MATH 285: Stochastic Processes

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Today: Hidden Markov chains

Homework 4 is due on Friday, February 11, 11:59 PM

Example: Occasionally Dishonest Casino Casino has two dice: fair (F) and loaded (L). • F: P(i)=6 • L: P(i) = 0.5, P(i) = 0.1 for $i \ge 2$ Casino switches the die: • F → L with probability 0.05 · L → F with probability 0.95 As a player you don't know which die is in use, you only observe the number that is rolled. Suppose you play the game (roll the die) 6 times and observe 1,1,1,1,1. Q: What is the most likely sequence of dice used by ?

Hidden Markov Model Def 16.2 A Hidden Markov Model (HMM) is a pair of stochastic processes (Xn, Yn), where (Yn) is a Markov chain with state space S, and (Xn) nzo has a possibly different state space R, and the vector valued process Zn = (Xn, Yn) is a Markov chain. For yes and xeR the conditional probabilities $e_{y}(x) = .$ are called the emission probabilities. Let p: SxS > [0,1] be the transition kernel for (Yn). It is taken as an assumption that the transition kernel for (Zn) is $P \mid Z_{n+1} = (x', y') \mid Z_n = (x, y) =$

Hidden Markov Model Remarks (1) In general (Xn) is not a Markov chain (2) Transition kernel for (Zn) does not depend on x; this is not true in general for Markov chains on SxR P[Xo=xo, Yo=yo, Xi=xi,..., Yn-i=yn-i, Xn=xn, Yn=yn]

⇒ P[Xo=xo, X,=x,,--, Xn=xn | Yo=yo, --, Yn=yn]=

Example: Occasionally Dishonest Casino (2)

Construct a HMM that models ODC $S = (Y_n)$ MC on S with transition probabilities $P(F_1L) = P(L_1F) = P(L_1F$

$$R = X_n \in R$$

Emission probabilities:
$$e_F(i) = for all i \in \mathbb{R}$$

 $e_L(i) = \begin{cases} i = 1 \\ i \in \{2,3,4,5,6\} \end{cases}$

$$Z_n = (X_n, Y_n)$$

The forward algorithm Let (Xn, Yn) be a HMM. Denote · X = (xo, x,..., xn) the observed sequence · y = (yo, y,, --, yN) the state sequence • $P[X] = P[X_0 = X_0, ..., X_N = X_N]$ • P[x,y] = P[X0 = x0,..., Xn = xn, Y0 = y0,..., Yn = yn] Q: What is the probability of (yo, y,,...,yn) given that we observe $(x_0, x_1, ..., x_N)$? Using the above notation, we have to compute P(y|x) =We know that P[x, y] =

The forward algorithm Direct way of computing $\mathbb{P}[\times]$ $\mathbb{P}[\times] =$ Problem: computationally infeasible ~ computations grows exponentially fast with N The forward algorithm allows to compute P[X] in polynomial time. Fix observed sequence x = (xo, x, ..., xn). For any yes and ne so, 1,..., N} define the probability dn (y) = that first n observations occured and the hidden state is y. The forward algorithm

Then
$$d_{n+1}(y') = \mathbb{P}[X_0 = x_0, X_1 = x_1, ..., X_n = x_n, X_{n+1} = x_{n+1}, Y_{n+1} = y']$$

Now condition on $X_0, X_1, ..., X_n, Y_n$

$$P[X_0 = x_0, X_1 = x_1, X_n = x_n, Y_n = y, X_{n+1} = x_{n+1}, Y_{n+1} = y']$$

$$P[Z_{n+1}(x_{n+1},y')|X_0=x_0,X_1=x_1,...,X_n=x_n,Y_n=y]$$

The forward algorithm Therefore, dn+1 (4') = and we can compute P[x] =Complexity of the forward algorithm: · < < (y) = P[X = x 0, Y = y] = By (★) we need ~ operations to compute &n(y) By (**) we have to compute dn(y) for all n,y

Proof of Lemma 16.5

$$P[Z_{n+1}(x_{n+1}, y') | X_0 = Z_0, X_1 = Z_1, ..., X_n = Z_n, Y_n = y]$$

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 $P[X_0 = Z_0, X_1 = Z_1, ..., X_n = Z_n, Y_n = y, X_{n+1} = Z_{n+1}, Y_{n+1} = y']$

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