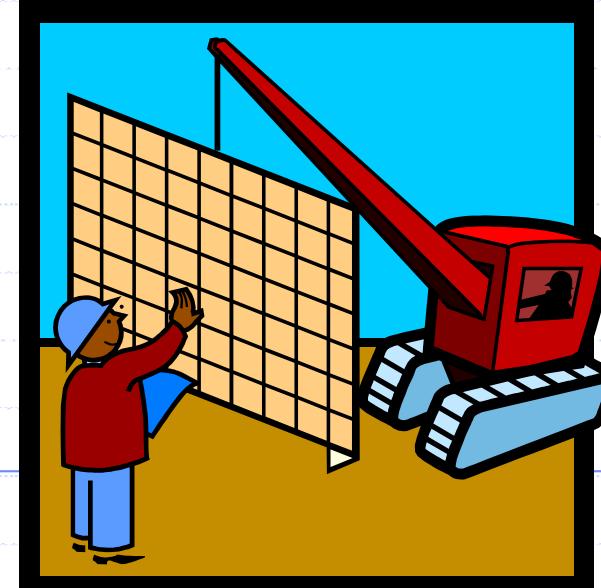
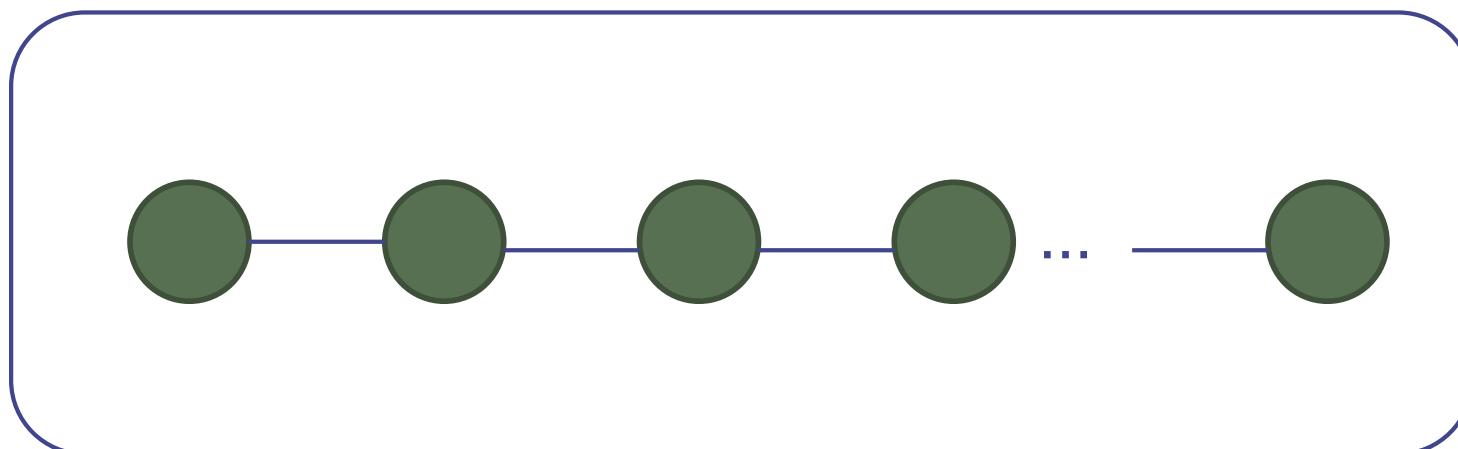


Vector, List and Sequence



Overview and Reading

- ◆ Reading: Chapters: 6.1, 6.2, and 6.3
- ◆ A data structure that stores n elements in a linear order
 - Called list or sequence
- ◆ Didn't we learn “array” and “linked list”?
 - We are talking about more abstract ADTs than them



Three ADTs

◆ Vector (also called Array List)

- Access each element using a notion of **index** in [0,n-1]
- Index of element e: the number of elements that are before e
- Typically we use the “**index**” (e.g., [])
- A more general ADT than “array”

◆ List

- Not using an index to access, but use a node to access
- Insert a new element e before some “**position**” p
- A more general ADT than “linked list”

◆ Sequence

- Can access an element as vector and list (using both **index** and **position**)

◆ (Note) Can implement the above ADTs using various ways

- array, singly linked list, doubly linked list, circular linked list

Vectors (or Array Lists)



The Array List ADT

❑ The Vector or Array List

ADT extends the notion of array by storing a sequence of objects

❑ An element can be accessed, inserted or removed by specifying its index (number of elements preceding it)

❑ An exception is thrown if an incorrect index is given (e.g., a negative index)

◆ Main methods:

- **at(integer i):** returns the element at index i without removing it
- **set(integer i, object o):** replace the element at index i with o
- **insert(integer i, object o):** insert a new element o to have index i
- **erase(integer i):** removes element at index i

◆ Additional methods:

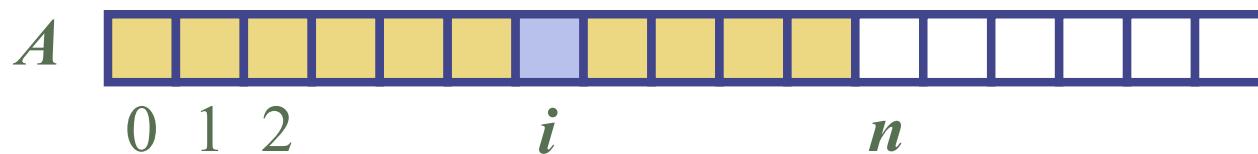
- **size()**
- **empty()**

Applications of Array Lists

- ◆ Direct applications
 - Sorted collection of objects (elementary database)
- ◆ Indirect applications
 - Auxiliary data structure for algorithms
 - Component of other data structures
- ◆ Basically, every place where you can use “array”.

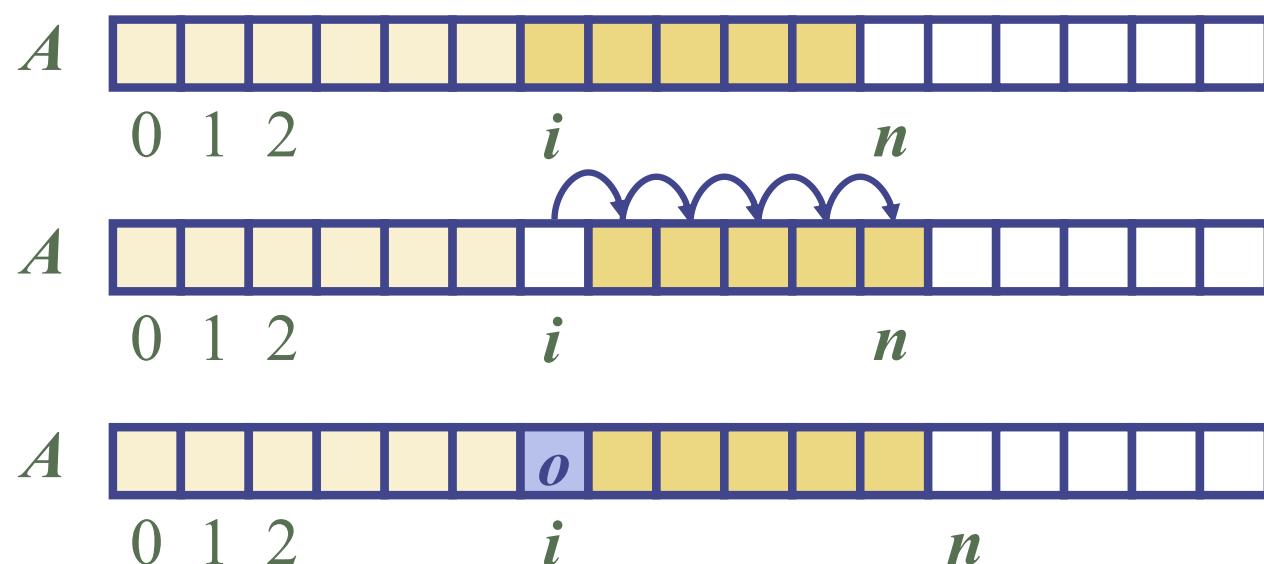
Array-based Implementation of Vector

- ◆ Use an array A of size N
- ◆ A variable n keeps track of the size of the array list (number of elements stored)
- ◆ Operation $\text{at}(i)$ is implemented in $O(1)$ time by returning $A[i]$
- ◆ Operation $\text{set}(i,o)$ is implemented in $O(1)$ time by performing $A[i] = o$



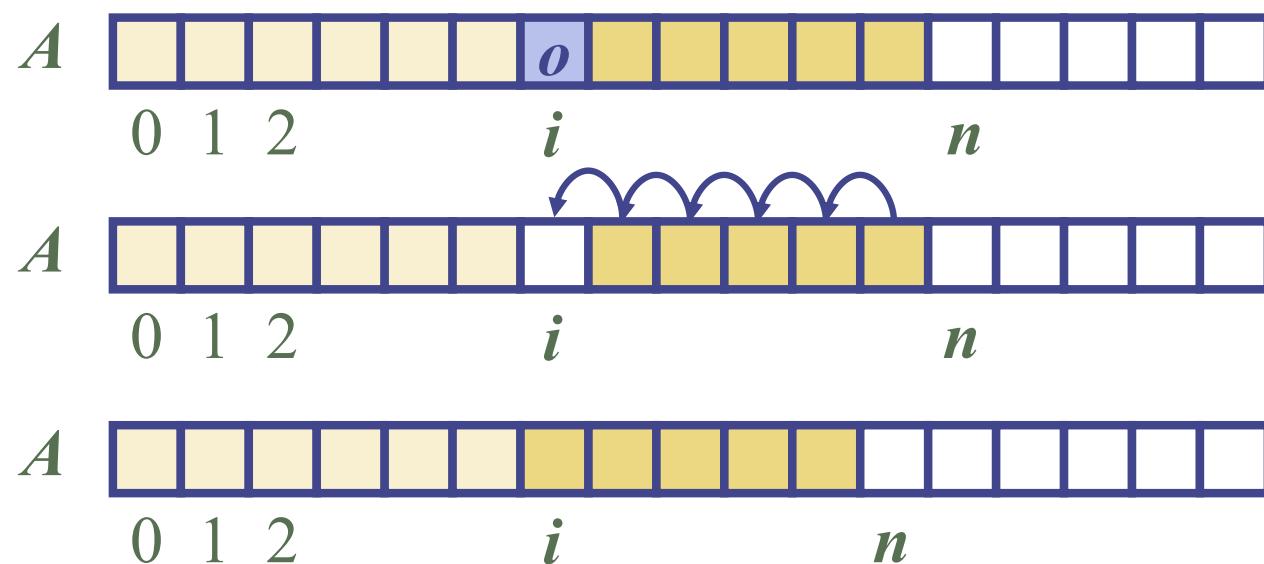
Insertion

- ◆ In operation $\text{insert}(i, o)$, we need to make room for the new element by shifting forward the $n - i$ elements $A[i], \dots, A[n - 1]$
- ◆ In the worst case ($i = 0$), this takes $O(n)$ time



Element Removal

- ◆ In operation $\text{erase}(i)$, we need to fill the hole left by the removed element by shifting backward the $n - i - 1$ elements $A[i + 1], \dots, A[n - 1]$
- ◆ In the worst case ($i = 0$), this takes $O(n)$ time



Performance

- ◆ In the array-based implementation of an array list:
 - The space used by the data structure is $O(n)$
 - *size*, *empty*, *at* and *set* run in $O(1)$ time
 - *insert* and *erase* run in $O(n)$ time in worst case
- ◆ If we use the array in a circular fashion, operations *insert*(0, x) and *erase*(0, x) run in $O(1)$ time
- ◆ In an *insert* operation, when the array is full, instead of throwing an exception, we can replace the array with a larger one

Growable Array-based Array List

- ❑ In an $\text{insert}(o)$ operation (without an index), we always insert at the end
- ❑ When the array is full, we replace the array with a larger one
- ❑ How large should the new array be?
 - Incremental strategy: increase the size by a constant c
 - Doubling strategy: double the size

Algorithm $\text{insert}(o)$

```
if  $t = S.length - 1$  then
     $A \leftarrow$  new array of
        size ...
    for  $i \leftarrow 0$  to  $n-1$  do
         $A[i] \leftarrow S[i]$ 
     $S \leftarrow A$ 
     $n \leftarrow n + 1$ 
     $S[n-1] \leftarrow o$ 
```

- ◆ For size n array, “re-grow” operation requires n copies

Which is better?: Incremental or Doubling

◆ Comparison Method 1

- Given the current size of $S = n$
- Worst-case running time
 - ◆ Incremental strategy: $O(1)$
 - ◆ Doubling strategy: $O(n)$

◆ Are you happy?

- Happy if your focus is really the worst-case
- Unhappy
 - ◆ For doubling strategy, the total number of resizing array size would be small

◆ Can we reconsider the analysis method?

Which is better?: Incremental or Doubling

◆ Comparison Method2

- Compute the total time $T(n)$ needed to perform a series of n insert(o) operations
- Assume that we start with an empty stack represented by an array of size 1

◆ We call amortized time of an insert operation **the average time taken by an insert over the series of operations**, i.e., $T(n)/n$

- This can be a fairer comparison in some cases

◆ Amortized analysis (분할상환분석 in Wiki)

Incremental Strategy Analysis

- ◆ We replace the old array with a new one $k = n/c$ times
- ◆ The total time $T(n)$ of a series of n insert operations is proportional to

$$\begin{aligned}n + c + 2c + 3c + 4c + \dots + kc &= \\n + c(1 + 2 + 3 + \dots + k) &= \\n + ck(k + 1)/2\end{aligned}$$

- ◆ Since c is a constant, $T(n)$ is $O(n + k^2)$, i.e., $O(n^2)$
- ◆ The amortized time of an insert operation is $O(n)$

Doubling Strategy Analysis

- ◆ We replace the old array with a new one

$$k = \log_2 n \text{ times}$$

- ◆ The total time $T(n)$ of a series of n insert operations is proportional to

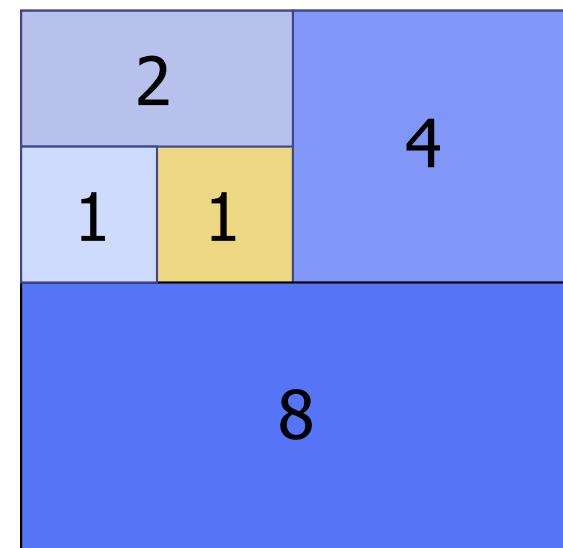
$$n + 1 + 2 + 4 + 8 + \dots + 2^k =$$

$$n + 2^{k+1} - 1 =$$

$$3n - 1$$

- ◆ $T(n)$ is $O(n)$

geometric series



- ◆ The amortized time of an insert operation is $O(1)$

Professor, I have a question

- ◆ In “computing spans”, why didn’t you do amortized analysis?
- ◆ Can we do it?
- ◆ Is it meaningful?
- ◆ Think about this!
 - I am ready to discuss if you get your version of answer ready.

Algorithm	<i>spans2(X, n)</i>	#
	$S \leftarrow$ new array of n integers	n
	$A \leftarrow$ new empty stack	1
	for $i \leftarrow 0$ to $n - 1$ do	n
	while ($\neg A.empty()$) \wedge	
	$X[A.top()] \leq X[i]$) do n	
	$A.pop()$	n
	if $A.empty()$ then	n
	$S[i] \leftarrow i + 1$	n
	else	
	$S[i] \leftarrow i - A.top()$	n
	$A.push(i)$	n
	return S	1

Vectors in C++ STL

```
#include <vector>           // provides definition of vector
using std::vector;          // make vector accessible

vector<int> myVector(100);   // a vector with 100 integers
```

- vector(*n*):** Construct a vector with space for *n* elements; if no argument is given, create an empty vector.
- size():** Return the number of elements in *V*.
- empty():** Return true if *V* is empty and false otherwise.
- resize(*n*):** Resize *V*, so that it has space for *n* elements.
- reserve(*n*):** Request that the allocated storage space be large enough to hold *n* elements.
- operator[i]:** Return a reference to the *i*th element of *V*.
 - at(*i*):** Same as *V[i]*, but throw an `out_of_range` exception if *i* is out of bounds, that is, if *i* < 0 or *i* ≥ *V.size()*.
- front():** Return a reference to the first element of *V*.
- back():** Return a reference to the last element of *V*.
- push_back(*e*):** Append a copy of the element *e* to the end of *V*, thus increasing its size by one.
- pop_back():** Remove the last element of *V*, thus reducing its size by one.

Difference between resize() and reserve()?

Logistics

- ◆ First programming assignment
 - Deadline: Sep, 19th
- ◆ Problem Solving Homework
 - Deadline: Oct, 1st
- ◆ You should keep reading the textbook
- ◆ Sep 24th, 26th : No class
 - Thanksgiving

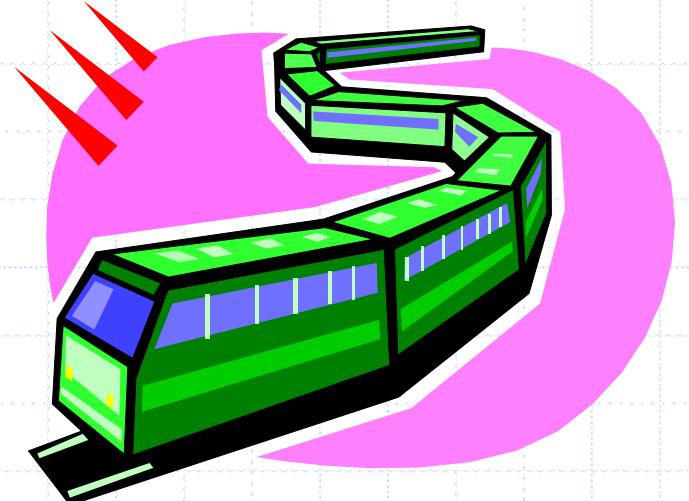
Last Class

◆ Vector and List

◆ Vector

- Access elements by “index”
- Incremental vs. Doubling Strategy
 - ◆ Amortized analysis

Lists



(Node) List ADT

- ◆ The Node List ADT models a sequence of positions storing arbitrary objects
- ◆ It establishes a before/after relation between positions
- ◆ Generic methods:
 - `size()`, `empty()`

□ Iterators:

- `begin()`, `end()`

□ Update methods:

- `insertFront(e)`,
`insertBack(e)`
- `removeFront()`,
`removeBack()`

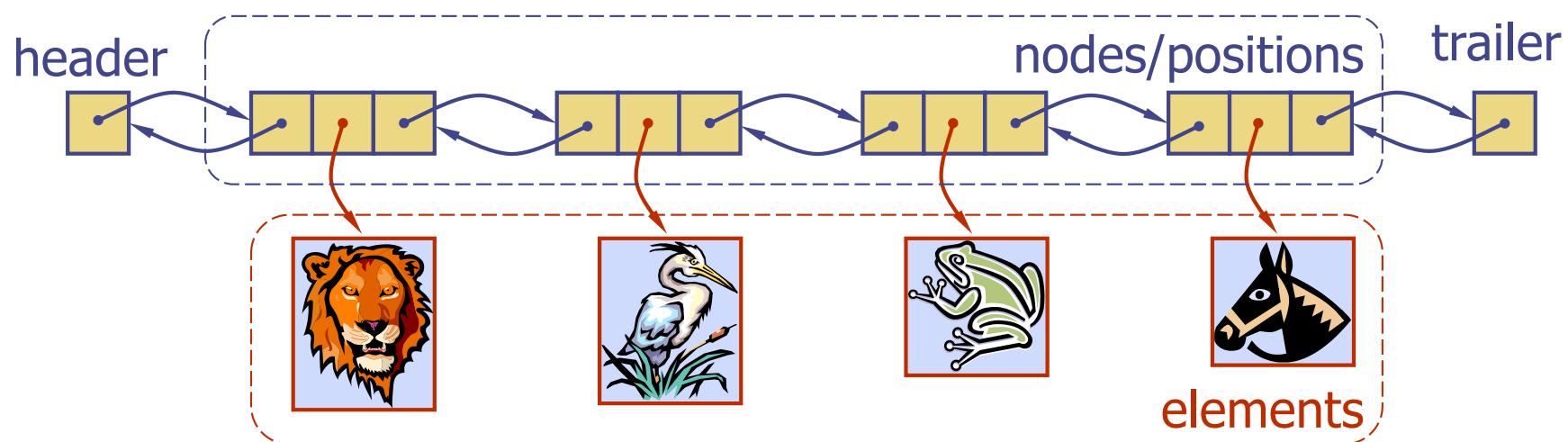
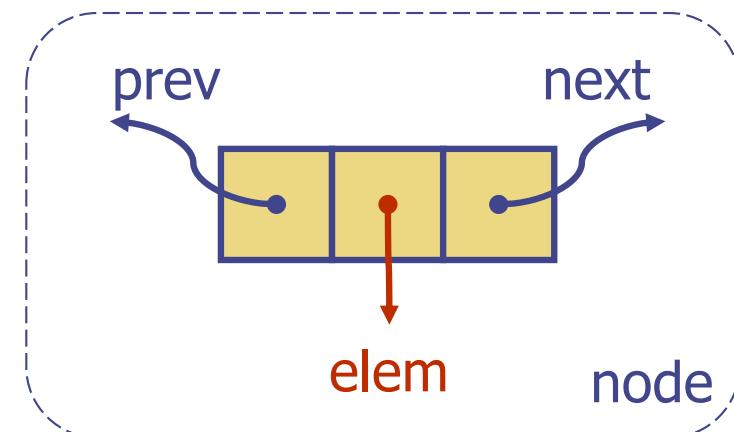
□ Iterator-based update:

- `insert(p, e)`
- `remove(p)`

(Question) No method for accessing a specific node?
We will talk about this later

Implementation based on DLL (covered this)

- ◆ A doubly linked list provides a natural implementation of the Node List ADT
- ◆ Nodes implement Position and store:
 - element
 - link to the previous node
 - link to the next node
- ◆ Special trailer and header nodes



Performance

- ◆ In the implementation of the List ADT by means of a doubly linked list
 - The space used by a list with n elements is $O(n)$
 - The space used by each position of the list is $O(1)$
 - All the operations of the List ADT run in $O(1)$ time

Lists in C++ STL

```
#include <list>
using std::list;                      // make list accessible
list<float> myList;                   // an empty list of floats
```

list(*n*): Construct a list with *n* elements; if no argument list is given, an empty list is created.

size(): Return the number of elements in *L*.

empty(): Return true if *L* is empty and false otherwise.

front(): Return a reference to the first element of *L*.

back(): Return a reference to the last element of *L*.

push_front(*e*): Insert a copy of *e* at the beginning of *L*.

push_back(*e*): Insert a copy of *e* at the end of *L*.

pop_front(): Remove the fist element of *L*.

pop_back(): Remove the last element of *L*.

Containers, Iterators, and Generic algorithms



Sorting: Vector and List

- ◆ I want to find “yiyung” in Vector or List objects

```
vector<string> V(100);  
list<string> L(100);  
// some data insertion to V and L
```

```
//Design 1: different function  
find_vector(&V);  
find_list(&L);
```

```
//Design 2: function overloading  
find(&V);  
find(&L);
```

Do you like these? Why? Why not?

This is how we can do in C++

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

It is cool. But why is it cool?

Mysterious things

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

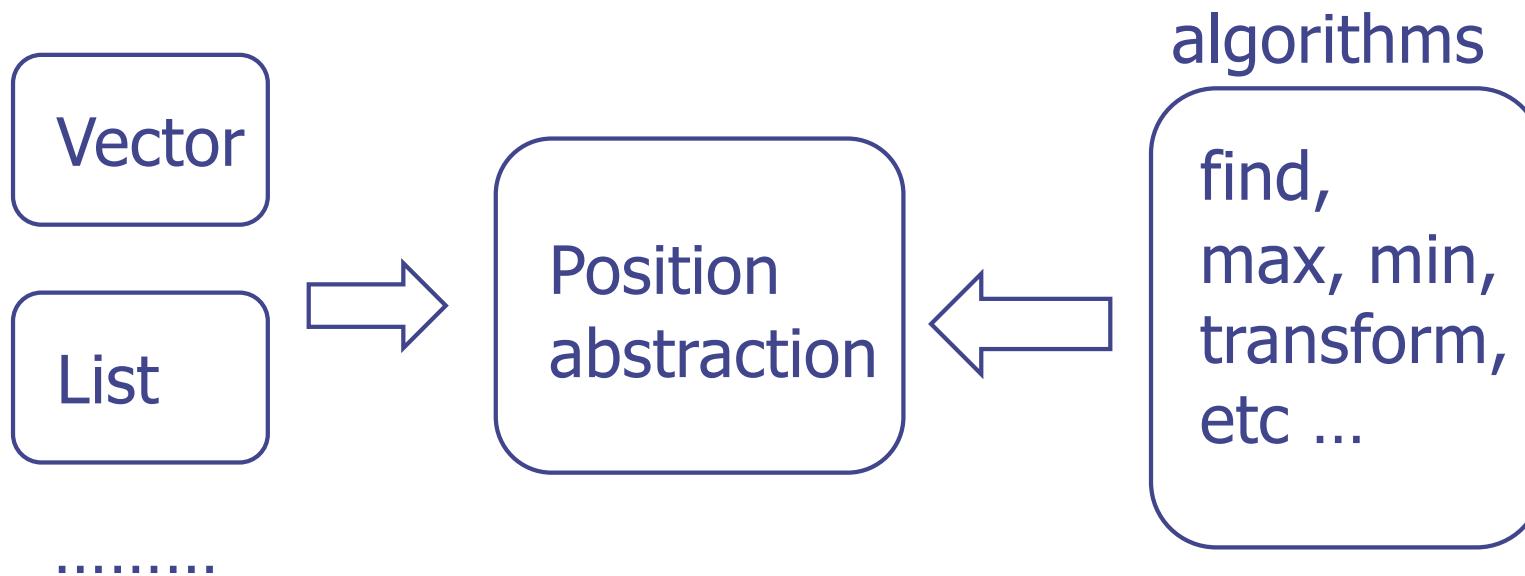
    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

iterator? Looks like a “position” of vector or list. Hmm....

Goal and Design Challenge

- ◆ Lots of data structures (or classes in C++) that can contain various types of elements
 - “Container”
 - Examples: Vector, List, deque, set, map, etc ...



- ◆ How are you going to design this concept?
 - Again, from C++ STL designer's perspective

Position ADT

- The Position ADT models the notion of place within a data structure where a single object is stored
- It gives a unified view of diverse ways of storing data, such as
 - a cell of an array
 - a node of a linked list
- “A” method of accessing the element at position **p**:
 - object **p.element()**: returns the element at position
 - In C++ it is convenient to implement this as ***p**
 - Operator overloading
- ◆ Implemented as “iterator” in C++

Containers and Iterators in C++

- ◆ An **iterator** abstracts the process of scanning through a collection of elements
- ◆ A **container** is an abstract data structure that supports element access through iterators
 - Data structures that support iterators
 - Examples include Stack, Queue, Vector, List
 - `begin()`: returns an iterator to the first element
 - `end()`: return an iterator to an imaginary position just after the last element
- ◆ An iterator behaves like a pointer to an element
 - `*p`: returns the element referenced by this iterator
 - `++p`: advances to the next element
- ◆ Extends the concept of **position** by adding a traversal capability

Example codes again

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

Ah-ha, it's an iterator!

Various Iterators

- ◆ (standard) iterator: allows read-write access to elements
- ◆ const iterator: provides read-only access to elements
- ◆ bidirectional iterator: supports both $++p$ and $-p$
- ◆ random-access iterator: supports both $p+i$ and $p-i$

STL Iterators in C++

- Each STL container type C supports iterators:

- `C::iterator` – read/write iterator type
- `C::const_iterator` – read-only iterator type
- `C.begin()`, `C.end()` – return start/end iterators

- This iterator-based operators and methods:

- `*p`: access current element
- `+p`, `-p`: advance to next/previous element
- `C.assign(p, q)`: replace C with contents referenced by the iterator range `[p, q]` (from p up to, but not including, q)
- `insert(p, e)`: insert e prior to position p
- `erase(p)`: remove element at position p
- `erase(p, q)`: remove elements in the iterator range `[p, q)`

Back to Iterator: STL Iterator-based Functions

- `vector(p, q)`: Construct a vector by iterating between *p* and *q*, copying each of these elements into the new vector.
- `assign(p, q)`: Delete the contents of *V*, and assigns its new contents by iterating between *p* and *q* and copying each of these elements into *V*.
- `insert(p, e)`: Insert a copy of *e* just prior to the position given by iterator *p* and shifts the subsequent elements one position to the right.
- `erase(p)`: Remove and destroy the element of *V* at the position given by *p* and shifts the subsequent elements one position to the left.
- `erase(p, q)`: Iterate between *p* and *q*, removing and destroying all these elements and shifting subsequent elements to the left to fill the gap.
- `clear()`: Delete all these elements of *V*.

STL Containers and Algorithms

```
#include <algorithm>
```

`sort(p,q)`: Sort the elements in the range from *p* to *q* in ascending order. It is assumed that less-than operator (“<”) is defined for the base type.

`random_shuffle(p,q)`: Rearrange the elements in the range from *p* to *q* in random order.

`reverse(p,q)`: Reverse the elements in the range from *p* to *q*.

`find(p,q,e)`: Return an iterator to the first element in the range from *p* to *q* that is equal to *e*; if *e* is not found, *q* is returned.

`min_element(p,q)`: Return an iterator to the minimum element in the range from *p* to *q*.

`max_element(p,q)`: Return an iterator to the maximum element in the range from *p* to *q*.

`for_each(p,q,f)`: Apply the function *f* the elements in the range from *p* to *q*.

STL Vector
STL deque

STL List

Example Code

```
#include <cstdlib> // provides EXIT_SUCCESS
#include <iostream> // I/O definitions
#include <vector> // provides vector
#include <algorithm> // for sort, random_shuffle

using namespace std; // make std:: accessible

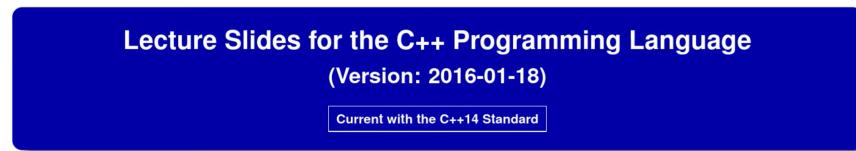
int main () {
    int a[] = {17, 12, 33, 15, 62, 45};
    vector<int> v(a, a + 6); // v: 17 12 33 15 62 45
    cout << v.size() << endl; // outputs: 6
    v.pop_back(); // v: 17 12 33 15 62
    cout << v.size() << endl; // outputs: 5
    v.push_back(19); // v: 17 12 33 15 62 19
    cout << v.front() << " " << v.back() << endl; // outputs: 17 19
    sort(v.begin(), v.begin() + 4); // v: (12 15 17 33) 62 19
    v.erase(v.end() - 4, v.end() - 2); // v: 12 15 62 19
    cout << v.size() << endl; // outputs: 4

    char b[] = {'b', 'r', 'a', 'v', 'o'};
    vector<char> w(b, b + 5); // w: b r a v o
    random_shuffle(w.begin(), w.end()); // w: o v r a b
    w.insert(w.begin(), 's'); // w: s o v r a b

    for (vector<char>::iterator p = w.begin(); p != w.end(); ++p)
        cout << *p << " "; // outputs: s o v r a b
    cout << endl;
    return EXIT_SUCCESS;
}
```

If you want to know more about iterators,

- ◆ Please watch this video



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For additional information and resources related to these lecture slides (including errata and *lecture videos* covering the material on many of these slides), please visit:
<http://www.ece.uvic.ca/~mdadams/cppbook>

If you like these lecture slides, *please show your support* by posting a review of them on Google Play:
<https://play.google.com/store/search?q=Michael%20D%20Adams%20C%2B%2B&c=books>



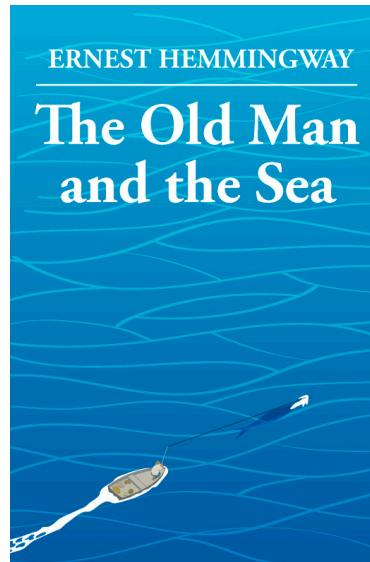
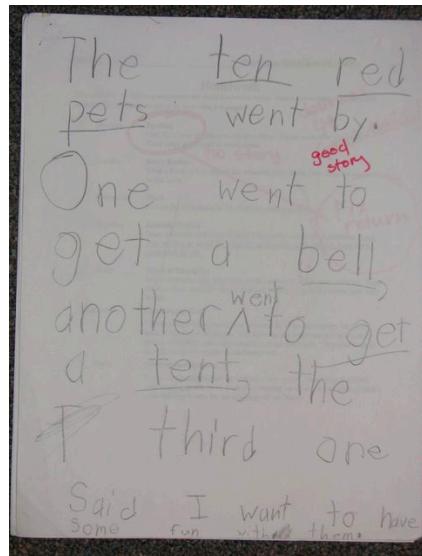
<https://www.youtube.com/watch?v=TxufBysSPK0>

- ◆ Please

- I hate to answer the question “Is this included in the exam?”

What should be your next question?

- ◆ Can I implement iterators in C++, in addition to just knowing how to use them?
 - Someone like the C++ STL designer
- ◆ Ch 6.2.3: Some level of explanation:
 - Beyond the topic of this class
- ◆ I will be happy to discuss this if you visit my office.



Sequences



Sequence ADT

- ◆ The Sequence ADT is the union of the Array List and Node List ADTs
- ◆ Elements accessed by
 - Index, or
 - Position
- ◆ Generic methods:
 - `size()`, `empty()`
- ◆ ArrayList-based methods:
 - `at(i)`, `set(i, o)`, `insert(i, o)`,
`erase(i)`
- ◆ List-based methods:
 - `begin()`, `end()`
 - `insertFront(o)`,
`insertBack(o)`
 - `eraseFront()`,
`eraseBack()`
 - `insert (p, o)`, `erase(p)`
- ◆ Bridge methods:
 - `atIndex(i)`, `indexOf(p)`

Applications of Sequences

- ◆ The Sequence ADT is a basic, general-purpose, data structure for storing an ordered collection of elements
- ◆ Direct applications:
 - Generic replacement for stack, queue, vector, or list
 - small database (e.g., address book)
- ◆ Indirect applications:
 - Building block of more complex data structures

Questions?