

# Lab08 & hw08

NJU SICP22 TAs

# Scheme

- Writing executable math equation,
- ... in linked list.

# Lab08p1: over or under

```
(define (over-or-under a b)
  -1 if a is less than b
  0 if a is equal to b
  1 if a is greater than b
)
```

# Lab08p1: over or under

```
(define (over-or-under a b)
  -1 if (< a b)
  0 if (= a b)
  1 else
)
```

# Lab08p1: over or under

```
(define (over-or-under a b)
  (cond ((< a b) -1)
        ((= a b) 0)
        (else 1))
)
```

# Lab08p2: make adder

- “closure”
- You already know this well in python.

# Lab08p3: composed

- “High-order function”
- You already know this well in python.

# Lab08p4: gcd

```
def gcd(a, b):  
    if b == 0:  
        return a  
    return gcd(b, a % b)
```



# Lab08p4: gcd

```
def gcd(a, b):  
    if (= b 0):  
        a  
    (gcd (remainder a b))
```

expression

# Lab08p4: gcd

```
def gcd(a, b):  
    (if (= b 0)  
         a  
         (gcd (remainder a b)))
```

control

# Lab08p5: make a list

```
(define lst
```

```
'YOUR-CODE-HERE ; not a comment!
```

```
)
```

# Lab08p5: ordered

Math: ordered(L)

If L is nil: return True

If L is a :: nil: return True

If L is a :: b :: L':

    return a < b /\ ordered (b :: L')

# Lab08p5: ordered

Math2Code: ordered(L)

If (null? L): #t

If (null? (cdr L)): #t

Else:

return ((car L) < (car (cdr L)) /\ (ordered (cdr L)))

; (car L) is a, (car (cdr L)) is b

# Lab08p5: ordered

Code: ordered(L)

```
(cond
  ((null? L) #t)
  ((null? (cdr L)) #t)
  (else
   ((car L) < (car (cdr L)) /\ (ordered (cdr L)))
  )
; (car L) is a, (car (cdr L)) is b
```

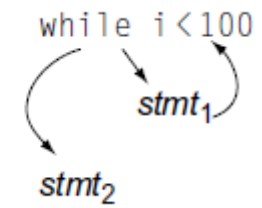
# Scheme

- Writing executable **math equation**,
- ... in linked list.
- Functional & Imperative
- Dataflow & Control flow

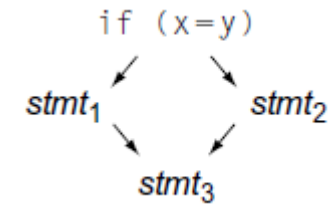
...

↓  
A = B + C  
↓  
D = C + E  
↓  
return D

```
while(i<100)
  begin
    stmt1
  end
  stmt2
```



```
if (x=y)
  then stmt1
  else stmt2
  stmt3
```



Control flow: how execution is performed?

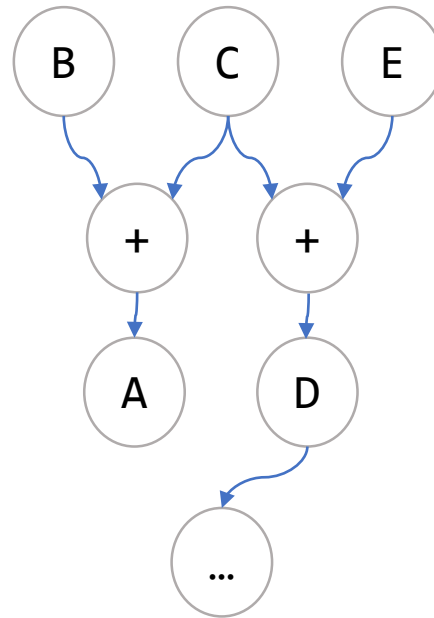


...

$A = B + C$

$D = C + E$

return D



How data is mathematically related?  
“Data dependences”

# Scheme

- Writing executable **math equation**,
  - ... in linked list.
  - Functional & Imperative
  - Dataflow & Control flow
- 
- Functional programming is in somewhat “higher-level” than the imperative:
    - Let the compiler & machine decide the tedious control flow,
    - Human only needs to know the math.
      - And think the math.

# Scheme

- Writing executable **math equation**,
- ... in linked list.

**You can visit last year's review  
for tygg's perspective.**

- Functional programming is in somewhat “higher-level” than the imperative:
  - Let the compiler & machine decide the tedious control flow,
  - Human only needs to know the math.
    - And think the math.

# Hw08p1: pow

Math: pow base exp =

If exp = 0: return 1

If exp % 2 = 0: return pow base (exp//2)

If exp % 2 = 1: return base \* (pow base (exp//2))

$$x^{2y} = (x^y)^2$$

$$x^{2y+1} = x(x^y)^2$$

Hw08p1: pow

**You definitely know  
how to do this!**

Math: pow base exp =

If exp = 0: return 1

If exp % 2 = 0: return pow base (exp//2)

If exp % 2 = 1: return base \* (pow base (exp//2))

$$x^{2y} = (x^y)^2$$

$$x^{2y+1} = x(x^y)^2$$

## Hw08p2: filter-lst

Math: filter-list fn L

If L = nil: return nil

If L = a::L':

    If (fn a): return a :: (filter-lst fn L')

    else: return (filter-lst fn L')

# Hw08p3: no repeats

Math: no-repeats L

If L = nil: return nil

If L = a::L':

    return a::(no-repeats (filter-lst (not eq a) L'))

# Hw08p4: substitute

Math: substitute s old new

If s = nil: return s

If s = a::L':   ;; a is not list

    If a = old: return new :: substitute L' old new

    Else: return old :: substitute L' old new

If s = L::L':   ;; L is list

    return (substitute L old new) :: (substitute L' old new)



# Hw08p5: Sub All

Math: sub-all s olds news

If olds = nil: return s

If olds = old::olds', news = new::news':

    sub-all (substitute s old new) olds' news'

# Hw08p6: Tree in Scheme

- How to represent a tree with linked list?
- Node: (label, branches)

# Hw08p7: Label Sum

Math: label-sum t

(label t) + (sum (map label-sum (branches t)))

# Hw09p1: symbol differentiation, +

Math: 
$$\frac{\partial(f(x) + g(x))}{\partial x} = \frac{\partial f(x)}{\partial x} + \frac{\partial g(x)}{\partial x}$$

```
(define (derive-sum expr var)
  (make-sum (derive (first-operand expr) var)
             (derive (second-operand expr) var)))
```

# Hw09p1: symbol differentiation, \*

Math: 
$$\frac{\partial(f(x) \times g(x))}{\partial x} = \frac{\partial f(x)}{\partial x} \times g(x) + f(x) \times \frac{\partial g(x)}{\partial x}$$

```
(define (derive-product expr var)
  (make-sum
    (make-product (derive (first-operand expr) var)
                  (second-operand expr))
    (make-product (derive (second-operand expr) var)
                  (first-operand expr))))
```

# Hw09p1: symbol differentiation, exp

Math:

$$\frac{\partial (f(x)^{g(x)})}{\partial x} = g(x) \times f(x)^{g(x)-1}$$

```
(define (derive-exp exp var)
  (make-product (second-operand exp)
    (make-exp (first-operand exp)
      (make-sum (second-operand exp) -1))))
```

# Congratulation!

Now you know the *most* elegant language in the world ☺