最小编辑距离算法 Minimum Edit Distance

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编辑距离

编辑前字符串s

编辑后字符串t

编辑操作p:插入、删除、替换

"编辑距离"定义为"编辑操作的次数"

源文: She is a star with the theatre company.

机器译文: 她是与剧院公司的一颗星。

参考译文: 她是剧团的明星。

编辑距离: 6

删除次数(4次): 与公司 一颗

替换次数(2次): 剧院→剧团 星→明星

计算机器译文 跟正确答案之 间的距离

如何计算最小编辑距离

原始串 s o t

目标串 stop

插入操作的权值

(insertCost): 1

删除操作的权值

(deleteCost) : 1

替换操作的权值

(substituteCost): 2

编辑操作1 s o t

→ stot (1.插入, 1分, 累计1分)

→ stop (2.替换,2分,累计3分)

编辑距离: 3

s o t

编辑操作2

 \rightarrow stt

(1.替换,2分,累计2分)

→ s t o (2.替换, 2分, 累计4分)

→ stop (3.插入,1分,累计5分)

编辑距离:5

最小编辑距离计算: 动态规划

```
i,目标串字符位置
D(0, 0) = 0
D(i, 0) = insertCost * i
                                                               i, 原始串字符位置
D(0, j) = deleteCost * j
D(i, j) = \min \begin{cases} D(i-1, j) + insertCost(target_i) \\ D(i-1, j-1) + substituteCost(source_j, target_i) \\ D(i, j-1) + deleteCost(source_j) \end{cases}
          substituteCost = 0 if target[i] = source[j]
                                2 otherwise
          insertCost = 1
```

deleteCost = 1

最小编辑距离算法描述

```
function Min-Edit_Distance (target, source)
n = length(target);
m = length(source);
create distance matrix d[n,m];
d[0,0]=0;
d[0,1]=1,...d[0,m]=m;
d[1,0]=1,...d[n,0]=n;
for each i from 1 to n do
   for each j from 1 to m do
         d[i, j] = min(d[i-1, j] + insertCost(target_i)),
                    d[i-1, j-1] + substituteCost(source_i, target_i),
                    d[i, j-1] + deleteCost(source<sub>i</sub>));
return d[n,m];
```

source: s o t

target: s t o p

n = length (target)

m = length (source)

Create matrix d [n, m];

 $i=0 \ j=0$

d[0,0] = 0; d[0,1] = 1; ...; d[0,m] = m;d[1,0] = 1; ...; d[n,0] = n;

$j \uparrow$	3	t				
	2	O				
	1	S				
	0	#	S	t	O	p
	#	0	1	2	3	4

source: s o t

target: s t o p

n = length (target)

m = length (source)

j	3	t				
	2	О				
	1	S	0			
	0	#	S	t	O	p
	#	0	1	2	3	4
Ī						

$$i=1$$
 $j=1$

$$d[1,1] = min \begin{cases} d[0,1] + insert(t[1]) = 2 \\ d[0,0] + substitute(s[1],t[1]) = 0 \\ d[1,0] + delete(s[1]) = 2 \end{cases} = 0$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

3	3	t				
2		O	1			
1	-	S	0			
C)	#	S	t	O	p
#	ŧ	0	1	2	3	4

$$i=1 j=2$$

$$d[1,2] = min \begin{cases} d[0,2] + insert(t[1]) = 3 \\ d[0,1] + substitute(s[2],t[1]) = 3 \\ d[1,1] + delete(s[2]) = 1 \end{cases} = 1$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

j	3	t	2			
	2	О	1			
	1	S	0			
	0	#	S	t	O	p
	#	0	1	2	3	4
Ī						

$$i=1 j=3$$

$$d[1,3] = min \begin{cases} d[0,3] + insert(t[1]) = 4 \\ d[0,2] + substitute(s[3],t[1]) = 4 \\ d[1,2] + delete(s[3]) = 2 \end{cases} = 2$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

j	3	t	2			
	2	О	1			
	1	S	0	1		
	0	#	S	t	O	p
	#	0	1	2	3	4
Ţ						

$$i=2 j=1$$

$$d[2,1] = min \begin{cases} d[1,1] + insert(t[2]) = 1 \\ d[1,0] + substitute(s[1],t[2]) = 3 \\ d[2,0] + delete(s[1]) = 3 \end{cases} = 1$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

j	3	t	2			
	2	О	1	2		
	1	S	0	1		
	0	#	S	t	O	p
	#	0	1	2	3	4
Ī						

$$i=2$$
 $j=2$

$$d[2,2] = min \left\{ \begin{array}{l} d[1,2] + insert(t[2]) = 2 \\ d[1,1] + substitute(s[2],t[2]) = 2 \\ d[2,1] + delete(s[2]) = 2 \end{array} \right\} = 2$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

j	3	t	2	1		
	2	О	1	2		
	1	S	0	1		
	0	#	S	t	O	p
	#	0	1	2	3	4

$$i=2 j=3$$

$$d[2,3] = min \left\{ \begin{array}{l} d[1,3] + insert(t[2]) = 3 \\ d[1,2] + substitute(s[3],t[2]) = 1 \\ d[2,2] + delete(s[3]) = 3 \end{array} \right\} = 1$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

3	t	2	1		
2	O	1	2		
1	S	0	1	2	
0	#	S	t	O	p
#	0	1	2	3	4

$$i=3 j=1$$

$$d[3,1] = min \left\{ \begin{array}{l} d[2,1] + insert(t[3]) = 2 \\ d[2,0] + substitute(s[1],t[3]) = 4 \\ d[3,0] + delete(s[1]) = 4 \end{array} \right\} = 2$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

$j \uparrow$	3	t	2	1		
	2	О	1	2	1	
	1	S	0	1	2	
	0	#	S	t	O	p
	#	0	1	2	3	4
_						

$$i=3 j=2$$

$$d[3,2] = min \left\{ \begin{array}{l} d[2,2] + insert(t[3]) = 3 \\ d[2,1] + substitute(s[2],t[3]) = 1 \\ d[3,1] + delete(s[2]) = 3 \end{array} \right\} = 1$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

j	3	t	2	1	2	
	2	О	1	2	1	
	1	S	0	1	2	
	0	#	S	t	O	p
	#	0	1	2	3	4
<u>!</u>	-					

$$i=3$$
 $j=3$

$$d[3,3] = min \begin{cases} d[2,3] + insert(t[3]) = 2 \\ d[2,2] + substitute(s[3],t[3]) = 4 \\ d[3,2] + delete(s[3]) = 2 \end{cases} = 2$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

]	3	t	2	1	2	
	2	O	1	2	1	
	1	S	0	1	2	3
	0	#	S	t	О	p
	#	0	1	2	3	4

$$i=4$$
 $j=1$

$$d[4,1] = min \left\{ \begin{array}{l} d[3,1] + insert(t[4]) = 3 \\ d[3,0] + substitute(s[1],t[4]) = 5 \\ d[4,0] + delete(s[1]) = 5 \end{array} \right\} = 3$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

$j \uparrow$	3	t	2	1	2	
	2	О	1	2	1	2
	1	S	0	1	2	3
	0	#	S	t	O	p
	#	0	1	2	3	4
<u></u>						

$$i=4$$
 $j=2$

$$d[4,2] = min \begin{cases} d[3,2] + insert(t[4]) = 2 \\ d[3,1] + substitute(s[2],t[4]) = 4 \\ d[4,1] + delete(s[2]) = 4 \end{cases} = 2$$

source: s o t

target: s t o p

n = length (target)

m = length (source)

j	3	t	2	1	2	3
	2	О	1	2	1	2
	1	S	0	1	2	3
	0	#	S	t	О	p
	#	0	1	2	3	4
<u>!</u>						-

$$i=4$$
 $j=3$

$$d[4,3] = min \left\{ \begin{array}{l} d[3,3] + insert(t[4]) = 3 \\ d[3,2] + substitute(s[3],t[4]) = 3 \\ d[4,2] + delete(s[3]) = 3 \end{array} \right\} = 3$$

```
s o t 编辑操作①
s t o t (1. 插入t, 1分, 累计1分)
s t o p (2. t替换p, 2分, 累计3分)
```

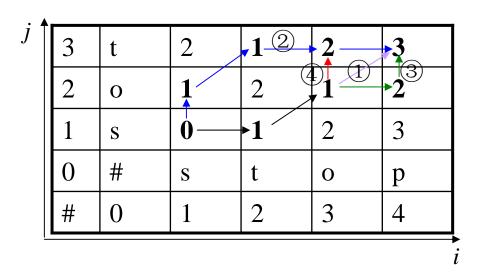
```
sot 编辑操作②

s <sup>↑</sup> t (1. 删除o, 1分, 累计1分)

s to (2. 插入o, 1分, 累计2分)

s top(3. 插入p, 1分, 累计3分)
```

3 t 0 p	(3. 1四/ c b)	1月,
s o t		编辑操作③
sto_t	(1. 插入t,	1分,累计1分)
stopt	(2. 插入p,	1分,累计2分)
stop	(3. 删除t,	1分,累计3分)



```
s o t 编辑操作④
s t o t (1. 插入t, 1分, 累计1分)
s t o (2. 删除t, 1分, 累计2分)
s t o p (3. 插入p, 1分, 累计3分)
```

最小编辑距离计算练习

intention → execution

```
intention inten* tion
\downarrow\downarrow\downarrow\downarrow\downarrow\downarrow
execution
ssss
2222=10
inten* tion
\downarrow\downarrow\downarrow\downarrow\downarrow
*execution
dsssi
12221=8
```

最小编辑距离计算练习

n	9	8	9	10	11	12	11	10	9	8
0	8	7	8	9	10	11	10	9	8	9
i	7	6	7	8	9	10	9	8	9	10
t	6	5	6	7	8	9	8	9	10	11
n	5	4	5	6	7	8	9	10	11	10
е	4	3	4	5	6	7	8	9	10	9
t	3	4	5	6	7	8	7	8	9	8
n	2	3	4	5	6	7	8	7	8	7
i	1	2	3	4	5	6	7	6	7	8
#	0	1	2	3	4	5	6	7	8	9
	#	е	X	е	С	u	t	i	0	n

参考文献

 Daniel Jurafsky & James H. Martin, 2000, Speech and Language Processing: An Introduction to Natural Language Processing, Computational Linguistics, and Speech Recognition, Chapter 5, section 5.6, pp153-156, Prentice-Hall Inc..