6.842 Randomness and Computation

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Lectures on Random Walks

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1 Definitions

Definition 1. Let Ω be a set of states (throughout this note, Ω is finite). A sequence of random walks $X_0, X_1, \ldots \in \Omega$ is a *Markov chain* if it satisfies the *Markovian property*, i.e., for each $t \in \mathbb{N}$ and for all $x_1, \ldots, x_t, y \in \Omega$,

$$\mathbb{P}[X_{t+1} \mid X_1 = x_1, \dots, X_t = x_t] = \mathbb{P}[X_{t+1} = y \mid X_t = x_t].$$

WLOG, we assume that transitions are independent of time. For $x, y \in \Omega$, let

$$P(x,y) = \mathbb{P}\left[X_{t+1} = y \mid X_t = x\right].$$

Interpreted as a matrix, P is called the *transition matrix* of the Markov chain. We can also interpret the transition matrix P as a weighted directed graph with vertex set Ω such that the weight on $(i,j) \in \Omega^2$ equals P(i,j).

A random walk on a directed graph is a special case of Markov chains.

Definition 2. A random walk on a directed graph G = (V, E) is a sequence $S_1, S_2, \ldots \in V$ such that S_{t+1} is picked uniformly in $N^+(S_t)$, i.e., the transition matrix P is defined so that for $x, y \in V$,

$$P(x,y) = \begin{cases} \frac{1}{d^+(x)}, & \text{if } (x,y) \in E, \\ 0, & \text{otherwise.} \end{cases}$$

Definition 3. An $n \times n$ matrix P is called a *stochastic matrix* if for all $i \in [n]$,

$$\sum_{i=1}^{n} P(i,j) = 1.$$

For each $t \in \mathbb{N}$, let $P_t(x, y)$ ve the transition probability from x to y for t steps. Then for all $x, y \in \Omega$ and $t \in \mathbb{N}$,

$$P^{t}(x,y) = \begin{cases} P(x,y), & \text{if } t = 1, \\ \sum_{z \in \Omega} P(x,z) P^{t-1}(z,y) & \text{if } t > 1. \end{cases}$$

Interpreted as matrix multiplication, for each $t \in \mathbb{N}$ with t > 1,

$$P^t = P \cdot P^{t-1}.$$

Let $\pi^{(0)} = (\pi_1^{(0)}, \dots, \pi_n^{(0)})$ be the initial distribution, where $\pi_i^{(0)}$ is the probability of starting at vertex i for each $i \in [n]$. Let $\pi^{(t)}$ be the distribution after t steps for each $t \in \mathbb{N}$. For each $t \in \mathbb{N}$,

$$\pi^{(t)} = \pi^{(0)} P^t.$$

¹WLOG, we assume $\Omega = [n]$.

Definition 4. A distribution π^* is called a *stationary distribution* of a Markov chain with state set Ω and transition matrix P if for all $x \in \Omega$,

$$\pi^*(x) = \sum_{y \in \Omega} \pi^*(y) P(y, x).$$

Definition 5. A Markov chain with state set Ω and transition matrix P is said to be *irreducible* if for all $x, y \in \Omega$, there exists $t \in \mathbb{N}$ such that $P^t(x, y) > 0$.

Definition 6. A Markov chain with state set Ω and transition matrix P is said to be *aperiodic* if for all $x \in \Omega$,

$$\gcd \{ t \in \mathbb{N} : p^t(x, x) > 0] \} = 1.$$

Definition 7. A Markov chain with state set Ω and transition matrix P is said to be *ergodic* if there exists $t^* \in \mathbb{N}$ such that for all $t \in \mathbb{N}$ with $t > t^*$ and for all $x, y \in \Omega$, we have $P^t(x, y) > 0$.

Theorem 8. Every ergodic Markov chain has a unique stationary distribution.

In the special case of a random walk on an undirected graph G = (V, E), the stationary distribution $\pi^* = (\pi_1^*, \dots, \pi_n^*)$ is given by $\pi_i^* = d(i)/(2|E|)$ for all $i \in [n]$. Therefore, for a random walk on a d-regular graph or on a directed graph with each in-degree and each out-degree equal to d, the stationary distribution is uniform; this is not true in general directed graphs.

2 Hitting Time, Cover Time and Commute Time

Definition 9. Consider a random walk on a graph G = (V, E). For $x, y \in V$, the hitting time $H_{x,y}$ is defined to be the expected number of steps to go from x to y. For each $x \in V$, we call $H_{x,x}$ the recurrence time for x.

Theorem 10. Consider a random walk on a graph G = (V, E) with stationary distribution π^* . For each $x \in V$,

$$h_{x,x} = \frac{1}{\pi_*(x)}.$$

Proof sketch. Consider a very long walk. Then a $\pi^*(x)$ fraction of the positions are x. Then the average gap between the occurrences of x is $h_{x,x} = \pi^*(x)^{-1}$.

Definition 11. Consider a random walk on a graph G = (V, E). For $u \in V$, the cover time $C_u(G)$ is defined to be the expected steps from u to visit all states in Ω . Define $C(G) = \max_{u \in V} C_u(G)$.

Following are several examples of the cover time:

- $C(K_n) = \Theta(n \log n)$, where K_n is the complete graph on n vertices with a self-loop at each vertex. This can be proved by a coupon collector argument.
- $C(L_n) = \Theta(n^2)$, where L_n is the *n*-vertex line graph with a self-loop at each vertex.
- $C(\text{lollipop}_n) = \Theta(n^3)$, where lollipop_n is an n-vertex lollipop vertex formed by $L_{n/2}$ and $K_{n/2}$ joined at a vertex. This is illustrated in Figure 1.

Theorem 12. Let G be an undirected graph. Then²

$$C(G) \leq O(mn)$$
.

When the context is clear, we denote m = |E| in a graph G = (V, E).

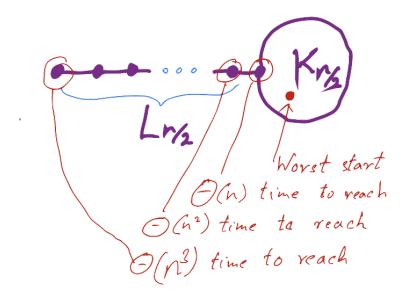


Figure 1: A lollipop graph $lollipop_n$ and its cover time.

Definition 13. Consider a random walk on a graph G = (V, E). For $x, y \in V$, the *commute time* $C_{x,y} = C_{x,y}(G)$ is defined to be the expected number of steps for the random walk to start at x, hit y and return to x.

Proposition 14. For $x, y \in V$,

$$C_{x,y} = h_{x,y} + h_{y,x}.$$

Proof. This is due to linearity of expectation.

Lemma 15. Consider a random walk on a connected undirected graph G = (V, E). For each $(x, y) \in E$,

$$C_{x,y} \leq O(m)$$
.

Proof. Construct a graph G' by adding a self-loop at each vertex with probability 1/2. Let $x, y \in V$. We claim that $C_{x,y}(G') = 2C_{x,y}(G)$. To see this, for each path from x to y in G', removing the self-loops in the path gives a path in G, and the expected fraction of self-loops in the path is 1/2. Then G' is ergodic. This implies that there exists a unique stationary distribution π^* .

Consider a walk u_1, u_2, \ldots , where $u_i \in V$ and $(u_i, u_{i+1}) \in E$ for each $i \in \mathbb{N}$. We look for commutes of the form

$$x \to y \to \ldots \to x \to y$$
.

For each $i \in \mathbb{N}$,

$$\mathbb{P}[u_i = x, u_{i+1} = y] = \mathbb{P}[u_i = x] \cdot \mathbb{P}[u_{i+1} = y \mid u_i = x] = \frac{d(x)}{2m} \cdot \frac{1}{d(x)} = \frac{1}{2m}.$$

Therefore, the expected fraction of $x \to y$ equals 1/(2m). This implies that the expected gap between the $(x \to y)$'s equals 2m. This proves that $C_{x,y}(G) = O(m)$.

Proof of Theorem 12. Let T be a spanning tree of G. Let $(v_0, v_1, \ldots, v_{2n-2})$ be a DFS traversal of T. For instance, (1, 2, 3, 2, 4, 2, 1, 5, 1) is a DFS traversal of the tree given in Figure 2. Then

$$C(G) \le \sum_{i=0}^{2n-3} h_{v_i, v_{i+1}} = \sum_{(x,y) \in E(T)} C_{x,y} \le (n-1) \cdot O(m) = O(mn).$$

This completes the proof.

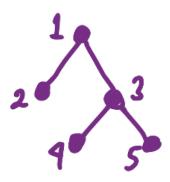


Figure 2: (1, 2, 3, 2, 4, 2, 1, 5, 1) is a DFS traversal of the tree in the figure.