多8单纯同调以5奇异同调 设长为单纯多形,LCK为了多形。将设:Hn(K,L)=>Hn(IKI,IH) 为些人的当前然的态射。C.(K,L)—>S.(IKI,ILI). 为此, 选定长的项点集的一个排序 记水中的n-单形金体为: { Challe In } () Challe K) 12 公的成点,集为 Vd<---< Vn (其中"<"从 K的顶点集的排序键 (Ush..., Un), d ∈ In. 生成的 了下来的). ∀n, Ź·×: Cn(K) —> Sn(IKI). by ‡lź. (V_n^2, \dots, V_n^2) \longrightarrow $(S_a: \Delta_n \rightarrow |K|, \sum_{i=n}^n t_i V_i^2$ "

定理4、设长为单纯多形、LCK为于多形、则历经态射 C.(K, L) -> S.(|K|, |L|). 流导的解闭态 Hn(K,L)→ Hn(IK1,1L1) 为13) 好, ∀n∈Z20. 证明: 发记上一声的情形 Case 1. $K = K^n$, for some $n \in \mathbb{Z}_{20}$ $\delta \hat{\chi} \stackrel{.}{\succeq}$. 对至对的任务的知知. n=0, 12 ski. 作文はマオ n ≤ k-1 1を11)3: Hi(K)=> H;(|K1), サショロ T 1/2 n = k. 表虑短正分别:0-)C.(Kk1)->C.(Kk)->C.(Kk,Kh) $0 \rightarrow 5.(|K^{k-1}|) \rightarrow 5.(|K^{k}|,|K^{k-1}|)$ ~~ 长正会引:

 $H_{i+1}(\mathbb{K}^{k},\mathbb{K}^{k-1}) \xrightarrow{S} H_{i}(\mathbb{K}^{k-1}) \longrightarrow H_{i}(\mathbb{K}^{k}) \longrightarrow H_{i}(\mathbb{K}^{k},\mathbb{K}^{k-1}) \longrightarrow H_{i-1}(\mathbb{K}^{k-1})$ Claim: $H:(K^k, K^{k-1}) \longrightarrow H:(IK^{k}, IK^{k-1}) \xrightarrow{3} 45), \forall i$ $H:(K^k, K^{k-1}) \cong \begin{cases} b K \otimes k - \mathring{\mathcal{A}} \mathring{\mathcal{K}} \mathring{\mathcal{$

有台换图包: $H_{i}\left(\prod_{d\in I_{k}}\widetilde{\Delta_{k}^{2}},\prod_{d\in I_{k}}\widetilde{\Delta_{k}^{2}}\right)$ \longrightarrow $H_{i}\left(\lceil k^{k}\rceil,\lceil k^{k-1}\rceil\right)$ $\frac{115}{H:\left(\frac{11}{261k}\right)^{2k}} \stackrel{(1)}{=} \frac{115}{2k} \stackrel{(1)}{=} \frac{115}{2k}$ 子型·· $H:(|K^{k}|,|K^{k-1}|) \cong H:(\stackrel{\downarrow}{\downarrow} \stackrel{$

因为当主权时 $H_{\hat{a}}(K^{k},K^{k-1}) \cong 0$, $H_{\hat{a}}(1K^{k}1,1K^{k-1}1) \cong 0$ 所以只需证: $H_k(K^k, K^{k-1}) \longrightarrow H_k(|K^k|, |K^{k-1}|) \stackrel{1}{\sim} 17) \stackrel{1}{\sim} 3$ (K., K.,) (K., K.,) (K.,) (+) Z la. 小人之(vò,···, Vn) 引(チZしが類的 的什:(说,…,以)一以,则有当换图表:

 $H_{k}(\mathbb{K}^{k},\mathbb{K}^{k-1}) \xrightarrow{\longrightarrow} H_{k}(\mathbb{K}^{k},\mathbb{K}^{k-1}) \xrightarrow{\parallel_{l}} 6a$ OF Hk (Sh, DSh) $\begin{array}{c|c}
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(\downarrow) & \searrow & (\downarrow$ 干岩 Claim 第分 由five lemma, H;(kk) -> H;(Kk1) 为13)构、自约完成. Case 2. K ≠ Kⁿ, Y n ∈ Z₂, o 公かる: H:(K) → H:(K1)为同场, Vi∈Z20 Fact, IK1中的任何一个原子集必包含在长的一个有限子鸟形 中(LIMunkres],[Hatcher], 专引了一节)

∀ | K | 上級-介奇异 i-cycle 6,由(*), ヨn∈Z₂0, st. 6为 | K" | 中的-介奇异 i-cycle.

H:(K") -> H:(K) > d

(ase 1 | 115] H; (|K"|) -> . (|K|) |

考虑交换图表:

于岩. ad 6 Hi(K), s.t.

d -> o在H:(IKI)中听代表写

· H:(K) -> H:(IKI)为单轴。

V K中的 i- cycle C, 设 c→> 6, 且 6=0 ≥, 其中之为 |K|上的 - 介奇异的(i+1)-chain,

由(x), ∃n∈Z20, s+06与互落在同意S((K)1)→ S.((K1)が) 像中、考虑交换图表。 团市。c=o in Hi(K). 最后治上丰中的情形。 考虑,每正今到的分类是图表: 0-) C.(L) -> C.(K) -> C.(K,L) -> ° 0-) C.(L) -> C.(K) -> S.(K,L) -> ° 0-) S.(H) -> S.(K,L) -> ° 一)长正含别: Hn(L) -> Hn(K) -> Hn(K,L) -> Hn-1(L) -> Hn-1(K) 15. If. Ifs. Ha(ILI) -> Ha(IKI) -> Ha(IKI, ILI) -> Har, (ILI) -> Har, (IKI)

由己的的急性的绝对版本, f.f.f.f.为同物。 12) five lemma => f3 % 13] \$5). 定理1、设义为介三南到分和超升空间,(K,h),(K',h')为 X的的个三面别分、则有历经的同韵: $\varphi_{kk'}: H_n(k) \stackrel{\mathcal{Z}}{=} H_n(k')$, $\forall n \in \mathbb{Z}_{\geq 0}$ 自然: 花(K,h), (K',h'), (K'',h'')为 X的三角别分, PKK" i.e. $q_{k'k''} \circ q_{kk'} = q_{kk''}$ 记明, 南流明治: $Q_{kk'} = Q_{k'}^{-1} \circ Q_{k}$ $\varphi_{k'}: H_n(k') \xrightarrow{u_{k'}} H_n(|k'|) \xrightarrow{h_{x}} H_n(x')$ $\varphi_{k'}: H_n(k') \xrightarrow{u_{k'}} H_n(|k'|) \xrightarrow{h_{x}} H_n(x')$ Hr (K") UK" > Hr (K")

对于有限单纯复形X,其奇异同调解Hn(X)均为 有性生成分别 推设: 花义为宗的流形。 => H(X) 为有限生成和日期 艺义, 花义为有限单纯多形, 从(x)的自由部分的直和项区 的个数、称为X的事个为是Hi基 多9. 应用. B"为n维闭球, 则从3"引B"的连 定理与(Brower不动点定理)设 镇映新必有不动点. 无不动点,构造映射 引 87一) 5~~ 元子 17月: ンタ 子: Bn 一) Bh

证明: 25 f. Bn 一) Bn 无不动点, 构造映射 g. Bn 一) Sn-1.

专规 笔 g(x) = 从 f(x) 连则 x 的 射线与 Sn-1的 支点。

*** 21 g 连续, g | 5n-1 = Id 5n-1

考虑多会 5-1 in Bn -\$ 5-1 9. i = Idsn-1 $Id = (90i)_{x} : H_{n-1}(5^{n-1}) \longrightarrow H_{n-1}(5^{n-1})$ []₂₋₁(B²) $=) g_{\chi}: H_{n-1}(B^n) -) H_{n-1}(S^{n-1}) 为 講 同意, (n = 1)$ n=1. gx: Z一文田区、矛盾 鸽上,9×不可能为满同态,矛盾。因而于必有不动点,

宇理 6. (Generalized Jordan carve theorem) ガカトモ Z+、ガ文 C 为 5 中 同配于5个的了空间,到5个人有两个连遍分支。 定理了 (invariance of domain) 没UCR"为开子集 f:UnR"为连续单轴,f(U)CR"也为开子集,归 f: U → f(U) 为同凡. 定理8. (Invariance of dimension).设UCRM, VCRM均为开子集 [2,] $U \cong V$, [2,] M = N. 京祖子 => 至祖多: V = 1 0 i, R = 1 R 由室理下. joi(U)为R"中开集,且joi(U)Cj(R")

定理6与党理了之话则伦毅于: (a) 花DC5°为一个闭的k-胞腔,则fi(5°\D)=o, Yi. (b) 花 C C S 同间下 5k, 0 ∈ k < n, 约 → 空间, 21. $H_i(S^r \setminus C) \cong \begin{cases} Z \\ 0 \end{cases}$ if i = n-k-1 else 为难以从(山),(6)。 宣祖6之谷明: CC5°, C=5°1, (b), \Rightarrow $H_{\circ}(S^{n}\setminus C) \cong \mathbb{Z}_{\rightarrow} H_{\circ}(S^{n}\setminus C) \cong \mathbb{Z} \oplus \mathbb{Z}_{\rightarrow}$ I STIC 有两个溢路连遍分支 5局部溢級连通 对于局部通路连通空间,适路连遍分支与连通分支吻合

=> 5°\c的有两个连通分支. 是理了之证明: 对R"做一点紧化:R"U(su) = S" f:U—>Sn为连续单新、UCRT.干话 f(U)为 s^n 中开集 (=> f(U)为同胚) YXEU、安はtx1为f(U)的内点 ∃870.5.t.开球 B(X;8) C U.D. 1.2.要任于(B(X;8)) B(x; S) CU, 为Sr中开集 flab(x:8) : aB(x:8) -> f(aB(x:8)). TI $\partial B(x;\delta)$ A B(x;\delta)

 $f|_{\overline{B(x;\delta)}}$, $\overline{B(x;\delta)} \longrightarrow f(\overline{B(x;\delta)}) \not \supset \overline{|}$

由宽理6, ST(C)的有两个(通路)连通分支 $S^n \setminus C = \left(S^n \setminus f(\overline{B(x;\delta)})\right) \coprod f(B(x;\delta)).$ 透路连通 Sn、f(aB(xis)). 透路達面 (a), $H_{o}(S^{n}) f(\overline{B(x;\delta)})) = 0$. $\Rightarrow S' \setminus C = (S') + (\overline{B(X;S)}) + (B(X;S)) 为连通分支分$ → f(B(X; 8)) 为 5 中日子条 # 下面看手子(四)、(b)的计算。 Mayer-Vietoris sequence: (M-V sequence). $\frac{1}{\sqrt{2}}$ \times ; top. Sp. A, B \subset \times , int(A) U int(B) = \times , 记见二个人的, 有正分别:

0 -> Sn (Anis) -> Sn(A) + Sn(B) -> Sn(X) -> 0 ~~ 链复形的还含引: 0-> S.(ANB)-> S.(A) (B) -> S.(X) -> o. ~) 长正分引: (Homological IM-V sequence). -.. -> H, (AnB) -> H, (A) + H, (B) -> H, (X) -> H, (AnB) -> H, (A) + 在(米)基础上加工 U) SI(ANB) -> SI(A) (B) (B) -> SI(X) -> 0 (后未的) 0-) So(ANB)-> So(A) (B) -> So(X) ->=

U→ SI(AOB) → SI(A) ⊕ SI(B) → SI(X) → 0 $0 \rightarrow S_{\bullet}(A \cap B) \rightarrow S_{\bullet}(A) \oplus S_{\bullet}(B) \rightarrow S_{\bullet}(X) \rightarrow S$ 对之应用超正台引诱与长正台引引强。 .. $\rightarrow \widetilde{H}_{n}(A \cap B) \rightarrow \widetilde{H}_{n}(A) \oplus \widetilde{H}_{n}(B) \rightarrow \widetilde{H}_{n}(X) \rightarrow \widetilde{H}_{n-1}(A \cap B) \rightarrow \widetilde{H}_{n-1}(A) \oplus \widetilde{H}_{n-1}(B)$ (a).Hi(S^\D)=o, Yi、 其中D为闭的k-胞腔、

好好的场。 k=0, S^1\D⇒R", 温恕成立

设有 k s l-1 时, 已记 门(s^\D) e o, \Vi, 其中 D为闭的胞腔 下设の为何如见一胞腔、造和闭胚h: II->D $A = S^{n} \setminus h(I^{l-1} \times I^{o}, \frac{1}{2}), B = S^{n} \setminus h(I^{l-1} \times I^{l-1}, 1)$ (I = [0,1]) $AUB = S^n \setminus h(I^{l-1} \times \{\frac{l}{l}\})$ r_{n=2}, l=2 ANB = S'\D closed (l-1)-ce () Tz 12) M-V segnence: $\widetilde{H}_{j+1}(S^n(h(\underline{I^{l}}X\S\{\underline{l}\}))\to \widetilde{H}_{j}(A\cap B)\to \widetilde{H}_{j}(A)\oplus \widetilde{H}_{j}(B)\to \widetilde{H}_{j}(S^n(h(\underline{I^{l}}X\S\{\underline{l}\}))\to \widetilde{H}_{j}(S^n(h(\underline{I^{l}}X\S\{\underline{l}\}))\to \widetilde{H}_{j}(B)\to \widetilde{H}_{j}(B)$ 川的假设 的由假设 $\widehat{H}_{j}(\underline{A}\underline{\cap B}) \stackrel{\mathcal{L}}{=} \widehat{H}_{j}(A) \oplus \widehat{H}_{j}(B)$.

 $T:\mathcal{F}:\widetilde{H}_{j}(S^{n}(D)) \cong 0$ 假设SND中的新年j-cycle C 不为boundary. 由河初 $\widehat{H}_{j}(S^{n}\setminus D)$ 一 $\widehat{H}_{j}(A)$ $\widehat{H}_{j}(B)$ C代表的同湖美一一、CC视为A中Cycle所代表的同间集, c视为B中cycle所代表的图记到集) 不好没c初为A中cycle所代表的同调集不为零。1.e. C不为A中的 boundary $12 A = S^{n} \setminus h(I^{l-1} \times I_{1})$, $I_{1} = [0, \frac{1}{2}]$ 再将工一等分,重复上述操作、设工、等分后得明的两个 (子) 足向为工之,工之, (5°\h(ヹ'×エィ)) => Ĥ;(5°\h(ヹ'xエュ))的Ĥ;(5°\ $h(I^{0-1}\times I^{2})$ 12 c ふちらかん(IP-1×I'l) ** boundary . 12 Il= I!. 重复流轨道,其口耳口耳口下。一、满足 $\left| \overline{L}_{n} \right| = \frac{1}{2^{n}}$

C 不为 Sⁿ \ h(I^{Q-1} × Im) it boundary, ∀ m ≥ 1 $(S^n \setminus h(I^{l-1} \times I_i) \subset S^n \setminus h(I^{l-1} \times I_i) \subset \cdots \subset S^n \setminus h(I^{l-1} \times I_m) \subset \cdots)$ 5m/ h(Il-1x{p}) $S^n \setminus h(I^l)$. closed (2-1) - (el) 12 (Im = {P}. 由归纳缩说。C沙为Smil(Il-1x/pl)中boundary。 $d = \sum_{j=1}^{n} m_j G_j$ $S^{m} \setminus h(I^{l-1} \times \{p\}) = \bigcup_{m=1}^{m} S^{n} \setminus h(I^{l-1} \times I_{m}) \rightarrow \bigcup_{j=1}^{m} G_{j}$ 7.43i 3° 5° $h(I^{l-1} \times I_{mo}) \supset \bigcup_{j=1}^{l} I_{m} 6_{j}$. $= d \in S_{jt1}(S^n \setminus h(1^{j-1} \times Im_{-})) = c + 5^n \setminus h(1^{j-1} \times Im_{-}) \neq 0$ boundary. 3th

i + i = n - k - 1 $\mathcal{A}_{i}(S^{r}\setminus C)\cong \begin{cases} \mathcal{A}_{i} \end{cases}$ else (其中 C = Sk, osk < n). 对人间间的话: k=0. $S^{n-1} \times R \stackrel{\sim}{\sim} S^{n-1}$. $I_{2}^{n-1} \stackrel{\sim}{\sim} I_{3}^{n-1} \stackrel{\sim}{\sim} I_{3}^$ 设对 k s e - 1 已治, 下设 C 至 St. 固笔一个同胜 $h: S^1 \longrightarrow C$ $S_{+}^{l} = \{(x^{*}, ..., x^{l}) \in S^{l} | x^{l} \ge 3^{l}, S_{-}^{l} = \{(x^{*}, ..., x^{l}) \in S^{l} | x^{l} \le 3^{l}\}$ $\frac{1}{2}A = S^{n} \setminus h(S^{l}), \quad B = S^{n} \setminus h(S^{l}).$ $A \cap B = S' \setminus C' A \cup B = S' \setminus h(S_{+}^{l} \cap S_{-}^{l}).$ ·<-1 M-V sequence: