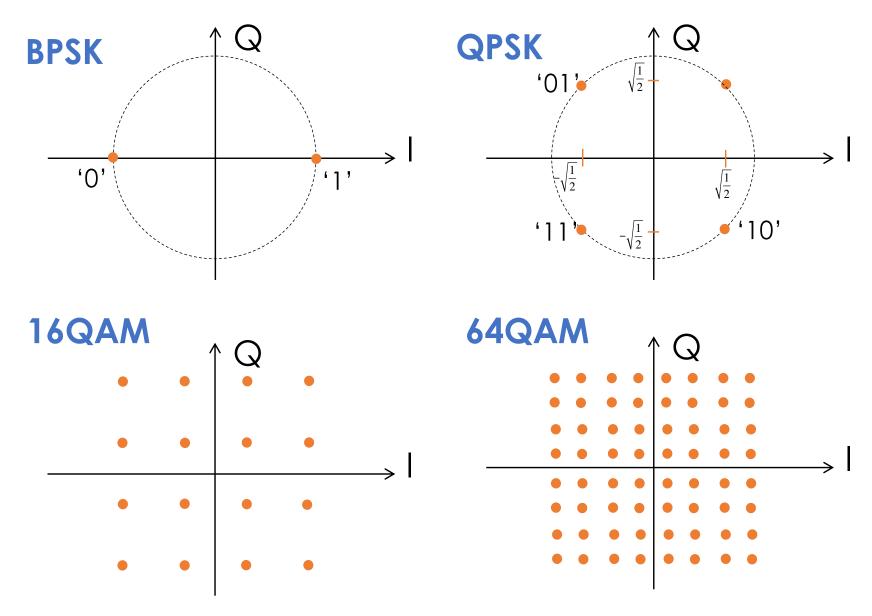
Wireless Communication Systems @CS.NCTU

Lecture 4: Rate Adaptation

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PSK and QAM



Agenda

- What is bit-rate adaptation?
- What are the challenges?
- Receiver-based bit-rate adaptation
- Transmitter-based bit-rate adaptation
- Bit-rate adaptation for multicast

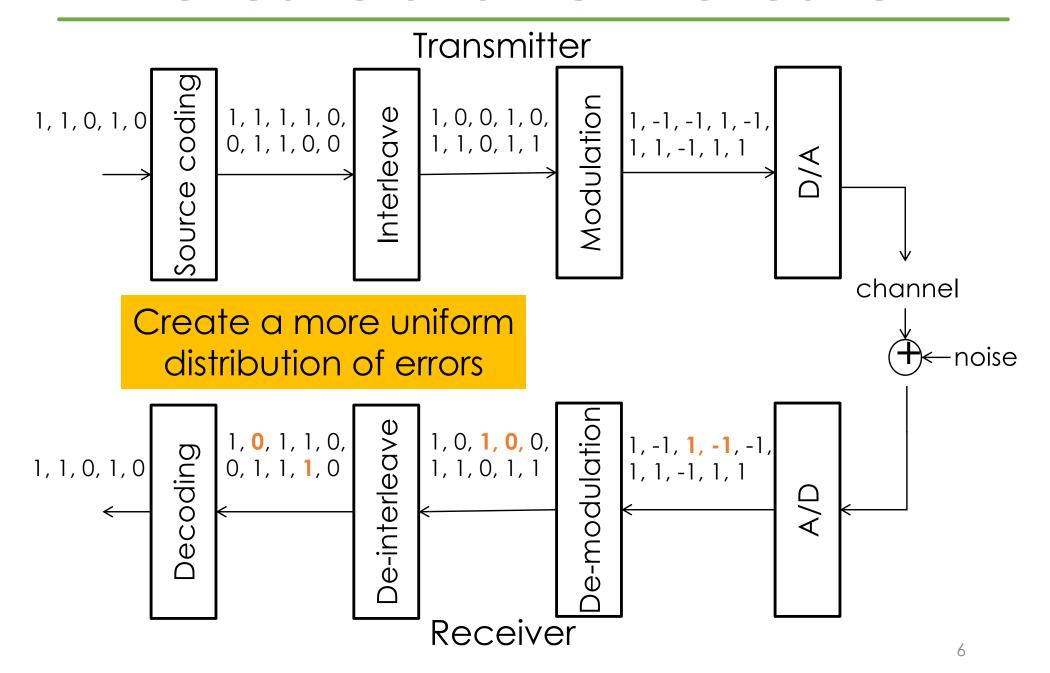
Bit-Rates in 802.11

Bit-	802.11	DSSS	Modulation	Bits	Coding	Mega-
rate	Stan-	or		per	Rate	Symbols
	dards	OFDM		Symbol		per
						second
1	b	DSSS	BPSK	1	1/11	11
2	b	DSSS	QPSK	2	1/11	11
5.5	b	DSSS	CCK	1	4/8	11
11	b	DSSS	CCK	2	4/8	11
6	a/g	OFDM	BPSK	1	1/2	12
9	a/g	OFDM	BPSK	1	3/4	12
12	a/g	OFDM	QPSK	2	1/2	12
18	a/g	OFDM	QPSK	2	3/4	12
24	a/g	OFDM	QAM-16	4	1/2	12
36	a/g	OFDM	QAM-16	4	3/4	12
48	a/g	OFDM	QAM-64	6	2/3	12
54	a/g	OFDM	QAM-64	6	3/4	12

Coding Rate

- Avoid random errors
 - 1/2: Add 1x redundant bits
 - -3/4: Add 1/3x redundant bits
- Haven't solved the problem yet
 - Data input: 1, 1, 0, 1, 0, 1, 1, 0, ...
 - After encoding: 1, 1, 1, 1, 0, 0, 1, 1, 1, 1, 0, 0,
 - Still one bit error → Suffer from burst errors

Interleave and De-interleave



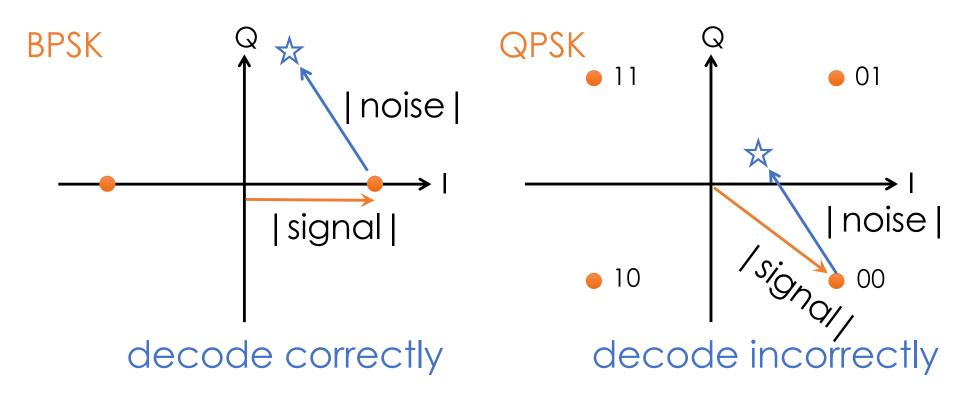
Channel Quality vs. Bit-Rate

- When channels are very good
 - Encode more digital bits as a symbol
- When channels are noisy
 - Encode fewer data bits as a sample

Why is it affected by the channel quality?

Error Probability vs. Modulations

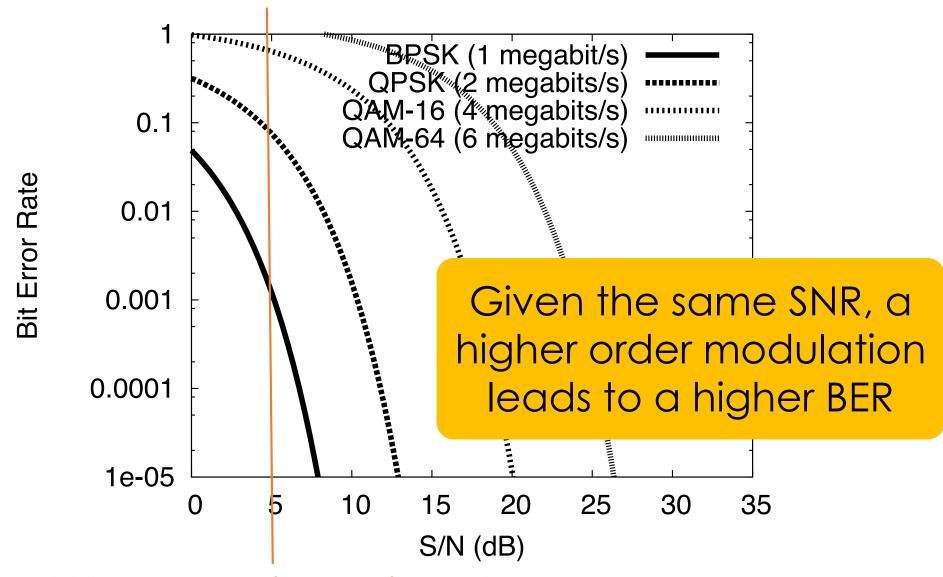
Given the same SNR



 $SNR = 10log10 (|signal|^2/|noise|^2)$

Given the same SNR, decodable for BPSK, but un-decodable for QPSK

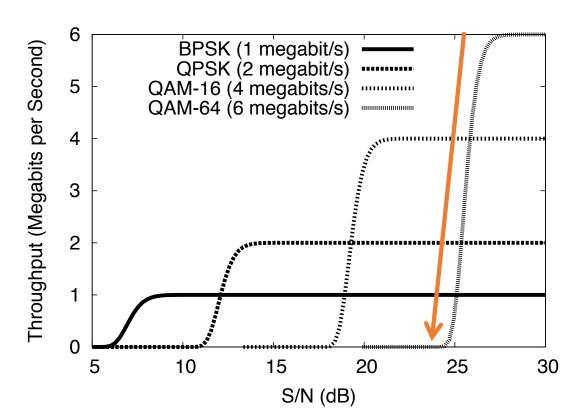
SNR vs. BER (Bit Error Rate)



802.11 operating region 5dB

SNR vs. PDR (Packet Delivery Ratio)

- In 802.11, a packet is received correctly if it passes the CRC check (all bits are correct)
 - Receive all or none
- Given an SNR value, BER and PDR change with bit-rates



 $PDR(r) = (1-BER(r))^n$

Throughput(r) = PDR(r) * r

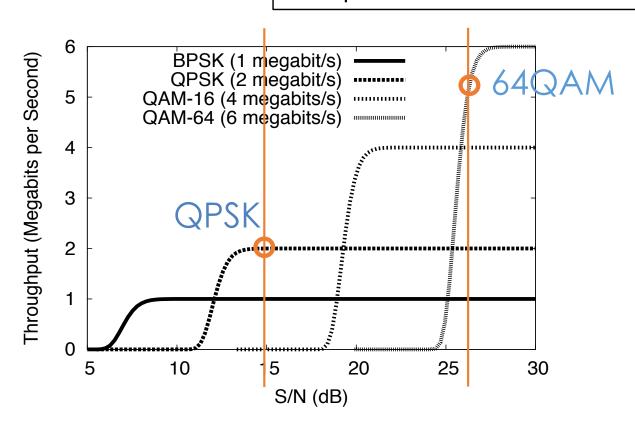
Throughput degrades quickly even with a small BER

Bit-Rate Selection

 Given the SNR, select the optimal bit-rate that achieves the highest throughput

$$r^* = \arg\min_r \mathrm{PDR}(r) * r$$

Ideal case without considering the protocol overhead



Difficulties with Rate Adaptation

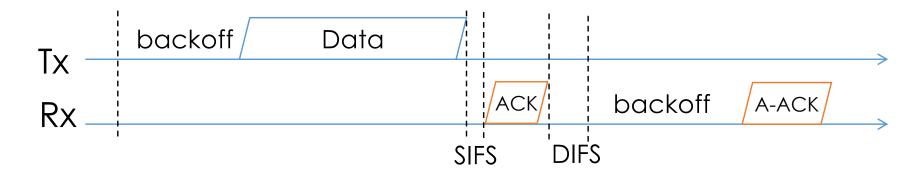
- Channel quality changes very quickly
 - Especially when the device is moving
- Can't tell the difference between
 - poor channel quality due to noise/interference/collision (high | noise |)
 - poor channel quality due to long distance (low | signal |)

Ideally, we want to decrease the rate due to low signal strength, but not interference/collisions

Types of Auto-Rate Adaptation

	Transmitter-based	Receiver-Based
SNR-based		RBAR, OAR, ESNR
ACK-based	ARF, AARF, ONOE	
Throughput-based	SampleRate, RRAA	
Partial packet		ZipTx
Soft information		SoftRate

Sync. ACK vs. Async ACK



Synchronous ACK

- Sent immediately after SIFS as a control frame (defined in 802.11)
- Cost the minimum overhead
- Only know whether the packet is transmitted correctly

Asynchronous ACK

- Sent as a data frame
- Cost additional overhead
- Can include more detailed information (e.g., error rate)

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Properties

Selected by Tx	Selected by Rx	
Sync. ACK	Async. ACK	
Less accurate	Higher overhead	

Rx-based Adaptation

Receiver Based Auto Rate (RBAR)

- The receiver measures the SNR of the RTS, and picks the optimal rate based on the SNR-to-rate lookup table
- Piggyback the selected rate in CTS

Opportunistic Auto Rate (OAR)

- Similar to RBAR, but consider the channel coherence time
- If the channel is good, opportunistically send more packets since the channel time of each frame is short

Pros

 More accurate since the Rx can measure the up-to-date channel condition

Cons

Rely on asynchronous ACK, causing a higher overhead

Tx-based Adaptation

SampleRate

- Default in Linux

RRAA

Robust Rate Adaption Algorithm

In common

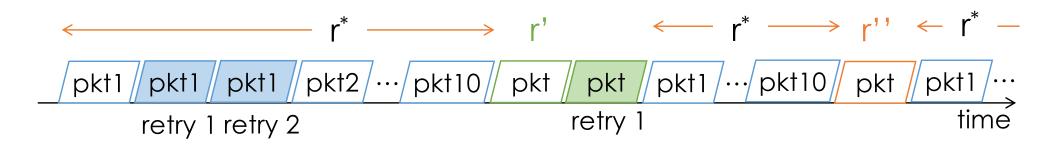
- Probe the packets at a rate not used currently
- See if switching to another rate gives a higher throughput

Differences

- Switch the rate by estimating the effective throughput
- Switch the rate by measuring the packet loss rate

SampleRate – Tx-based Adaptation

- Default in Linux
- Periodically send packets at a randomlysampled bit-rate other than the current bit-rate
 - Let r* be the current best rate
 - After sending 10 packets at the best rate, send a packet at a randomly-sampled rate
 - Estimate the achievable throughput of the sampled rates



SampleRate – Throughput Estimation

- How to estimate the effective throughput of a rate?
 - Calculate the transmission time of a L-bit packet
 - Consider packet length (I), bit-rate (r), number of retries (n), backoff time

$$T_{tx}(r, n, l) = T_{\text{DIFS}} + T_{\text{back off}}(n) + (n+1)(T_{\text{SIFS}} + T_{\text{ACK}} + T_{\text{header}} + l/r)$$

 Select the rate that has the smallest measured average transmission time to deliver a L-bit packet

$$r^* = \min_r T_{tx}(r, n, L)$$

SampleRate

- Do not sample the rates that
 - Have failed four successive times
 - Are unlikely to be better than the current one
- Is thought of the most efficient scheme for static environments
 - SNR, and thereby BER and best rate, do not change rapidly over time
- Waste channel time for sampling if the channel is very stable

RRAA – Tx-based Adaptation

- Robust Rate Adaption Algorithm
- Root causes of packet failures
 - Channel fading: mainly determined by the link distance
 - Random events: collisions, cross-technique interference (e.g., bluetooth or microwave)

Goal

- Robust against random loss: Should not switch the rate due to random channel variation
- Responsive to drastic channel changes: Should respond quickly to significant channel changes

S. Wong, H. Yang, S. Lu, V. Bharghavan, "Robust Rate Adaptation for 802.11 Wireless Networks," ACM MOBICOM, 2006

RRAA

- Use short-term loss ratio to assess the channel
 - Probe a window of N frames at a bit-rate
 - Estimate the loss ratio

$$P = \frac{\text{\# lost frames}}{\text{\# transmitted frame}}$$



- Stay unchanged if the loss ratio is acceptable
 - $-P_{min} < P < P_{max}$
- Switch the rate to
 - A higher one if $P < P_{min}$: imply that the channel is good enough to try the higher rate
 - A lower one if P > P_{max}: imply that the channel is too bad to use the current rate

RRAA – Parameter Configuration

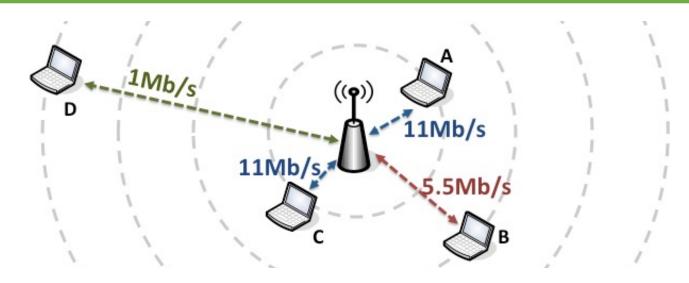
- P_{max}: Maximum tolerable loss threshold
 - the effective throughput of the current rate should be no worse than the loss-free throughput at a lower rate

$$(1-P_{\max}^r)\frac{l}{T_{rx}(r,n,l)}=\frac{l}{T_{rx}(r-1,n=1,l)}$$

$$\Rightarrow P_{\max}^r=1-\frac{T_{rx}(r,n,l)}{T_{rx}(r-1,n=1,l)}$$
 • P_{min}: Opportunistic rate Increase threshold

- - Harder to predict because we do not know how good is good enough
 - Heuristic: $P_{\min} = P_{\max}^{r+1}/\beta, \beta = 2$
- Window size N
 - Long enough to capture the minimum probability P_{min}

Rate Adaptation for Multicast



- Why it is difficult?
 - Can only assign a single rate to each packet
 - But the channel conditions of clients are different
- Possible Solutions
 - For reliable transmission: select the rate based on the worst node
 - For non-reliable transmission: provide clients heterogeneous throughput

Reliable Multicast Protocol

- Before rate adaptation, we should first ask:
 - How to efficiently collect ACK from multicast clients?
- Leader-based Protocol (LBP)
 - Select one of the receivers as the leader to reply ACK
 - Leader

if receive successfully, send ACK otherwise, send NACK

- Others

if receive successfully, do nothing otherwise, send NACK

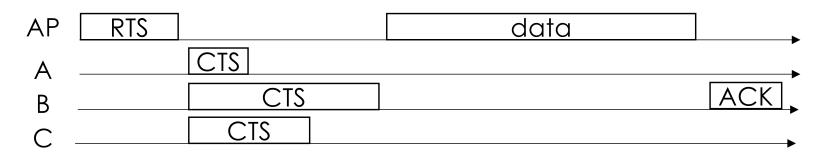
- Retransmit if the AP receives any NACK

J. Kuri and S. Kasera, "Reliable Multicast in Multi-Access Wireless LANs," IEEE INFOCOM, Mar. 1999.

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Rate Adaptation for <u>Data Multicast</u>

- Rate Adaptive Reliable Multicast (RAM)
 - Should pick the bit-rate based on the channel of the worst receiver
- Say we have three receivers A, B, and C
 - Each receiver feedbacks CTS at its optimal rate chosen based on its SNR
 - The AP detects the lowest rate by measuring the longest channel time occupied by CTS



Rate Adaptation for Video Multicast

- Video codec usually allows some losses
 - Receive more frames -> better video quality
 - Receive less frame → lower video quality
- No need to receive everything
 - No need to be constrained by the channel of the worst receiver
- One would expect a video quality proportional to its channel condition, i.e., differential QoS
 - Higher SNR → better video quality
 - Lower SNR → lower video quality

J. Villalon et. Al., "Cross-Layer Architecture for Adaptive Video Multicast Streaming over Multirate Wireless LANs," IEEE JSAC, vol. 25, no. 4, pp. 699-711, May 2007.

Rate Adaptation for Video Multicast

- H-ARSM (Hybrid Auto Rate Selection Mechanism)
- Mainly consider two video layers: base layer and

enhancement layer

Heuristic; not really optimizing for QoS/QoE

Design principles

- Guarantee a minimum video quality
 - Ensure that everyone reliably gets the base layer
 - Again, send at the rate according to the worst receiver
- Pick a more aggressive rate for the enhancement layer
 - Use the next higher rate if there exist one (or more)
 receivers with an SNR above the threshold of that rate

Recent Proposals

ZipTx

K. Lin, N. Kushman and D. Katabi, "Harnessing Partial Packets in 802.11 Networks," ACM MOBICOM, 2008

Exploit partial packets with consideration of bit-rate adaptation

SoftRate

M. Vutukuru, H. Balakrishnan and K. Jamieson, "Cross-Layer Wireless Bit Rate Adaptation," ACM SIGCOMM, 2009

Exploit soft information to improve selection accuracy

FARA

H. Rahul, F. Edalat, D. Katabi and C. Sodini, "Frequency-Aware Rate Adaptation and MAC Protocols," ACM MOBICOM, 2009

Adapt the bit-rate for every OFDM subcarrier

• ESNR

D. Halperin, W. Hu, A. Sheth and D. Wetherall, "Predictable 802.11 Packet Delivery from Wireless Channel Measurements", ACM SIGCOMM, 2010

Consider frequency selective fading