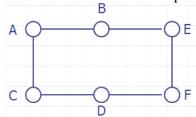
- 1. Answer the following questions about the graph G having n = 6 vertices, below.
  - a. Is G Hamiltonian?
  - b. Can you find two non-adjacent vertices the sum of whose degrees is less than 6?
  - c. Do these facts contradict Ore's Theorem? Explain.



- 2. Show that TSP is NP-complete. (Hint: use the relationship between TSP and HamiltonianCycle discussed in the slides. You may assume that the HamiltonianCycle problem is NP-complete.)
- 3. Below is another variation of the Knapsack problem. Given a set  $S = \{s_0, s_1, ..., s_{n-1}\}$  of items, weights  $\{w_0, w_1, ..., w_{n-1}\}$ , values  $\{v_0, v_1, ..., v_{n-1}\}$ , a max weight W, and a min value V, find a subset T of S whose total value is no less than V and total weight is at most W.

Show that the SubsetSum problem is polynomial reducible to this Knapsack problem.

- 4. Show that the worst case for VertexCoverApprox can happen by giving an example of a graph G which has these properties:
  - a. G has a smallest vertex cover of size s
  - b. VertexCoverApprox outputs size 2\*s as its approximation to optimal size.
- 5. Find an O(n) algorithm that does the following: Given a size n input array of integers, output the first numbers in the array (from left to right) whose sum is exactly 10 (or indicate that no such numbers can be found).
- 6. The decision problem formulation of the Vertex Cover problem is this: Given a positive integer k, and a graph G, is there a vertex cover for G having size  $\leq k$ ? Show that this decision problem belongs to NP.