

CUBE: A MEMORY-LITE COSMOLOGICAL N -BODY ALGORITHM

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Abstract

Cosmological large scale structure N -body simulations are computation-light, memory-heavy problems in supercomputing. Traditional N -body simulation algorithms use at least 24-byte memory per particle, of which six 4-byte single precision floating point numbers keep track phase space coordinates of each particle. Here we present the algorithm and accuracy of a new parallel, memory-lite, Particle-Mesh based N -body code, where each particle can occupy as low as 6-byte memory. This is accomplished by storing relative position and relative velocity of each particle, in the format of 1-byte-integer, respect to their averaged value of a mesh-grid. The remaining information is given by complimentary density and velocity fields, which are negligible in memory space, and proper ordering of particles, which gives no extra memory. Our numerical experiments show that this integer based N -body algorithm provides acceptable accuracy compared to traditional algorithm in cosmological simulations. This significant lowering of memory-to-computation ratio breaks the bottleneck of scaling up and speeding up large cosmological N -body simulations on multi-core and heterogenous computing systems.

1. INTRODUCTION

N -body simulation is a powerful tool to solve the highly nonlinear dynamic problems (Hockney & Eastwood 1988). It is widely used in cosmology to model the formation and evolution of the large scale structure (LSS). With the fast development of parallel supercomputers, we are able to simulate a system of more than a trillion (10^{12}) N -body particles. To date the largest N -body simulation in application is the “TianNu” simulation (Yu et al. 2017; Emberson et al. 2017) run on the TianHe-2 supercomputer by cosmological simulation code CUBEP3M (Harnois-Déraps et al. 2013). It uses nearly 3×10^{12} particles to simulate the cold dark matter (CDM) and cosmic neutrino evolution through the cosmic age.

The N -body simulations use considerable amount of memory, because the phase space coordinates (x, y, z, v_x, v_y, v_z) of each N -body particle must be stored as, at least, six single precision floating point numbers (24 bytes). Contrarily, their computing workload can be alleviated by many algorithms, like Particle-Mesh [cite] and Tree [cite], scale as $o(N \log N)$ or even $o(N)$. On the other hand, modern supercomputer systems use multi cores, many integrated cores (MIC) and even densely parallelized GPUs, bringing orders of magnitude higher computing power, whereas these architectures usually have limited memory allocation. Thus, the computation-light but memory-heavy applications, compared to matrix multiplication and decomposition calculations, are less suitable for fully usage of the computing power of modern supercomputers. For example, although native and offload modes of CUBEP3M are able to run on the Intel Xeon-PHI MIC architectures, with the require ment

of enough memory, TianNu simulation were done on TianHe-2 with only its CPUs – 73% of the total memory but only 13% of the total computing power.

We present a new N -body simulation code CUBE, using as low as 6 byte per particle (here after we use “bpp” referring to “byte per particle”). We show that it gives accurate results in cosmological LSS simulations. The method is presented in §2, and a comparison between this method and traditional method is shown in §3. Discussions and conclusions are in §??.

2. METHOD

The most memory consuming part of a N -body simulation is usually the phase space coordinates of N -body particles – 24 bpp (4 single precision floating numbers) must be used to store each particles’ 3D position and velocity vectors. CUBEP3M, an example of a memory-efficient parallel N -body code, can use as low as 40 bpp in sacrificing computing speed (Harnois-Déraps et al. 2013). This includes the phase coordinates (24 bpp) for particles in physical domain and buffered region, a linked list (4 bpp), and a global coarse mesh and local fine mesh. 4-byte real numbers are not necessarily adequate in representing the *global* coordinates in simulations. If the box size is many orders of magnitude larger than interactive distance between particles, especially in the field of resimulation of dense subregions, double precision (8-byte) coordinates are needed to avoid truncation errors. Another solution is to record *relative* coordinates for both position and velocity. CUBE replaces the coordinates and linked list 24+4=28 bpp memory usage with an integer based storage, reduces the basic memory usage from 28 bpp down to 6 bpp, described as following 2.1 and 2.2. The algorithm is described in 2.3.

2.1. Particle position storage

We construct a uniform mesh throughout the space and each particle belongs to its parent cell of the mesh. Instead of storing global coordinates of each particle, we store its offset relative to its parent cell which contains the particle. We divide the cell, in each dimension d , evenly into $2^8 = 256$ bins, and use a 1-byte (8 bits) integer $\chi_d \in \{-128, -127, \dots, 127\}$ to indicate which bin it locates in this dimension. The global locations of particles are given by cell-ordered format in memory space, and a complimentary number count of particle number in this mesh (density field) will give complete information of particle distribution in the mesh. Then the global coordinate in d th dimension x_d is given by $x_d = (n_c - 1) + (\chi_d + 128 + 1/2)/256$, where $n_c = 1, 2, \dots, N_c$ is the index of the coarse grid. The mesh is chosen to be coarse enough such that the density field takes negligible memory. This coarse density field can be further compressed into 1-byte integer format, such that a 1-byte integer show the particle number in this coarse cell in range 0 to 255. In the densest cells (rarely happened) where there are ≥ 255 particles, we can just write 255, and write the actual number as a 4-byte integer in another file.

In a simulation with volume L^3 and N_c^3 coarse cells, particle positions are stored with a precision of $L/(256N_c)$. The force calculation (e.g. softening length) should be configured much finer than this resolution, discussed in later sections. On the other hand, particle position can also be stored as 2-byte (16 bits) integers to increase the resolution. In this case, each coarse cell is divided into $2^{16} = 65536$ bins and the position precision is $L/(65536N_c)$, precise enough compared to using 4-byte global coordinates, see later results. We denote this case “x2” and denote using 1-byte integers for positions “x1”.

We collectively write the general position conversion formulae

$$\chi_d = [2^{8n_x}(x_d - [x_d])] - 2^{8n_x-1}, \quad (1)$$

$$x_d = (n_c - 1) + 2^{-8n_x} (\chi_d + 2^{8n_x-1} + 1/2), \quad (2)$$

where $[]$ is the operator to take the integer part. $n_x \in \{1, 2\}$ is the number of bytes used for each integer, x_d and χ_d are floating and integer version of the coordinate. The velocity counterpart of them are $n_v = 1, 2$, v_d and v_c . $n_c = 1, 2, \dots$ is the coarse grid index, given by the ordering of particles and a integer based particle density field (see 2.3.1). The position resolution for n_x -byte integer, “ xn_x ”, is $2^{-8n_x} L/N_c$.

As a $n_x = 1$, 1D ($d = 1$), 4-coarse-cell ($N_c = 4$) example, if

$$\chi_1 = (-128, 127, 0, 60),$$

particle number density

$$\rho_c^{1D} = (1, 0, 2, 1),$$

then in unit of coarse cells,

$$x_1 = (0.001953125, 2.998046875, 2.501953125, 3.736328125).$$

2.2. Particle velocity storage

Similarly, actual velocity in d th dimension v_d is decomposed into an averaged velocity field on the same coarse

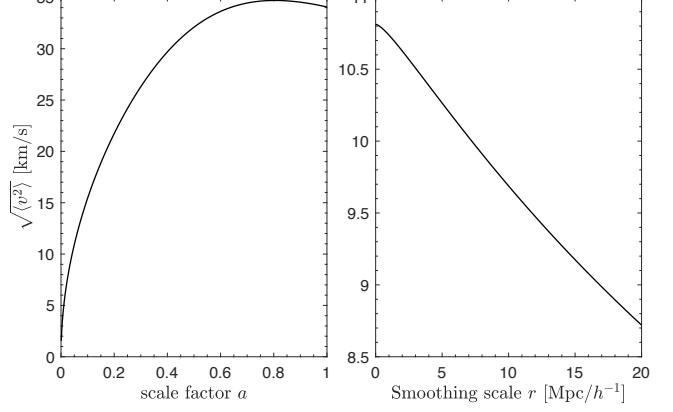


FIG. 1.— Xin, could you please check the units of both y -axis?

grid v_c and a residual Δv relative to this field:

$$v_d = v_c + \Delta v. \quad (3)$$

v_c is always recorded and kept updated, and occupies negligible memory. We then divide velocity space Δv into uneven bins, and use a n_v -byte integer to indicate which Δv bin the particle is located.

The reason why we use uneven bins is that, slower particles are more abundant compared to faster ones, and one should better resolve slower particles tracing at least linear evolution. On the other hand, there could be extreme scattering particles (in case of particle-particle force), and we can safely ignore or less resolve those non-physical particles. One of the solution is that, if we know the probability distribution function (PDF) $f(\Delta v)$ we divide its cumulative distribution function (CDF) $F(\Delta v) \in (0, 1)$ into 2^{8n_v} bins to determine the boundary of Δv bins, and particles should evenly distribute in the corresponding uneven Δv bins. Practically we find that either $f(v_d)$ or $f(\Delta v)$ is close to Gaussian, so we can use Gaussian CDF, or any convenient analytic functions which are close to Gaussian, to convert velocity between real numbers and integers.

The essential parameter of the velocity distribution is its variance. On non-linear scale, the velocity distribution function is non-Gaussian. However, to the first order approximation, we simply assume it as Gaussian and characterized by the variance

$$\sigma_v(a, k) = \frac{1}{3}(aHfD)^2 \int_0^k d^3 q \frac{P_L(q)}{q^2}, \quad (4)$$

where $a(z)$ is the scale factor, $H(z)$ the Hubble parameter, D is the linear growth factor, $f = d \ln D / d \ln a$, and $P_L(k)$ is the linear power spectrum of density contrast at redshift zero. $\sigma_v(a, k)$ is a function of cosmic evolution a and a smoothing scale k , or r (see Figure 1). Δv is the velocity dispersion relative to the coarse grid, so we approximate its variance as

$$\sigma_\Delta^2(a) = \sigma_v^2(a, r_c) - \sigma_v^2(a, r_p), \quad (5)$$

where r_c is the scale of coarse grid, and r_p is the scale of average particle separation. In each dimension of 3D velocity field, we use $\sigma_\Delta^2(a)/3$ according to the equipartition theorem. On different scales, we measure the statistics of v_d , v_c and Δv and find good agreement with the above model.



FIG. 2.— Spacial decomposition in CUBE in a 2D analogy. In this example, there are 2 images per dimension ($\text{nn} = 2$), and 2 tiles per image per dimension ($\text{nnt} = 2$). The orange boxes show the overlapped physical+buffer regions, inside of which one physical regions is indicated with green.

The simulation results are very insensitive if we manually tune the variance of the model σ_Δ within an order of magnitude. However, in $n_\nu = 1$ case, the method of using uneven bins gets much better results than simply using equal bins between minimum and maximum values [$\min(\Delta v)$, $\max(\Delta v)$]. So, one can safely use a standard Λ CDM (cold dark matter with a cosmological constant as dark energy) for slightly different cosmological models, in equation (4). In CUBE, the velocity conversion takes the formula

$$\nu_d = \left\lfloor (2^{8n_\nu} - 1)\pi^{-1} \tan^{-1} \left((v_d - v_c) \sqrt{\pi/2\sigma_\Delta^2} \right) \right\rfloor, \quad (6)$$

$$v_d = v_c + \tan \left(\frac{\pi\nu_d}{2^{8n_\nu} - 1} \right) \sqrt{2\sigma_\Delta^2/\pi}, \quad (7)$$

where $\lfloor \cdot \rfloor$ is the operator to take the nearest integer. Tangent functions are convenient and compute very fast. Compared to error functions used in Gaussian case, they take the same variance at $\nu_d = 0$ but resolve high velocities relatively better. Note again that proper choice of conversion formulae and σ_Δ only optimizes the velocity space sampling, but does not affect the physics.

Initially, particles are generated by initial condition generator, at a higher redshift. The coarse grid density field v_c is also generated at this step by averaging all particles in the coarse cell. A global σ_Δ is calculated by equation (5), where linear approximation is hold. Then velocities are stored by equation (6). During the simulation, v_c is updated every time step, and a nonlinear σ_Δ is measured directly from the simulation, and can be simply used in the next time step, after scaled by the ratio of growth factors between two adjacent time steps. More details see section 2.3.2.

2.3. Code overview

CUBE uses a 2-level PM force calculation, same as CUBEP3M. However, in order to apply the integer based

format to the N -body simulation, substantial structural changes need to be done. CUBE is written in Coarray Fortran, where Coarray features implement MPI communication between computation nodes/images¹. The algorithm is described in this language.

2.3.1. Spacial decomposition

CUBE uses cubic decomposition structures. The global simulation volume is decomposed into nn^3 cubic sub-volumes with nc coarse grids per side, and each of these are assigned to a coarray *image*. Inside of an image, the sub-volume is further decomposed into nnt^3 cubic *tiles*, with $\text{nt}=\text{nc}/\text{nnt}$ coarse grids per side, which is a essential unit in the calculations. Each tile is surrounded by a *buffer* region which is ncb coarse cells thick. The buffer is designed for two reasons: (1) computing the fine mesh force, whose cut-off length $\text{nforce.cutoff} \leq \text{ncb}$, and (2) collecting all possible particles travelling from a tile's buffer region to its center, *physical* region. Figure 2 shows the spacial decomposition in a 2-dimensional analogy, with $\text{nn} = 2$ and $\text{nnt} = 2$. According to this spatial decomposition, and as discussed in the last two subsections, we declare $\{\rho_c, \chi_d, v_c, \nu_d\}$, or in variable expression $\{\text{rho_c}, \text{xp}, \text{vfield}, \text{vp}\}$ as follows:

```
integer(1) rho_c(nex,nex,nex,nnt,nnt)[nn,nn,*]
integer(nx) xp(npmax)[nn,nn,*]
real(4) vfield(nex,nex,nex,nnt,nnt)[nn,nn,*]
integer(nv) vp(npmax)[nn,nn,*]
```

where $\text{nex} = \text{nt} + 2 \times \text{ncb}$ covers the buffer region on both sides, nnt is the tile dimensions, and nn is the image codimensions². We denote actual number of particles in a given image nplocal , and $\text{npmax} > \text{nplocal}$ is a value (discussed in Section 2.4) large enough to store particles. The particle position and velocity arrays xp and vp (equivalent to χ_d and ν_d in 2.1 and 2.2 respectively) are required to be sorted according to the same memory layout of rho_c , such that n_c and thus global positions of particles x_d can be obtained by equation (2). $\{\text{rho_c}, \text{xp}, \text{vfield}, \text{vp}\}$ provides a complete information of positions and velocities of particles, and we call it a snapshot, or checkpoint.

2.3.2. Algorithm

Figure 3 shows the overall structure of the code and these subroutines are described in following paragraphs.

initialize and **read_particles** – The subroutine **initialize** creates necessary FFT plans and read in configuration files telling the program at which redshifts we need to do checkpoints, halofinds, or stop the simulation. Force kernels **kern_c**, **kern_f** are also created or read in here. In **read_particles**, for each image, we read in all particles in *physical* regions of every tile (indicated in Figure 2), i.e., $\{\text{rho_c}, \text{xp}, \text{vfield}, \text{vp}\}$. These are obtained by the initial condition generator (see Appendix A). Because physical regions of tile are *complete* and *disjoint* in space, particles at this stage are also complete and disjoint. We call it “disjoint state”. At this stage,

¹ Images are the concept of computing nodes or MPI tasks in Coarray Fortran. We use this terminology in this paper.

² Coarray Fortran concept. Codimensions can do communications between images.

```

program CUBE
  call initialize
  call read_particles
  call buffer_density
  call buffer_xp
  call buffer_vp
  do
    call timestep
    call update_xp
    call buffer_density
    call buffer_xp
    call update_vp
    call buffer_vp
    if(checkpoint_step) then
      call update_xp
      call checkpoint
      if (final_step) exit
      call buffer_density
      call buffer_xp
      call buffer_vp
    endif
  enddo
  call finalize
end

```

FIG. 3.— Overall structure of CUBE.

`rho_c`'s buffer regions of each tile are 0, and the elements of `xp` and `vp` beyond physical number $-xp(nplocal+1:)$ and $vp(nplocal+1:)$ can be arbitrary and are not used.

`buffer_density`, `buffer_x` and `buffer_v` – In order to use the integer based format, `xp` and `vp` must always be ordered, and their number density field `rho_c` must always be present. In updating arrays of `xp` and `vp` (not simultaneously) of a local tile, a vicinity (buffer) region of the physical region is also needed. First, buffer regions of `rho_c` is synchronized between tiles and images by subroutine `buffer_density`. Then, by subroutines `buffer_x` and `buffer_v` respectively, `xp` and `vp` are updated to contain common, buffered particles, in an order according to the new, buffered `rho_c`. We call this stage “buffered state”. After `particle_initialize` is done by all images (synchronized by a “`sync all`”, which is equivalent to a `mpi_barrier`), these three subroutines are called and particles are converted from disjoint state to buffered state.

`timestep` – We operate a Runge-Kutta 2 method for time integration. i.e., we update position (D =drift) and velocity (K =kick) every half time step. For n time steps, the operation would be $(DKKD)^n$ which is 2nd order accurate. The actual simulation applies varied time steps. In each iteration of the main loop, we firstly call `timestep`, where a increment of time `dt` is controlled by particles' maximum velocities, accelerations, cosmic expansion and any other desired conditions.

`update_xp` – According to `dt`, subroutine `update_xp` updates the positions of particles in a “gather” algorithm (in contrast, CUBEP3M uses “scatter” algorithm) tile by tile. For each particle, `xp` and `vp` are converted to x_d and v_d by Equations(2,7). Because each tile is in the buffered state with buffer depth `ncb`, we are able to collect all possible particles whose $v_d \times dt < ncb$ from physical+buffer region to its physical region.

In order to keep particles ordered, for each tile, we first do $x_d = x_d + v_d \times dt$ on all particles to obtain an updated density and velocity field on the tile, `rho_c_new` and `vfield_new`. Then, this calculation is done on

```

subroutine update_xp
  do (each tile)
    do (each coarse grid)
      do (each particle)
        calculate particle's new coarse grid position
        update rho_c_new
        update vfield_new
      enddo
    enddo
    do (each coarse grid)
      do (each particle)
        calculate particle's new accurate position
        calculate particle's index
        update xp_new(index) & vp_new(index)
      enddo
    enddo
    do (each coarse grid)
      discard buffer information
      replace xp and vp for this tile
    enddo
  enddo
  sync all
  update velocity dispersion
  check nplocal & npglobal
end

```

FIG. 4.— Pseudocode for subroutine `update_xp`.

the same tile again³ to generate a new, local particle list `xp_new` and `vp_new` on the tile by Equations(1,6). Here, the ordering of particles depends on `rho_c_new`, and `vfield_new` is used as v_c in Equation(6). Then, another iteration is done on this tile to discard buffer regions of `{rho_c_new, xp_new, vfield_new, vp_new}`, converting buffer state to disjoint state, and it replaces corresponding part of `{rho_c, xp, vfield, vp}`. When all tiles are iterated, the entire image is in disjoint state, and `nplocal` is updated. Then all images are synchronized by `sync all` and we sum up `nplocal` over images to check if total number of particles `npglobal` is conserved. These steps are summarized in Figure 4.

`update_vp` – After `update_xp`, or drift D , we call `buffer_density` and `buffer_xp` in order that particle positions are in buffered state, based on which, we update velocities (kick K) of particles in physical region of each tile. CUBE uses a 2-level particle mesh scheme (Harnois-Déraps et al. 2013). Local fine forces have a force cutoff `nforce_cutoff` $\leq ncb$, i.e., $F_{\text{fine}}(r > ncb) = 0$. So, if we apply a fine-grid particle-mesh on an extended tile with buffered-depth `ncb` (see Figure 2 and regard it periodic), the force on physical regions does not depend on the false periodic boundary assumption, and the resulting fine force F_{fine} and velocity update on physical region is correct.

The compensating coarse grid force F_{coarse} is globally computed by using a coarser (usually by factor of 4) mesh by dimensional splitting – a distributed-memory cubic decomposed 3D coarse density field is interpolated by particles, and we Fourier transform data in consecutive three dimensions with global transposition in between (known as the pencil decomposition). After the multiplication of force kernels, the inverse transform takes place to get the cubic distributed coarse force field F_{coarse} , upon which velocities are updated again.

An optional particle-particle (PP) force F_{pp} can be called to increase the force resolution and the velocities

³ This repetition scales as $o(N)$ and is computational inexpensive.

```

subroutine update_vp
  do (each tile)
    calculate fine mesh density
    calculate fine mesh force: local FFT
    update fine mesh velocity
    update maximum acceleration
    if (PP force) then
      PP calculation & velocity update
      update maximum acceleration
    endif
  enddo
  sync all
  calculate coarse mesh density
  calculate coarse mesh force: global pencil FFT
  update coarse mesh velocity
  update maximum acceleration
  update maximum velocity
end

```

FIG. 5.— Pseudocode for subroutine `update_vp`.
are updated again. The collective operations of velocity update is Equation(7), $v_d = v_d + F_{\text{total}}$ and Equation(6).

The maximum accelerations $\max(\dot{v}_{\text{fine}})$, $\max(\dot{v}_{\text{coarse}})$, $\max(\dot{v}_{\text{pp}})$ and maximum of velocities $\max(v_d)$ are collected controlling dt for the `timestep` in the next iteration. We also update σ_{Δ}^2 according to the new $v_d - v_c$.

By far, `vp` in physical regions are correctly updated. Remember that particle locations in the buffer regions has updated before PM and remained unchanged. So we simply call `buffer_v` again such that the `update_x` in the next iteration will be done correctly. We call `update_vp` to convert `vp` into buffered state. These steps are summarized in Figure 5.

`checkpoint` – If a desired redshift is reached, we execute the last drift step in the $(DKKD)^n$ operation by `update_xp`, and call `checkpoint` to save the disjoint state of $\{\text{xp}, \text{vp}, \text{rho}_c, \text{vfield}\}$ on disk. Related operations, like run-time halo finder, projections are also done at this point. If final desired redshift is reached, `final_step` let us exit the main loop. These corresponding logical variables are controlled in `timestep`.

`finalize` – Finally, in `finalize` subroutine we destroy all the FFT plans and finish up any timing or statistics taken in the simulation.

2.4. Memory layout

Here we list the memory-consuming arrays and how they scale with different configurations of the simulation. We classify them into (1) arrays of particles, (2) coarse mesh arrays and (3) fine mesh arrays.

2.4.1. Arrays of particles

The arrays of particles contains `xvp(6,npmax)` and a temporary `xvp_new(6,np.tile)`⁴. `npmax` is the number of particles `xvp` can store:

$$\text{npmax} = \langle \text{nplocal} \rangle \left(1 + \frac{2 \times \text{ncb}}{\text{nt}} \right)^3 (1 + \epsilon_{\text{image}}), \quad (8)$$

where $\langle \text{nplocal} \rangle$ is the average number of particles per image. The second term above let us store particles in buffer regions, and the third term $1 + \epsilon_{\text{image}}$ takes into account the inhomogeneity of `nplocal` on different images. When each image models smaller physical scales, ϵ_{image} should be set larger.

⁴ We use “`xvp`” to represent $\{\text{xp}, \text{vp}\}$, and use “`xvp_new`” to represent $\{\text{xp_new}, \text{vp_new}\}$.

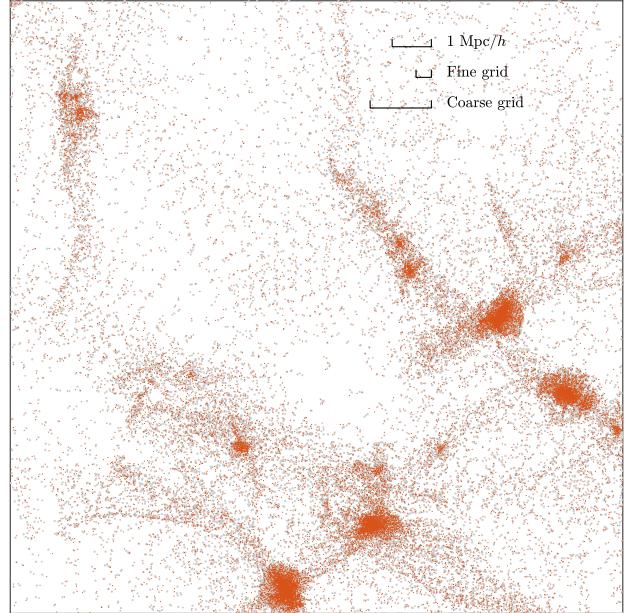


FIG. 6.— particles.

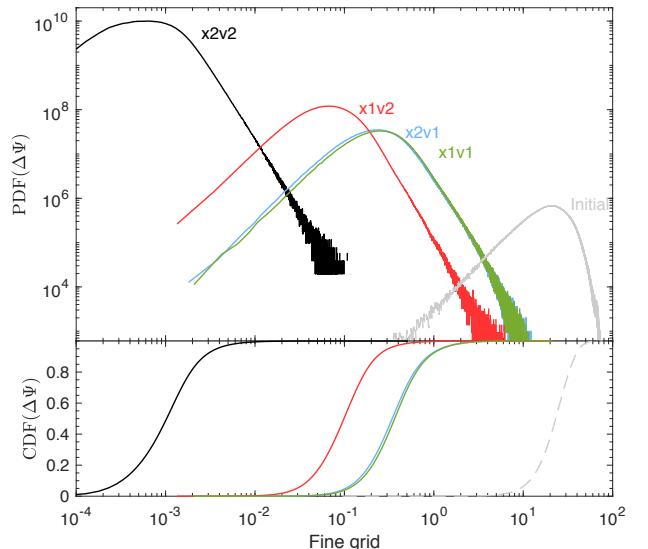


FIG. 7.— Displacement.

`xvp_new` stores temporary particles only on tiles:

$$\text{nptile} = \langle \text{nplocal} \rangle \left(\frac{1}{\text{nnt}} \right)^3 (1 + \epsilon_{\text{tile}}), \quad (9)$$

where, recall that `nnt` is the number of tiles per image per dimension, and similarly ϵ_{tile} controls the inhomogeneity on scale of tiles. Larger `nnt` causes more inhomogeneity on smaller tiles, and ϵ_{tile} can be much larger than ϵ_{image} , however the term nnt^{-3} decreases much faster. So, the majority memory is occupied by `xvp`.

2.4.2. Coarse mesh arrays

On coarse mesh, `rho_c`, `vfield` and force kernel `kern_c` should always be kept. They are usually configured to be 4 times coarser than fine grids and particle number density. In this case, each coarse grid, 3D `vfield` takes only 12 bytes, and 64 particles will at least take 384

TABLE 1
MEMORY LAYOUT FOR A CERTAIN CONFIGURATION

Type	Array	Memory usage		
		/GB	/bpp	Percentage
Particles	xvp	29.9	8.24	87.4%
	xvp_new	1.69	0.465	4.93% ^c
	PID	0	0	0%
	Subtotal	31.5	8.71	92.4%
Coarse mesh	rho_c	0.296	0.0818	0.870%
	vfield	0.889	0.245	2.60%
	kern_c	0.340	0.0939	0.996%
	force_c	(0.690)	(0.190)	(2.02%)
	rho_c_new	0.0140	0.00388	0.0411%
	Pencil-FFT	(0.454)	(0.125)	(1.33%)
	Subtotal	1.55	0.429	4.55%
Fine mesh	kern_f	1.06	0.292	3.09%
	force_f	(1.63)	(0.450)	(4.77%)
	Fine-FFT	(1.41)	(0.389)	(4.13%)
	Subtotal	1.06	0.292	3.09%
Total		34.2	9.43	100%
Optimal limit		21.7	6	63.6%

bytes ($n_\chi = n_\nu = 1$ case). Similarly, pencil-FFT arrays, coarse force arrays etc. are memory-light and temporary. The total coarse mesh arrays take less than 10% of the memory of xvp.

2.4.3. Fine mesh arrays

On the local fine mesh, only a force kernel array needs to be kept. The fine mesh density arrays, force field arrays are temporary. Since xvp_new, coarse force arrays, pencil-FFT arrays are also temporary and are not used in any calculation simultaneously, they can be overlapped in memory by using equivalent statements. The total fine mesh arrays take less than 5% of the memory of xvp.

2.4.4. Compare with traditional algorithms

To show the improvement of memory usage, we give an example of TianNu simulation (Yu et al. 2017) on the Tianhe-2 supercomputer, which used a traditional CUBEP3M code. TianNu’s particles number is limited by memory per computing node – per computing node, an average of 576^3 neutrino particles and 288^3 CDM particles are used, and consumes about 40 GB of memory⁵, or about 186 bpp. A memory efficient case of CUBEP3M by using large physical scales and at costs of speed, still uses about 40 bpp.

By using CUBE, with parameters $n_\chi = n_\nu = 1$, nc = 384, nnt = 4, ncb = 6, $\epsilon_{\text{image}} = 10\%$, $\epsilon_{\text{tile}} = 100\%$ and $\langle \text{nplotal} \rangle = 1536^3$, we use about 38.3 GB of memory, corresponding to 10.6 bpp. This can be done on most of the supercomputers, even laptops.

3. RESULTS

We use the same initial conditions, and use CUBEP3M and CUBE to simulate the large scale structure formation separately and compare their results.

Results of `izipx=1`

Results of `izipx=2`

Cosmological simulation, PM method.

PID.

Summarize memory.

⁵ Additional memory is used for OpenMP parallelization and for particle IDs to differentiate different particle species.

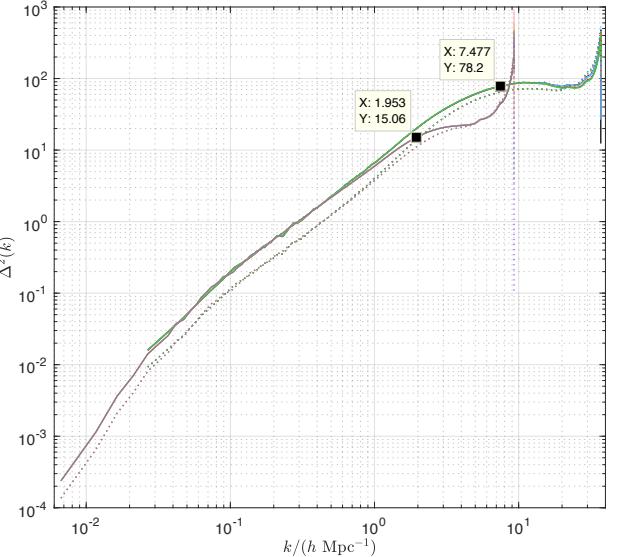


FIG. 8.— power spectrum.

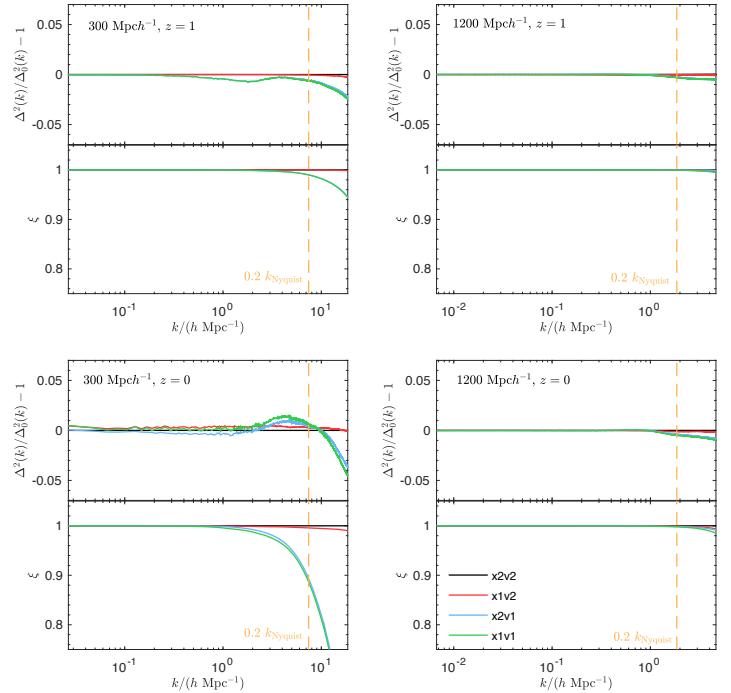


FIG. 9.— power spectrum and cross correlation.

Phi, GPU.

APPENDIX
A. INITIAL CONDITION GENERATOR

REFERENCES

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