

Homework 5

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Exer.1

Denote the original r-s-flow and s-t-flow as f_{rs} and f_{st} respectively, we use them to construct a new r-t-flow f_{rt} , thus prove the transitivity of flow.

Let

$$f_{rt} = V \times V \rightarrow \mathbb{R} = \begin{cases} f_{rs}(u, v) & f_{rs}(u, v) \neq 0 \\ f_{st}(u, v) & f_{st}(u, v) \neq 0 \\ 0 & otherwise \end{cases}$$

Here V is the set of all vertices.

We just need to prove f_{rt} is a flow.

- Capacity constraint:

Case1: $f_{rs}(u, v) \neq 0$, then $f_{rt}(u, v) = f_{rs}(u, v) \leq c(u, v)$

Case2: $f_{st}(u, v) \neq 0$, then $f_{rt}(u, v) = f_{st}(u, v) \leq c(u, v)$

Case3: Otherwise, $f_{rt}(u, v) = 0 \leq c(u, v)$, since $c(u, v) \geq 0$

- Skew symmetry:

Case1: $f_{rs}(u, v) \neq 0$, then $f_{rt}(u, v) = f_{rs}(u, v) = -f_{rs}(v, u) = -f_{rt}(v, u)$

Case2: $f_{st}(u, v) \neq 0$, then $f_{rt}(u, v) = f_{st}(u, v) = -f_{st}(v, u) = -f_{rt}(v, u)$

Case3: Otherwise, $f_{rs}(u, v) = 0, f_{st}(u, v) = 0 \Rightarrow f_{rt}(u, v) = 0$

$f_{rs}(v, u) = -f_{rs}(u, v) = 0, f_{st}(v, u) = 0 \Rightarrow f_{rt}(v, u) = 0$

$f_{rt}(u, v) = 0 = -f_{rt}(v, u)$

- Flow conservation:

For all $u \in V/\{r, t\}$,

$$\begin{aligned}
 \sum_{v \in V} f_{rt}(u, v) &= \sum_{v \in V, f_{rs}(u, v) > 0} f_{rt}(u, v) + \sum_{v \in V, f_{rs}(u, v) > 0} f_{rt}(u, v) + \sum 0 \\
 &= \sum f_{rs}(u, v) + \sum f_{st}(u, v) + 0 \\
 &= 0 + 0 + 0 \\
 &= 0
 \end{aligned}$$

Exer.2

First we prove the easy part, these two cases will not occur simultaneously. Prove it by contradiction.

Suppose both cases are true simultaneously, then there are k vertex disjoint paths p_1, \dots, p_k from s to t . and there exists v_1, \dots, v_{k-1} such that $G - \{v_1, \dots, v_{k-1}\}$ contains no s-t-path.

For any v_1, \dots, v_{k-1} , the can take place in at most $k-1$ paths in p_1, \dots, p_k since p_1, \dots, p_k are disjoint vertex paths. Then we know there must be at least one path from p_1, \dots, p_k left, which connects s and t , leading to a contradiction.

Next we prove that one of them must occur.

The second case is obvious since we can just remove all vertices except s, t from the V , then obviously there're no s-t-path now.

For the first case, we construct such network-graph. Edges connecting with s, t are assigned with ∞ capacity so that it will not affect the inner flow. All other edges are assigned with capacity 1.

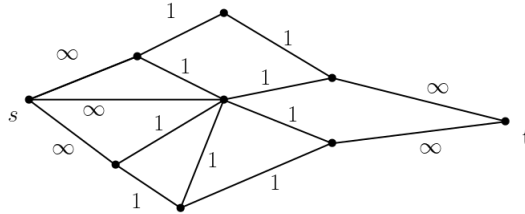


图 1: network

In such a network, the value of a flow is the number of vertex disjoint

paths in such a flow. Since the capacity of each edge is 1, we can know for sure that no two s-t-path cross with each other, otherwise there must be a vertex with units bigger than 1. Based on that, we see k , the value of a flow is the number of paths in it. Thus leading to k disjoint paths.