

CSC 252: Computer Organization

Spring 2025: Lecture 8

Instructor: Yuhao Zhu

Department of Computer Science
University of Rochester

Announcement

- You might still have three slip days.
- Read the instructions before getting started!!!
 - You get 1/4 point off for every wrong answer
 - Maxed out at 10
- TAs are best positioned to answer your questions about programming assignments!!!
- Programming assignments do NOT repeat the lecture materials. They ask you to synthesize what you have learned from the lectures and work out something new.
- Logics and arithmetics problem set: <https://www.cs.rochester.edu/courses/252/spring2025/handouts.html>.
 - Not to be turned in.

Announcement

- Email me to schedule an office hour.
- Or come to my office at around 3 on Wednesday or Friday.

Today: Control Instructions

- Control: Conditional branches (**if... else...**)
- Control: Loops (**for, while**)
- Control: Switch Statements (**case... switch...**)

Conditional Branch Example

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq    %rsi,%rdi # x:y
    jle     .L4
    movq    %rdi,%rax
    subq    %rsi,%rax
    ret

.L4:      # x <= y
    movq    %rsi,%rax
    subq    %rdi,%rax
    ret
```

| Register | Use(s) |
|----------|--------------|
| %rdi | x |
| %rsi | y |
| %rax | Return value |

cmpq sets ZF, SF, OF

jle checks $ZF \mid (SF \wedge OF)$

| | | |
|----|----|----|
| 0 | 0 | 0 |
| ZF | SF | OF |

Conditional Branch Example

```
unsigned long absdiff
(unsigned long x, unsigned
long y)
{
    unsigned long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq    %rsi,%rdi # x:y
    jbe     .L4
    movq    %rdi,%rax
    subq    %rsi,%rax
    ret
.L4:       # x <= y
    movq    %rsi,%rax
    subq    %rdi,%rax
    ret
```

| Register | Use(s) |
|----------|--------------|
| %rdi | x |
| %rsi | y |
| %rax | Return value |

| | | |
|----|----|----|
| 0 | 0 | 0 |
| ZF | SF | OF |

Conditional Jump Instruction

```
cmpq    %rsi, %rdi
jle     .L4
```

Jump to label if less
than or equal to

- Semantics:
 - If **%rdi** is less than or equal to **%rsi** (both interpreted as **signed value**), jump to the part of the code with a label **.L4**
- Under the hood:
 - **cmpq** instruction sets the condition codes
 - **jle** reads and checks the **condition codes**
 - If condition met, modify the Program Counter to point to the address of the instruction with a label **.L4**

How Should `cmpq` Set Condition Codes?

`cmpq %rsi, %rdi`

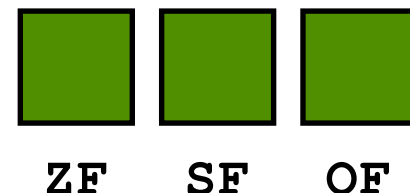
- Essentially, how do we know `%rdi <= %rsi`?
- Calculate `%rdi - %rsi`
- `%rdi == %rsi` if and only if `%rdi - %rsi == 0`
- `%rdi < %rsi` if and only if: ~~`%rdi - %rsi < 0`~~ (is it correct??)
 - `%rdi - %rsi < 0` and the result doesn't overflow, or
 - `%rdi - %rsi > 0` and the result does overflow

| | | |
|----------------|--|--|
| No Overflow | $\begin{array}{r} 001 \\ -) 010 \\ \hline 111 \end{array}$ | $\begin{array}{r} 1 \\ -) 2 \\ \hline -1 \end{array}$ |
| | 111 | -1 |
| Overflow | $\begin{array}{r} 101 \\ -) 011 \\ \hline 010 \end{array}$ | $\begin{array}{r} -3 \\ -) 3 \\ \hline -6 \end{array}$ |
| | 010 | -6 |

ZF Zero Flag (result is zero)

SF Sign Flag (result is negative)

OF Overflow Flag (result overflow)



How Should `cmpq` Set Condition Codes?

`cmpq %rsi, %rdi`

- Essentially, how do we know `%rdi <= %rsi`?
- Calculate `%rdi - %rsi`
- `%rdi == %rsi` if and only if `%rdi - %rsi == 0`
- `%rdi < %rsi` if and only if: ~~`%rdi - %rsi < 0`~~ (is it correct??)
 - `%rdi - %rsi < 0` and the result doesn't overflow, or
 - `%rdi - %rsi > 0` and the result does overflow

- `%rdi <= %rsi` if and only if
 - ZF is set, or
 - SF is set but OF is not set, or
 - SF is not set, but OF is set
- or simply: **ZF | (SF ^ OF)**

ZF Zero Flag (result is zero)

SF Sign Flag (result is negative)

OF Overflow Flag (result overflow)

| | | |
|-----------|-----------|-----------|
| 0 | 1 | 0 |
| ZF | SF | OF |

Conditional Branch Example

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq    %rsi,%rdi # x:y
    jle     .L4
    movq    %rdi,%rax
    subq    %rsi,%rax
    ret

.L4:      # x <= y
    movq    %rsi,%rax
    subq    %rdi,%rax
    ret
```

| Register | Use(s) |
|----------|--------------|
| %rdi | x |
| %rsi | y |
| %rax | Return value |

cmpq sets ZF, SF, OF

jle checks $ZF \mid (SF \wedge OF)$

| | | |
|----|----|----|
| 0 | 0 | 0 |
| ZF | SF | OF |

How Does the Hardware Check Overflow?

- ZF and SF are easily set by just examining the bits
- How about OF? How do we know A-B leads to overflow (A and B are treated as signed)
 - If $A < 0$ & $B > 0$, but the result > 0 , or
 - If $A > 0$ & $B < 0$, but the result < 0
 - So again, just have to check the bits

| | | | | |
|----------------|--|---|--|--|
| No Overflow | $\begin{array}{r} 001 \\ -) 010 \\ \hline 111 \end{array}$ | $\begin{array}{r} 1 \\ -) 2 \\ \hline -1 \end{array}$ | | |
| | | | | |
| Overflow | $\begin{array}{r} 101 \\ -) 011 \\ \hline 010 \end{array}$ | $\begin{array}{r} -3 \\ -) 3 \\ \hline 2 \end{array}$ | $\begin{array}{r} 011 \\ -) 100 \\ \hline 111 \end{array}$ | $\begin{array}{r} 3 \\ -) -4 \\ \hline -1 \end{array}$ |

Conditional Jump Instruction

```
cmpq    %rsi, %rdi
jbe     .L4
```

Jump to label if
below or equal to

- Semantics of **jbe**:
 - Treat the data in **%rdi** and **%rsi** as **unsigned values**.
 - If **%rdi** is less than or equal to **%rsi**, jump to the part of the code with a label **.L4**
- Under the hood:
 - **cmpq** instruction sets the condition codes
 - **jbe** reads and checks the condition codes
 - If condition met, modify the Program Counter to point to the address of the instruction with a label **.L4**

How Should `cmpq` Set Condition Codes?



`cmpq %rsi, %rdi`

- How do we know `%rdi <= %rsi`? This time for unsigned values
- Calculate `%rdi - %rsi`
- `%rdi == %rsi` if and only if `%rdi - %rsi == 0`
- `%rdi < %rsi` if a carry is generated during subtraction

| | | | |
|-------|------|---|---|
| | 001 | ← | 1 |
| -) | 111 | ← | 7 |
| <hr/> | | | |
| | C010 | | |

ZF Zero Flag (result is zero)

CF Carry Flag (for unsigned)

| | |
|---|---|
|  |  |
| CF | ZF |

How Should `cmpq` Set Condition Codes?

`cmpq %rsi, %rdi`

- How do we know `%rdi <= %rsi`? This time for unsigned values
- Calculate `%rdi - %rsi`
- `%rdi == %rsi` if and only if `%rdi - %rsi == 0`
- `%rdi < %rsi` if a carry/borrow is generated during subtraction

11111111 10000000

- `%rdi <= %rsi` (as unsigned) if and only if:
 - ZF is set, or
 - CF is set
- or simply: **ZF | CF**
- This is what `jbe` checks

ZF Zero Flag (result is zero)

CF Carry Flag (for unsigned)

| | |
|-----------|-----------|
| 1 | 0 |
| CF | ZF |

Putting It All Together

```
cmpq    %rsi,%rdi
jle     .L4
```

```
cmpq    %rsi,%rdi
jbe     .L4
```

- `cmpq` sets all 4 condition codes simultaneously
- ZF, SF, and OF are used when comparing signed value (e.g., `jle`)
- ZF, CF are used when comparing unsigned value (e.g., `jbe`)

```
11111111 10000000
cmpq 0xFF, 0x80
```

```
      10000000
- )  11111111
-----
c10000001
```

ZF Zero Flag

CF Carry Flag

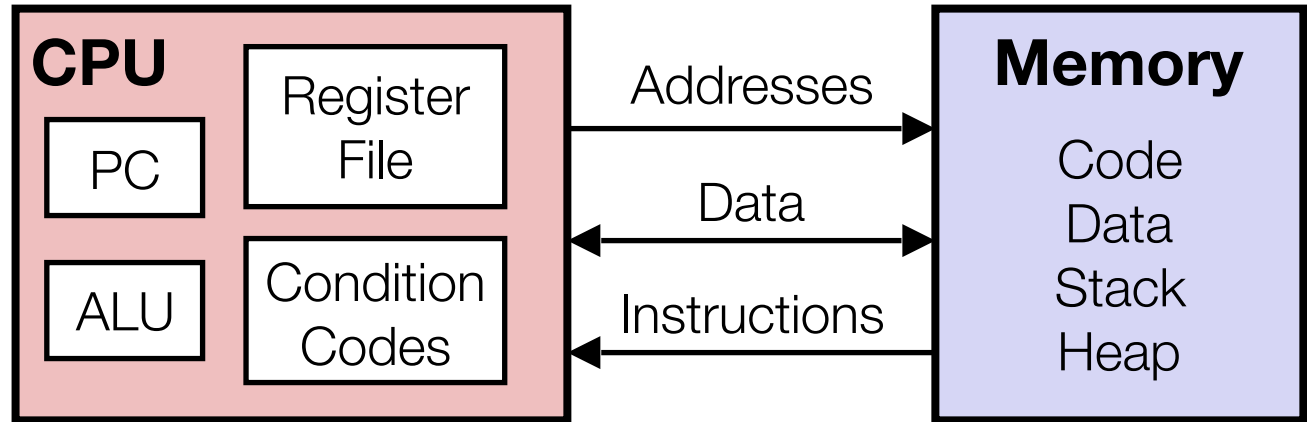
SF Sign Flag

OF Overflow Flag (for signed)

| | | | |
|----|----|----|----|
| 1 | 0 | 1 | 0 |
| CF | ZF | SF | OF |

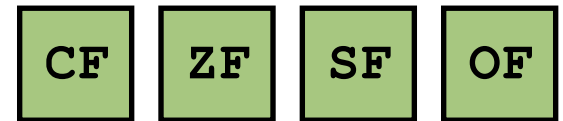
Condition Codes Hold Test Results

Assembly
Programmer's
Perspective
of a Computer



- Condition Codes

- Hold the status of most recent test
- 4 common condition codes in x86-64
- A set of special registers (more often: bits in one single register)
- Sometimes also called: Status Register, Flag Register



CF Carry Flag

ZF Zero Flag

SF Sign Flag

OF Overflow Flag (for signed)

Jump Instructions

- Jump to different part of code (designated by a label) depending on condition codes

| | | |
|------------|--------------------------|------------------------|
| jle | $(SF \wedge OF) \mid ZF$ | Less or Equal (Signed) |
|------------|--------------------------|------------------------|

| | | |
|------------|--------------|---------------------------|
| jbe | $CF \mid ZF$ | Below or Equal (unsigned) |
|------------|--------------|---------------------------|

Implicit Set Condition Codes

`addq %rax, %rbx`

- Arithmetic instructions implicitly set condition codes (think of it as side effect)
 - CF** set if `%rax + %rbx` generates a carry (i.e., unsigned overflow)
 - ZF** set if `%rax + %rbx == 0`
 - SF** set if `%rax + %rbx < 0`
 - OF** set if `%rax + %rbx` overflows when `%rax` and `%rbx` are treated as signed numbers
 - `%rax > 0, %rbx > 0, and (%rax + %rbx) < 0`, or
 - `%rax < 0, %rbx < 0, and (%rax + %rbx) >= 0`

```
if ((x+y)<0) {  
    ...  
}
```

`addq 0xFF, 0x80`

`jle 1000000`

`+) 11111111`

C01111111



CF **ZF** **SF** **OF**

Today: Control Instructions

- Control: Conditional branches (`if... else...`)
- Control: Loops (**`for`**, **`while`**)
- Control: Switch Statements (`case... switch...`)

“Do-While” Loop Example

- Popcount: Count number of 1's in argument x


do-while version

```
long pcount_do
(unsigned long x) {
    long result = 0;
    do {
        result += x & 0x1;
        x >>= 1;
    } while (x);
    return result;
}
```

goto Version


```
long pcount_goto
(unsigned long x) {
    long result = 0;
    loop:
        result += x & 0x1;
        x >>= 1;
        if(x) goto loop;
    return result;
}
```

“Do-While” Loop Assembly



```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

| Register | Use(s) |
|----------|---------------|
| %rdi | Argument x |
| %rax | result |



```
    movl    $0, %rax    # result = 0
.L2:                                # loop:
    movq    %rdi, %rdx
    andl    $1, %rdx    # t = x & 0x1
    addq    %rdx, %rax   # result += t
    shrq    $1, %rdi    # x >>= 1
    jne     .L2          # if (x) goto loop
    ret
```

General “Do-While” Translation

do-while version

```
<before>;  
do {  
    body;  
} while (A < B) ;  
<after>;
```



goto Version

```
<before>  
.L1: <body>  
    if (A < B)  
        goto .L1  
<after>
```

Replace with a

conditional jump
instruction

```
<before>  
.L1: <body>  
    cmpq B, A  
    jl .L1  
<after>
```

Assembly
Version

General “While” Translation

while version

```
<before>;  
while (A < B) {  
    body;  
}  
<after>;
```



goto Version

```
<before>  
goto .L2  
.L1: <body>  
.L2: if (A < B)  
      goto .L1  
<after>
```



Assembly
Version

```
<before>  
jmp .L2  
.L1: <body>  
.L2: cmpq A, B  
      jg .L1  
<after>
```

“While” Loop Example

while version

```
long pcount_while
(unsigned long x) {

    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

goto Version

```
long pcount_goto_jtm
(unsigned long x) {
    long result = 0;
    goto test;
loop:
    result += x & 0x1;
    x >>= 1;
test:
    if(x) goto loop;
    return result;
}
```


“For” Loop Example

```
for (init; test; update) {  
    body  
}
```

```
//assume unsigned int is 4 bytes  
long pcount_for (unsigned int x)  
{  
  
    size_t i;  
    long result = 0;  
    for (i = 0; i < 32; i++)  
    {  
        result += (x >> i) & 0x1;  
    }  
    return result;  
}
```

init

i = 0

test

i < 32

update

i++

body

```
{  
    result += (x >> i)  
    & 0x1;  
}
```

Convert “For” Loop to “While” Loop

For Version

```
before;  
for (init; test; update) {  
    body;  
}  
after
```

While Version

```
before;  
init;  
while (test) {  
    body;  
    update;  
}  
after;
```

Assembly Version

```
before  
init  
jmp .L2  
.L1: body  
      update  
.L2: cmpq A, B  
      jg .L1  
after
```

Today: Control Instructions

- Control: Conditional branches (**if... else...**)
- Control: Loops (**for, while**)
- Control: Switch Statements (**case... switch...**)

Switch Statement Example

```
long switch_eg (long x, long y, long z)
{
    long w = 1;
    switch(x) {
        case 1:
            w = y*z;
            break;
        case 2:
            w = y/z;
        case 3:
            w += z;
            break;
        case 5:
        case 6:
            w -= z;
            break;
        default:
            w = 2;
    }
    return w;
}
```

Fall-through case

Multiple case labels

For missing cases, fall back to default

Converting to a cascade of if-else statements is simple, but cumbersome with too many cases.

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
  case val_0:  
    Block 0  
  case val_1:  
    Block 1  
  
  ....  
  case val_n-1:  
    Block n-1  
}
```

Jump Table

JTab:

| |
|---------|
| Targ0 |
| Targ1 |
| Targ2 |
| • |
| • |
| • |
| Targn-1 |

Jump Targets

Targ0: Code Block 0

Targ1: Code Block 1

Targ2: Code Block 2

•
•
•

Targn-1: Code Block n-1

- Each code block starts from a unique address (Targ0, Targ1, ...)
- Jump table stores all the target address
- Use the case value to index into the jump table to find where to jump to

Assembly Directives (Pseudo-Ops)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

- Directives:

- Not real instructions, but assist assembler. Think of them as messages to help the assembler in the assembly process.

- **.quad**: tells the assembler to set aside the next 8 bytes in memory and initialize with the value of the operand (a label here, which itself is an address)
- **.align**: tells the assembler that addresses of the the following data will be aligned to 8 bytes
- **.section**: denotes different parts of the object file
- **.rodata**: read-only data section

Jump Table and Jump Targets

Jump Table

```
.section .rodata
.align 8
.L4:
    .quad .LD # x = 0
    .quad .L1 # x = 1
    .quad .L2 # x = 2
    .quad .L3 # x = 3
    .quad .LD # x = 4
    .quad .L5 # x = 5
    .quad .L5 # x = 6
```

jmp .L3 will go
to .L3 and start
executing from there

Jump Targets

```
.L1:                                # Case 1
    movq    %rsi, %rax
    imulq   %rdx, %rax
    jmp     .done

.L2:                                # Case 2
    movq    %rsi, %rax
    cqto
    idivq   %rcx

.L3:                                # Case 3
    addq    %rcx, %rax
    jmp     .done

.L5:                                # Case 5,6
    subq    %rdx, %rax
    jmp     .done

.LD:                                # Default
    movl    $2, %eax
    jmp     .done
```

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
  case val_0:  
    Block 0  
  case val_1:  
    Block 1  
  
  ....  
  case val_n-1:  
    Block n-1  
}
```

Jump Table

| | |
|-------|-----------|
| JTab: | Target0 |
| | Target1 |
| | Target2 |
| | • |
| | • |
| | • |
| | Targetn-1 |

Jump Targets

Target0: Code Block 0

Target1: Code Block 1

Target2: Code Block 2

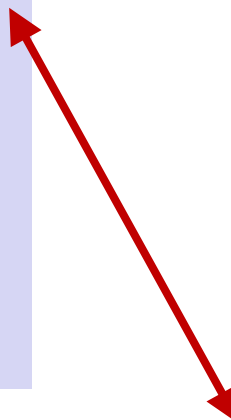
•
•
•

Targetn-1: Code Block n-1

.Done:

Code Blocks (x == 1)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```



```
switch(x) {
case 1:      // .L1
    w = y*z;
    break;
    ...
}
```

| Register | Use(s) |
|----------|--------------|
| %rdi | Argument x |
| %rsi | Argument y |
| %rdx | Argument z |
| %rax | Return value |

```
.L1:
    movq    %rsi, %rax    # y
    imulq   %rdx, %rax    # y*z
    jmp     .done
```

Code Blocks (x == 2, x == 3)

```
.section .rodata
.align 8
.L4:
    .quad .L2 # x = 2
    .quad .L3 # x = 3
    .quad .L2 # x = 4
    .quad .L5 # x = 5
    .quad .L5 # x = 6
```

```
switch(x) {
...
case 2:      // .L2
    w = y/z;
    /* Fall Through */
case 3:      // .L3
    w += z;
    break;
...
}
```

```
.L2:                                # Case 2
    movq    %rsi, %rax
    cqto
    idivq   %rcx                    # y/z

.L3:                                # Case 3
    addq    %rcx, %rax              # w += z
    jmp     .done
```

| Register | Use(s) |
|----------|--------------|
| %rdi | Argument x |
| %rsi | Argument y |
| %rdx | Argument z |
| %rax | Return value |

Code Blocks (x == 5, x == 6, default)

```
.section .rodata
.align 8
.L4:
.quad .LD # x = 0
.quad .L1 # x = 1
.quad .L2 # x = 2
.quad .L3 # x = 3
.quad .LD # x = 4
.quad .L5 # x = 5
.quad .L5 # x = 6
```

```
switch(x) {
...
case 5: // .L5
case 6: // .L5
    w -= z;
    break;
default: // .LD
    w = 2;
}
```

| Register | Use(s) |
|----------|--------------|
| %rdi | Argument x |
| %rsi | Argument y |
| %rdx | Argument z |
| %rax | Return value |

```
.L5:                                # Case 5,6
    subq %rdx, %rax # w -= z
    jmp  .done
.LD:                                # Default:
    movl $2, %eax  # 2
    jmp  .done
```

Implementing Switch Using Jump Table

Switch Form

```
switch(x) {  
  case val_0:  
    Block 0  
  case val_1:  
    Block 1  
  
  ....  
  case val_n-1:  
    Block n-1  
}
```

Jump Table

| | |
|------|-----|
| .LJ: | .LD |
| | .L1 |
| | .L2 |
| | • |
| | • |
| | • |
| | .L5 |

Jump Targets

.LD: Code Block 0

.L1: Code Block 1

.L2: Code Block 2

•
•
•

.L5: Code Block n-1

.Done:

- The only thing left...
 - How do we jump to different locations in the jump table depending on the case value?

Indirect Jump Instruction

The address we want to jump to is stored at $.LJ + 8 * x$

```
.section .rodata
.align 8
.LJ:
    .quad .LD # x = 0
    .quad .L1 # x = 1
    .quad .L2 # x = 2
    .quad .L3 # x = 3
    .quad .LD # x = 4
    .quad .L5 # x = 5
    .quad .L5 # x = 6
```

```
# assume x in %rdi
movq  .LJ(,%rdi,8), %rax
jmp   *%rax
```

- Indirect Jump: **jmp *%rax**
 - %rax specifies the address to jump to (PC = %rax)
- Direct Jump (**jmp .LJ**), directly specifies the jump address
- Indirect Jump specifies where the jump address is located

An equivalent syntax in x86:

```
jmp    *.LJ(,%rdi,8)
```

Summary

Switch Form

```
switch(x) {  
  case val_0:  
    Block 0  
  case val_1:  
    Block 1  
  
  ....  
  case val_n-1:  
    Block n-1  
}
```

Jump Table

JTab:

| |
|---------|
| Targ0 |
| Targ1 |
| Targ2 |
| • |
| • |
| • |
| Targn-1 |

Jump Targets

Targ0: Code Block 0

Targ1: Code Block 1

Targ2: Code Block 2

•
•
•

Targn-1: Code Block n-1

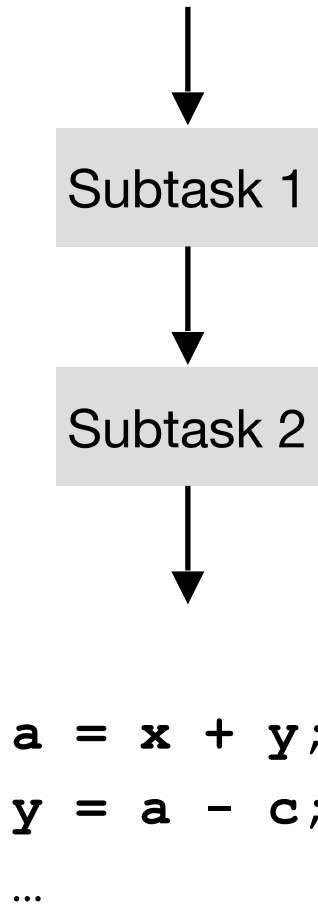
- Each code block starts from a unique address (Targ0, Targ1, ...)
- Jump table stores all the target address
- Use the case value to index into the jump table to find where to jump to

Not The Only Way

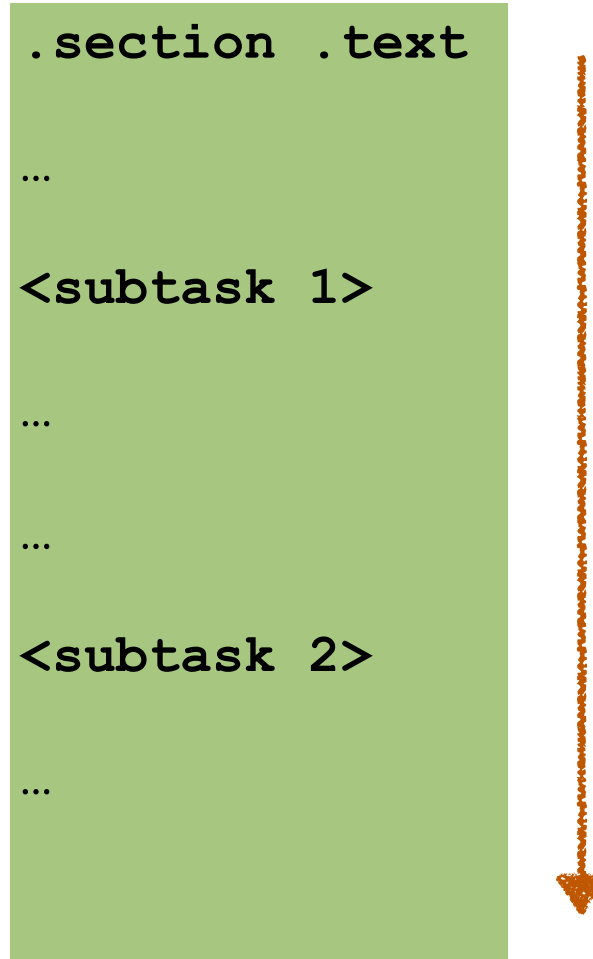
- Jump table might not be the most efficient implementation; certainly not the only way to implement switch-case.
- What if x can take a very large value range. Do we need to have a giant jump table?
- Let's say x can be any integer from 1 to 1 million, but anything between 8 and 1 million fall back to the default case. Can we avoid a 1 million entry jump table (which isn't too bad if you calculate the size)?
 - Have an if-else check first followed by an 8-entry table.

Summary

Sequential

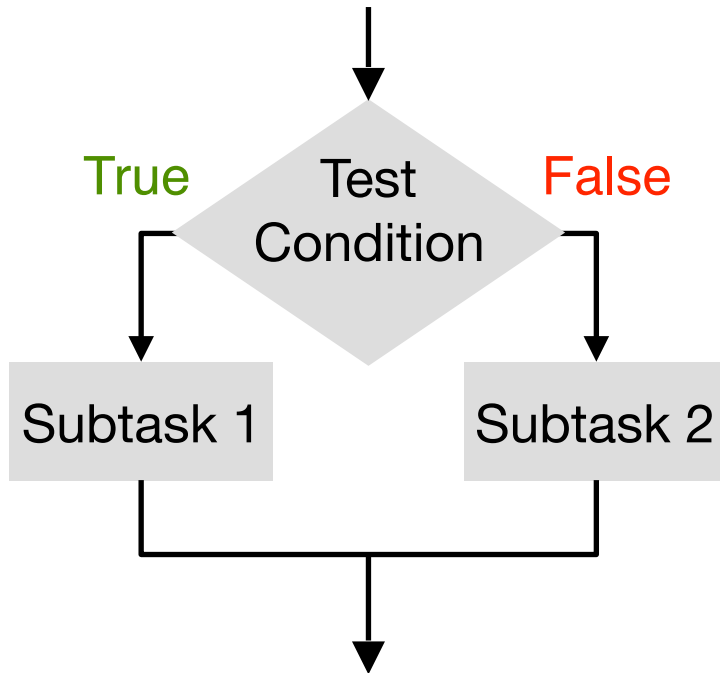


Memory



Summary

Conditional



```
if (x > y) r = x - y;  
else r = y - x;
```

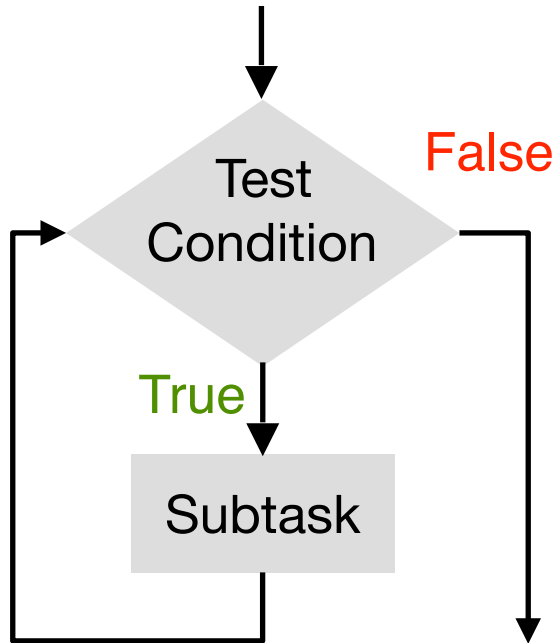
Memory

```
.section .text  
  
...  
cmpq  
jle .L2  
.L1 <subtask 1>  
  
...  
  
jmp .done  
.L2 <subtask 2>  
  
...  
.done
```



Summary

Iterative



```
while (x > 0) {  
    x-- ;  
}
```

Memory

```
.section .text
```

```
...
```

```
addq  
jmp .L2  
.L1:
```

```
...
```

```
<subtask>
```

```
...
```

```
.L2:  
    cmpq A, B  
    jg .L1
```

```
...
```

```
...
```

