prod_{i=1}^n \lambda_i \end{aligned}$$

This supports the conclusion.

## Problem 2

```
z=1e10
expit<-function(z){
  return(exp(z)/(1+exp(z)))
}
expit(z)
```

```
## [1] NaN
```

It returns NaN, which means that it is not computable in R when  $z$  is large.

```
z=1e10
exp(z)
```

```
## [1] Inf
```

Because when  $z$  is large and  $\exp(z)$  is difficult to represent, R returns ‘Inf’. As we all know,  $\text{Inf}/(\text{Inf}+1)$  is obviously not computable in R.

To fix that, we can implement the following transformation:

$$\frac{\exp(z)}{1 + \exp(z)} = \frac{1}{\exp(-z) + 1}$$

Now even  $z$  becomes large,  $\exp(-z) \rightarrow 0$ .  $\frac{1}{0+1}$  is still computable in R. We can validate it using the following code:

```
z=1e10
expit_trans<-function(z){
  return(1/(1+exp(-z)))
}
expit_trans(z)
```

```
## [1] 1
```

This shows that the numerical issue has been successfully fixed.

### Problem 3

```
set.seed(1)
z <- rnorm(10, 0, 1)
x <- z + 1e12
formatC(var(z), 20, format = 'f')
```

```
## [1] "0.60931443706111987346"
```

```
formatC(var(x), 20, format = 'f')
```

```
## [1] "0.60931216345893013386"
```

Firstly, we claim that  $Var(X) = Var(Z + a) = Var(Z)$  where  $a$  is a fixed number. i.e, the variance of  $z$  and  $x$  mentioned above should be same mathematically.

Then we look into  $x$  and  $z$  vector to get some insights about what happens:

```
formatC(z[4], 20, format = 'f')
```

```
## [1] "1.59528080213779155372"
```

```
formatC(x[4], 20, format = 'f')
```

```
## [1] "1000000000001.59533691406250000000"
```

It seems that the accuracy of difference between corresponding element of  $x$  and  $z$  limited to  $10^{-3}$ . Theoretically,  $x$  is around  $10^{12}$ . Then the absolute error in representing it is  $x\epsilon \approx 2 * 10^{-4}$ . In other words, the machine epsilon between numbers around  $10^{12}$  and the next significant number is quite big compared with the base number( the element of  $z$ ).

This means that we have accuracy of variance to the order  $10^{-4}$ , which explains these two estimates agree to only 4 or 5 digits.

```
set.seed(12)
z <- rnorm(10, 0, 1)
x <- z + 1e12
formatC(var(z), 20, format = 'f')
```

```
## [1] "1.00131389554051986046"
```

```
formatC(var(x), 20, format = 'f')
```

```
## [1] "1.00132647818989228838"
```

## Problem 4

(a)

Since we have exactly  $p$  cores to do the computation, even if break up  $Y$  into  $n$  individual column-wise, we can only carry out  $p$  processes at the same time, which makes the former division meaningless.

(b)

- Approach A

For each task, we have two input matrices with dimensions of  $m * n$  and  $n * n$ . Besides, we have calculated matrix with dimension of  $m * n$ . As a result, we need  $2 * m * n + n * n$  of memory for each task. Since we have  $p$  tasks carried out at the same time, we need  $2n^2 + n^2p$  of memory. Then the total communication cost is  $2n^2 + n^2p$ .

- Approach B

For each task, we have two input matrices with dimensions of  $m * n$  and  $n * m$ . Besides, we have calculated matrix with dimension of  $m * m$ . As a result, we need  $2 * m * n + m * m$  of memory for each task. Since we have  $p$  tasks carried out at the same time, we need  $2n^2 + m^2p = 2n^2 + mn$  of memory at the same time. Then the total communication cost is  $(2n^2 + mn) * p = 2n^2p + n^2$ .

In conclusion, approach B is better for minimizing memory use and approach A is better for minimizing communication cost.