

# NJUPT XCPC Templates

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November 26, 2020

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# 1 数据结构

## 1.1 树状数组

```
1 // 注意树状数组不能处理下标0开始
2 // 一维
3 int c[N];
4
5 inline int lowbit(int x) {
6     return x & -x;
7 }
8
9 int add(int x, int y) {
10     for (int i = x; i <= n; i += lowbit(i)) c[i] += y;
11 }
12
13 int sum(int x) {
14     int res = 0;
15     for (int i = x; i; i -= lowbit(i)) res += c[i];
16     return res;
17 }
18
19 // 二维
20 LL c[N][N];
21
22 inline int lowbit(int x) {
23     return x & -x;
24 }
25
26 void add(int x, int y, LL v) {
27     for (int i = x; i <= n; i += lowbit(i))
28         for (int j = y; j <= m; j += lowbit(j))
29             c[i][j] += v;
30 }
31
32 LL query(int x, int y) {
33     LL ans = 0;
34     for (int i = x; i; i -= lowbit(i))
35         for (int j = y; j; j -= lowbit(j))
36             ans += c[i][j];
37     return ans;
38 }
```

## 1.2 线段树

```
1 // 例子：维护区间加法区间求和
2 #include <bits/stdc++.h>
3
4 using namespace std;
5
6 typedef long long LL;
7
```

```
8  const int N = 100010;
9
10 int n, m;
11 LL a[N];
12 struct Tree {
13     int l, r;
14     LL sum, add;
15 } tr[N << 2];
16
17 void pushup(int u) {
18     tr[u].sum = tr[u << 1].sum + tr[u << 1 | 1].sum;
19 }
20
21 void build(int u, int l, int r) {
22     if (l == r) {
23         tr[u] = {l, r, a[r], 0};
24     } else {
25         tr[u] = {l, r};
26         int mid = tr[u].l + tr[u].r >> 1;
27         build(u << 1, l, mid), build(u << 1 | 1, mid + 1, r);
28         pushup(u);
29     }
30 }
31
32 void pushdown(int u) { // 先下方标记, 再清空标记
33     Tree &left = tr[u], &right = tr[u << 1 | 1];
34     left.add += root.add, left.sum += (LL)(left.r - left.l + 1) * root.add;
35     right.add += root.add, right.sum += (LL)(right.r - right.l + 1) * root.add;
36     ;
37     root.add = 0;
38 }
39
40 void modify(int u, int l, int r, LL y) {
41     if (tr[u].l >= l && tr[u].r <= r) {
42         tr[u].sum += (tr[u].r - tr[u].l + 1) * y;
43         tr[u].add += y;
44     } else {
45         pushdown(u);
46         int mid = tr[u].l + tr[u].r >> 1;
47         if (l <= mid) modify(u << 1, l, r, y);
48         if (r > mid) modify(u << 1 | 1, l, r, y);
49         pushup(u);
50     }
51 }
52
53 LL query(int u, int l, int r) {
54     if (tr[u].l >= l && tr[u].r <= r) {
55         return tr[u].sum;
56     } else {
57         pushdown(u);
58         int mid = tr[u].l + tr[u].r >> 1;
59         LL res = 0;
60         if (l <= mid) res += query(u << 1, l, r);
```

```
60     if (r > mid) res += query(u << 1 | 1, l, r);
61     return res;
62 }
63 }
64
65 int main() {
66     scanf("%d%d", &n, &m);
67     for (int i = 1; i <= n; i++) scanf("%lld", &a[i]);
68     build(1, 1, n);
69     char op[2];
70     int l, r;
71     while (m--) {
72         scanf("%s%d%d", op, &l, &r);
73         if (*op == 'C') {
74             LL d;
75             scanf("%lld", &d);
76             modify(1, l, r, d);
77         } else {
78             printf("%lld\n", query(1, l, r));
79         }
80     }
81     return 0;
82 }
```

### 1.3 ST 表

```
1 // ST表可维护区间最值/区间gcd
2 int n, a[N];
3 int f[N][M]; // f[i][j]表示区间[i, i+2^j-1]区间的最大值
4 int Log2[N];
5
6 void ST_pre() {
7     Log2[2] = 1;
8     for (int i = 3; i < N; i++) Log2[i] = Log2[i >> 1] + 1;
9     for (int i = 1; i <= n; i++) f[i][0] = a[i];
10    for (int j = 1; j < M; j++) {
11        for (int i = 1; i + (1 << j) - 1 <= n; i++)
12            f[i][j] = max(f[i][j-1], f[i + (1 << j - 1)][j - 1]);
13    }
14 }
15
16 int query(int l, int r) {
17     int k = Log2[r - l + 1];
18     return max(f[l][k], f[r - (1 << k) + 1][k]);
19 }
```

### 1.4 Splay

## 2 图论

### 2.1 最短路

### 2.2 最小生成树

### 2.3 次小生成树

### 2.4 有向图的强连通分量

### 2.5 无向图的双连通分量

### 2.6 最近公共祖先

### 2.7 2-SAT

### 2.8 网络流

#### 2.8.1 Dinic

test

```
1  #include <bits/stdc++.h>
2
3  using namespace std;
4
5  const int N = 100010, M = 200010, INF = 1e9;
6
7  int n, m, S, T;
8  int h[N], e[M], w[M], ne[M], idx;
9  int q[N];
10 int d[N], cur[N]; // d[i]表示点i的层次, cur[i]表示i的当前弧
11
12 void add(int a, int b, int c) {
13     e[idx] = b, w[idx] = c, ne[idx] = h[a], h[a] = idx++;
14     e[idx] = a, w[idx] = 0, ne[idx] = h[b], h[b] = idx++;
15 }
16
17 bool bfs() { // 判断残留网络是否存在增广路 (即从S到T存在边全大于0的路径)
18     int hh = 0, tt = -1;
19     memset(d, -1, sizeof d);
20     q[++tt] = S, d[S] = 0, cur[S] = h[S];
21     while (hh <= tt) {
22         int t = q[hh++];
23         for (int i = h[t]; ~i; i = ne[i]) {
24             int j = e[i];
25             if (d[j] == -1 && w[i]) {
26                 d[j] = d[t] + 1;
27                 cur[j] = h[j];
28                 q[++tt] = j;
29                 if (j == T) return true; // 找到了增广路, 此时增广路的流量即f[T]
30             }
31         }
32     }
33     return false; // 残留网络不存在增广路, 那么此时原图的可行流流量就是最大流
```

```

34 }
35
36 // 从起点到u, 流量最大值为limit
37 int find(int u, int limit) {
38     if (u == T) return limit;
39     int flow = 0; // u->T的流量
40     for (int i = cur[u]; ~i && flow < limit; i = ne[i]) {
41         int j = e[i];
42         cur[u] = i; // 更新当前弧
43         if (d[j] == d[u] + 1 && w[i]) {
44             int t = find(j, min(w[i], limit - flow));
45             if (!t) d[j] = -1; // 删点
46             w[i] -= t, w[i ^ 1] += t, flow += t;
47         }
48     }
49     return flow;
50 }
51
52 int dinic() {
53     int max_flow = 0;
54     while (bfs()) while (int flow = find(S, INF)) max_flow += flow;
55     return max_flow;
56 }
57
58 int main() {
59     cin >> n >> m >> S >> T;
60     memset(h, -1, sizeof h);
61     while (m--) {
62         int a, b, c;
63         scanf("%d%d%d", &a, &b, &c);
64         add(a, b, c); // 初始流量为0, 对于原图的残留网络, 正向边容量为c-0=c, 反向边容量
                       // 为0+0=0
65     }
66     printf("%d\n", dinic());
67     return 0;
68 }

```

### 2.8.2 EK

```

1  #include <bits/stdc++.h>
2
3  using namespace std;
4
5  const int N = 1010, M = 20010, INF = 1e9;
6
7  int n, m, S, T;
8  int h[N], e[M], w[M], ne[M], idx;
9  bool st[N];
10 int q[N], f[N]; // f[i]表示以i结尾的增广路径的流量 (即路径上容量的最小值)
11 int pre[N]; // pre[i]表示i的前驱边的编号 (即指向i的边)
12
13 void add(int a, int b, int c) {

```

```
14     e[idx] = b, w[idx] = c, ne[idx] = h[a], h[a] = idx++;
15     e[idx] = a, w[idx] = 0, ne[idx] = h[b], h[b] = idx++;
16 }
17
18 bool bfs() { // 判断残留网络是否存在增广路 (即从S到T存在边全大于0的路径)
19     int hh = 0, tt = -1;
20     memset(st, false, sizeof st);
21     q[++tt] = S, st[S] = true, f[S] = INF;
22     while (hh <= tt) {
23         int t = q[hh++];
24         for (int i = h[t]; ~i; i = ne[i]) {
25             int j = e[i];
26             if (!st[j] && w[i]) {
27                 q[++tt] = j;
28                 st[j] = true;
29                 pre[j] = i; // 记录j的前驱边
30                 f[j] = min(f[t], w[i]);
31                 if (j == T) return true; // j是终点, 即找到了增广路, 此时增广路的流量即f
                                     [T]
32             }
33         }
34     }
35     return false; // 残留网络不存在增广路, 那么此时原图的可行流流量就是最大流
36 }
37
38 int EK() {
39     int flow = 0;
40     while (bfs()) { // 如果残留网络存在增广路f', 就将他的流量加到原网络的流量上
41         flow += f[T];
42         for (int i = T; i != S; i = e[pre[i] ^ 1]) { // 更新残留网络
43             w[pre[i]] -= f[T], w[pre[i] ^ 1] += f[T];
44         }
45     }
46     return flow; // 最大流
47 }
48
49 int main() {
50     cin >> n >> m >> S >> T;
51     memset(h, -1, sizeof h);
52     while (m--) {
53         int a, b, c;
54         scanf("%d%d%d", &a, &b, &c);
55         add(a, b, c); // 初始流量为0, 对于原图的残留网络, 正向边容量为c-0=c, 反向边容量
                        为0+0=0
56     }
57     printf("%d\n", EK());
58     return 0;
59 }
```

## 2.9 二分图

## 3 其他

### 3.1 高精度

```
1 struct hll {
2     int num[4010], len, sign;
3     hll() { len = 0, sign = 1; }
4     hll(int x) { *this = x; }
5     hll(long long x) { *this = x; }
6     hll(char *ss) { *this = ss; }
7     hll(string ss) { *this = ss; }
8     hll& operator = (const int &x) {
9         int val = x;
10        if (val >= 0) sign = 1, len = 0;
11        else if (val < 0) sign = -1, val = -val, len = 0;
12        do {
13            num[len++] = val % 10, val /= 10;
14        } while (val);
15        return *this;
16    }
17    hll& operator = (const long long &x) {
18        long long val = x;
19        if (val >= 0) sign = 1, len = 0;
20        else if (val < 0) sign = -1, val = -val, len = 0;
21        do {
22            num[len++] = val % 10, val /= 10;
23        } while (val);
24        return *this;
25    }
26    hll& operator = (const string &ss) {
27        len = ss.size();
28        int start;
29        if (ss[0] == '-') sign = -1, start = 1;
30        else sign = 1, start = 0;
31        for (int i = len - 1; i >= start; i--) num[len - i - 1] = ss[i] - '0';
32        if (sign == -1) len--;
33        return *this;
34    }
35    hll& operator = (const char *ss) {
36        len = strlen(ss);
37        int start;
38        if (ss[0] == '-') sign = -1, start = 1;
39        else sign = 1, start = 0;
40        for (int i = len - 1; i >= start; i--) num[len - i - 1] = ss[i] - '0';
41        if (sign == -1) len--;
42        return *this;
43    }
44    hll& operator = (const hll &t) {
45        len = t.len, sign = t.sign;
46        for (int i = 0; i < len; i++) num[i] = t.num[i];
47        return *this;
```



```

48     }
49     int abs_cmp(const hll &a, const hll &b) const { // |a|>|b|时返回1, 相等返回0
        , 小于返回-1
50         if (a.len > b.len) return 1;
51         else if (a.len < b.len) return -1;
52         else {
53             for (int i = a.len - 1; i >= 0; i--) {
54                 if (a.num[i] < b.num[i]) return -1;
55                 if (a.num[i] > b.num[i]) return 1;
56             }
57             return 0;
58         }
59     }
60     int cmp(const hll &t) const { // *this与t比较, 小于返回-1, 等于返回0, 大于返回1
61         if (sign != t.sign) {
62             if (sign == 1) return 1;
63             else return -1;
64         } else {
65             if (abs_cmp(*this, t) == 1) return sign;
66             else if (abs_cmp(*this, t) == 0) return 0;
67             else return -sign;
68         }
69     }
70     hll abs_plus(const hll &a, const hll &b) { // |a|+|b|, ans的符号与a和b原来的
        符号相同
71         hll ans;
72         ans.sign = a.sign;
73         for (int i = 0, carry = 0; i < a.len || i < b.len || carry; i++) {
74             if (i < a.len) carry += a.num[i];
75             if (i < b.len) carry += b.num[i];
76             ans.num[ans.len++] = carry % 10;
77             carry /= 10;
78         }
79         return ans;
80     }
81     hll abs_minus(const hll &a, const hll &b) { // ||a|-|b||, ans的符号为|a|-|b|
        的符号
82         hll ans, c, d;
83         if (abs_cmp(a, b) >= 0) ans.sign = 1, c = a, d = b;
84         else ans.sign = -1, c = b, d = a;
85         for (int i = 0, borrow = 0; i < c.len; i++) {
86             borrow = c.num[i] - borrow;
87             if (i < d.len) borrow -= d.num[i];
88             ans.num[ans.len++] = (borrow + 10) % 10;
89             if (borrow >= 0) borrow = 0;
90             else borrow = 1;
91         }
92         while (ans.len > 1 && ans.num[ans.len - 1] == 0) ans.len--; // 去除前导0
93         return ans;
94     }
95     bool operator == (const hll &t) const { return cmp(t) == 0; }
96     bool operator != (const hll &t) const { return !(cmp(t) == 0); }
97     bool operator < (const hll &t) const { return cmp(t) == -1; }

```

```

98     bool operator > (const hll &t) const { return cmp(t) == 1; }
99     bool operator <= (const hll &t) const { return !(cmp(t) == 1); }
100    bool operator >= (const hll &t) const { return !(cmp(t) == -1); }
101    hll operator + (const hll &t) {
102        hll ans;
103        if (sign == t.sign) { // 同号 直接相加, 符号不变
104            ans = abs_plus(*this, t);
105        } else { // 异号
106            if (sign == 1) { // 前正 + 后负 == 前绝对值 - 后绝对值
107                ans = abs_minus(*this, t);
108            } else { // 前负 + 后正 == 后绝对值 - 前绝对值
109                ans = abs_minus(t, *this);
110            }
111        }
112        return ans;
113    }
114    hll operator - (const hll &t) {
115        hll ans;
116        if (sign == t.sign) { // 同号
117            ans = abs_minus(*this, t);
118            if (sign == 1) { // 前正 - 后正
119                ; // 不用做了
120            } else { // 前负 - 后负
121                ans.sign *= -1;
122            }
123        } else { // 异号
124            if (sign == 1) { // 前正 - 后负 == 前绝对值 + 后绝对值
125                ans = abs_plus(*this, t);
126            } else { // 前负 - 后正 == -(前绝对值 + 前绝对值)
127                ans = abs_plus(t, *this);
128                ans.sign = -1;
129            }
130        }
131        return ans;
132    }
133    hll operator * (const hll &t) { // 高精度*高精度
134        hll ans;
135        memset(ans.num, 0, len + t.len << 2);
136        ans.sign = sign * t.sign;
137        ans.len = len + t.len - 1; // a位数乘以b位数, 得到的结果是a+b-1位数, 或a+b位数
138        for (int i = 0; i < len; i++) {
139            for (int j = 0; j < t.len; j++)
140                ans.num[i + j] += num[i] * t.num[j];
141        }
142        for (int i = 0; i < ans.len - 1; i++) {
143            if (ans.num[i] >= 10) {
144                ans.num[i + 1] += ans.num[i] / 10;
145                ans.num[i] %= 10;
146            }
147        }
148        // 看最高位是否需要进位, 如果有进位, 答案最终是a+b位数, 否则是a+b-1位数
149        if (ans.num[ans.len - 1] >= 10) {

```

```
150         ans.num[ans.len] = ans.num[ans.len - 1] / 10;
151         ans.num[ans.len - 1] %= 10;
152         ans.len++;
153     }
154     while (ans.len > 1 && ans.num[ans.len - 1] == 0) ans.len--; // 去除前导0
155     return ans;
156 }
157 hll operator += (const hll &t) {
158     return *this + t;
159 }
160 hll operator *= (const hll &t) {
161     return *this * t;
162 }
163 void print() {
164     if (sign == -1) putchar('-');
165     for (int i = len - 1; i >= 0; i--) putchar(num[i] + '0');
166 }
167 };
```

### 3.2 莫队