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Problem 1

Answer the following questions about the graph G having $n = 6$ vertices, below. a. Is G Hamiltonian? Yes b. Any two non adjacent vertices. c. No

Problem 2

To show Any Problem Q in NP , $Q \xrightarrow{\text{poly}} \text{TSP}$.

Just show Hamiltonian $\xrightarrow{\text{poly}} \text{TSP}$.

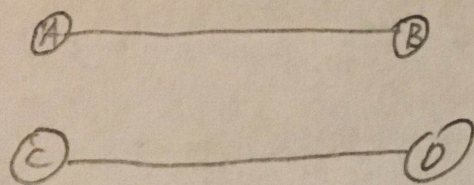
We can use input of the Hamiltonian problem, after s of the input, ($O(p(n))$), then use the new input as n input. then we can see after $O(p(n))$, the so Hamiltonian could be processed to be the solution with much work. so, Hamiltonian $\xrightarrow{\text{poly}} \text{TSP}$.

Problem 3

To show subset sum poly \rightarrow new Knapsack Problem.

We can set V to be 0, in new Knapsack and W to be subset problem, so the solution to the subset is solution of knapsack problem in occasion with
 \Rightarrow Subset Problem poly \rightarrow new knapsack Problem

Problem 4



Smallest vertex size
VertexCoverApprox:

Problem 5

Use dynamic programming, $B[i, j]$ denotes whether there is a subset with previous i integers in the set whose sum is j . We can build such an array from $B[0, 1]$ to $B[n, 10]$, then iterate over the $B[0, 10]$ to $B[n, 10]$ and return the i with the first i that makes $B[i][10]$ to be true.

Problem 6

6. Just ^{to} verify the solution in $p(n)$ time complexity.

$G = (V, E)$.

1. To verify that every edge in G , at least one vertex solution, this takes $O(m)$. ($O(n^2)$ in the worst case)

2. To show the size of vertex cover is less or this takes $O(1)$.

$\Rightarrow Q \subseteq NP$.