

Monthly Meeting on July

Yuichiro Honda

Morita lab. M1

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1 Introduction

2 Progress

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matroid

definition

Let $M = (S, \mathcal{I})$ where S is ground set, and \mathcal{I} is a family satisfying $\mathcal{I} \subseteq 2^S$. M is called a **matroid** when \mathcal{I} satisfies:

$$\emptyset \in \mathcal{I} \tag{1}$$

$$I_1 \subset I_2, I_2 \in \mathcal{I} \Rightarrow I_1 \in \mathcal{I} \tag{2}$$

$$I_1, I_2 \in \mathcal{I}, |I_1| < |I_2| \Rightarrow \exists i_2 \in I_2 ; I_1 \cup \{i_2\} \in \mathcal{I} \tag{3}$$

\mathcal{I} is called an **independent set family**.

A maximal element in \mathcal{I} with respect to order “ \subseteq ” is called a **base**.

Every base is the same size. This size is called **rank**.

matroid examples

vector matroid

Linearly independent sets construct matroid:

$$\mathbf{a}_1 = \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix}, \mathbf{a}_2 = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}, \mathbf{a}_3 = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}, \mathbf{a}_4 = \begin{pmatrix} 1 \\ 0 \\ -1 \end{pmatrix}$$

ground set S is $\{\mathbf{a}_1, \mathbf{a}_2, \mathbf{a}_3, \mathbf{a}_4\}$, and \mathcal{I} consists of linearly independent combinations, that:

$$\begin{aligned} \mathcal{I} = & \{\emptyset, \{\mathbf{a}_1\}, \{\mathbf{a}_2\}, \{\mathbf{a}_3\}, \{\mathbf{a}_4\}, \\ & \{\mathbf{a}_1, \mathbf{a}_2\}, \{\mathbf{a}_1, \mathbf{a}_3\}, \{\mathbf{a}_1, \mathbf{a}_4\}, \{\mathbf{a}_2, \mathbf{a}_3\}, \\ & \{\mathbf{a}_2, \mathbf{a}_4\}, \{\mathbf{a}_3, \mathbf{a}_4\}, \{\mathbf{a}_1, \mathbf{a}_2, \mathbf{a}_3\}, \{\mathbf{a}_1, \mathbf{a}_3, \mathbf{a}_4\}, \{\mathbf{a}_2, \mathbf{a}_3, \mathbf{a}_4\}\} \end{aligned}$$

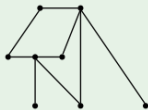
Then \mathcal{I} satisfies (1), (2), (3).

matroid examples

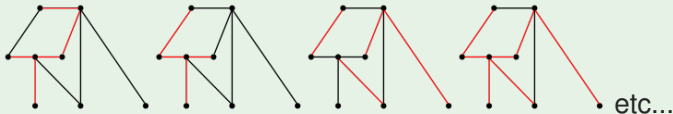
graphic matroid

In graph theory, forests construct matroid:

Let $G = (V, E)$ be a graph where V is a set of vertices and E is a set of edges:



When \mathcal{I} is a family of forests, such as:



Then \mathcal{I} satisfies (1), (2), (3).

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ON DISJOINT COMMON BASES IN TWO MATROIDS

problem 1

Let $M_1 = (S, \mathcal{I}_1)$ and $M_2 = (S, \mathcal{I}_2)$ be matroids on the ground set S , where \mathcal{I}_1 and \mathcal{I}_2 are the respective families of independent sets. A set $B \subseteq S$ that is both a base of M_1 and M_2 is called a common base. The problem is to decide whether S can be partitioned into common bases.

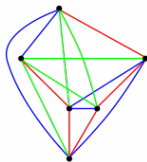


Figure: partitioned into bases

ON DISJOINT COMMON BASES IN TWO MATROIDS

conjecture 2 (Rota's conjecture)

Let $M = (T, \mathcal{I})$ be a matroid of rank n . Let A_1, \dots, A_n be a partition of T into bases of M . Then there are disjoint bases B_1, \dots, B_n such that $|A_i \cap B_j| = 1$ for every $i = 1, \dots, n$ and $j = 1, \dots, n$.

Next conjecture is the generalization of Rota's conjecture:

conjecture 3 (Chow's conjecture)

Let $M = (T, \mathcal{I})$ be a matroid of rank n with the property that T can be partitioned into b bases where $3 \leq b \leq n$. Let $I_1, \dots, I_n \in \mathcal{I}$ be disjoint independent sets, each size at most b . Then there exists a partition of T into sets A_1, \dots, A_n such that $I_i \subseteq A_i$ and $|A_i| = b$ for every $i = 1, \dots, n$, and there exist disjoint bases B_1, \dots, B_b such that $|A_i \cap B_j| = 1$ for every $i = 1, \dots, n$ and $j = 1, \dots, b$.

ON DISJOINT COMMON BASES IN TWO MATROIDS

theorem 4

Problem 1 can be solved in polynomial time if and only if this is under the additional assumption that one of the matroids is a direct sum of uniform matroids.

definition

Let $M_1 = (E_1, \mathcal{I}_1)$, $M_2 = (E_2, \mathcal{I}_2)$ be matroids where $E_1, E_2 \neq \emptyset, E_1 \cap E_2 = \emptyset$.

M is called a **direct sum** of 2 matroids, M_1 and M_2 when:

$$M = (E_1 \cup E_2, \mathcal{I}_1 \oplus \mathcal{I}_2)$$

$$\text{where } \mathcal{I}_1 \oplus \mathcal{I}_2 = \{X_1 \cup X_2 \mid X_1 \in \mathcal{I}_1, X_2 \in \mathcal{I}_2\}$$

uniform matroid

Let $U = (S, \mathcal{I})$ be a matroid where $\mathcal{I} \in 2^S$ satisfies:

$$\mathcal{I} = \{I \mid |I| \leq k\}$$

Then U is a matroid of rank k . This U is called a **uniform matroid**.

claim 5

Conjecture 3 is false for every b such that $2 \leq b \leq \frac{n}{3}$.

corollary 6

Problem 1 can be solved in polynomial time if and only if this is true under the additional assumption that M_2 is a direct sum of uniform matroids whose blocks are each independent in M_1 .

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next month

Yet these methods show that Chow's conjecture is false when $3 \leq b \leq \frac{n}{3}$, Rota's conjecture is still alive since it's Chow's conjecture in the case of $b = n$. There's still room for research about Chow's conjecture.

Also, computational complexity of problem 1 is still open. This paper gives necessary and sufficient condition to solve problem 1 in polynomial time only when problem 1 can be solved. It's yet to be discovered exactly when problem 1 can be solved.

What I'm going to do next month is to think about these problems.

Thank you for your attention.