6.842 Randomness and Computation April 11, 2022

Homework 4

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1. *Collaborators and sources:* Guanghao Ye, Zixuan Xu.

   Proof. □

2. (a) *Collaborators and sources:* none.

Proof. Note that $\mathbb{1}_{\text{test accepts}} = (1 + f(x)f(y)f(z))/2$. By the Fourier transform of f and by linearity of expecation,

$$\begin{aligned}
\mathbb{E}[f(x)f(y)f(z)] &= \mathbb{E}\left[\left(\sum_{S\subset[n]} \hat{f}(S)\chi_S(x)\right)\left(\sum_{T\subset[n]} \hat{f}(T)\chi_T(y)\right)\left(\sum_{U\subset[n]} \hat{f}(U)\chi_U(z)\right)\right] \\
&= \sum_{S,T,U\subset[n]} \hat{f}(S)\hat{f}(T)\hat{f}(U)\,\mathbb{E}\left[\chi_S(x)\chi_T(y)\chi_U(x\circ y\circ w)\right].
\end{aligned}$$

Let $S, T, U \subset [n]$. For all $i \in [n]$, since $x_i, y_i \in \{\pm 1\}$, then $x_i^2 = y_i^2 = 1$. Hence,

$$\begin{aligned}
\chi_S(x)\chi_T(y)\chi_U(x\circ y\circ w) &= \left(\prod_{i\in S} x_i\right)\left(\prod_{i\in T} y_i\right)\left(\prod_{i\in U} x_i y_i w_i\right) \\
&= \left(\prod_{i\in S\cap U} x_i^2\right)\left(\prod_{i\in T\cap U} y_i^2\right)\left(\prod_{i\in S\triangle U} x_i\right)\left(\prod_{i\in T\triangle U} y_i\right)\left(\prod_{i\in U} w_i\right) \\
&= \chi_{S\triangle U}(x)\chi_{T\triangle U}(y)\chi_U(w).
\end{aligned}$$

If $S = T = U$, since $w_1, \ldots, w_n$ are all chosen independently and since $\mathbb{E}[w_i] = (-1)\cdot\delta + 1\cdot(1-\delta) = 1 - 2\delta$ for all $i \in [m]$, then

$$\mathbb{E}\left[\chi_{S\triangle U}(x)\chi_{T\triangle U}(y)\chi_U(w)\right] = \mathbb{E}\left[\prod_{i\in S} w_i\right] = \prod_{i\in S}\mathbb{E}\left[w_i\right] = (1-2\delta)^{|S|}.$$

Now, suppose that either $S \neq U$ or $T \neq U$. WLOG assume that $S \neq U$. Then $S\triangle U \neq \emptyset$. Let $j \in S\triangle U$. For $x \in \{\pm 1\}^n$, let $x^{\oplus j}$ be the vector obtained by flipping the j^{th} bit in x. Then we can partition $\{\pm 1\}^n$ into (unordered) pairs $(x, x^{\oplus j})$. Therefore,

$$\begin{aligned}
\mathbb{E}\left[\chi_{S\triangle U}(x)\right] &= \frac{1}{2^n}\sum_{x\in\{\pm1\}^n}\chi_{S\triangle U}(x) = \frac{1}{2^n}\sum_{\text{pairs }(x,x^{\oplus j})}\left(\chi_{S\triangle U}(x) + \chi_{S\triangle U}\left(x^{\oplus j}\right)\right) \\
&= \frac{1}{2^n}\sum_{\text{pairs }(x,x^{\oplus j})}\left(x_j\prod_{i\in(S\triangle U)\setminus\{j\}} x_i + (-x_j)\prod_{i\in(S\triangle U)\setminus\{j\}} x_i\right) = 0.
\end{aligned}$$

Since x, y and w are chosen independently, then for all $S, T, U \subset [n]$ such that either $S \neq U$ or $T \neq U$,

$$\mathbb{E}\left[\chi_{S\triangle U}(x)\chi_{T\triangle U}(y)\chi_U(w)\right] = \mathbb{E}\left[\chi_{S\triangle U}(x)\right]\mathbb{E}\left[\chi_{T\triangle U}(y)\right]\mathbb{E}\left[\chi_U(w)\right] = 0.$$

Therefore,

$$\begin{aligned}
\mathbb{P}[\text{test accepts}] &= \mathbb{E}\left[\mathbb{1}_{\text{test accepts}}\right] = \mathbb{E}\left[\frac{1+f(x)f(y)f(z)}{2}\right] = \frac{1}{2} + \frac{1}{2}\mathbb{E}[f(x)f(y)f(z)] \\
&= \frac{1}{2} + \frac{1}{2}\sum_{S,T,U\subset[n]}\hat{f}(S)\hat{f}(T)\hat{f}(U)\,\mathbb{E}\left[\chi_{S\triangle U}(x)\chi_{T\triangle U}(y)\chi_U(w)\right] \\
&= \frac{1}{2} + \frac{1}{2}\sum_{S\subset[n]}(1-2\delta)^{|S|}\hat{f}(S)^3.
\end{aligned}$$

This completes the proof. □

(b) *Collaborators and sources:* none.

Proof. Let $f : \{\pm 1\}^n \to \{\pm 1\}$ be a dictator function. Then $f = \chi_{\{j\}}$ for some $j \in [n]$. Therefore, $\hat{f}(\{j\}) = 1$ and $\hat{f}(S) = 0$ for all $S \subset [n]$ with $S \neq \{j\}$. By part (a),

$$\begin{aligned}
\mathbb{P}[\text{test accepts}] &= \frac{1}{2} + \frac{1}{2} \sum_{S \subset [n]} (1-2\delta)^{|S|} \hat{f}(S)^3 \\
&= \frac{1}{2} + \frac{1}{2} \left((1-2\delta)^{|\{j\}|} \hat{f}(\{j\})^3 + \sum_{\substack{S \subset [n] \\ S \neq \{j\}}} (1-2\delta)^{|S|} \hat{f}(S)^3 \right) \\
&= \frac{1}{2} + \frac{1}{2} \left((1-2\delta)^1 \cdot 1^3 + \sum_{\substack{S \subset [n] \\ S \neq \{j\}}} (1-2\delta)^{|S|} \cdot 0^3 \right) \\
&= \frac{1}{2} + \frac{1}{2}(1-2\delta) = 1 - \delta.
\end{aligned}$$

This completes the proof. □

(c) *Collaborators and sources:* none.

Proof. Let $f : \{\pm 1\}^n \to \{\pm 1\}$ be such that f passes with probability at least $1 - \varepsilon$ for some $\varepsilon \in (0, 1/2)$. By part (a),

$$1 - \varepsilon \leq \mathbb{P}[\text{test accepts}] = \frac{1}{2} + \frac{1}{2} \sum_{S \subset [n]} (1 - 2\delta)^{|S|} \hat{f}(S)^3.$$

Rearranging the above inequality and applying Parseval's identity yield

$$\begin{aligned} 1 - 2\varepsilon &\leq \sum_{S \subset [n]} (1 - 2\delta)^{|S|} \hat{f}(S)^3 \leq \left(\max_{S \subset [n]} (1 - 2\delta)^{|S|} \hat{f}(S) \right) \sum_{S \subset [n]} \hat{f}(S)^2 \\ &= \left(\max_{S \subset [n]} (1 - 2\delta)^{|S|} \hat{f}(S) \right) \cdot 1 = \max_{S \subset [n]} (1 - 2\delta)^{|S|} \hat{f}(S). \end{aligned}$$

Hence, there exists $S \subset [n]$ such that $(1-2\delta)^{|S|} \hat{f}(S) \geq 1-2\varepsilon$. Set $\delta = \varepsilon$ in the test. Then $(1 - 2\varepsilon)^{|S|} \hat{f}(S) \geq 1 - 2\varepsilon$. Since $\varepsilon \in (0, 1/2)$, then $1 - 2\varepsilon \in (0, 1)$, so $(1 - 2\varepsilon)^{|S|} \in (0, 1]$. Therefore,

$$\hat{f}(S) \geq \frac{1 - 2\varepsilon}{(1 - 2\varepsilon)^{|S|}} \geq \frac{1 - 2\varepsilon}{1} = 1 - 2\varepsilon.$$

This completes the proof. $\square$

(d) *Collaborators and sources:* none.

By part (c), if f passes with probability at least $1-\varepsilon$ for some $\varepsilon \in (0,1/2)$, then there exists $S \subset [n]$ such that $(1-2\varepsilon)^{|S|}\hat{f}(S) \geq 1-2\varepsilon$ by setting $\delta = \varepsilon$ in the test. Since $\operatorname{dist}(f,\chi_S) \in [0,1]$, then $\hat{f}(S) = 1 - 2\operatorname{dist}(f,\chi_S) \in [-1,1]$. Since $\varepsilon \in (0,1/2)$, then $1-2\varepsilon \in (0,1)$. If $|S| \geq 2$, then $0 < (1-2\varepsilon)^{|S|} < 1-2\varepsilon$, so $(1-2\varepsilon)^{|S|}\hat{f}(S) < 1-2\varepsilon$, a contradiction. Therefore, one of the following two cases holds:

(i) $|S| = 1$ and $\hat{f}(S) = 1$ (so $\operatorname{dist}(f,\chi_S) = 0$, and $f = \chi_S$ is a dictator function);

(ii) $|S| = 0$ and $\hat{f}(S) \geq 1-2\varepsilon$ (so $\operatorname{dist}(f,\chi_\emptyset) \leq \varepsilon$).

Hence, if f is ε-close to $\chi_\emptyset \equiv 1$ (a non-dictator function), then f also passes with probability at least $1-\varepsilon$.

Note that for any dictator function, say $\chi_{\{j\}}$ for some $j \in [n]$,

$$\mathop{\mathbb{P}}_{x\in\{\pm1\}^n}\left[\chi_{\{j\}}(x) = 0\right] = \mathop{\mathbb{P}}_{x\in\{\pm1\}^n}[x_j = 0] = \frac{|\{x \in \{\pm1\}^n : x_j = 0\}|}{2^n} = \frac{2^{n-1}}{2^n} = \frac{1}{2}.$$

In other words, any dictator function equals 0 for half of the inputs, and 1 for the other half. We give a simple fix to the test by applying the following new test before the original test. For any sufficiently small $\eta > 0$, we independently and uniformly sample $\Theta(\log(1/\eta))$ random inputs from $\{\pm1\}^n$, and reject if and only if more than $3/4$ of the values are 1. If f is ε-close to $\chi_\emptyset \equiv 1$ for some $\varepsilon \in (0,1/8)$, then by the Chernoff bound,

$$\mathbb{P}[\text{new test rejects } f] = 1 - \mathbb{P}[\leq 3/4 \text{ of the values are } 1] \geq 1 - e^{-\Theta(\log(1/\eta))} = 1 - \Theta(\eta).$$

On the other hand, if f is a dictator function, then by the Chernoff bound,

$$\mathbb{P}[\text{new test accepts } f] = 1 - \mathbb{P}[> 3/4 \text{ of the values are } 1] \geq 1 - e^{-\Theta(\log(1/\eta))} = 1 - \Theta(\eta).$$

Hence, if f passes the combination of the new test and the original test with probability at least $1-\varepsilon$ and with $\delta = \varepsilon$ in the original test for some sufficiently small $\varepsilon > 0$, then f is a dictator function with probability at least $1-\Theta(\eta)$; on the other hand, if f is a dictator function, then the union bound implies that f passes the combined test with probability at least $1-\Theta(\eta)-\delta$. This shows that the combined test is a dictator test.