CS 217 – Algorithm Design and Analysis

Shanghai Jiaotong University, Fall 2019

Handed out on Monday, 2019-10-21
First submission and questions due on Monday, 2019-10-28
You will receive feedback from the TA.
Final submission due on Monday, 2019-11-04

5 More on Network Flows

Exercise 1. Let $G = (V, c)$ be a flow network. Prove that flow is "transitive" in the following sense: if r, s, t are vertices, and there is an r–s-flow of value k and an s–t-flow of value k, then there is an r–t-flow of value k.

Proof. Denote the original r-s-flow and s-t-flow as f_{rs} and f_{st} respectively. We prove $f_{rt} = k$ by contradiction.

Suppose $f_{rt} < k$, denote $f_{rt} = p$. By Max-Flow-Min-Cut THM, there is a cut C with $cap(C) = p < k$. If $s \in C$, we know that C is also a r-s-cut. But by Max-Flow-Min-Cut THM, min $cap(r - s - cut) = k > cap(C) = p$, which leads to a contradiction. If $s \notin C$, the proof is similar. So $f_{rt} \geq k$, so there is a flow of value k in between r, t.

□

5.1 Vertex Disjoint Paths

Let G be a directed graph. Two paths p_1, p_2 from s to t are called *vertex disjoint* if they don't share any vertices except s and t.

Theorem 2 (Menger's Theorem)**.** *Let G be a graph and $s \neq t$ two vertices therein. Let $k \in \mathbf{N}_0$. Then exactly one of the following is true:*

1. *There are k vertex disjoint paths $p_1, \ldots, p_k$ from s to t; that is, no two p_i, p_j share any vertex besides s and t.*

2. *There are vertices $v_1, \ldots, v_{k-1} \in V \setminus \{s,t\}$ such that $G - \{v_1, \ldots, v_{k-1}\}$ contains no s–t-path.*

Exercise 3. Prove Menger's Theorem. You have to prove two things: first, not both cases above can occur (this is rather easy); second, one of them must occur (this requires a tool from the lecture).

Proof. First we prove the easy part, these two cases will not occur simultaneously.Prove it by contradiction.

Suppose both cases are true simultaneously, then there are k vertex disjoint paths $p_1, ..., p_k$ from s to t. and there exists $v_1, ..., v_{k-1}$ such that $G - \{v_1, ..., v_{k-1}\}$ contains no s-t-path.

For any $v_1, ..., v_{k-1}$, the can take place in at most $k-1$ paths in $p_1, ..., p_k$ since $p_1, ..., p_k$ are disjoint vertex paths. Then we know there must be at least one path from $p_1, .., p_k$ left, which connects s and t, leading to a contradiction.

Next we prove that one of them must occur.

The second case is obvious since we can just remove all vertices except s, t from the V, then obviously there're no s-t-path now.

For the first case, we construct such a network-graph with all edge in the original graph assigned capacity 1. In such a network, the value of a flow is

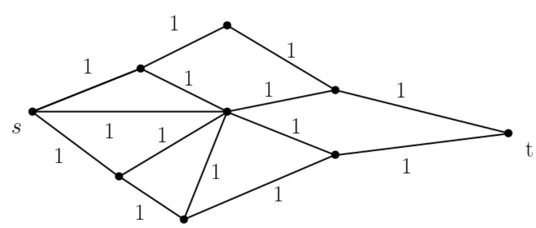


Figure 1: network

the number of vertex disjoint paths in such a flow. Since the capacity of each edge is 1, we can know for sure that no two s-t-path cross with each other, otherwise there must be a vertex with units bigger than 1. Based on that, we see k, the value of a flow is the number of paths in it. Thus leading to k disjoint paths.

In fact, we can prove that the maximum number of disjoint paths is equal to the minimum number of vertex set which separates s, t.

We construct a new network from original graph. For a vertex v other than s, t, we split it to v^+, v^-, and draw a new directed edge v^-v^+. For an edge (s, v), replace it with sv^-. For an edge (v, t), replace it with v^+t. For other edges (u, v), replace it with (u^+v^-, v^+u^-). All new edges are assigned capacity 1. Here's an example.

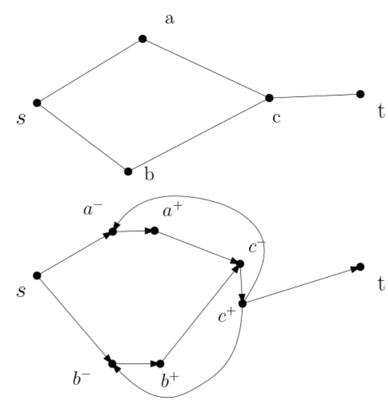


Figure 2: network

In such a network G', we split the nodes and put directed edge to make sure no two flow cross each other at some nodes. Also, any flow f in this network can produce a set P of disjoint (s, t) paths, $|P| = val(f)$. We add an edge (u, v) to a path if (u^+, v^-) is in the flow. Any cut C has a sub set C' with all arcs like (v^-, v^+), and C' still a cut in G'. Then we can produce a set of vertices V in G with all nodes $v((v^-, v^+) \in C')$.

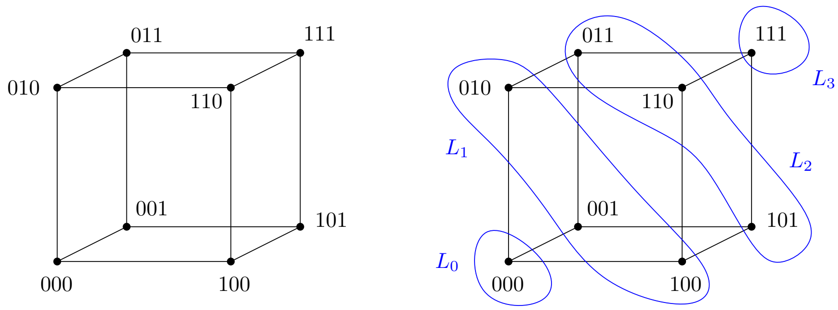
Then we know that for a maximum set $P'$of disjoint paths in G, $|P'| \geq |P| = val(f) = |C| \geq |C'| \geq |V| \geq |P'|$ So all of these values are equal, which finishes the proof.

$\square$

Let $V = \{0, 1\}^n$. The n-dimensional Hamming cube H_n is the graph (V, E) where $\{u, v\} \in E$ if u, v differ in exactly one coordinate. Define the i^{th} level of H_n as

$$L_i := \{u \in V \mid \|u\|_1 = i\} ,$$

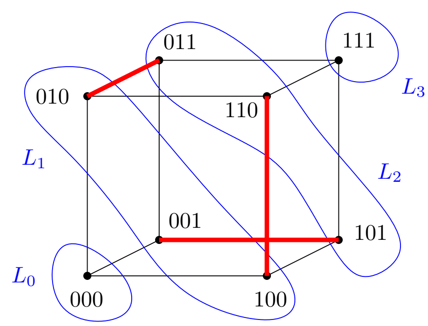
i.e., those vertices u having exactly i coordinates which are 1.



The 3-dimensional Hamming cube and the four sets L_0, L_1, L_2, L_3.

Exercise 4. [Matchings in H_n] Consider the induced bipartite subgraph $H_n[L_i \cup L_{i+1}]$. This is the graph on vertex set $L_i \cup L_{i+1}$ where two edges are connected by an edge if and only if they are connected in H_n.

Show that for $i \leq n/2$ the graph $H_n[L_i \cup L_{i+1}]$ has a matching of size $|L_i| = \binom{n}{i}$.



A matching of size 3 between L_1 and L_2.

Proof. □

Exercise 5. Let H_n be the n-dimensional Hamming cube. For $i < n/2$ consider L_i and L_{n-i}. Note that $|L_i| = \binom{n}{i} = \binom{n}{n-i} = L_{n-i}$, so the L_i and L_{n-i} have the same size. Show that there are $\binom{n}{i}$ paths $p_1, p_2, \dots, p_{\binom{n}{i}}$ in H_n such that (i) each p_i starts in L_i and ends in L_{n-i}; (ii) two different paths p_i, p_j do not share any vertices. **Hint 1.** Model this problem as a network flow with vertex capacities. What would the maximum flow be in this network? **Hint 2.** It's not *that* easy. If you try to work from both sides towards the middle by combining matchings between levels, you will certainly run into problems as how to glue things together in the middle. I have never seen any “meet in the middle” proof that works. **Hint 3.** There is a “direct” proof by induction that does not require anything about network flows.

Proof. □

5.2 Matchings and Vertex Covers

The following exercise was on the final exam of CS 499 (mathematical foundations of computer science) in spring 2019.

Exercise 6. Let $\nu(G)$ denote the size of a maximum matching of G. Show that a bipartite graph G has at most $2^{\nu(G)}$ minimum vertex covers.

Proof. Suppose the number of minimum vertex covers of graph G is $m\left(G\right)$. We can easily Now we prove that there could be no more than $2^{\nu(G)}$ minimum

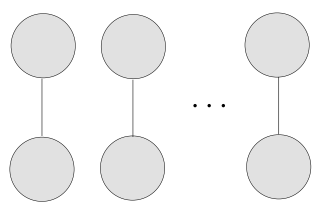


Figure 3: example

vertex cover. Let $G = A \cup B$, since G is a bipartite graph. Assume that the maximum matching is $e_1, e_2, \ldots, e_t$, in which $t = \nu\left(G\right)$. By **Konig's Theorem** we know that the number of vertices in a minimum vertex cover is exactly t.

Plus we can prove that every vertex in K, the vertex cover, is an endpoint of a matched edge. Hence the result is obvious. □

Obviously, this is not true for general (non-bipartite) graphs: the triangle K_3 has $\nu(K_3) = 1$ but it has three minimum vertex covers. The five-cycle C_5 has $\nu(C_5) = 2$ but has five minimum vertex covers.

Exercise 7. Is there a function $f : \mathbf{N}_0 \to \mathbf{N}_0$ such that every graph with $\nu(G) = k$ has at most $f(k)$ minimum vertex covers? How small a function f can you obtain?

Proof. Since this is for every graph. First consider K_{2r+1}, we have $\frac{k}{r}$ K_{2r+1}. Hence $\nu(G) = \frac{k}{r} \cdot r = k$

$$f(k) \geq (2r+1)^{\frac{k}{r}}.$$

For example, if we choose $r = 1$, there are k K_3, we have $f(k) \geq 3^k$. And let $r \to \infty$ we have

$$f(k) \geq e^{2k}.$$

Next we prove that the number of minimum vertex cover is at most e^{2k}. Consider $\forall$ minimum vertex cover K. We prove

$$|K| \leq e^{2k}.$$

Let $\forall u, v \in K$. If u, v are connected, then at least one of them is on a edge of maximum matching. $\square$