

Lecture 3 : Model categories

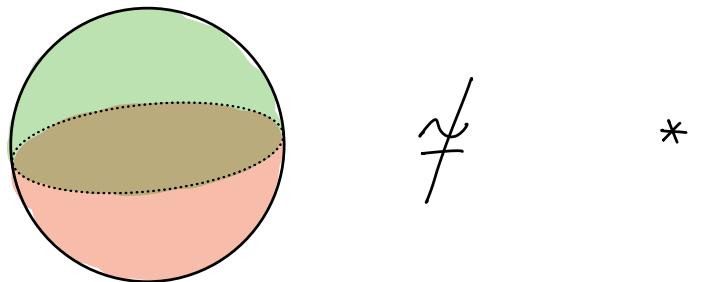
CW-complexes and Kan complexes

support equivalent homotopy theories.

$$D^2 \leftarrow S^1 \longrightarrow D^2$$

Consider : $\begin{array}{ccc} ? & \downarrow & ? \\ ? & \downarrow & ? \\ * & \longleftarrow S^1 & \longrightarrow * \end{array}$ in Top .

The vertical maps are all (weak) htpy eq., but if we take the pushout of each row, we get non-equivalent spaces :



Using the language of model categories, we can deal with these (seemingly) ill-behaved equivalences and make precise e.g. S^2 is the "right" pushout.

Def: Let M be a category with small limits & colimits.

A model structure on M consists of three classes of morphisms called:

- weak equivalences W ($\xrightarrow{\sim}$) } determine the homotopy theory.

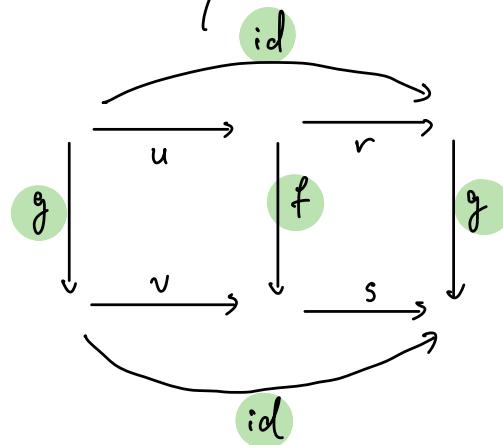
- cofibrations C (\rightarrowtail) }

- fibrations F (\rightarrowtail) }

satisfying (1 - 4).

Weak equivalences are the equivalences we care about (e.g. weak htpy eq. in $\text{Top}/\underline{sSet}$) and the axioms ensure that they indeed behave like equivalences.

(1) If



and $f \in W$ then $g \in W$.

if $W = \{ \text{id} \circ \text{id} \}$ then $g^{-1} = \begin{matrix} \xrightarrow{v} \\ \uparrow f^{-1} \\ \xrightarrow{v} \end{matrix}$

(2) Given $\xrightarrow{f} \xrightarrow{g}$, if two of f, g, gf are in W then so is the last.

2-out-of-3 property

Although we are only interested in CW-complexes / Kan complexes, these objects are NOT closed under the sort of things we want to do (e.g. taking (co)limits in $\text{Top}/\underline{\text{sSet}}$).

So we'll work in $\text{Top}/\underline{\text{sSet}}$ & use the model structure to capture the "niceness" of CW-complexes.

- CW-complexes are built as colimits (of disks),

so it's easy to construct maps out of them.

- Kan complexes have the dual property;

it's easy to construct maps into them (by definition).

$$\begin{array}{ccc} \Delta_k^n & \xrightarrow{\alpha} & X \\ \downarrow & \nearrow & \nearrow \beta \\ \Delta^n & & \end{array}$$

In the language of model categories,

- CW-complexes are cofibrant, and

- Kan complexes are fibrant.

precise definitions in a few slides

Cofibrations are maps with a relative version of easy-to-map-out-of property.

i.e. it's easy to extend maps along them.



Dually, fibrations have a relative version of easy-to-map-into property.

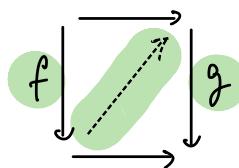
i.e. it's easy to lift maps through them.

can extend a map defined
on a subset to the whole set

Cofibrations are "like injections" and fibrations are "like surjections".

can always pick a point in the preimage

Def: Let f, g be morphisms in \mathcal{M} . If every commutative square of the form

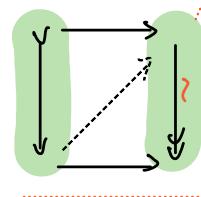


simultaneously extending the unique composite along f and lifting it through g .

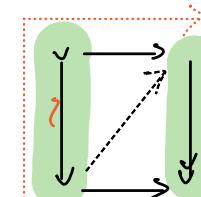
admits a **diagonal lift**, we say:

- f has the **left lifting property** (LLP) w.r.t. g , and
- g has the **right lifting property** (RLP) w.r.t. f .

(3) $\mathcal{C} = \{\text{maps with LLP w.r.t. all maps in } \mathcal{F} \circ \mathcal{W}\}$.



these lifts "trivially" exist up to htpy.



$\tilde{\mathcal{F}} = \{\text{maps with RLP w.r.t. all maps in } \mathcal{C} \circ \mathcal{W}\}$.

(It follows from (1) & (3) that $\mathcal{C} \circ \mathcal{W} \neq \tilde{\mathcal{F}} \circ \mathcal{W}$ can be similarly characterised by lifting properties.)

(4) Any map f in \mathcal{M} admits factorisations of the form :

$$\begin{array}{ccc} e_{\text{W}} & \nearrow & e_{\text{F}} \\ & \sim & \\ & \searrow & \\ & f & \end{array}$$

and

$$\begin{array}{ccc} e_{\text{W}} & \nearrow & e_{\text{F} \circ \text{W}} \\ & \searrow & \\ & f & \end{array}$$

approximation of f
by co/fibration.

In most cases, these factorisations can be chosen functorially : $\mathcal{M}^{\{ \cdot \rightarrow \cdot \}} \rightarrow \mathcal{M}^{\{ \cdot \rightarrow \cdot \rightarrow \cdot \}}$

Examples :

\mathcal{M}	<u>Top</u> classical / Quillen	<u>sSee</u> classical / Kan / Quillen	<u>Cart</u> canonical / fork
\mathcal{W}	weak htpy eq.	weak htpy eq.	equivalences
\tilde{f}	RLP wrt $D^n \times \{0\}$ \downarrow $D^n \times I$ Serre fibrations	RLP wrt Λ_k^n \downarrow Δ^n Kan fibrations	RLP wrt $\{ \cdot \cong \cdot \}$ \downarrow isofibrations
\mathcal{C}	LLP wrt $\tilde{f} \circ \mathcal{W}$	monomorphisms	injective on objects

Def: $X \in M$ is

- **cofibrant** if $0 \xrightarrow{!} X$ is a **cofibration**
e.g. CW-ex in Top

easy to map out of

- **fibrant** if $X \xrightarrow{!} I$ is a **fibration**.
e.g. Kan ex in $s\text{Set}$

easy to map into

A **cofibrant replacement** of $X \in M$ is **cofibrant**

$QX \in M$ t/w $QX \xrightarrow{\sim} X$.

A **fibrant replacement** _____ **fibrant**

$RX \in M$ t/w $X \xrightarrow{\sim} RX$.

Observation: For any $X \in M$, by (4) we can always construct

$$\begin{array}{ccc} & QX & \\ & \nearrow \downarrow & \\ 0 & \xrightarrow{!} & X \end{array}$$

and

$$\begin{array}{ccc} & RX & \\ & \nearrow \downarrow & \\ X & \xrightarrow{!} & I \end{array}$$

- • Can extend Q & R to **morphisms**, e.g.

$$\begin{array}{ccccc} 0 & \xrightarrow{!} & QY & & \\ \downarrow & \nearrow Qf & \downarrow f & & \\ QX & \xrightarrow{\sim} & X & \xrightarrow{f} & Y \end{array}$$

In general, Q & R are functorial only up to htpy.

- Any $X \in M$ is connected by a **zigzag** of weak eq. to

an object that is both **fibrant** & **cofibrant**:

$$\begin{array}{ccccc} 0 & \longrightarrow & QX & \xrightarrow{\sim} & X \\ & & \downarrow f & & \\ & & RQX & & \\ & & \downarrow & & \\ & & I & & \end{array}$$

In a model category, the notion of homotopy makes sense.

Def: Let $X, Y \in \mathcal{M}$.

- A **cylinder object** for X is a factorisation
e.g. $X \times I$ in Top, $X \times \Delta^1$ in sSet.

- A **path object** for Y is a factorisation

$$\begin{array}{ccc} & Cyl(X) & \\ \langle i_0, i_1 \rangle \nearrow & \downarrow & \\ X \amalg X & \xrightarrow{\quad} & X \\ & \langle id, id \rangle & \end{array}$$

$$\begin{array}{ccc} & Path(Y) & \\ \nearrow & \downarrow & \\ Y & \xrightarrow{(id, id)} & Y \times Y \\ & \downarrow & \\ & (p_0, p_1) & \end{array}$$

Def: Let $f, g: X \rightarrow Y$ in \mathcal{M} .

- A **left homotopy** from f to g is (a choice of $Cyl(X)$ t/w)
a map $Cyl(X) \xrightarrow{H} Y$ s.t. $H i_0 = f$ & $H i_1 = g$. $f \stackrel{l}{\sim} g$
- A **right homotopy** from f to g is (a choice of $Path(Y)$ t/w)
a map $X \xrightarrow{K} Path(Y)$ s.t. $p_0 K = f$ & $p_1 K = g$. $f \stackrel{r}{\sim} g$

Observation: $f \stackrel{l}{\sim} g \Rightarrow hf \stackrel{l}{\sim} hg$ for $\cdot \xrightarrow{k} \cdot \xrightarrow{f} \cdot \xrightarrow{g} \cdot \xrightarrow{h}$.

$$f \stackrel{r}{\sim} g \Rightarrow fk \stackrel{r}{\sim} gk$$

Fact: If $X \in \mathcal{M}$ is cofibrant and $Y \in \mathcal{M}$ is fibrant, then \sim^l & \sim^r agree on $\mathcal{M}(X, Y)$, and moreover it's an equivalence relation. We'll just write \sim in this case.

Whitehead's theorem:

Let $f: X \rightarrow Y$ with $X \& Y$ both fibrant & cofibrant.

Then $f \in \mathcal{W}$ iff f is a homotopy eq. $\exists g: Y \rightarrow X$ s.t. $gf \sim \text{id}_X$ & $fg \sim \text{id}_Y$.

Def: The homotopy category $\text{Ho}(\mathcal{M})$ is obtained from \mathcal{M} by

- restricting to the fibrant & cofibrant objects, and
- quotienting the hom-sets by \sim .

Fact: • There is a functor $\Upsilon: \mathcal{M} \rightarrow \text{Ho}(\mathcal{M})$ given by
 $\Upsilon(X \xrightarrow{f} Y) = (RQX \xrightarrow{[RQf]} RQY)$ htpy class containing RQf .

• $f \in \mathcal{W} \Leftrightarrow \Upsilon(f)$ is an isomorphism.

• Υ is universal among those functors $\mathcal{M} \rightarrow \mathcal{C}$ that send weak equivalences to isomorphisms:

$$\begin{array}{ccc} \mathcal{M} & \xrightarrow{\Upsilon} & \text{Ho}(\mathcal{M}) \\ & \cong \downarrow & \downarrow \text{essentially unique} \\ & & \mathcal{C} \end{array}$$

$\text{Ho}(\mathcal{M})$ is the localisation of \mathcal{M} at \mathcal{W} .

What's the right notion of **morphism** between **model categories**?

If we just want $F: M \rightarrow N$ to induce $\text{Ho}(M) \rightarrow \text{Ho}(N)$

then F simply needs to preserve **weak eq.**, but such F can mess up **computations**.

Here's a hint.

Recall: co/fibrations facilitate computations

Exercise: Consider an adjunction $\mathcal{C} \rightleftarrows \mathcal{D}$ and morphisms $f: A \rightarrow B$ in \mathcal{C} &

$g: X \rightarrow Y$ in \mathcal{D} . Then there is a **bijection**

$$\begin{array}{ccc} \begin{matrix} A \\ \downarrow f \\ B \end{matrix} & \xrightarrow{\quad} & \begin{matrix} GX \\ \downarrow Gg \\ GY \end{matrix} \quad \text{in } \mathcal{C} & \longleftrightarrow & \begin{matrix} FA \\ \downarrow Ff \\ FB \end{matrix} & \xrightarrow{\quad} & \begin{matrix} X \\ \downarrow g \\ Y \end{matrix} \quad \text{in } \mathcal{D} \end{array}$$

and either admits a **diagonal lift** iff the other does. Consequently,

$$f \text{ has LLP wrt. } Gg \iff Ff \text{ has LLP wrt. } g.$$

Def: An adjunction between model categories $\mathcal{M} \rightleftarrows \mathcal{N}$ is a Quillen adjunction if

- G preserves triv. fib. \Leftrightarrow • F preserves cofibrations ; and
- G preserves fibrations.

F preserves triv. cof. $\Leftrightarrow F$ is left Quillen / G is right Quillen

Quillen adjunctions present homotopically meaningful universal properties.

Fact: A Quillen adjunction $\mathcal{M} \rightleftarrows \mathcal{N}$ induces an adjunction

$$\mathcal{M} \rightleftarrows \mathcal{N}$$

$$\text{Ho}(\mathcal{M}) \rightleftarrows \text{Ho}(\mathcal{N})$$

derived functors

e.g. $s\text{Set} \rightleftarrows \text{Top}$. In fact, it's an example of a Quillen equivalence.

Def: A Quillen equivalence is a Quillen adjunction $F \dashv G$ s.t.

$$f: X \rightarrow GY \in \mathcal{W}_{\mathcal{M}} \iff \bar{f}: FX \rightarrow Y \in \mathcal{W}_{\mathcal{N}}$$

(transpose of f)

Fact: $LF \dashv RG$ is an adjoint equivalence iff $F \dashv G$ is a Quillen equivalence.