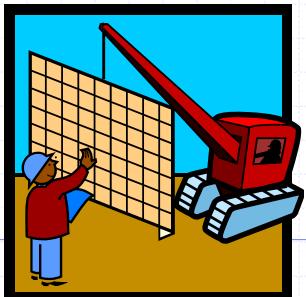


Vector, List and Sequence



1

Three ADTs

◆ Vector (also called Array List)

- Access each element using a notion of **index** in $[0, n-1]$
- Index of element e: the number of elements that are before e
- Typically we use the “**index**” (e.g., `[]`)
- A more general ADT than “array”

◆ List

- Not using an index to access, but use a node to access
- Insert a new element e before some “**position**” p
- A more general ADT than “linked list”

◆ Sequence

- Can access an element as vector and list (using both **index** and **position**)

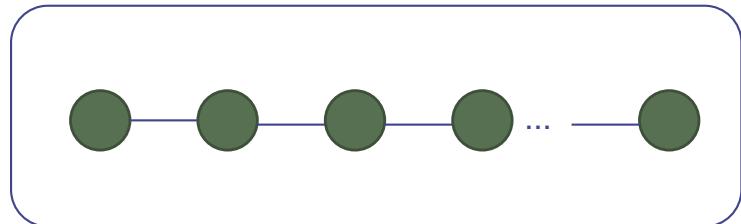
◆ (Note) Can implement the above ADTs using various ways

- array, singly linked list, doubly linked list, circular linked list

3

Overview and Reading

- ◆ Reading: Chapters: 6.1, 6.2, and 6.3
- ◆ A data structure that stores n elements in a linear order
 - Called list or sequence
- ◆ Didn’t we learn “array” and “linked list”?
 - We are talking about more abstract ADTs than them



2

Vectors (or Array Lists)



4

The ArrayList ADT

□ The Vector or ArrayList

ADT extends the notion of array by storing a sequence of objects

□ An element can be accessed, inserted or removed by specifying its index (number of elements preceding it)

□ An exception is thrown if an incorrect index is given (e.g., a negative index)

◆ Main methods:

- **at(integer i):** returns the element at index i without removing it
- **set(integer i, object o):** replace the element at index i with o
- **insert(integer i, object o):** insert a new element o to have index i
- **erase(integer i):** removes element at index i

◆ Additional methods:

- **size()**
- **empty()**

5

Applications of Array Lists

◆ Direct applications

- Sorted collection of objects (elementary database)

◆ Indirect applications

- Auxiliary data structure for algorithms
- Component of other data structures

◆ Basically, every place where you can use “array”.

6

Array-based Implementation of Vector

◆ Use an array A of size N

◆ A variable n keeps track of the size of the array list (number of elements stored)

◆ Operation $at(i)$ is implemented in $O(1)$ time by returning $A[i]$

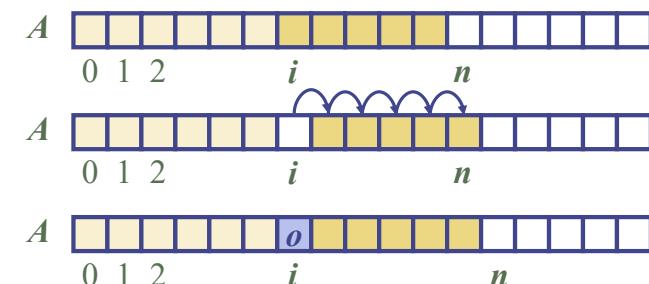
◆ Operation $set(i, o)$ is implemented in $O(1)$ time by performing $A[i] = o$



7

Insertion

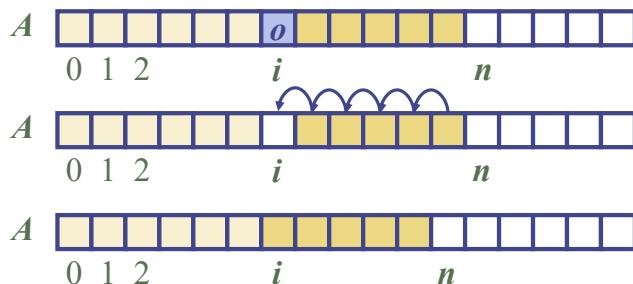
- ◆ In operation $insert(i, o)$, we need to make room for the new element by shifting forward the $n - i$ elements $A[i], \dots, A[n - 1]$
- ◆ In the worst case ($i = 0$), this takes $O(n)$ time



8

Element Removal

- ◆ In operation $\text{erase}(i)$, we need to fill the hole left by the removed element by shifting backward the $n - i - 1$ elements $A[i + 1], \dots, A[n - 1]$
- ◆ In the worst case ($i = 0$), this takes $O(n)$ time



9

Performance

- ◆ In the array-based implementation of an array list:
 - The space used by the data structure is $O(n)$
 - size , empty , at and set run in $O(1)$ time
 - insert and erase run in $O(n)$ time in worst case
- ◆ If we use the array in a circular fashion, operations $\text{insert}(0, x)$ and $\text{erase}(0, x)$ run in $O(1)$ time
- ◆ In an insert operation, when the array is full, instead of throwing an exception, we can replace the array with a larger one

10

Growable Array-based Array List

- In an $\text{insert}(o)$ operation (without an index), we always insert at the end
- When the array is full, we replace the array with a larger one
- How large should the new array be?
 - Incremental strategy: increase the size by a constant c
 - Doubling strategy: double the size

Algorithm $\text{insert}(o)$

```
if  $t = S.length - 1$  then
     $A \leftarrow$  new array of
        size ...
    for  $i \leftarrow 0$  to  $n-1$  do
         $A[i] \leftarrow S[i]$ 
     $S \leftarrow A$ 
     $n \leftarrow n + 1$ 
     $S[n-1] \leftarrow o$ 
```

- ◆ For size n array, “re-grow” operation requires n copies

Which is better?: Incremental or Doubling

- ◆ Comparison Method 1
 - Given the current size of $S = n$
 - Worst-case running time
 - Incremental strategy: $O(1)$
 - Doubling strategy: $O(n)$
- ◆ Are you happy?
 - Happy if your focus is really the worst-case
 - Unhappy
 - For doubling strategy, the total number of resizing array size would be small
- ◆ Can we reconsider the analysis method?

11

12

Which is better?: Incremental or Doubling

◆ Comparison Method2

- Compute the total time $T(n)$ needed to perform a series of n insert(o) operations
- Assume that we start with an empty stack represented by an array of size 1

◆ We call amortized time of an insert operation **the average time taken by an insert over the series of operations**, i.e., $T(n)/n$

- This can be a fairer comparison in some cases

◆ Amortized analysis (분할상환분석 in Wiki)

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Doubling Strategy Analysis

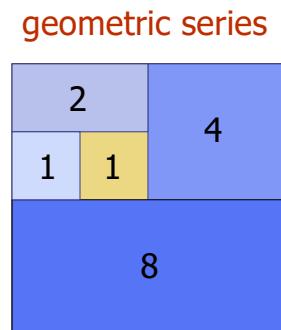
◆ We replace the old array with a new one $k = \log_2 n$ times

◆ The total time $T(n)$ of a series of n insert operations is proportional to

$$\begin{aligned}n + 1 + 2 + 4 + 8 + \dots + 2^k &= \\n + 2^{k+1} - 1 &= \\3n - 1\end{aligned}$$

◆ $T(n)$ is $O(n)$

◆ The amortized time of an insert operation is $O(1)$



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Incremental Strategy Analysis

- ◆ We replace the old array with a new one $k = n/c$ times
- ◆ The total time $T(n)$ of a series of n insert operations is proportional to

$$\begin{aligned}n + c + 2c + 3c + 4c + \dots + kc &= \\n + c(1 + 2 + 3 + \dots + k) &= \\n + ck(k + 1)/2\end{aligned}$$

- ◆ Since c is a constant, $T(n)$ is $O(n + k^2)$, i.e., $O(n^2)$
- ◆ The amortized time of an insert operation is $O(n)$

14

Professor, I have a question

◆ In “computing spans”, why didn’t you do amortized analysis?

◆ Can we do it?

◆ Is it meaningful?

◆ Think about this!

- I am ready to discuss if you get your version of answer ready.

Algorithm	<i>spans2(X, n)</i>	#
$S \leftarrow$	new array of n integers	n
$A \leftarrow$	new empty stack	1
for	$i \leftarrow 0$ to $n - 1$ do	n
while	$(\neg A.\text{empty}) \wedge X[A.\text{top}()] \leq X[i]$ do	n
	$A.pop()$	n
if	$A.\text{empty}()$ then	n
	$S[i] \leftarrow i + 1$	n
else		
	$S[i] \leftarrow i - A.\text{top}()$	n
	$A.push(i)$	n
return	S	1

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Vectors in C++ STL

```
#include <vector>           // provides definition of vector
using std::vector;          // make vector accessible

vector<int> myVector(100);  // a vector with 100 integers

vector(n): Construct a vector with space for  $n$  elements; if no argument is given, create an empty vector.
size(): Return the number of elements in  $V$ .
empty(): Return true if  $V$  is empty and false otherwise.
resize(n): Resize  $V$ , so that it has space for  $n$  elements.
reserve(n): Request that the allocated storage space be large enough to hold  $n$  elements.
operator[i]: Return a reference to the  $i$ th element of  $V$ .
at(i): Same as  $V[i]$ , but throw an out_of_range exception if  $i$  is out of bounds, that is, if  $i < 0$  or  $i \geq V.size()$ .
front(): Return a reference to the first element of  $V$ .
back(): Return a reference to the last element of  $V$ .
push_back(e): Append a copy of the element  $e$  to the end of  $V$ , thus increasing its size by one.
pop_back(): Remove the last element of  $V$ , thus reducing its size by one.
```

Difference between resize() and reserve()?

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Last Class

◆ Vector and List

◆ Vector

- Access elements by “index”
- Incremental vs. Doubling Strategy
 - ◆ Amortized analysis

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Logistics

◆ First programming assignment

- Deadline: Sep, 19th

◆ Problem Solving Homework

- Deadline: Oct, 1st

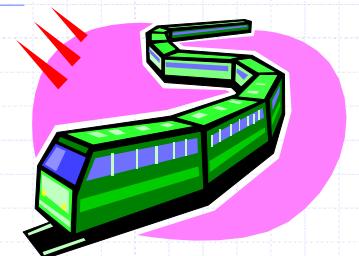
◆ You should keep reading the textbook

◆ Sep 24th, 26th : No class

- Thanksgiving

18

Lists



20

(Node) List ADT

- ◆ The Node List ADT models a sequence of positions storing arbitrary objects
- ◆ It establishes a before/after relation between positions
- ◆ Generic methods:
 - `size()`, `empty()`

- Iterators:
 - `begin()`, `end()`
- Update methods:
 - `insertFront(e)`,
`insertBack(e)`
 - `removeFront()`,
`removeBack()`
- Iterator-based update:
 - `insert(p, e)`
 - `remove(p)`

(Question) No method for accessing a specific node?
We will talk about this later

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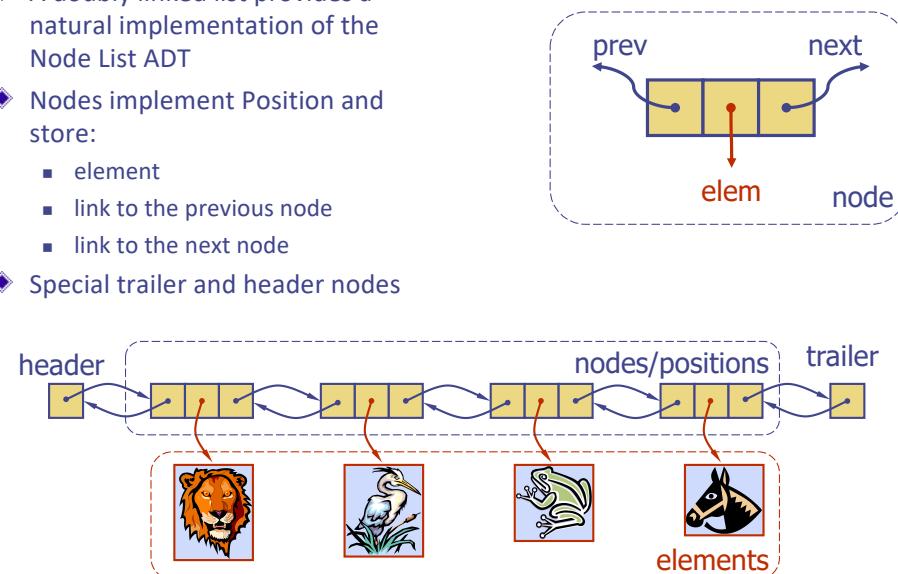
Performance

- ◆ In the implementation of the List ADT by means of a doubly linked list
 - The space used by a list with n elements is $O(n)$
 - The space used by each position of the list is $O(1)$
 - All the operations of the List ADT run in $O(1)$ time

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Implementation based on DLL (covered this)

- ◆ A doubly linked list provides a natural implementation of the Node List ADT
- ◆ Nodes implement Position and store:
 - element
 - link to the previous node
 - link to the next node
- ◆ Special trailer and header nodes



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Lists in C++ STL

```
#include <list>
using std::list;
list<float> myList; // make list accessible
// an empty list of floats
```

- `list(n)`: Construct a list with n elements; if no argument list is given, an empty list is created.
- `size()`: Return the number of elements in L .
- `empty()`: Return true if L is empty and false otherwise.
- `front()`: Return a reference to the first element of L .
- `back()`: Return a reference to the last element of L .
- `push_front(e)`: Insert a copy of e at the beginning of L .
- `push_back(e)`: Insert a copy of e at the end of L .
- `pop_front()`: Remove the first element of L .
- `pop_back()`: Remove the last element of L .

24

Containers, Iterators, and Generic algorithms



25

This is how we can do in C++

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

It is cool. But why is it cool?

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Sorting: Vector and List

- ◆ I want to find “yiyung” in Vector or List objects

```
vector<string> V(100);
list<string> L(100);
// some data insertion to V and L
```

```
//Design 1: different function
find_vector(&V);
find_list(&L);
```

```
//Design 2: function overloading
find(&V);
find(&L);
```

Do you like these? Why? Why not?

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Mysterious things

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

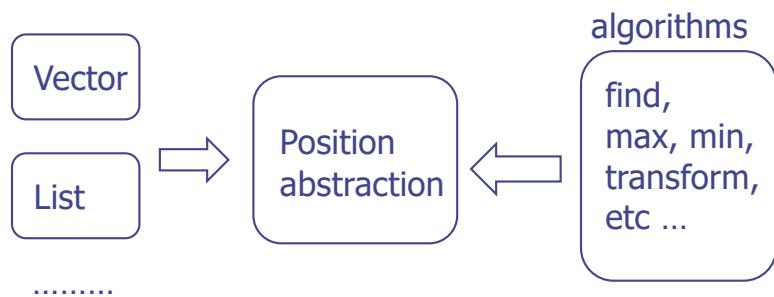
    return 0;
}
```

iterator? Looks like a “position” of vector or list. Hmm.....

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Goal and Design Challenge

- ◆ Lots of data structures (or classes in C++) that can contain various types of elements
 - “Container”
 - Examples: Vector, List, deque, set, map, etc ...



- ◆ How are you going to design this concept?

- Again, from C++ STL designer's perspective

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Containers and Iterators in C++

- ◆ An **iterator** abstracts the process of scanning through a collection of elements
- ◆ A **container** is an abstract data structure that supports element access through iterators
 - Data structures that support iterators
 - Examples include Stack, Queue, Vector, List
 - **begin()**: returns an iterator to the first element
 - **end()**: return an iterator to an imaginary position just after the last element
- ◆ An iterator behaves like a pointer to an element
 - ***p**: returns the element referenced by this iterator
 - **++p**: advances to the next element
- ◆ Extends the concept of **position** by adding a traversal capability

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Position ADT

- ◆ The **Position** ADT models the notion of place within a data structure where a single object is stored
- ◆ It gives a unified view of diverse ways of storing data, such as
 - a cell of an array
 - a node of a linked list
- ◆ “A” method of accessing the element at position **p**:
 - object **p.element()**: returns the element at position
 - In C++ it is convenient to implement this as ***p**
 - **Operator overloading**

- ◆ Implemented as “iterator” in C++

30

Example codes again

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it = find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;
    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it = find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;
    return 0;
}
```

Ah-ha, it's an iterator!

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Various Iterators

- ◆ **(standard) iterator:** allows read-write access to elements
- ◆ **const iterator:** provides read-only access to elements
- ◆ **bidirectional iterator:** supports both ++p and -p
- ◆ **random-access iterator:** supports both p+i and p-i

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Back to Iterator: STL Iterator-based Functions

`vector(p, q):` Construct a vector by iterating between p and q , copying each of these elements into the new vector.

`assign(p, q):` Delete the contents of V , and assigns its new contents by iterating between p and q and copying each of these elements into V .

`insert(p, e):` Insert a copy of e just prior to the position given by iterator p and shifts the subsequent elements one position to the right.

`erase(p):` Remove and destroy the element of V at the position given by p and shifts the subsequent elements one position to the left.

`erase(p, q):` Iterate between p and q , removing and destroying all these elements and shifting subsequent elements to the left to fill the gap.

`clear():` Delete all these elements of V .

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STL Iterators in C++

- Each STL container type C supports iterators:
 - $C::iterator$ – read/write iterator type
 - $C::const_iterator$ – read-only iterator type
 - $C.begin()$, $C.end()$ – return start/end iterators
- This iterator-based operators and methods:
 - $*p$: access current element
 - ++p , -p : advance to next/previous element
 - $C.assign(p, q)$: replace C with contents referenced by the iterator range $[p, q)$ (from p up to, but not including, q)
 - $\text{insert}(p, e)$: insert e prior to position p
 - $\text{erase}(p)$: remove element at position p
 - $\text{erase}(p, q)$: remove elements in the iterator range $[p, q)$

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STL Containers and Algorithms

`#include <algorithm>`

`sort(p, q):` Sort the elements in the range from p to q in ascending order. It is assumed that less-than operator (“ $<$ ”) is defined for the base type.

`random_shuffle(p, q):` Rearrange the elements in the range from p to q in random order.

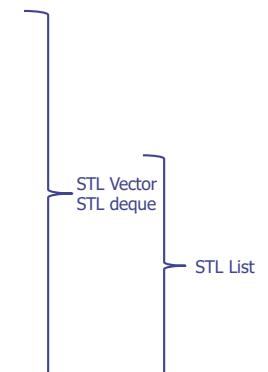
`reverse(p, q):` Reverse the elements in the range from p to q .

`find(p, q, e):` Return an iterator to the first element in the range from p to q that is equal to e ; if e is not found, q is returned.

`min_element(p, q):` Return an iterator to the minimum element in the range from p to q .

`max_element(p, q):` Return an iterator to the maximum element in the range from p to q .

`for_each(p, q, f):` Apply the function f to the elements in the range from p to q .



Example Code

```
#include <cstdlib>
#include <iostream>
#include <vector>
#include <algorithm>

using namespace std;

int main () {
    int a[] = {17, 12, 33, 15, 62, 45};
    vector<int> v(a, a + 6);
    cout << v.size() << endl;
    v.pop_back();
    cout << v.size() << endl;
    v.push_back(19);
    cout << v.front() << " " << v.back() << endl; // outputs: 17 19
    sort(v.begin(), v.begin() + 4); // v: (12 15 17 33) 62 19
    v.erase(v.end() - 4, v.end() - 2); // v: 12 15 62 19
    cout << v.size() << endl; // outputs: 4

    char b[] = {'b', 'r', 'a', 'v', 'o'};
    vector<char> w(b, b + 5);
    random_shuffle(w.begin(), w.end());
    w.insert(w.begin(), 's');
    for (vector<char>::iterator p = w.begin(); p != w.end(); ++p)
        cout << *p << " "; // outputs: s o v r a b
    cout << endl;
    return EXIT_SUCCESS;
}
```

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If you want to know more about iterators,

◆ Please watch this video

Lecture Slides for the C++ Programming Language
(Version: 2016-01-18)
Current with the C++14 Standard

Michael D. Adams
Department of Electrical and Computer Engineering
University of Victoria
Victoria, British Columbia, Canada
University of Victoria

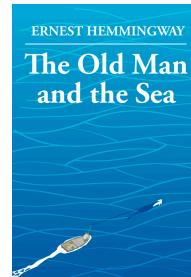
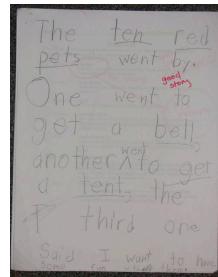
For additional information and resources related to these lecture slides (including errata and lecture videos covering the material on many of these slides), please visit:
<http://www.ece.uvic.ca/~mdadams/cppbook>

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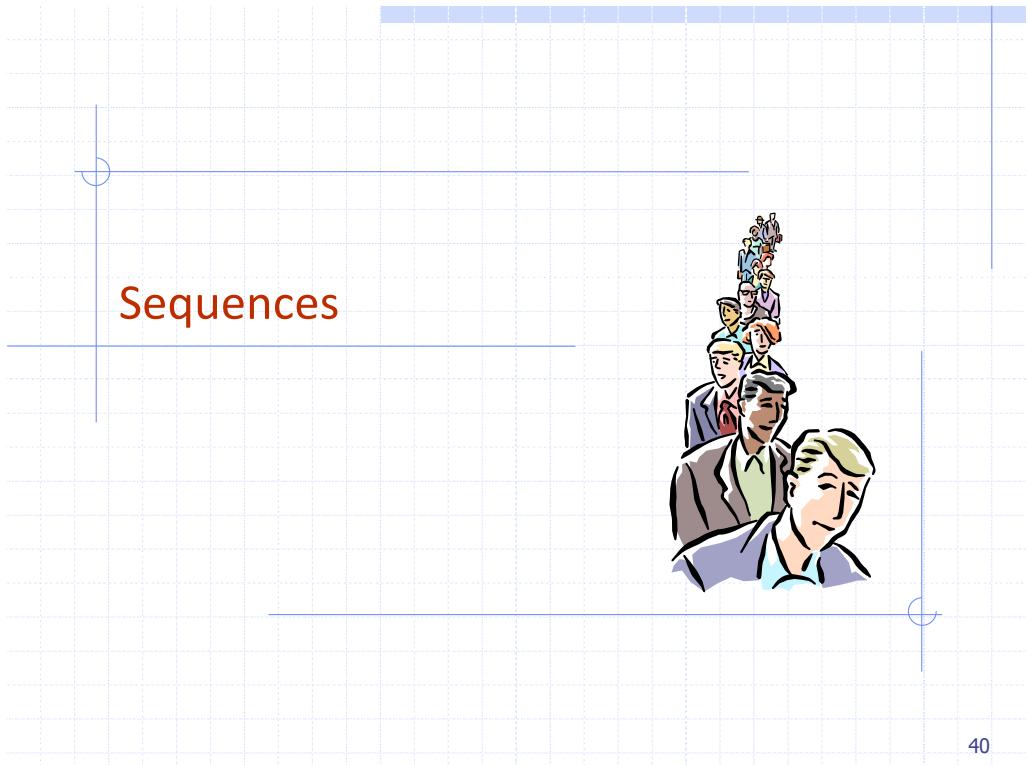
What should be your next question?

- ◆ Can I implement iterators in C++, in addition to just knowing how to use them?
 - Someone like the C++ STL designer
- ◆ Ch 6.2.3: Some level of explanation:
 - Beyond the topic of this class
- ◆ I will be happy to discuss this if you visit my office.



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Sequences



40

- ◆ The Sequence ADT is the union of the Array List and Node List ADTs
- ◆ Elements accessed by
 - Index, or
 - Position
- ◆ Generic methods:
 - `size()`, `empty()`
- ◆ ArrayList-based methods:
 - `at(i)`, `set(i, o)`, `insert(i, o)`,
`erase(i)`

- ◆ List-based methods:
 - `begin()`, `end()`
 - `insertFront(o)`,
`insertBack(o)`
 - `eraseFront()`,
`eraseBack()`
 - `insert(p, o)`, `erase(p)`
- ◆ Bridge methods:
 - `atIndex(i)`, `indexOf(p)`

- ◆ The Sequence ADT is a basic, general-purpose, data structure for storing an ordered collection of elements
- ◆ Direct applications:
 - Generic replacement for stack, queue, vector, or list
 - small database (e.g., address book)
- ◆ Indirect applications:
 - Building block of more complex data structures

Questions?