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#### **Problem 1**

- A. Is the graph G connected? If not, what are the connected components for G? No, connected components are: (A, B, C, F, H, G),(D, E, I)
- B. Draw a spanning tree/forest for G
- C. Is G a Hamiltonian graph? No, because there is no Hamiltonian cycle in this graph.
- D. Is there a Vertex Cover of size less than or equal to 5 for G? If so, what is the Vertex Cover? Yes, the Vertex cover is (D, E, F, A, G).

#### **Problem 2**

Hamiltonian Graphs. The following graph has a Hamiltonian cycle. Find it.

### **Problem 3**

Vertex Covers. Create an algorithm for computing the smallest size of a vertex cover for a graph. The input of your algorithm is a set V of vertices along with a set E of edges. Assume you have the following functions available (no need to implement these): ● computeEndpoints(edge) − returns the vertices that are at the endpoints of the input edge ● belongsTo(vertex, set) − returns true if the input vertex is a member of the given

Algorithm: conjute venera Cover

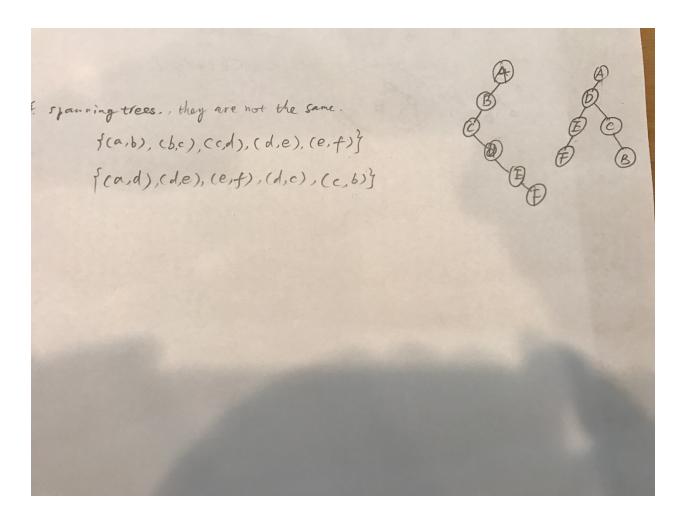
Input: G(V, E)

Output: a vertex cover with smallest size of vertices
initialize a vertext list,
while (E. size > 0). do

compute the vertex appears most times in oull E,
add the vertext to the list,

remove all edges with the vertext in the edge.

### **Problem 4**



# **Problem 5**

Write the pesudo-code for compute connected components algorithm discussed in class. Your algorithm can be built on top of DFS discussed in the slides.

Algorithm: compute Connected Congoneurs

Input: a graph G=(V, E),

Output: a List of Component

Befine a List of Vertex,

While (there is vertex not visited)

Single Component Loop, (add all reachable vertex to the component

Veturn List;

### **Problem 6**

Write the pesudo-code for the algorithm, discussed in class, that computes the shortest path length between two vertices in a graph. You can assume that: a. The graph is connected. b. A version of BFS is provided that accepts a specified starting

Algorithm find Shortest flooth

Input: Two vertices in a G=(V, E),

Output: List of Vertices from either VI to VE or Vato V,

choose VI or V2 as root, a List, a infound flog,

Individual a queue Q, erqueue the root,

while ('infound) {

N = Q. Size(); part

(Neare a List of Vertices as parths.

for (i=1, to n)

V v = dequeue ()

mark v as visited;

clifts (v equals Vs) {

Neturn list(i);