

Solution:

(1) Choose an undecidable problem that is not stated in terms of Turing Machines and provide a one-paragraph, less than half-page sketch of its undecidability proof.

Reference: "How is proving a context free language to be ambiguous undecidable?" *StackExchange*. 2011. <<https://cstheory.stackexchange.com/questions/4352/how-is-proving-a-context-free-language-to-be-ambiguous-undecidable>>

The undecidable problem I want to sketch the proof is, it is undecidable whether a context-free grammar is ambiguous.

The general idea is that we could reduce it from Post's Correspondence Problem (PCP). Suppose we can decide the language $L(G)$ generated by CFG G is ambiguous, we somehow construct a CFG and if the language is ambiguous then there is a derivation of some string in two different ways, but this leads to the result that we have a solution to PCP; similarly, if there is no ambiguity, then PCP cannot be solved. Therefore, we've reduced the original problem to PCP, and since PCP is undecidable, we can prove that it is undecidable whether a context-free grammar is ambiguous.

(2) Prove by induction that in a rooted tree, every node is reached by a DFS starting from the root.

A tree is a **connected** undirected graph with no simple circuits, and a rooted tree is a tree in which one vertex has been designated as the root and **every edge is directed away from the root**. A preorder traversal is a DFS that visits all tree nodes, given a rooted tree T and its root r (this rooted tree is not empty), preorder traversal works as follows:

def PreorderTraversal(T, r):

- 1 : visit r
- 2 : **for** each child c of r from left to right:
- 3 : T_c = subtree with c as its root
- 4 : PreorderTraversal(T_c, c)

Proof of Correctness

- Basis Step: when the rooted tree T only has one node, r , then preorder traversal visits r .
- Inductive Step: Suppose preorder traversal visits all nodes of a rooted tree T' , if T' is a subtree of another rooted tree T . Since T can be expressed as a root r with n children, each child i forms a subtree T_i . For any subtree, T_1, T_2, \dots, T_n , preorder traversal visits all nodes, therefore, PreorderTraversal(T, r) first visits r and then calls PreorderTraversal(T_i, r_i) for every subtree i , and PreorderTraversal(T, r) visits r and every node in its subtrees.

So every node is reached by a DFS starting from the root in a rooted tree. ■