Presentation CS595

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Spectre Attacks: Exploiting Speculative Execution

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Spectre attack

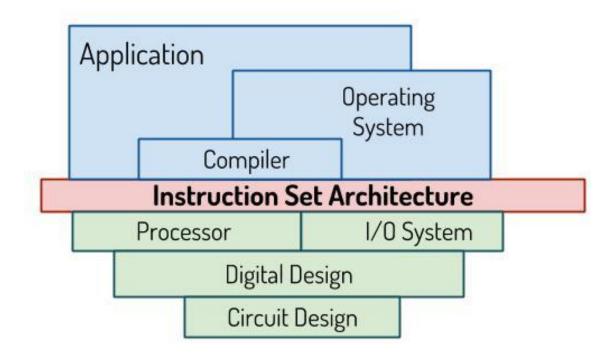
• Vulnerabilities in modern computers





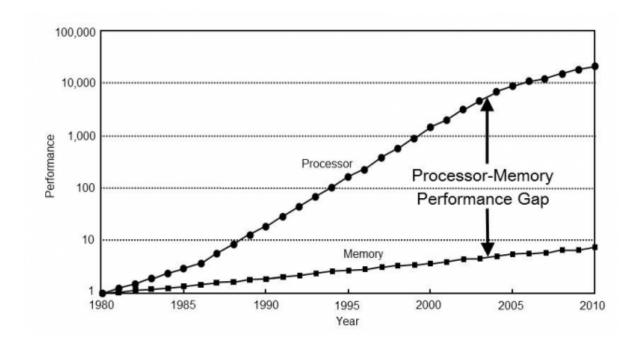
Background: Computer architecture

- Instruction set architecture (ISA)
- Execution order
 - In-order
 - Out-of-order
 - Re-order buffer (ROB)



Background: To boost CPU performance

- Improve hardware
- Design optimization
 - Speculative execution
- Memory hierarchy
 - CPU-memory
 - Caches



Spectre attack overview

- Conditional branch example
 - array1 = [1,2]
 - array2 = [x, y, z]
 - x is input under the attacker's control

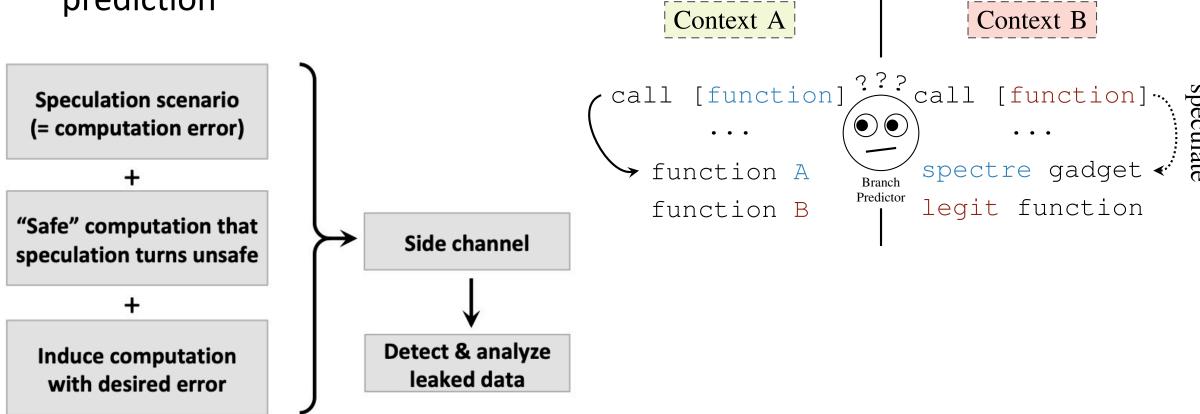
Spectre attack overview

- Attacker read array2[i*4096]
 - Find i=3 is fast
 - array[x] = k has been cached
 - Then secret k = 3 is revealed
- Side channel attack

array2[0 * 4096]	
array2[1* 4096]	
array2[2 * 4096]	
array2[3 * 4096]	cached
array2[4 * 4096]	

Spectre attack: Brach prediction

Attacker can misdirect the prediction



Mitigation options

- Not do speculative execution
 - Trade-off between performance and security
- Preventing access to secret data
 - Add new data dependencies
- Add hardware to hide speculative execution
- Not all speculative loads leak secret

Summary

- Spectre attacks leverage the speculative execution.
- Software security depends on having a clear common understanding between hardware and software.
- Trade-off between security and performance is always a problem.

Thank you