

## Electronics 101:

### Switches:

Consider a simple circuit :

Picture



When no power on 1,2 floats

⇒ We see noise when switch is open

When switch is open the input goes high

- When switch is closed we get a direct connect between  $V_{cc}$  & Gnd
  - Short, smoke
    - Too much current from V to ground

R1 limits flow to a small amount

- A pull-up resistor

(Resistant)  $R = (\text{Voltage}) V = SV$

### LEDs:

### Buttons:

- Fundamental form of input
- usually tied to general purpose I/O
  - sometimes on an ext interrupt
- we must sample the i/p to detect presses
  - is the button up or down?

- keypads are little more complex:



- A key press connect 2 wires
  - pressing A connects  $I_3$  &  $I_6$
  - We activate  $I_5$  thru  $I_8$  one at a time
  - On each, we sample  $I_1$  thru  $I_4$ 
    - o An active input indicates a connection => press
- how often do we need to sample?
- (Constraint)how fast can a human press/release?
  - sampling 100x a second is more than enough
  - 10x a sec is common
  - we have to worry about microcontroller loading issues
  - spending too much cycles on waiting user input
- problems :

- (a) when a button is pressed, the signal goes:



- (b) this means we can't use a single sample to decide on state

Note: what if it's an interrupt?

(c) we also need to stop declaring a transition.

Hysteresis : Solution to (holding button vs quick clicking button multiple times):

- We accumulate state to decide when something is on or off

- if we sample with on, increment

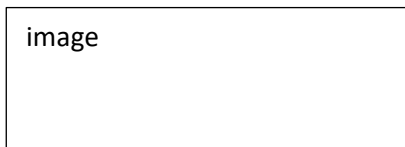
- if we sample with off, decrement

- we then define on/off values and min/max values

- avoid runaway counting when stable.

e.g. on=10 , off=4, max=12, min=0

- init to 0 & off.



- we tune our settings to ensure that we handle noise

- based on the h/w

- experimentation

Note:

- (1) We can use this to filter any noisy i/p(input)
  - a. With buttons, we call this debouncing
- (2) Easy to implement for multiple inputs with a table
  - a. Stores our constants & accumulator
  - b. Sampling simply iterates over the table

- c. Easy to change and add inputs

Ring:



Button issues:

- 1) do we response to a keypress or a keyrelease?
  - what do people expect vs what we can actually do
- 2) interrupts
  - if we have an interrupt that fires whenever we see a level, how do we actually deal with each interrupt?
    - stay in the ISR until release (against our design prinicle , simple)
    - disable that interrupt
      - when do we re-enable?

Solution: trigger interrupt on edges

- o Check pin within ISR to determine state
- 3) what about key repeats?
  - on for x milliseconds, trigger another key press
- 4) delays between key presses
  - filtering hardware or software can incur a delay before triggering an on.

Result:

- Fast keypresses are ignored
  - But if we shorten the debounce, we will get false positive due to the fact that noise getting thru
- 4.5) lots of hardware filtering will give us occasional bounces
  - need to determine how you development board behave
- 5) multiple key presses
  - what if we want to detect:
    - 1) button "A" is pressed , do i
    - 2) button "B" is pressed, do j
    - 3) button "A" then "B" is pressed, do k
    - 4) button "A" and "B" are pressed, do l
  - we must measure how closely spaced the presses are:
    - time space to declare separate

- time space to declare consecutive
- time space to declare together
- say we're sampling 10times/s
- people won't actually press at exactly the same time.
- we have to assume there is always a gap

Solution: need a state machine

### Event processing

- debouncing is low-level state \*\*\*\* to give us clean on/off states
- buttons have semantic meaning that we must interpret (app specific)
- we use state machines to track input & trigger actions
- state machines

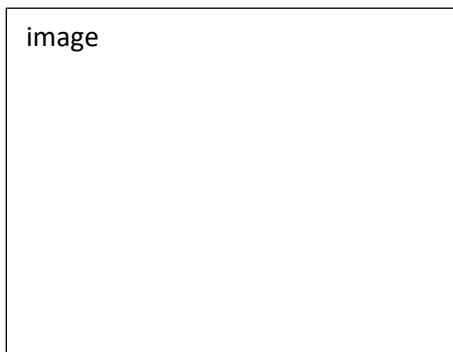
States => record previous input chain and their order

Transitions => define the next state based on the next input values

Actions => use states & transition to generate outputs

- Can be in state or on transition

e.g.



### Notes:

- our states have meaning
  - A,B,C don't cut it
- we don't have terminating states
  - run forever

- we don't have a stream of characters
  - we can sample inputs and not transition
  - happened since nothing changed
- we don't trigger an event to transition
  - we sample and see if we need to change
    - we are doing this on a heartbeat
  - with interrupts, ISR changes a value that our finite state machine code checks at next heartbeat.
- outputs of state machine : can be a sequence of operations

## Real-time FSMs

- time adds another input to FSMs(state machine)
- we can transition based on amount of time passed
  - e.g. timeout to return to start state
  - e.g. timeout to another state if no input change
  - e.g. multiple transition for the same input value input value based on how much time has passed while in current state

e.g.



- We use our heartbeat to timestamp when things happen
  - Or
- Have a count in FSM for or baseline
  - o Can be per state and/or per input

```
while(1)
{
    sleep_mode();
    sample_inputs();
    processFSM();
}
```

## FSM Implementation:

- a) Variable to track your current state
  - a. Use an enum to define all possible states
    - i. Enum X

```
{
    S1,
    S2,
    S3,
```



```
NUM_STATES
}
```

- Array of states is easy [NUM\_STATES]

b) Transitions

a. Look-up table

i. A row for each state & column for each input values

1. Entries indicate next state

– or –

2. Use function pointers to call for each state

a. Routine takes inputs & update state & generate output

while(1)

{

State = table[state](.....);

}

3. Just use a big switch

a. A case for each state

b. Could call a routine

c.