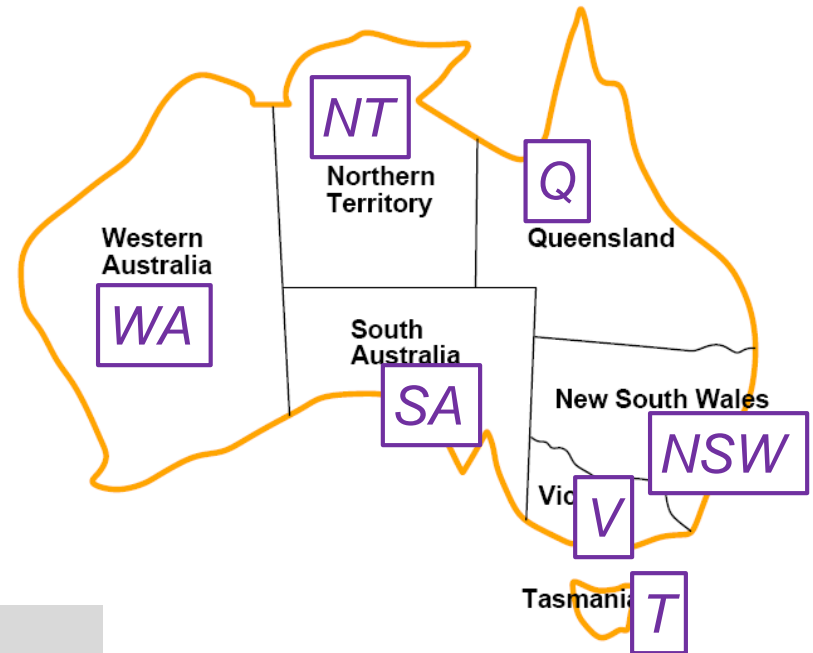
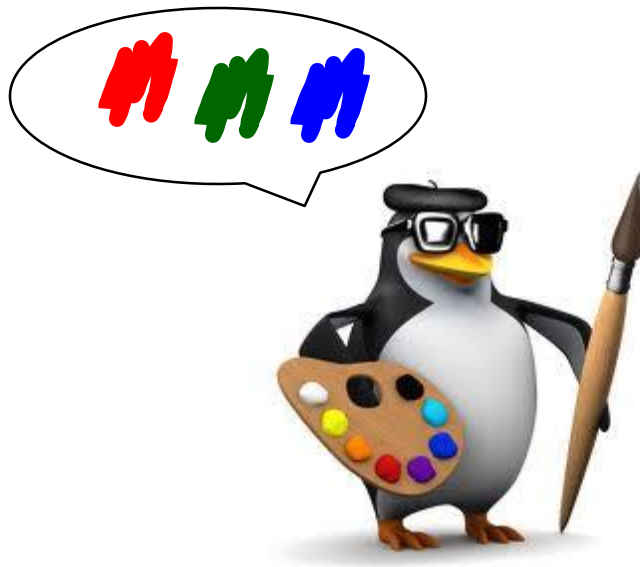


Constraint Satisfaction Problems (CSP)

Example: Map-Coloring Problems

Goal: To color the provinces of Australia as red, green, or blue, with no adjacent provinces having the same color.



Variables: *WA, NT, Q, NSW, V, SA, T*

Domains: *{red, green, blue}*

Constraints: Adjacent regions must have different colors: *WA ≠ NT*, etc.

Example: Cryptarithmic Puzzles

Goal: To assign (all different) digits to the symbols.



$$\begin{array}{r} \text{T W O} \\ + \text{T W O} \\ \hline \text{F O U R} \end{array}$$

Variables: $F, T, U, W, R, O, C_1, C_2, C_3$

Domains: $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$ for F, T, U, W, R, O

$\{0, 1\}$ for C_1, C_2, C_3

Constraints: $F > 0$

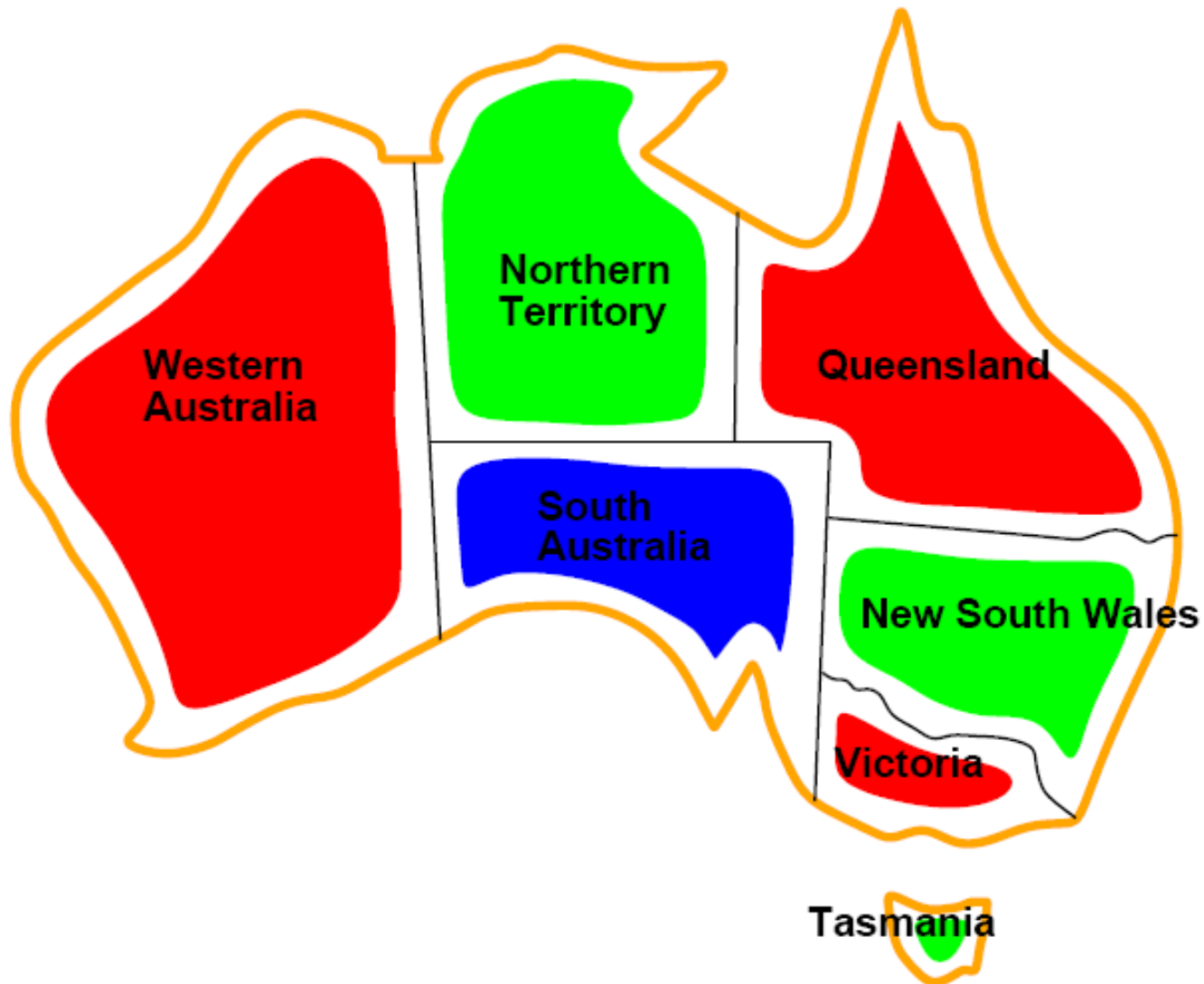
$\text{alldiff}(F, T, U, W, R, O)$

$O + O = R + 10 * C_1$, etc.

Constraint Satisfaction Problems

- Information about the state space:
 - Variables: X_i
 - Domains (allowed values for the variables): D_i (one for each variable)
- Goal test: Whether the variable assignments satisfy the given set of constraints C .
- Solution: Variable assignments that satisfy all the constraints.
- Search algorithms can be used to find solutions. (Only the final state, not the path, is needed).

Example Solution for Map-Coloring



Varieties of CSPs

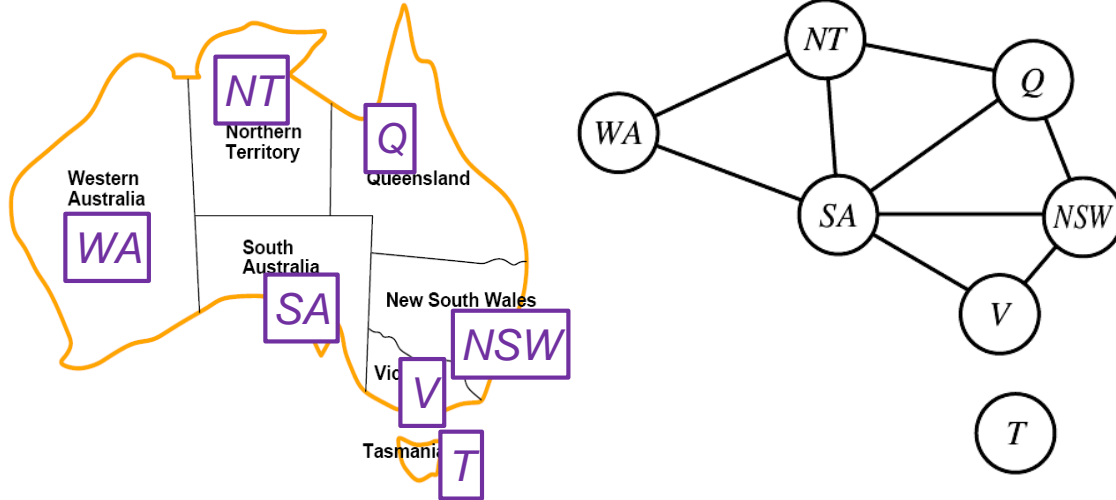
- Discrete variables with finite domains (covered here):
 - n variables, domain size = d : $O(d^n)$ complete assignments
- Discrete variables with infinite domains:
 - Example: job scheduling with the variables being the start/end days for the jobs.
 - Such CSPs with linear constraints are solvable. For such CSPs with nonlinear constraints, the problem is undecidable.
- Continuous variables:
 - Example: job scheduling with the variable being the start/end times for the jobs.
 - Such CSPs with linear constraints are solvable in polynomial time by linear programming methods.

Types of Constraints

- Unary constraints: involving a single variable
 - e.g., *SA* \neq *green*
- Binary constraints: involving pairs of variables
 - e.g., *SA* \neq *WA*
- Higher-order constraints: involving 3 or more variables
 - e.g., cryptarithmic column constraints
- Global constraints: involve an arbitrary number of variables
 - e.g., *alldiff*, *atmost*, etc.
 - Some are better solved with specialized methods.
- Preferences (soft constraints)
 - e.g., *SA* likes *red* more than *green*
 - often represented by a cost for each variable assignment
 - ➔ optimization problems

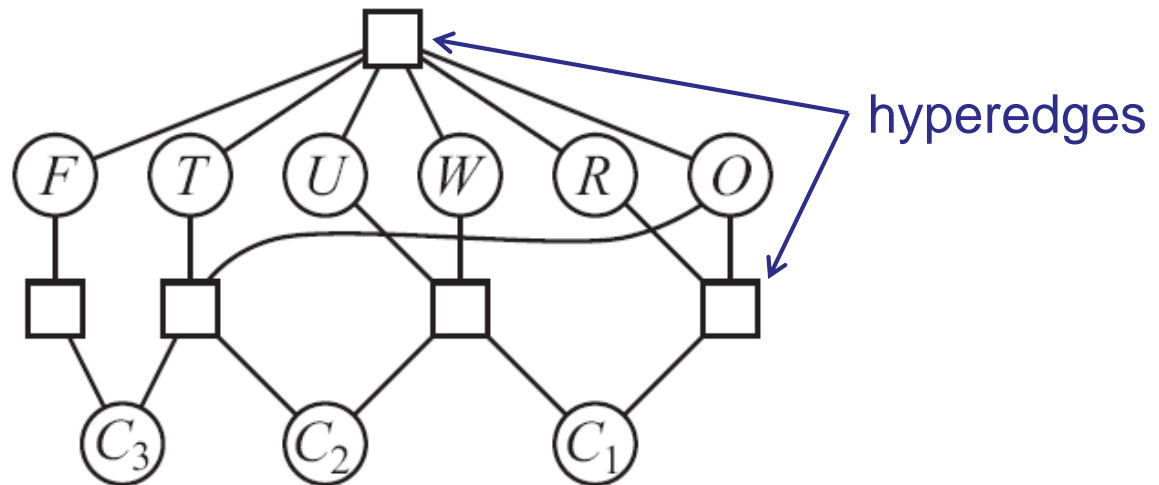
Constraint Graph

- Binary CSPs: Variables as vertices and edges as constraints



- Higher-order CSPs: Represented as **hypergraphs**

$$\begin{array}{r}
 T \ W \ O \\
 + \ T \ W \ O \\
 \hline
 F \ O \ U \ R
 \end{array}$$



Constraint Propagation

Inference in CSP: Given the constraints and possible assigned values of some variables, determine the domains (allowed values) of the other variables.

- **Node consistency:** A variable is node-consistent if each possible value satisfies its unary constraints.
- **Arc consistency:** An arc $(X_i \rightarrow X_j)$ is arc-consistent if each possible value of X_i leaves some valid values for X_j according to the binary constraints for $\{X_i, X_j\}$.
- **Path consistency:** A path $(X_i \rightarrow X_m \rightarrow X_j)$ is path-consistent if each pair of possible values for X_i and X_j leaves some valid values for X_m according to the binary constraints for $\{X_i, X_m\}$ and $\{X_j, X_m\}$.

AC-3 Algorithm

- Goal: To make every arc arc-consistent. (It attempts to reduce the domains by removing values that cause violation of arc-consistency.)
- Initialization: A set containing every arc in a CSP.
- In each step, an arc $(X_i \rightarrow X_j)$ is popped off the set for consideration.
 - Make $(X_i \rightarrow X_j)$ arc-consistent by removing from D_i those values that would leave no valid value for X_j according to the binary constraints for $\{X_i, X_j\}$.
 - If D_i is changed, then for each neighbor X_k of X_i , add the arc $(X_k \rightarrow X_i)$ to the set.
 - If D_i becomes empty, terminate the processing; there is no solution.

AC-3 Example

Example: X, Y are digits, and $Y=X^2$.

■ Initial domains:

- $D_X = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$

- $D_Y = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$

■ After checking the arc $(X \rightarrow Y)$:

- $D_X = \{0, 1, 2, 3\}$

■ Next, after checking the arc $(Y \rightarrow X)$:

- $D_Y = \{0, 1, 4, 9\}$

■ No more change; the process terminates.

AC-3 Example (Sudoku)

						5		
3		2		7		9	1	
6			9					
							2	6
	2		3			1	5	9
7	9		6		5		8	
1		9	7					
4	5					2	3	
	3	8	4	5		6		

An example problem:

- 81 variables
- 32 unary constraints
- Q: How many binary constraints?

Initially, all the unassigned blanks have domain $\{1,2,\dots,9\}$.

AC-3 Example (Sudoku)

						5		
3		2		7		9	1	
6			9					
							2	6
	2		3			1	5	9
7	9		6		5		8	
1	6	9	7			48	4	458
4	5	67				2	3	178
2	3	8	4	5		6	79	127

Some updated domains after checking arcs involving already assigned blanks.

						5		
3		2		7		9	1	
6			9					
							2	6
	2		3			1	5	9
7	9		6		5		8	
1	6	9	7			48	4	458
4	5	67				2	3	178
2	3	8	4	5		6	79	127

Further reduction of some domains.

AC-3 Example (Sudoku)

						5		
3		2		7		9	1	
6			9					
							2	6
	2		3			1	5	9
7	9		6		5		8	
1	6	9	7			48	4	458
4	5	67				2	3	178
2	3	8	4	5		6	79	127

						5		
3		2		7		9	1	
6			9					
							2	6
	2		3			1	5	9
7	9		6		5		8	
1	6	9	7			48	4	458
4	5	67				2	3	178
2	3	8	4	5		6	79	127

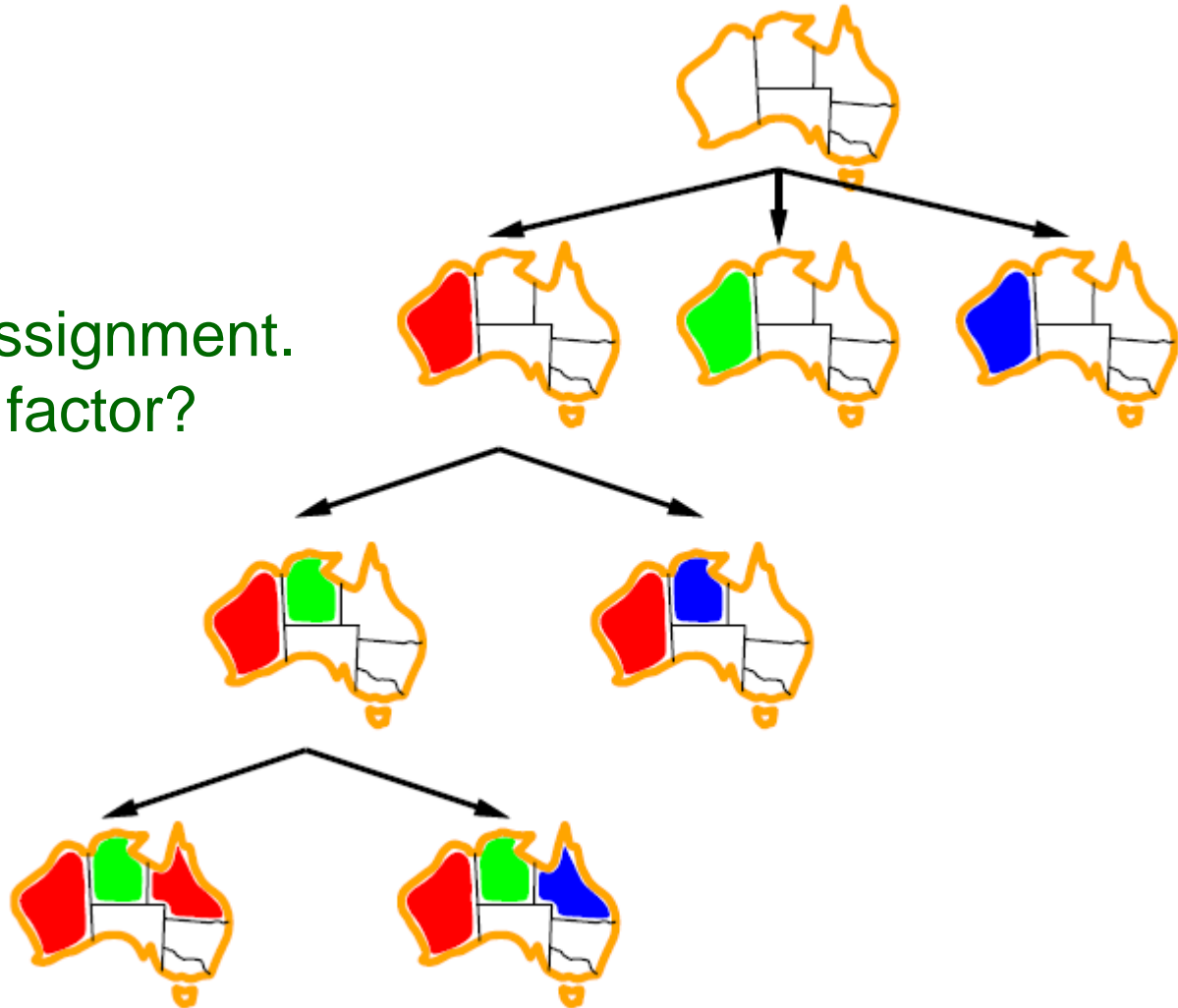
AC-3 can occasionally solve a problem by itself.

Solving CSPs by Searching

Backtracking search: depth-first, incremental formulation (assigning one variable in each step).

Example:

Q: Consider the first assignment.
What is the branching factor?



Backtracking Search

General-purpose methods can give huge gains in speed:

- Which variable should be assigned next?
 - **Minimum remaining values (MRV) heuristic**
 - **Degree heuristic**
- In what order should its values be tried?
 - **Least constraining value (LCV) heuristic**
- Can we detect an inevitable failure early?
- Can we take advantage of the problem structure?

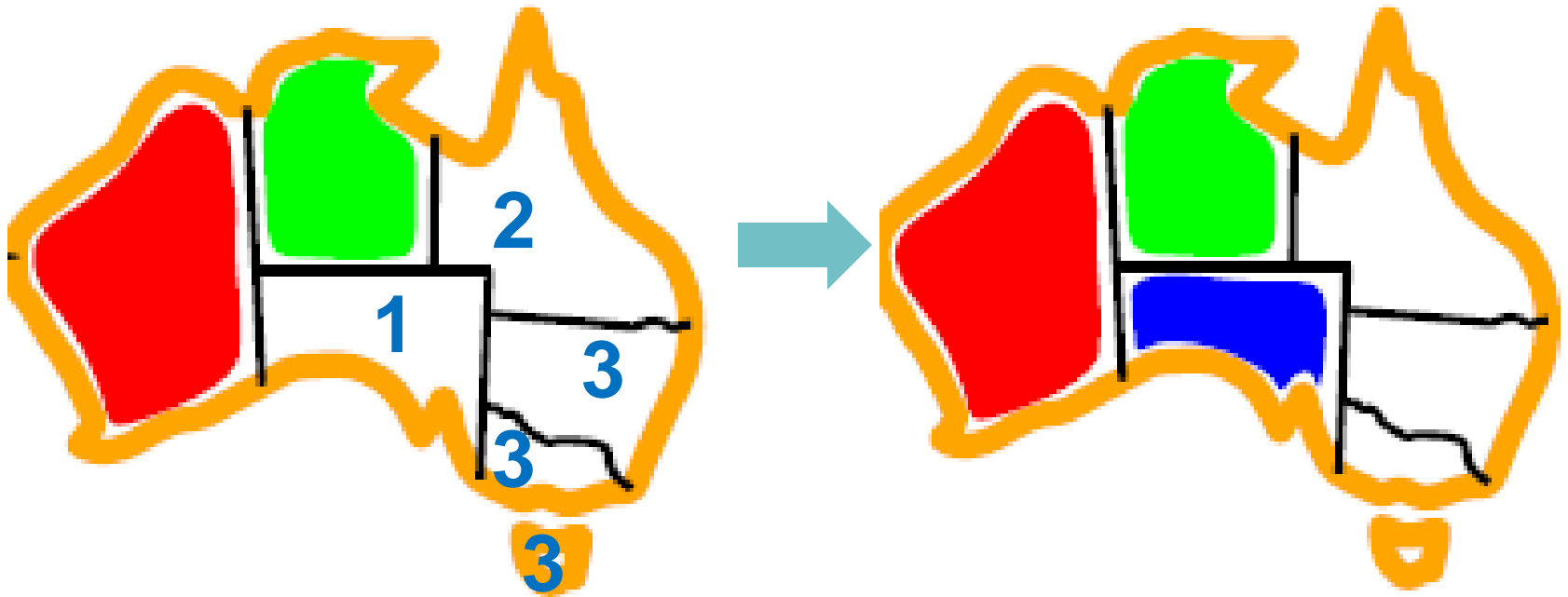


Usually applied in the order of MRV→degree→LCV.

Minimum Remaining Values (MRV)

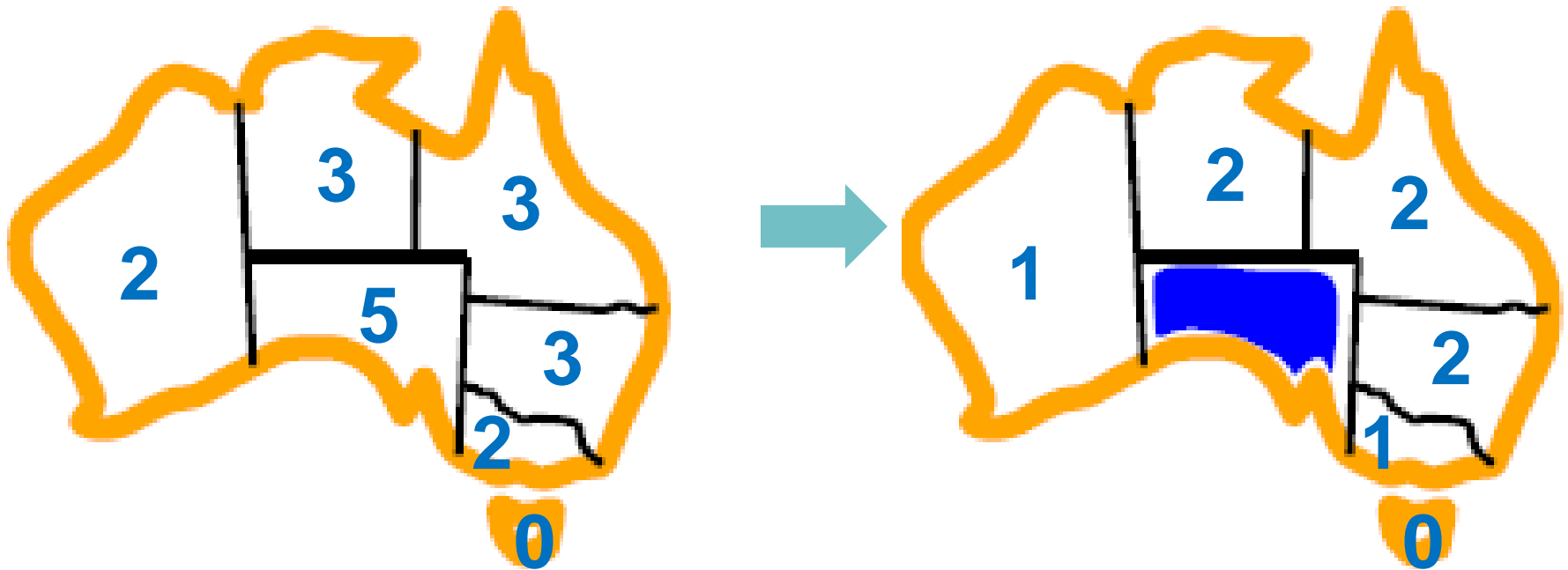
Idea: Choose the variable with the fewest legal values

Example: Which region to color next?



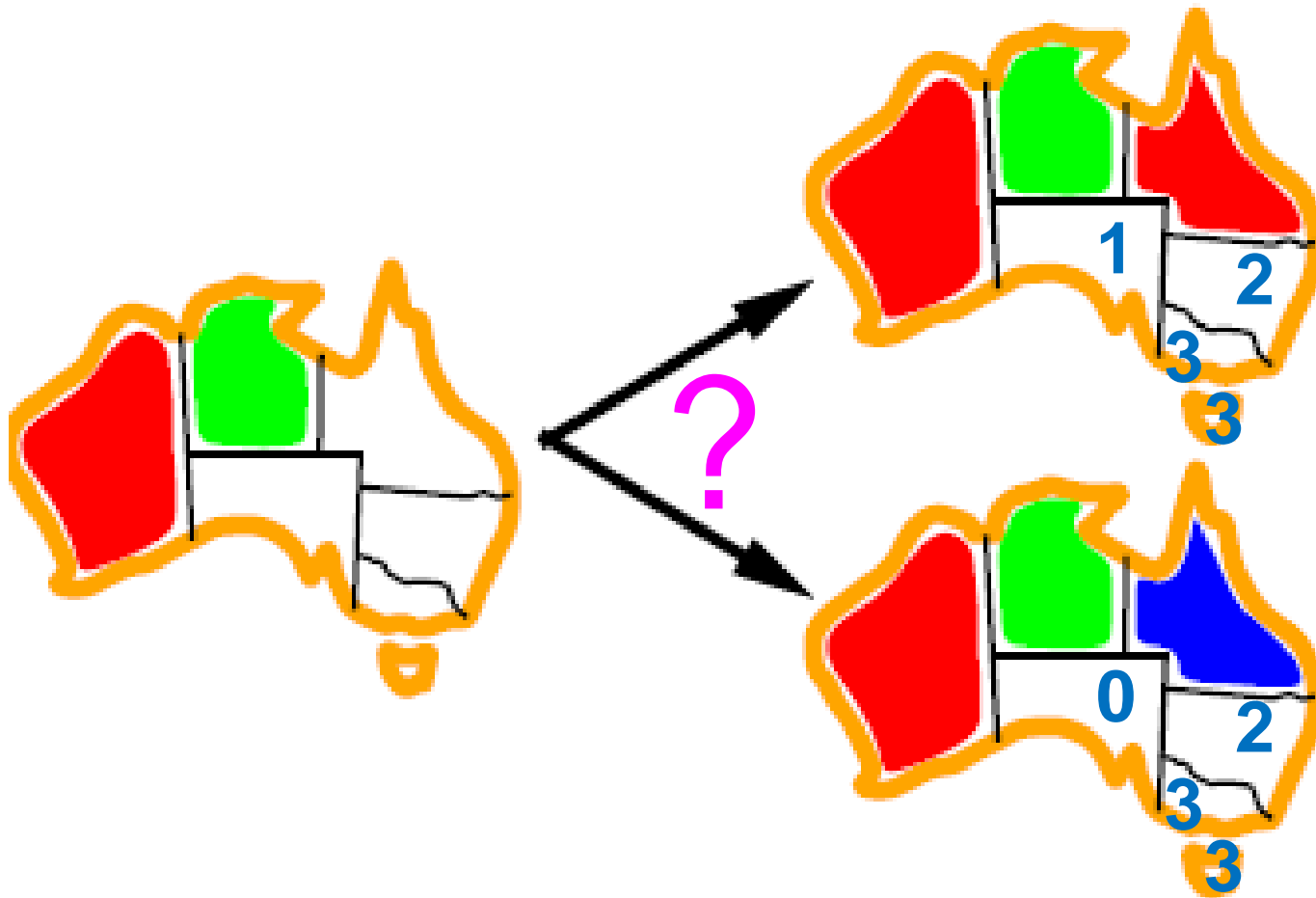
Degree Heuristic

Idea: Choose the variable with the most constraints on the remaining variables



Least Constraining Value (LCV)

Given a variable, choose the least constraining value: the one that rules out the fewest values in the remaining variables.



Interleaving Searching and Inference

Standard backtracking detects a failure only when it is assigning a variable with an empty domain.

Q: Can we detect an inevitable failure earlier?

■ **Forward checking:**




























- Idea: After each assignment, update the domains of all the neighbors of the assigned variable.
- If any domain becomes empty, then backtrack.

■ **Maintaining arc consistency:**

- Detects more upcoming failures than forward checking.
- A common method is to run **AC-3** (only updating the domains of unassigned variables) after each assignment.

Forward Checking Example



WA	NT	Q	NSW	V	SA	T
						
						
						
					Failure!	

Min-Conflicts Local Search for CSPs

- Start from a complete configuration (each variable assigned by randomly selecting a value from its domain).
- Iteratively select a variable (randomly), and reassign its value to minimize conflicts.
- Example (4 queens):

1			
2	Q	Q	
2			Q
1			

Q	3		
	2	Q	
	2		Q
	0		

Q		1	
		1	
		2	Q
	Q	1	

Q		Q	2
			2
			0
	Q		1

1		Q	
0			
1			Q
1	Q		

		Q	
Q			
			Q
	Q		

- This can solve million-queen problems!



CSPs and SATs

- SATs represent satisfiability problems. (In logic, satisfiability problems are concerned with whether a set of logical assignments (True or False) exists to make a Boolean formula evaluate to True.
- Many theoretical and practical computer science problems are related to SATs (e.g., logical circuit design).
- We will touch on SATs when we talk about logical reasoning. However, it is nice to know the connection first.
- Many approaches of CSPs, including backtracking search and local search, are commonly applied to solve SATs.