Monthly Meeting on October

Yuichiro Honda Morita lab. M1

2016/10/05

- 1 Previous work
- Progress
- 3 Next step

Last month

- searched an (polynomial) algorithm to partition two matroids into their common bases
- 2 found an algorithm to find all common bases in two matroids in $O(n(n^2 + t)\lambda)$

- Previous work
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References

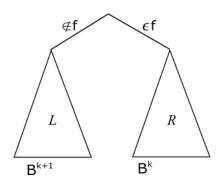


Komei Fukuda, Makoto Namiki: "Finding all common bases in two matroids", Discrete Applied Mathmatics 56 (1995) 231-243

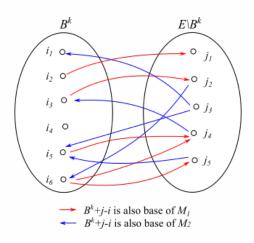
Main result

Given two matroids $M_1=(E,\mathcal{B}_1)$, $M_2=(E,\mathcal{B}_2)$, and a common base $B^1\in\mathcal{B}_1\cap\mathcal{B}_2$, there is an algorithm finding all common bases of them in $O(n(n^2+t)\lambda)$ where λ is number of the bases and t is the time to make one pivot operation.

- The algorithm use enumeration tree whose height is |E| = n
- In the beginning



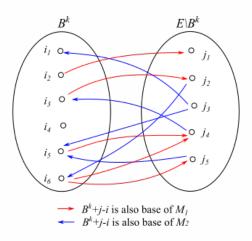
bipartite graph G



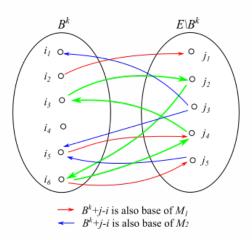
Theorem 1

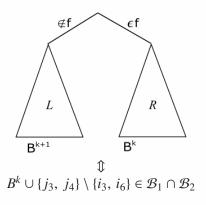
There exists a common base B^{k+1} in L if and only if there exists a directed cycle C in G which contains f and consists of elements in E less than or equal to f.

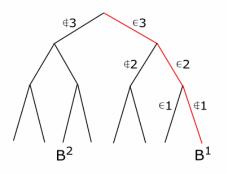
bipartite graph G when $f = i_6...$



bipartite graph G when $f = i_6...$

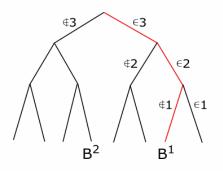






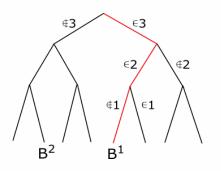
$$P = (R R R)$$

 $P_1 = R \text{ then } f = 1$
 $B^1 = \{2, 3\}$



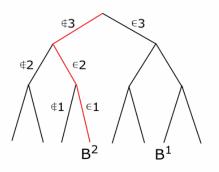
$$P = (LRR)$$

 $P_2 = R \text{ then } f = 2$
 $B^1 = \{2, 3\}$



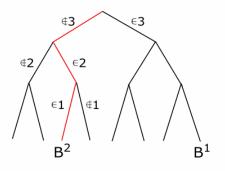
$$P = (LLR)$$

 $P_3 = R \text{ then } f = 3$
 $B^1 = \{2, 3\}$



$$P = (R R L)$$

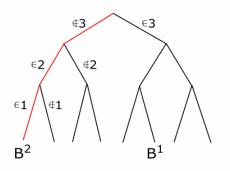
 $P_1 = R \text{ then } f = 1$
 $B^2 = \{1, 2\}$



$$P = (LRL)$$

$$P_2 = R \text{ then } f = 2$$

$$B^2 = \{1, 2\}$$



$$P = (L L L)$$

 f doesn't exist end

alogrithm "Enumerate"

input: $M_1 := (E, \mathcal{B}_1), M_2 := (E, \mathcal{B}_2), B^1 \in \mathcal{B}_1 \cap \mathcal{B}_2$ output: all common bases of M_1 and M_2

- 1 k := 1, output B^k . P := (R, ..., R). (Here, $P \in \{L, R\}^n$)
- **2** $f := min\{j \mid P_j = R\}$
- 3 If *f* doesn't exist, stop. Else, create *G* and find directed cycle *C* satisfying theorem 1 with breadth first search.
- 4 If C doesn't exist, goto step 2. Else, k := k+1, make pivot operation along C and obtain new B^k , $T_{M_1}(B^k)$, and $T_{M_2}(B^k)$. $P_f := L$, $\forall j < k; \ P_j := R$, goto step 2.

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next month

TODO:

- 1 learn about matroid intersection
- 2 learn how to partition a matoid into bases
- tackle two different partition matroids (can be partitioned into their common bases)
- generalized partition matroid of two different uniform matroids and any matroid