

## Chap 4 Roots of Nonlinear Equations

Nonlinear problem

$$f(x) = 0$$

Ex:  $J_m(\omega_{k,m} r) \cos(m\theta) \cos(c\omega_{k,m} t)$

Idea: Replace a problem you can't solve with an approximate one you can

Conn: Infinite iteration even in exact arithmetics

Linear problem

$$g(x) = 0$$

$$Ax = b \quad (3x + 2 = 0)$$

Replace a problem you can't solve with one that you can solve  
(e.g.  $Ax = b \Rightarrow LUx = b$ )

Finite step

## § 4.1 The rootfinding problem

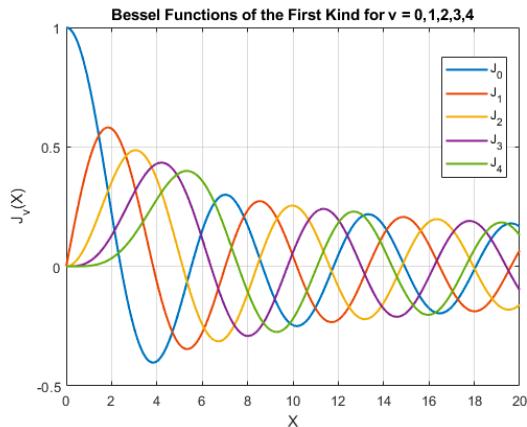
- Goal: Find  $r$ , such that  $f(r) = 0$   
 $\uparrow$   
root                     $\uparrow$   
a continuous scalar function ( $R'$ )

- MATLAB: "fzero"

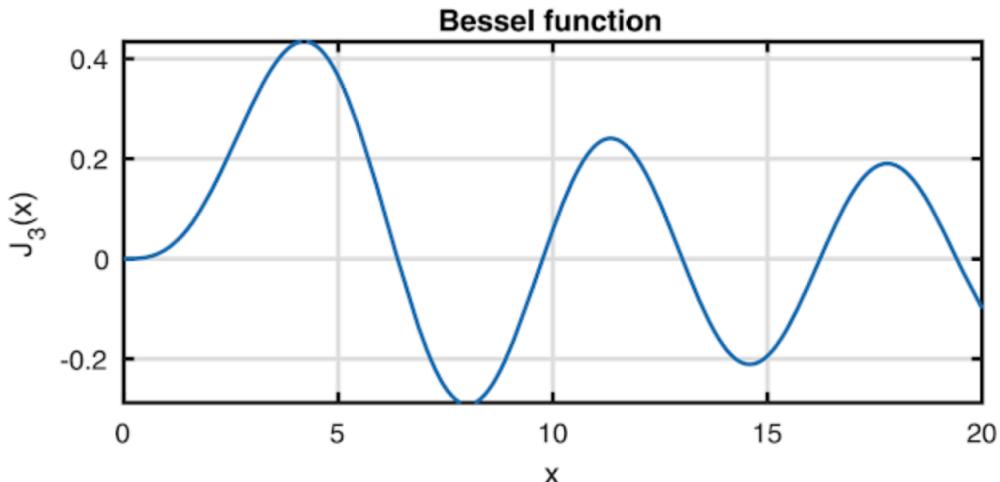
- Example: 4.1.1

- Displacement of the drumhead

- Bessel function of the first kind  
(Hard to find the root!)



```
J3 = @(x) besselj(3,x);  
fplot(J3,[0 20]), grid on
```



From the graph we see roots near 6, 10, 13, 16, and 19. We use `fzero` to find these roots accurately.

```
omega = [];  
for guess = [6,10,13,16,19]  
    omega = [omega; fzero(J3,guess)];  
end  
omega
```

## Conditioning, Error, Residual

- The perturbated problem

$$\text{Eq. } f(x)$$

$$\Rightarrow \tilde{f}(x) = f(x) + \varepsilon$$

perturbation

$$\text{Root } r \text{ s.t. } f(r) = 0$$

$$\Rightarrow \tilde{r} = r + \delta \text{ s.t. } \tilde{f}(\tilde{r}) = 0$$

- condition number

$$\kappa_r = \left| \frac{\delta}{\varepsilon} \right| \text{ as } \varepsilon \rightarrow 0$$

change in output  
change in input

By Taylor series

$$= 0$$

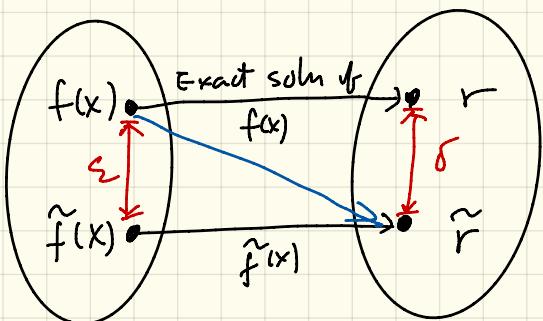
$$0 = f(r + \delta) + \varepsilon = f(r) + f'(r)\delta + \varepsilon + O(\delta^2)$$

Small number  
can be approximated  
by 0

与  $r$   
有關

$$\Rightarrow \kappa_r = \left| \frac{\delta}{\varepsilon} \right| \approx \left| \frac{1}{f'(r)} \right|$$

Recall that the absolute condition number for the evaluation of  $f$  at  $x$  is simply  $|f'(x)|$ .  
*The condition number of the rootfinding problem is equivalent to that of evaluating the inverse function.* The influence of  $|f'(r)|$  on rootfinding is easily illustrated.



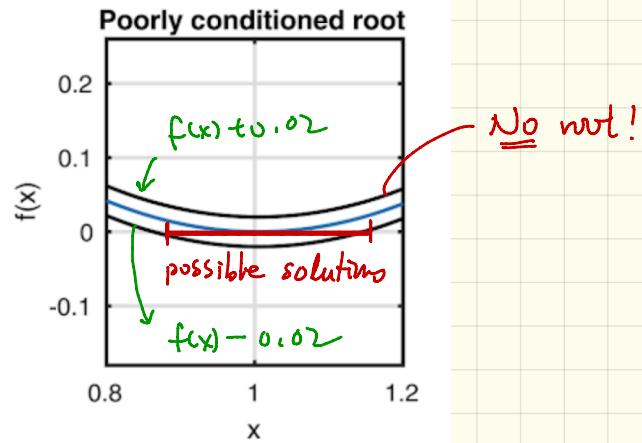
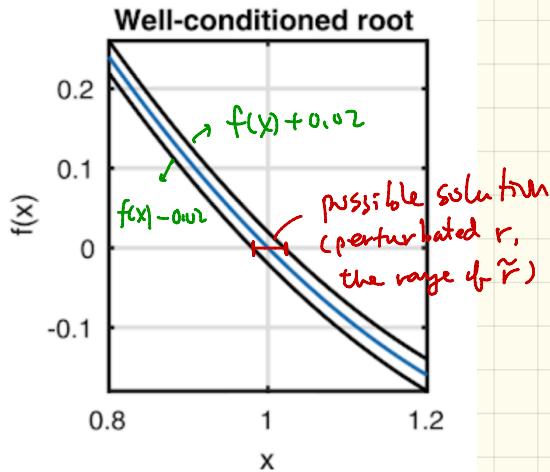
## Example 4.1.2

$$(1) f(x) = (x-1)(x-2)$$

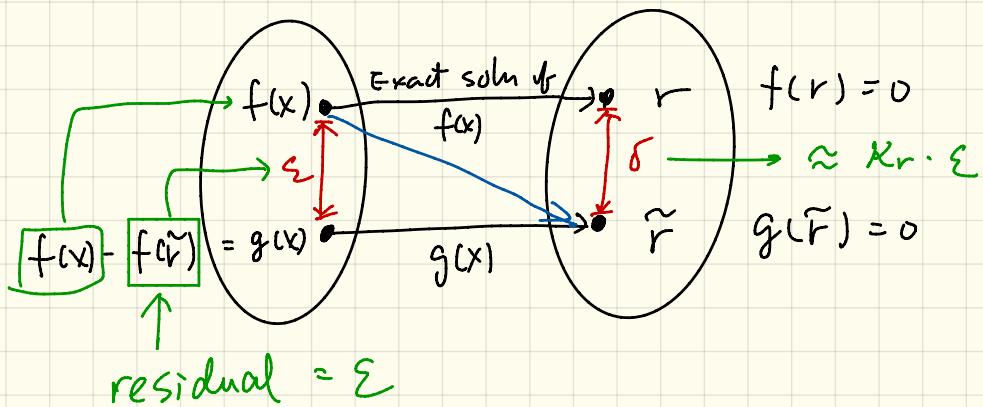
$$\Rightarrow r=1, f'(r)=-1, \kappa_r=1$$

$$(2) f(x) = (x-1)(x-1.01)$$

$$\Rightarrow r=1, f'(1) = -0.01, \kappa_r=100$$



- When  $|f'(r)|$  is small at the **root** (the cond. # is big)  
we may not get small computed error.
- Residual:  $f(\tilde{r}) \rightarrow$  Computable!



Small residual (i.e. small backward error)

## Multiple roots

Let  $r$  be a root of  $f(x)$  and  $g_f(x) = \frac{f(x)}{(x-r)}$ .

We have  $f(x) = (x-r) \cdot g_f(x)$

$$\Rightarrow f'(x) = (x-r) g'_f(x) + g_f(x)$$

$$\Rightarrow f'(r) = g(r) \quad (\text{By L'Hopital's rule, } g(x) \text{ is well-defined at } x=r \text{ as long as } f'(x) \text{ is well-defined})$$

[H(W)]

Claim:  $r$  is a simple root ( $\Leftrightarrow f'(r) \neq 0$ )

(i.e.  $g(r) \neq 0$ )

claim: Suppose  $f'(r) = g_f(r) = 0$  and  $g(x)$  is differentiable.

We have  $r$  is not a simple root of  $g_f(x) \Leftrightarrow g'_f(r) \neq 0$

since  $f''(x) = (x-r) f''(x) + 2f'(x)$

$$\Rightarrow f''(r) = 2f'(r)$$

$\Rightarrow f''(r) = 0 \Leftrightarrow r$  is a multiple root of  $f$

Thm Let  $r$  be a root of multiplicity  $m$  if

$$f(r) = f'(r) = \dots = f^{(m-1)}(r) = 0 \text{ and } f^{(m)}(r) \neq 0$$

Another view

$$f(x) = \underbrace{f(r)}_{r \text{ is simple root}} + \underbrace{f'(r)(x-r)}_{r \text{ is a double root}} + \frac{f''(r)}{2} (x-r)^2 + \dots + \frac{f^{(n)}(r)}{n!} (x-r)^n + \dots$$

$\Rightarrow f'(r) \neq 0$

$r$  is a double root

$$\Rightarrow \frac{f''(r)}{2} \neq 0$$

Recall  $\lambda_r = \frac{1}{|f'(r)|}$

If  $r_1$  is simple ( $m=1$ ): Ill-conditioned if  $f'(r_1)$  small

$$r_2 \approx r_1$$

$$\Leftrightarrow g(r_1) \approx 0$$

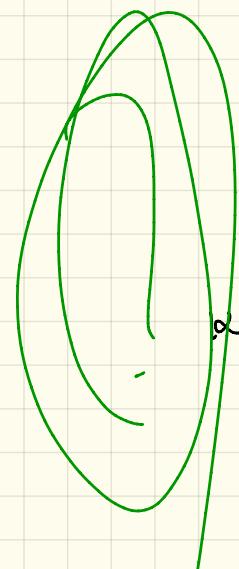
$\Leftrightarrow$  another root  $r_2$  is close to  $r_1$

$$\therefore g(x) = \frac{f(x)}{x - r_1} \Rightarrow$$

$$g(r_2) = \frac{f(r_2)}{r_2 - r_1}$$

$$f'(r_2) = g(r_2) = \frac{f(r_2)}{r_2 - r_1}$$

$$r_2 - r_1 = g(r_2) \cdot \underbrace{f(r_2)}$$



$$\alpha_{r_2} = \frac{1}{f'(r_2)}$$

## § 4.2 Fixed point iteration

Idea: Replace  $f(x) = 0$  with  $g(x) = x$  (fixed-point problem)

let  $g(x) = x - f(x)$  and  $r$  be the fixed point of  $g(x)$ .

(i.e.  $g(r) = r$ ). Then  $g(r) = r - f(r) = r$  or  $f(r) = 0$ .

Example 4.2.1

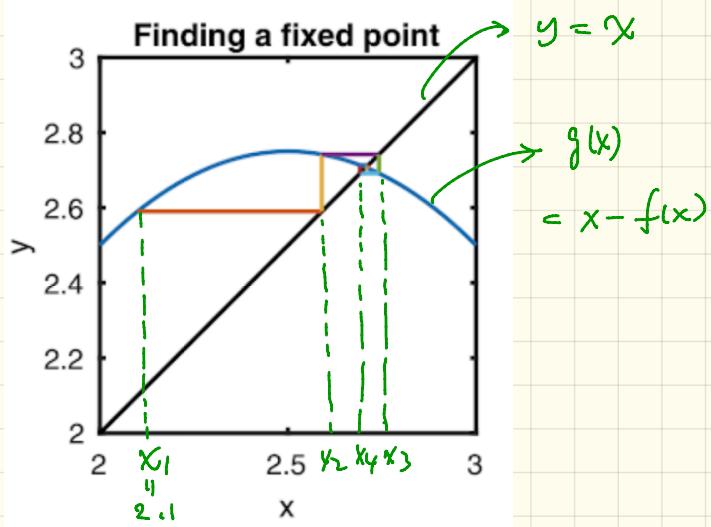
$$f(x) = x^2 - 4x + 3.5$$

MATLAB

```
>> f = @(x) x.^2 - 4*x + 3.5
```

```
>> r = roots([1 -4 3.5])
```

$$r = 2.7071, 1.2929$$

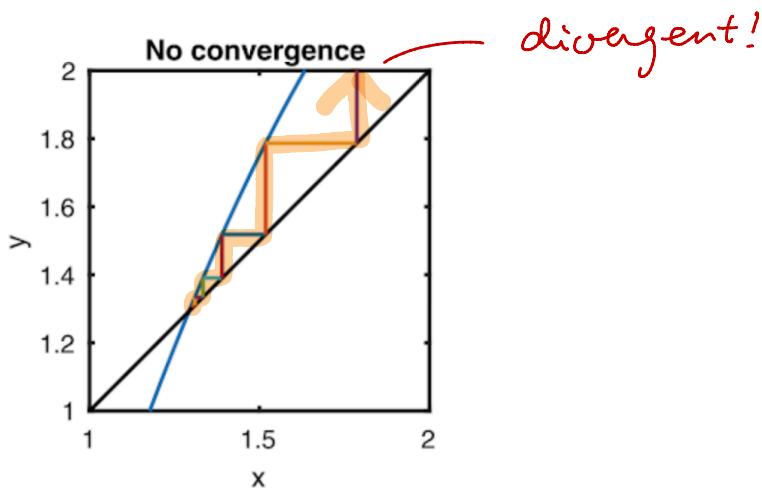


Now let's try to find the other fixed point in the same way. We'll use 1.3 as a starting approximation.

$$x_1 = 1.3$$

$x_k$  divergence!

```
cla  
fplot(g,[1 2])  
hold on, plot([1 2],[1 2], 'k'), ylim([1 2])  
x = 1.3; y = g(x);  
for k = 1:5  
    plot([x y],[y y], '-'), x = y;      % y --> new x  
    y = g(x);  plot([x x],[x y], '-')  % g(x) --> new y  
end
```



We started near the true fixed point, which is at about 1.293. But this time, the iteration is pushing us *away* from the correct answer.

## Series Analysis

- Let  $\varepsilon_k = x_k - r$  and  $x_{k+1} \leftarrow g(x_k)$

$$\text{Since } \varepsilon_{k+1} = x_{k+1} - r$$

$$\Rightarrow x_{k+1} = \varepsilon_{k+1} + \cancel{r} = g(x_k)$$

$$= g(\varepsilon_k + r) = g(r) + g'(r) \varepsilon_k + \frac{1}{2} g''(r) \varepsilon_k^2 + \dots$$

$$\Rightarrow \varepsilon_{k+1} = g'(r) \varepsilon_k + O(\varepsilon_k^2)$$

Small  $r$  if  $\varepsilon_k$  is small

$$\Rightarrow \varepsilon_{k+1} \approx g'(r) \cdot \varepsilon_k$$

$$\Rightarrow \varepsilon_{k+1} \approx g'(r)^{k+1} \cdot (\varepsilon_k \cdot \varepsilon_{k-1} \cdots \varepsilon_0)$$

$$\Rightarrow (1) |g'(r)| < 1, \varepsilon_{k+1} \rightarrow 0$$

$$(2) |g'(r)| > 1, \varepsilon_{k+1} \rightarrow \infty$$

Fixed point iteration

In Example 4.2.1

$$g(x) = -x^2 + 5x - 3.5$$

$$g'(x) = -2x + 5$$

$$|g'(\sim 2.71)| \approx |-0.42| < 1$$

$$|g'(\sim 1.29)| \approx |2.42| > 1$$

## Linear Convergence (Rate of convergence)

$$|\varepsilon_{k+1}| \approx |\varepsilon_k| \cdot \sigma, \quad 0 < \sigma < 1$$

$$|\varepsilon_k| = |x_k - r| \approx C |g'(r)|^k = C \sigma^k$$

$$\begin{aligned} \Rightarrow \log |\varepsilon_k| &\approx \log |C \sigma^k| = \underbrace{\log |C|}_{\beta} + \underbrace{k \log \sigma}_{\alpha} \\ &= \alpha k + \beta \end{aligned}$$

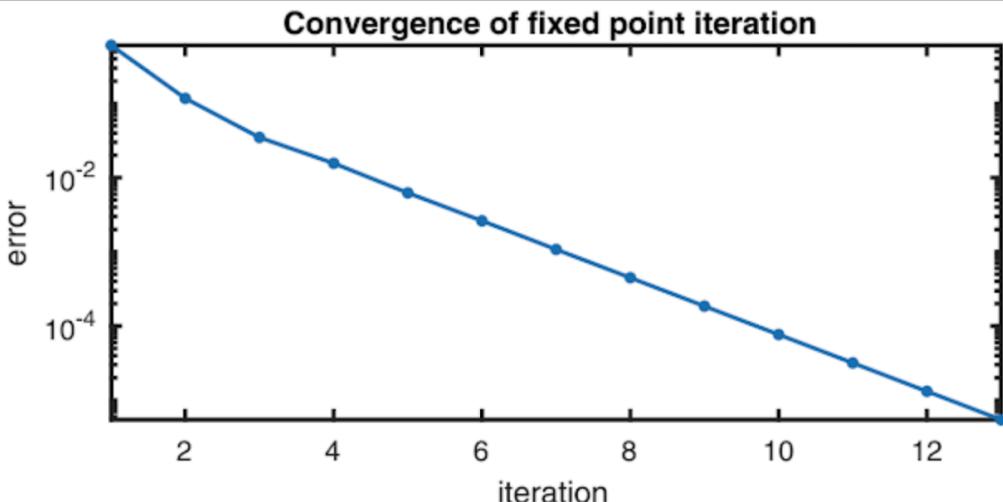
## Linear convergence

$$\lim_{k \rightarrow \infty} \frac{|x_{k+1} - r|}{|x_k - r|} = \sigma \quad \text{where } 0 < \sigma < 1$$

Practically speaking, linear convergence is identified by two different observations about the errors  $\epsilon_k$ :

1. The errors lie on a straight line on a log-linear graph, as implied by  $\log|\epsilon_k| \approx \alpha k + \beta$ .
2. The error is reduced by a constant factor in each step, as implied by  $|\epsilon_k| \approx C\sigma^k$ .

Both statements are approximate and only apply for sufficiently large values of  $k$ , so a certain amount of judgment has to be applied.



The first value is the slope, which must be negative. We can exponentiate it to get the convergence constant  $\sigma$ .

```
sigma = exp(p(1))
```

```
sigma =
0.4145
```

The numerical values of the error should decrease by a factor of  $\sigma$  at each iteration. We can check this easily with an elementwise division.

```
err(9:12) ./ err(8:11)
```

```
ans =
0.4138      0.4144      0.4141      0.4142
```

The methods for finding  $\sigma$  agree well.

## Contraction Maps

**Definition 4.2.1.** A function  $g$  is said to satisfy a **Lipschitz condition** with constant  $L$  on the interval  $S \subset \mathbb{R}$  if

$$|g(s) - g(t)| \leq L|s - t| \quad \text{for all } s, t \in S. \quad (4.2.5)$$

It can be shown that a function satisfying (4.2.5) is continuous in  $S$ . If  $L < 1$ , we call  $g$  a **contraction map**, because distances between points decrease after an application of  $g$ .

From the Fundamental Theorem of Calculus, which asserts  $g(s) - g(t) = \int_s^t g'(x) dx$ , it's not difficult to conclude that an upper bound of  $|g'(x)| \leq L$  for all  $x$  results in (4.2.5). But the weaker Lipschitz condition alone is enough to guarantee the success of fixed point iteration.

### Theorem 4.2.1: Contraction Mapping Theorem

Suppose that  $g$  satisfies (4.2.5) with  $L < 1$  on an interval  $S$ . Then  $S$  contains exactly one fixed point  $r$  of  $g$ . If  $x_1, x_2, \dots$  are generated by the fixed point iteration (4.2.1), and  $x_1, x_2, \dots$  all lie in  $S$ , then  $|x_k - r| \leq L^{k-1} |x_1 - r|$  for all  $k > 1$ .

*Partial proof.* First we show there is at most one fixed point in  $S$ . Suppose  $f(r) = r$  and  $f(s) = s$  in  $S$ . Then by (4.2.5),  $|r - s| = |g(r) - g(s)| \leq L|r - s|$ , which for  $L < 1$  is possible only if  $|r - s| = 0$ , so  $r = s$ .

Now suppose that for some  $r \in S$ ,  $g(r) = r$ . By the definition of the fixed point iteration and the Lipschitz condition,

$$|x_{k+1} - r| = |g(x_k) - g(r)| \leq L|x_k - r|,$$

which shows that  $x_k \rightarrow r$  as  $k \rightarrow \infty$ . To show that  $r$  must exist and complete the proof, one needs to apply the Cauchy theory of convergence of a sequence, which is beyond the scope of this book.  $\square$

## § 4.3 Newton's Method in One Variable

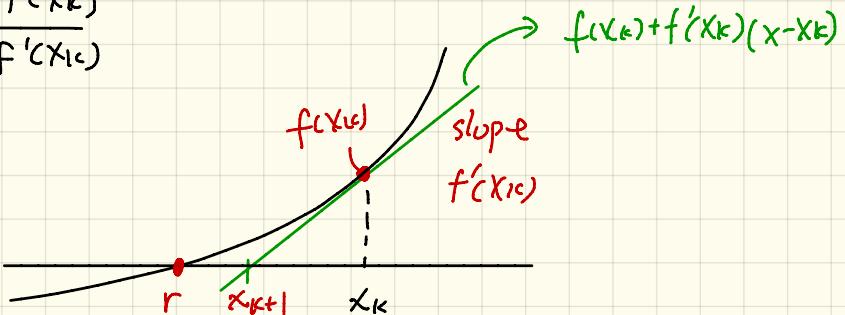
Idea: Replace  $f(x)$  with a linear model

$$f(x) = f(x_k) + f'(x_k)(x - x_k) + O(\beta^2)$$

$$\Rightarrow f(x) \approx f(x_k) + f'(x_k)(x - x_k)$$

$$\Rightarrow \text{Let } f(x_k) + f'(x_k)(x - x_k) = 0$$

$$\Rightarrow x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$



Example 4.3.1

## Quadratic Convergence

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)} \quad \varepsilon_k = r - x_k$$

$$\Rightarrow \varepsilon_{k+1} = \varepsilon_k - \frac{f(r - \varepsilon_k)}{f'(r - \varepsilon_k)}$$

$$\varepsilon_{k+1} = \varepsilon_k + \frac{f(r) - \epsilon_k f'(r) + \frac{1}{2} \epsilon_k^2 f''(r) + O(\epsilon_k^3)}{f'(r) - \epsilon_k f''(r) + O(\epsilon_k^2)}.$$

We use the fact that  $f(r) = 0$  and additionally assume now that  $f'(r) \neq 0$ . Then

$$\varepsilon_{k+1} = \varepsilon_k - \epsilon_k \left[ 1 - \frac{1}{2} \frac{f''(r)}{f'(r)} \epsilon_k + O(\epsilon_k^2) \right] \left[ 1 - \frac{f''(r)}{f'(r)} \epsilon_k + O(\epsilon_k^2) \right]^{-1}.$$

The series in the denominator is of the form  $(1+z)^{-1}$ . Provided  $|z| < 1$ , this is the limit of the geometric series  $1-z+z^2-z^3+\dots$ . Keeping only the lowest-order terms, we derive

$$\begin{aligned} \varepsilon_{k+1} &= \varepsilon_k - \epsilon_k \left[ 1 - \frac{1}{2} \frac{f''(r)}{f'(r)} \epsilon_k + O(\epsilon_k^2) \right] \left[ 1 + \frac{f''(r)}{f'(r)} \epsilon_k + O(\epsilon_k^2) \right] \\ &= -\frac{1}{2} \frac{f''(r)}{f'(r)} \epsilon_k^2 + O(\epsilon_k^3). \end{aligned} \tag{4.3.4}$$

Equation (4.3.4) suggests that, *eventually, each iteration of Newton's method roughly squares the error*. This behavior is called **quadratic convergence**. The formal definition of quadratic convergence is that there exists a number  $\alpha > 0$  such that

$$\lim_{k \rightarrow \infty} \frac{|x_{k+1} - r|}{|x_k - r|^2} = \alpha. \quad (4.3.5)$$

Recall that linear convergence is identifiable by trending toward a straight line on a log-linear plot of the error. When the convergence is quadratic, no such straight line exists—the convergence keeps getting steeper. Alternatively, note that (neglecting high-order terms)

$$\log(|\epsilon_{k+1}|) \approx 2 \log(|\epsilon_k|) + \text{constant},$$

which is equivalent to saying that the number of accurate digits approximately doubles at each iteration, once the errors become small enough.

Example 4.3.2

Example 4.3.3

## § 4.4 Interpolation - based

### Secant method

- $\rightarrow$   $\Rightarrow$

- Secant method
- Inverse interpolation
- Bracketing

Example 4.4.1 demonstrates the secant method. *In the secant method, one finds the root of the linear approximation through the two most recent root estimates.* That is, given previous approximations  $x_1, \dots, x_k$ , define the linear model function as the line through  $(x_{k-1}, f(x_{k-1}))$  and  $(x_k, f(x_k))$ :

$$q(x) = f(x_k) + \frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}}(x - x_k). \quad (4.4.1)$$

Solving  $q(x_{k+1}) = 0$  for  $x_{k+1}$  gives the formula

$$x_{k+1} = x_k - \frac{f(x_k)(x_k - x_{k-1})}{f(x_k) - f(x_{k-1})}, \quad n = 1, 2, \dots \quad (4.4.2)$$

Example 4.4.1

Function 4.4.1

## Superlinear Convergence

$$\epsilon_{k+1} \approx -\frac{1}{2} \frac{f''(r)}{f'(r)} \epsilon_k \epsilon_{k-1}. \quad (4.4.3)$$

If we make an educated guess that

$$\epsilon_{k+1} = c(\epsilon_k)^\alpha, \quad \epsilon_k = c(\epsilon_{k-1})^\alpha, \quad \alpha > 0,$$

then (4.4.3) becomes

$$[\epsilon_{k-1}^\alpha]^\alpha \approx C \epsilon_{k-1}^{\alpha+1} \quad (4.4.4)$$

for an unknown constant  $C$ . Treating the implied approximation as an equality, this becomes solvable if and only if the exponents match, i.e.,  $\alpha^2 = \alpha + 1$ . The only positive root of this equation is the golden ratio,

$$\alpha = \frac{1 + \sqrt{5}}{2} \approx 1.618.$$

Hence the errors in the secant method converge like  $\epsilon_{k+1} = c(\epsilon_k)^\alpha$  for  $1 < \alpha < 2$ , a situation called **superlinear convergence**.

Example 4.4.2

P.14' Bottom

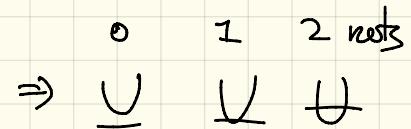
P.14' top

## Inverse Interpolation

Idea : Newton's method : 1 point

Secant method : 2 points

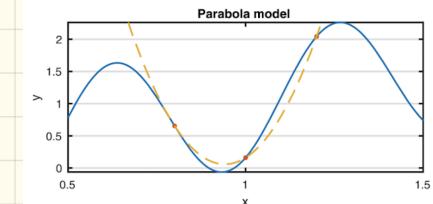
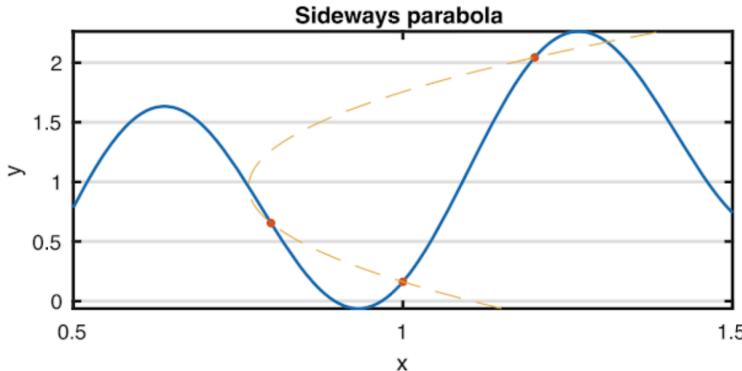
higher-order =  $\Rightarrow$  2 points ?



$\Rightarrow$  Inverse interpolation

To do inverse interpolation, we swap the roles of  $x$  and  $y$  in the interpolation.

```
c = polyfit(y,x,2); % coefficients of interpolating polynomial  
q = @(y) polyval(c,y);  
fplot(q,@(y) y,ylim,'--') % plot x=q(y), y=y
```



We repeat the process a few more times.

```
for k = 4:8  
    c = polyfit(y(k-2:k),x(k-2:k),2);  
    x(k+1) = polyval(c,0);  
    y(k+1) = f(x(k+1));  
end
```

Check the convergence rate on P1(j)

## Bracketing

- Idea: Intermediate Value Theorem

If  $f(x)$  is continuous on  $[a, b]$  and  $f(a)f(b) < 0$ ,  
 $f(x)$  have at least one root in  $[a, b]$

- Convergence:  $\left| \frac{b_k - a_k}{2} \right|^k$ . a guarantee method

## MATLAB

- To blend fast-converging methods w/ the guarantee provided by a bracket
- For example, bracketing  $\Rightarrow$  sub-intervals
  - $\Rightarrow$  inverse quadratic estimates
  - $\Rightarrow$  Secant methods

## § 4.5 Newton for Nonlinear Systems

Idea: Replace the nonlinear system on  $\mathbb{R}^n$  by a linear model

$$f: \mathbb{R}^n \rightarrow \mathbb{R}^n$$

$$f_1(x_1, \dots, x_n) = 0,$$

$$f_2(x_1, \dots, x_n) = 0,$$

⋮

$$f_n(x_1, \dots, x_n) = 0.$$

$$f(x+h) = f(x) + J(x)h + O(\|h\|^2), \quad (4.5.1)$$

where  $J$  is called the **Jacobian matrix** of  $f$  and is defined by

$$J(x) = \begin{bmatrix} \frac{\partial f_1}{\partial x_1} & \frac{\partial f_1}{\partial x_2} & \cdots & \frac{\partial f_1}{\partial x_n} \\ \frac{\partial f_2}{\partial x_1} & \frac{\partial f_2}{\partial x_2} & \cdots & \frac{\partial f_2}{\partial x_n} \\ \vdots & \vdots & & \vdots \\ \frac{\partial f_n}{\partial x_1} & \frac{\partial f_n}{\partial x_2} & \cdots & \frac{\partial f_n}{\partial x_n} \end{bmatrix} = \left[ \frac{\partial f_i}{\partial x_j} \right]_{i,j=1,\dots,n}. \quad (4.5.2)$$

- Notes: If  $f(x)$  is linear (e.g.  $f(x) = Ax = b$ ), then the Jacobian

$J(x)$  is a constant (e.g.  $A$ ) and the higher-order term

$O(\|h\|^2)$  disappear

## The Multidimensional Newton Iteration

With a method in hand for constructing a linear model for the vector system  $\mathbf{f}(\mathbf{x})$ , we can generalize Newton's method. Specifically, at a root estimate  $\mathbf{x}_k$ , we set  $\mathbf{h} = \mathbf{x} - \mathbf{x}_k$  in (4.5.1) and get

$$\mathbf{f}(\mathbf{x}) \approx \mathbf{q}(\mathbf{x}) = \mathbf{f}(\mathbf{x}_k) + \mathbf{J}(\mathbf{x}_k)(\mathbf{x} - \mathbf{x}_k).$$

We define the next iteration value  $\mathbf{x}_{k+1}$  by requiring  $\mathbf{q}(\mathbf{x}_{k+1}) = \mathbf{0}$ ,

$$\mathbf{0} = \mathbf{f}(\mathbf{x}_k) + \mathbf{J}(\mathbf{x}_k)(\mathbf{x}_{k+1} - \mathbf{x}_k),$$

which can be rearranged into

$$\mathbf{x}_{k+1} = \mathbf{x}_k - [\mathbf{J}(\mathbf{x}_k)]^{-1} \mathbf{f}(\mathbf{x}_k). \quad (4.5.3)$$

Note that  $\mathbf{J}^{-1} \mathbf{f}$  now plays the role that  $f/f'$  had in the scalar case; in fact the two are the same in one dimension. In computational practice, however, we don't compute matrix inverses. Instead, define the *Newton step*  $\mathbf{s}_k$  by the linear  $n \times n$  system

$$\mathbf{J}(\mathbf{x}_k) \mathbf{s}_k = -\mathbf{f}(\mathbf{x}_k), \quad (4.5.4)$$

so that  $\mathbf{x}_{k+1} = \mathbf{x}_k + \mathbf{s}_k$ . *Computing the Newton step is equivalent to solving a linear system using the Jacobian matrix and the function value.*

An extension of our series analysis of the scalar Newton method shows that the vector version is also quadratically convergent in any vector norm, under suitable circumstances.

[LAB] Example 4.5.2

Function 4.5.1

Example 4.5.3

## § 4.6 Quasi-Newton Methods

Issues of Newton's methods:

- Programming nuisance and computational expense of calculating the Jacobian matrix
- Divergence of many starting points

Jacobian by finite differences

Idea

- In 1-D,  $f'(x_k) \Rightarrow \left[ \frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}} \right] \rightarrow$  Conv. to  $f'(x_k)$  as  $x_k \rightarrow r$

- In n-D, jth column of the Jacobian:

$$J(x)e_j = \begin{bmatrix} \frac{\partial f_1}{\partial x_j} \\ \frac{\partial f_2}{\partial x_j} \\ \vdots \\ \frac{\partial f_n}{\partial x_j} \end{bmatrix} \approx \frac{f(x + \delta e_j) - f(x)}{\delta}, \quad j = 1, \dots, n.$$

-  $\delta = \sqrt{\sum \text{mach}}$

- Function 4.6.1

## Broyden's Update

### Issues of the finite-difference

- Needs  $n$  extra function evaluations ( $\vec{f}(\vec{x} + \delta e_j), j=1:n$ )
- Jacobian changes little as iteration converges

### Idea

- "Cheap-and-good-enough" way to update the Jacobian

### Derivation

Recall

$$0 \approx f(x_{k+1}) \approx f_k + J(x_k) \underbrace{(x_{k+1} - x_k)}$$

Notations  $\Rightarrow$

$$f_{k+1} \quad f_k \quad A_k \quad s_k \text{ (Newton step)}$$

Newton iteration:  $A_k s_k = -f_k$

After get  $x_{k+1}$ , we need  $A_{k+1}$

Idea: the "secant" method

new Jacobian

$$A_{k+1} = \boxed{\frac{f_{k+1} - f_k}{x_{k+1} - x_k}} \rightarrow \text{向量, 不能算, But}$$

$$\Rightarrow A_{k+1} (x_{k+1} - x_k) = (f_{k+1} - f_k) \quad \text{3) 算!}$$

$$\Rightarrow \underbrace{A_{k+1} s_k}_{=} = f_{k+1} - f_k$$

what's  $A_{k+1}$ ?

Idea:  $A_{k+1} = A_k + [\text{rank-1 matrix}]$

$$A_{k+1} = A_k + \underbrace{\frac{1}{s_k^T s_k} (f_{k+1} - f_k - A_k s_k)}_{\text{scalar}} s_k^T \quad (+)$$

①  $(+)*s_k$

$$A_{k+1} \cdot s_k = A_k \cdot s_k$$

$$+ \frac{1}{s_k^T s_k} (f_{k+1} - f_k - A_k s_k) s_k^T \cdot s_k$$

$$\Rightarrow A_{k+1} s_k = f_{k+1} - f_k = (+)$$

②  $\downarrow$  is rank-1 matrix!

$s_k \cdot s_k^T$  (outer product)

③ No extra function evaluations  
Needs only  $f_k, f_{k+1}$

- Note: With some assumptions,  $x_k \rightarrow r$  superlinearly  
( $A_k$  may not converge to the Jacobian)
- Note: Use finite-difference in the beginning.  
If the Broyden update fails, reinitialize  $A_k$  by finite-difference.

## Levenberg's Method

### Issues

- To choose starting points
- The Jacobian matrices (① derivative, ② finite-difference, ③ Broyden's iteratively updated) become increasingly inaccurate

### Ideas

- Detailed accuracy analysis (too expensive!)
- Checking whether take  $x_{k+1}$  or find an alternative

## Derivations

- Balance between two ideas

$$(A_k^T A_k + \boxed{\lambda} I) s_k = -A_k^T f_k \quad s_k = x_{k+1} - x_k \quad (+)$$

(1)  $\lambda = 0$

$$(+ \Rightarrow A_k^T A_k s_k = -A_k^T f_k$$

|||

$$A_k s_k = -f_k$$

(Regular linear model)

(2)  $\lambda \neq 0$

$$(+ \Rightarrow \lambda s_k = -A_k^T f_k$$

$\Rightarrow$   $s_k$  is the steepest descent dir.  
to min  $\|f(x)\|_2^2$

To solve  $f(x) = 0$

$$\Leftrightarrow \min \phi(x) = \|f(x)\|_2^2 = f(x)^T f(x)$$

$$\nabla \phi(x) = 2 \boxed{J(x)^T f(x)} \approx 2 \boxed{A_k^T f(x)}$$

$$= -2 \lambda s_k$$

$$\Rightarrow s_k \left( \approx -\frac{1}{2\lambda} \nabla \phi(x) \right)$$

$A_k = J(x)$  or  $A_k \approx J(x)$

## Implementation

Function 4.6.2

Example 4.6.1

## § 4.7 Nonlinear Least Squares

Recall:  $f(x) = f(x_k) + A_k(x - x_k) + O(\|h\|^2)$

- Jacobian
- Finite-difference
- Broyden's update

$\approx g(x)$  (linear model)

$$\Rightarrow g(x) \approx 0 \Rightarrow A_k s_k = -f_k \quad (s_k = x - x_k \text{ or } x = x_k + s_k)$$

search direction

	linear	nonlinear
square	$Ax = b$	$f(x) = 0$
overdetermined	$\min \ Ax - b\ _2$	$\min \ f(x)\ _2$

Square We can solve  $A_k s_k = -f_k$

Overdetermined

$$\begin{matrix} n \\ m \\ m > n \end{matrix}$$

$$f: \mathbb{R}^n \rightarrow \mathbb{R}^m$$

$$\begin{cases} f_1(x_1, \dots, x_n) = 0 \\ \vdots \\ f_m(x_1, \dots, x_n) = 0 \end{cases}$$

$$m > n$$

$$\left\{ \begin{array}{l} \min \|g(x)\|_2^2 \approx \|A_k s_k + f(x_k)\|_2 \\ x_{k+1} = \arg \min \|A_k s_k + f(x_k)\|_2 \end{array} \right.$$

Gauss-Newton Method

- { Function 4.5.1 newtonsys }
- { Function 4.6.2 levenberg }

## Nonlinear Data Fitting

We can apply the Gauss-Newton method to perform nonlinear data fitting.

$$\min_{c \in \mathbb{R}^n} \sum_{i=1}^m |g(t_i; c) - y_i|^2 = \min_{c \in \mathbb{R}^n} \|f(c)\|^2$$

↑  
coefficients to be determined

$$\text{Linear case: } g(t; c) = c_1 g_1(t) + c_2 g_2(t) + \dots + c_m g_m(t)$$

[LAB] Example 4.7.1