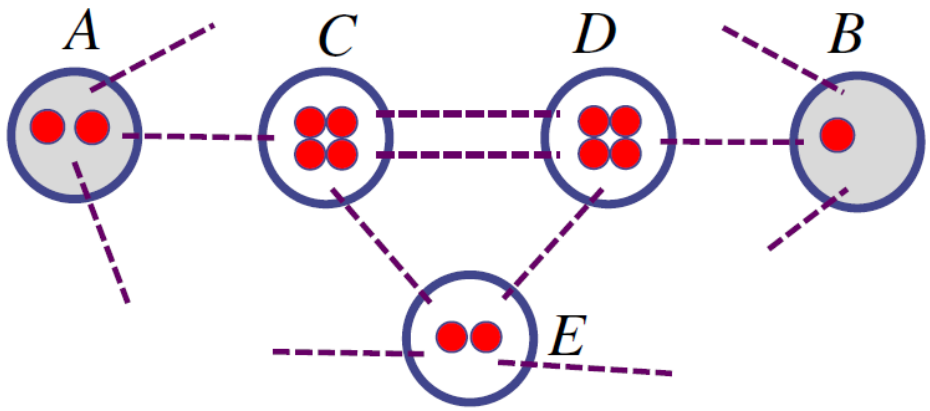
Data Structures
Programming Project #3

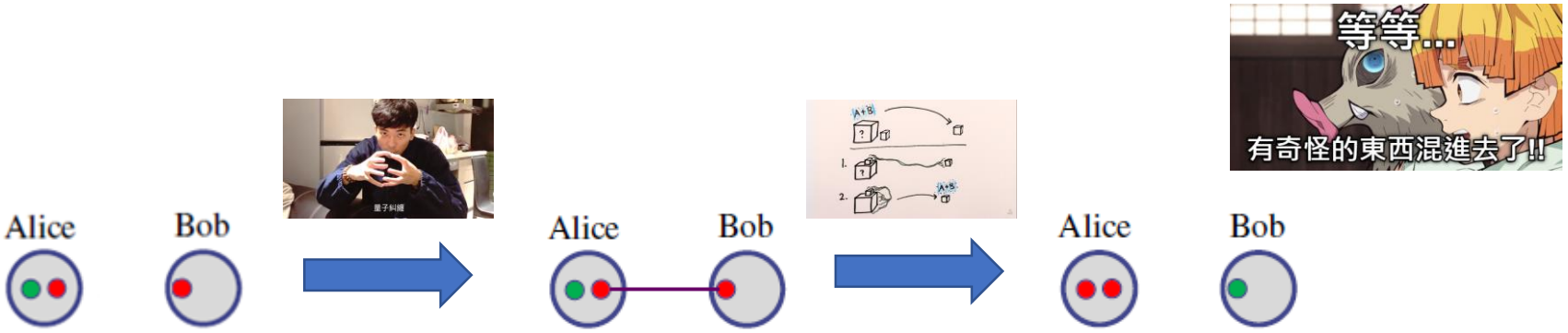
Data Transmission in Quantum Networks

- A quantum network:
- Nodes has a limited units of quantum memory
- Nodes are interconnected with a limited number of quantum channels (e.g., optical fiber)



Data Transmission in Quantum Networks

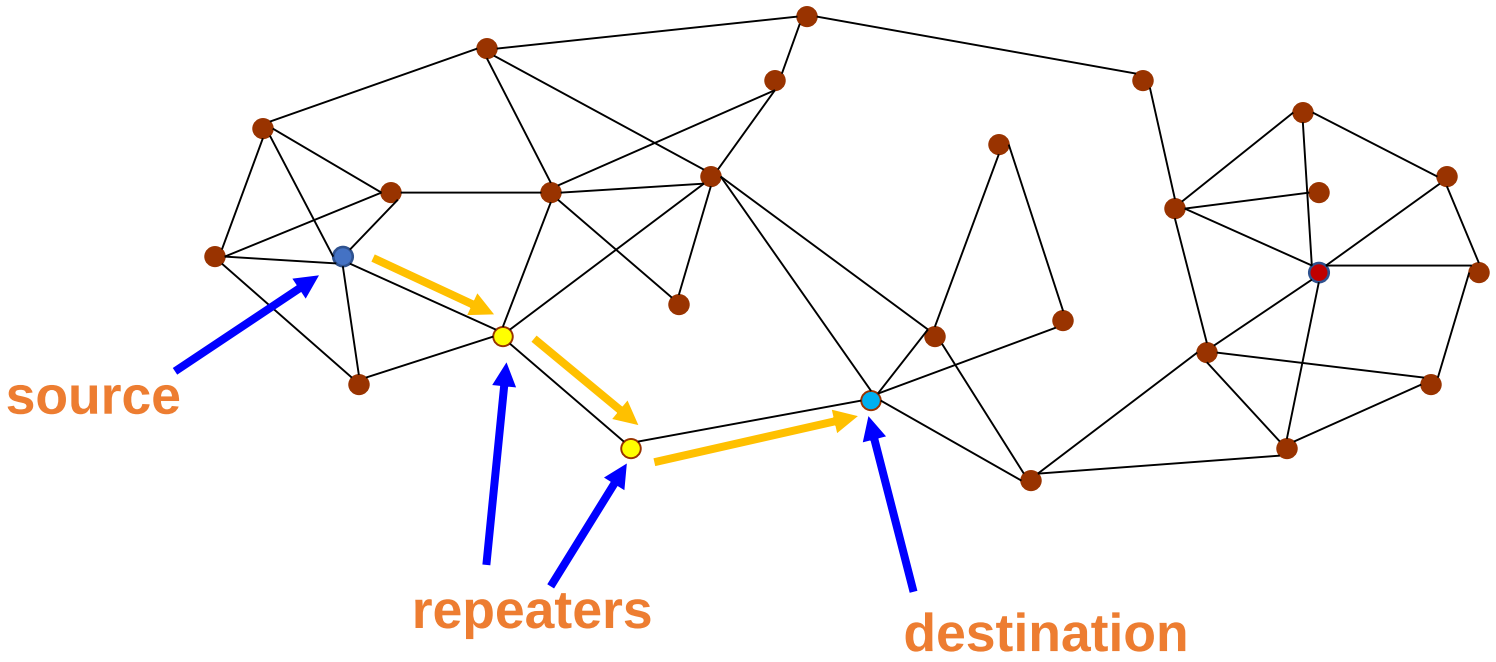
- Entangling (building an entangled link):
  Create an entangled pair between two nodes
- Precondition:
  Two nodes each with a quantum memory are interconnected with a quantum channel



https://www.popularmechanics.com/science/a25699/how-quantum-teleportation-works/

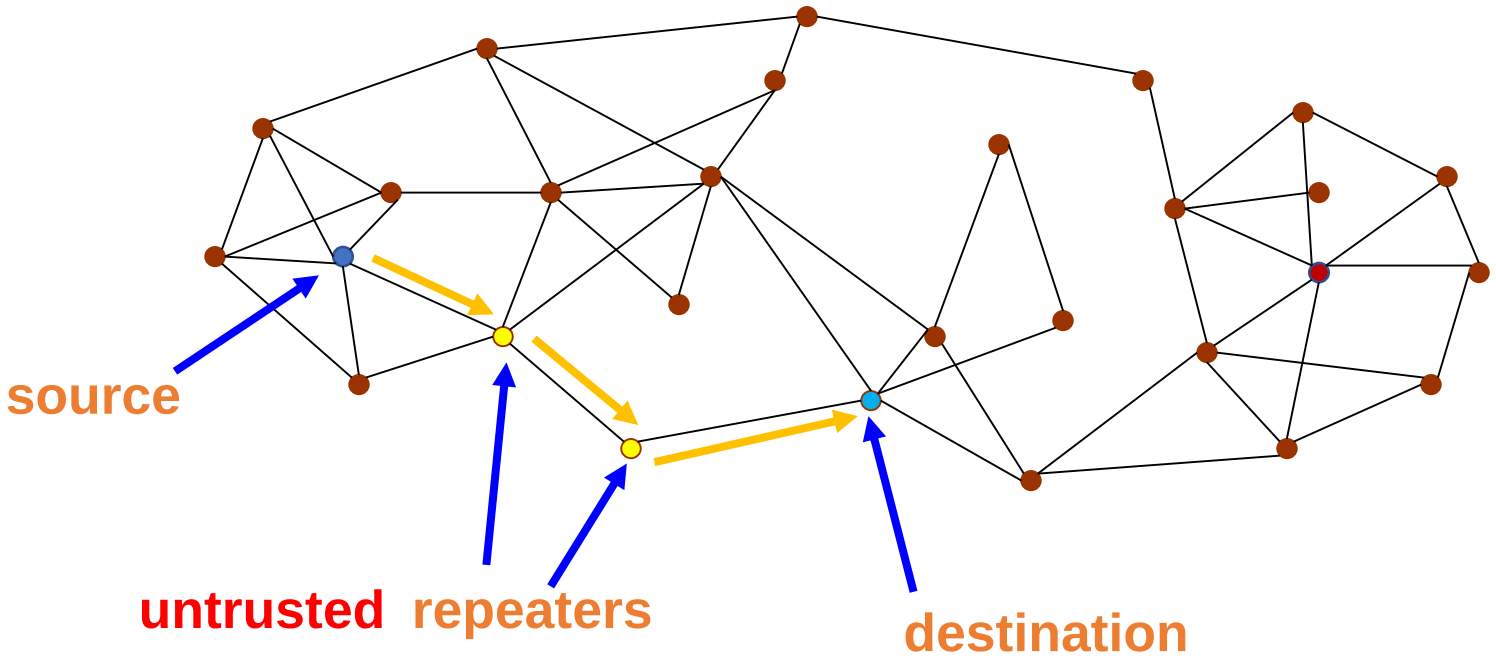
Data Transmission in Quantum Networks

- The two nodes may be distant from each other
- Classical networks:
  Repeaters use store and forward to transmit packets from a source to a destination



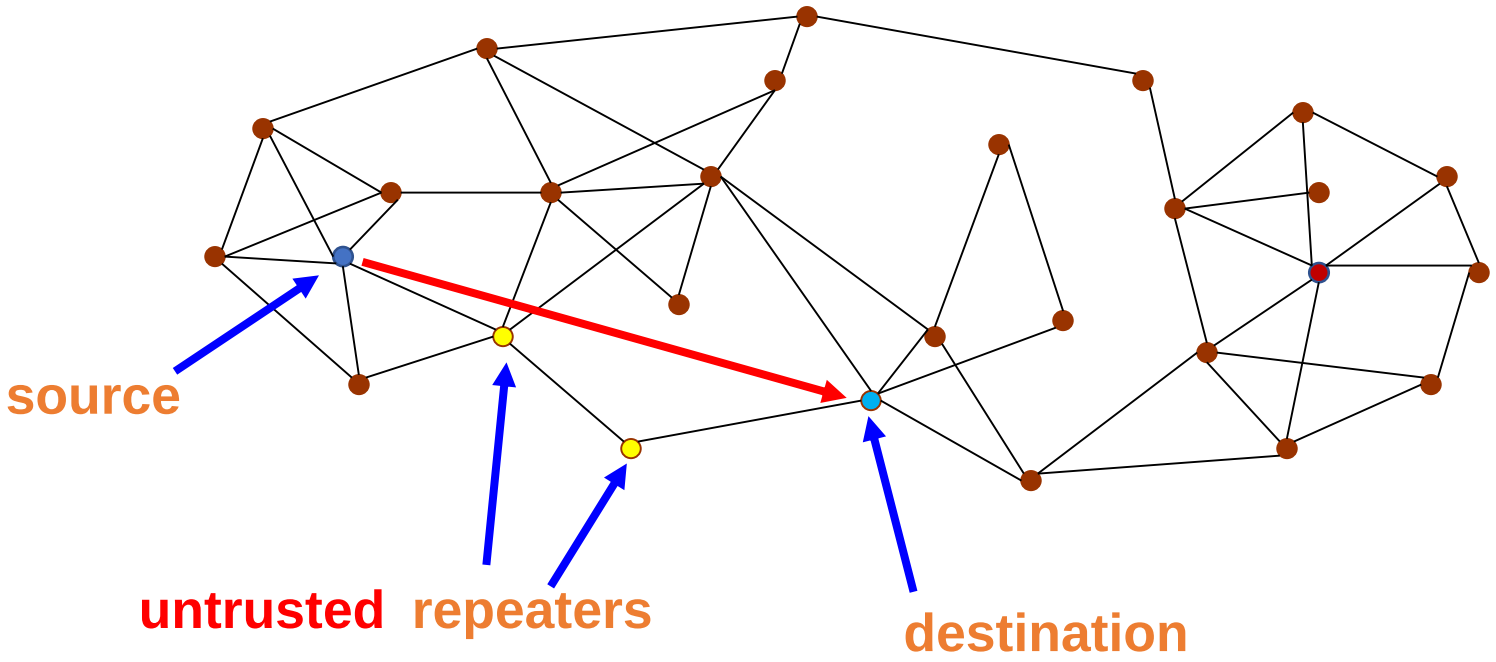
Data Transmission in Quantum Networks

- However, the data qubit may visit untrusted repeaters
- It could be destroyed, peeked at, or faked



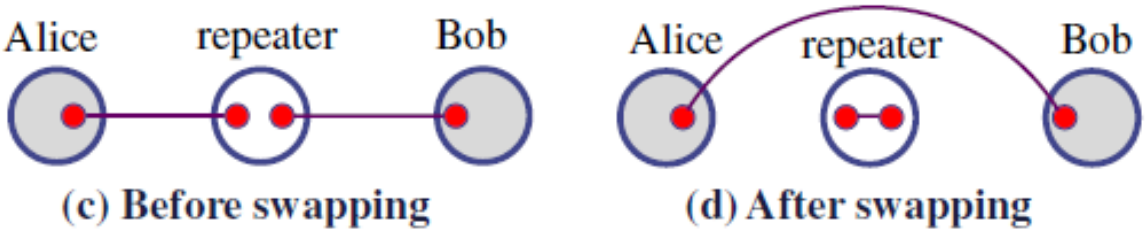
Data Transmission in Quantum Networks

- Can we send the data qubits without letting repeaters know?
- Yes, via Entanglement Swapping



Entanglement Swapping

- Alice has a data qubit for Bob
- Entangling: build the links between Alice and the repeater and between Bob and the repeater
- Swapping: build a long-distance entanglement

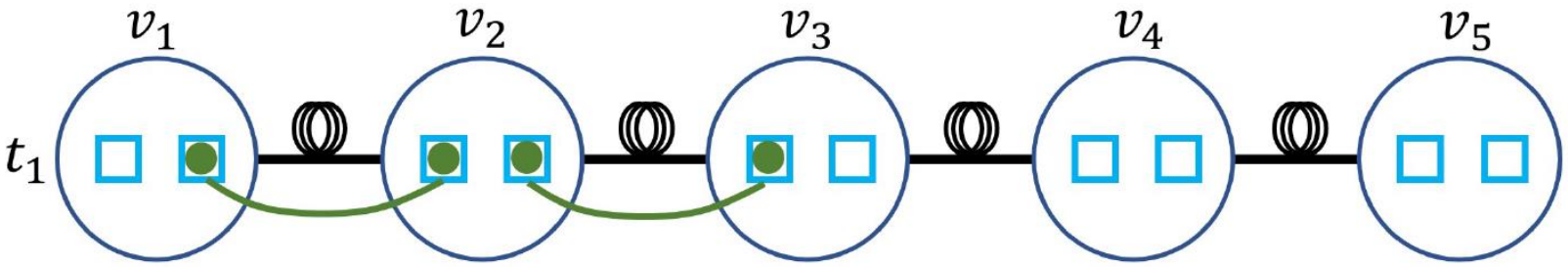


(c) Before swapping (d) After swapping

Scheduling of Entangling and Swapping

- To construct an entanglement between v_1 and v_5

- One possible scheduling is linear:
- Conduct swapping one hop by one hop from the source to the destination



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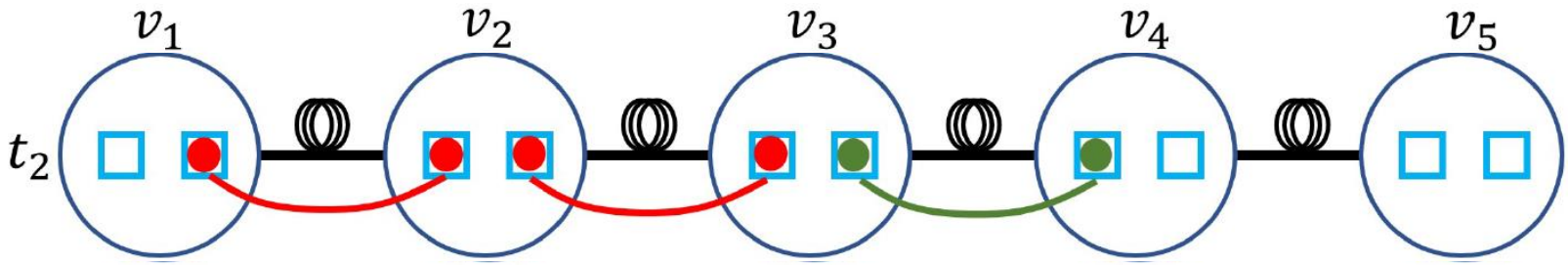

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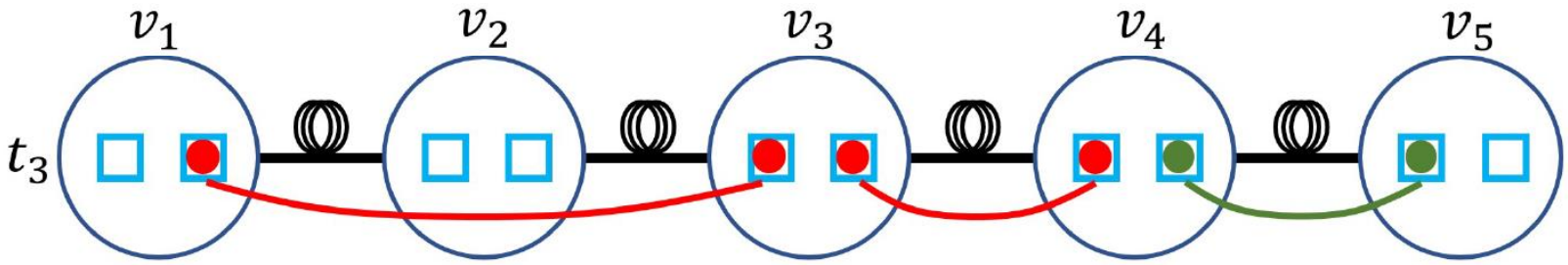

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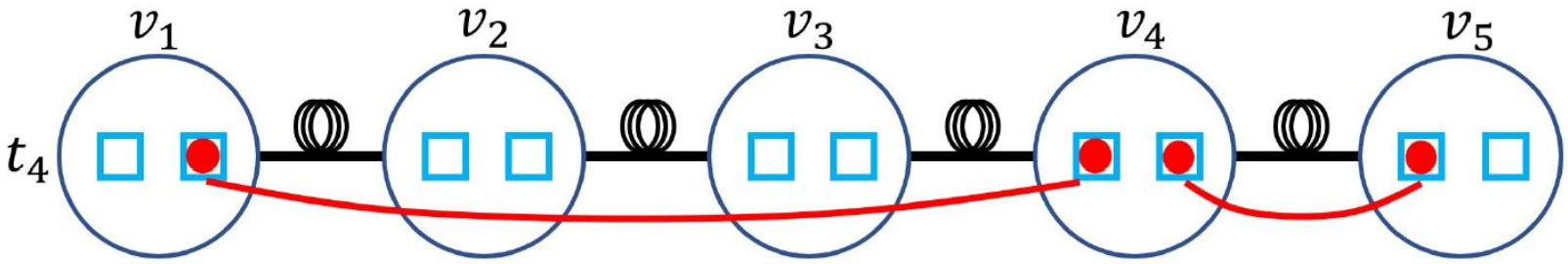

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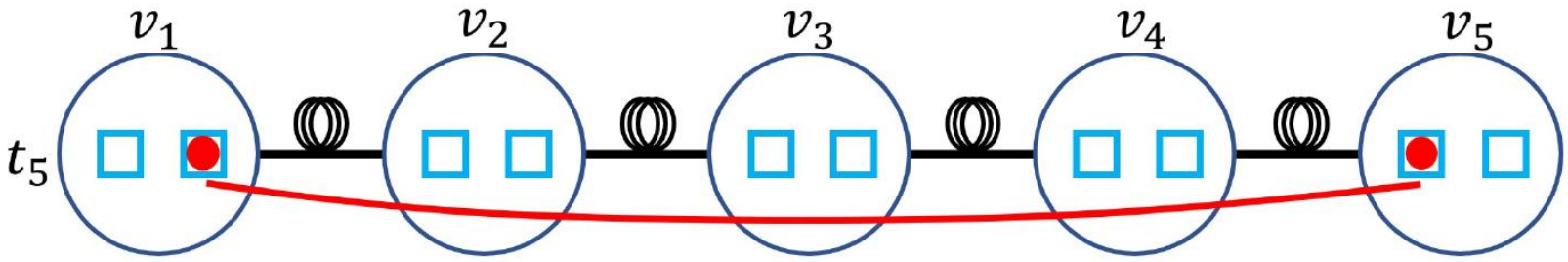
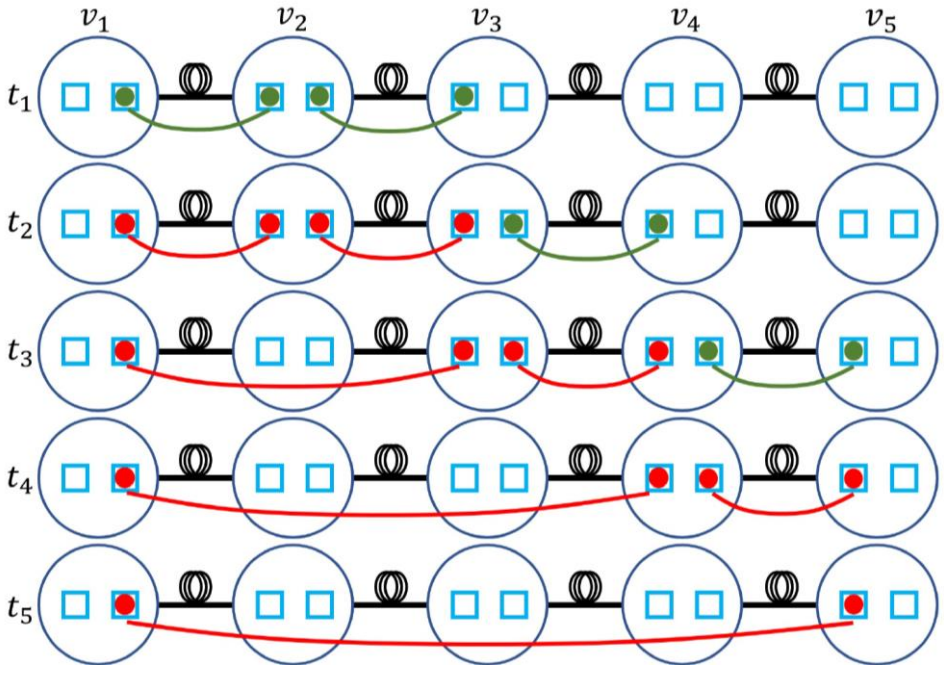
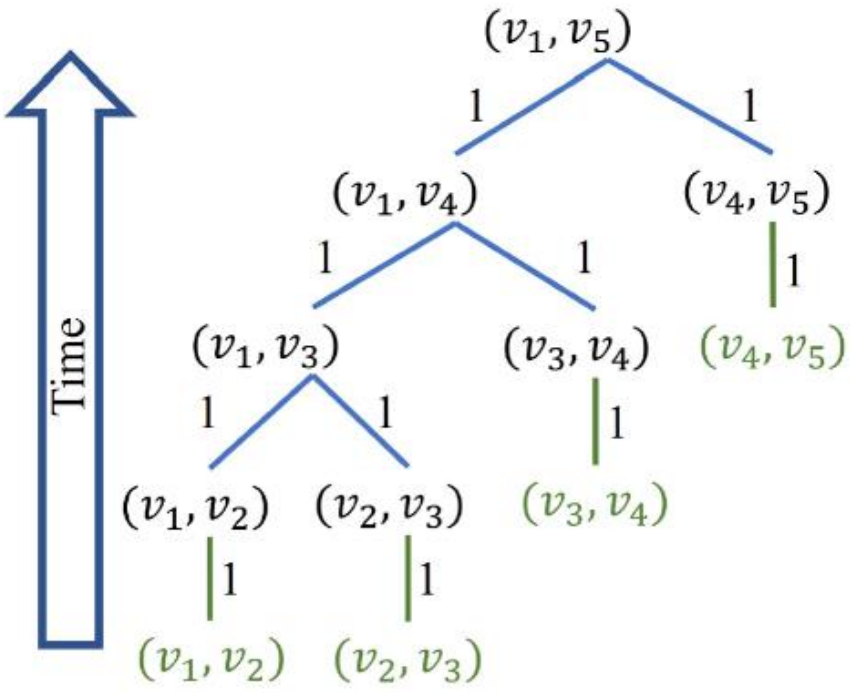


Scheduling vs Strategy Tree (Binary Tree)

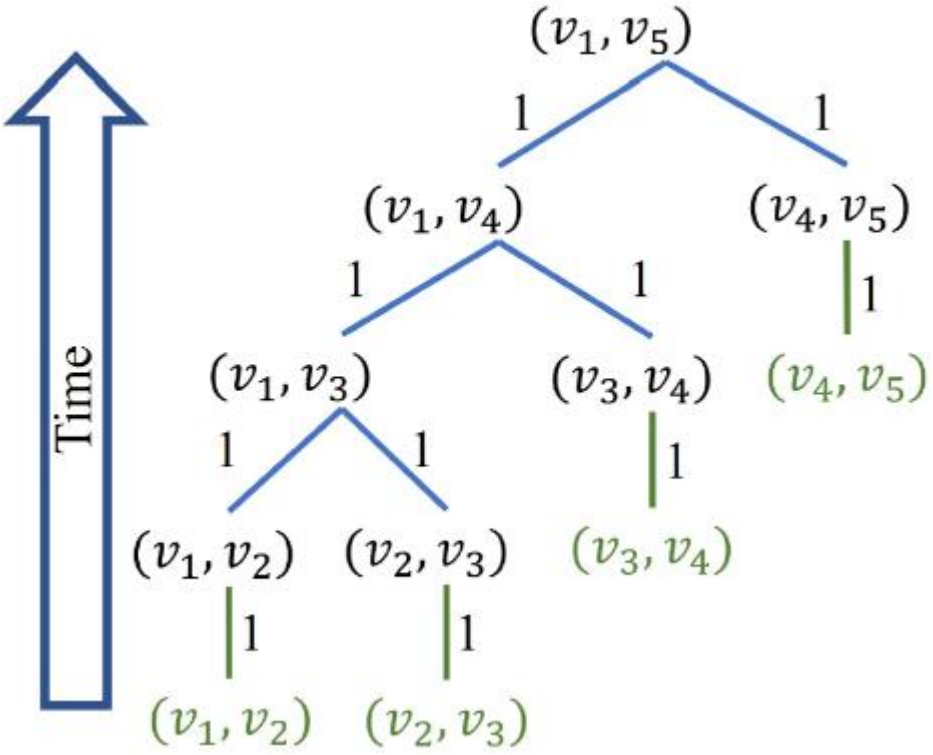
- Any scheduling can be represented by a binary tree

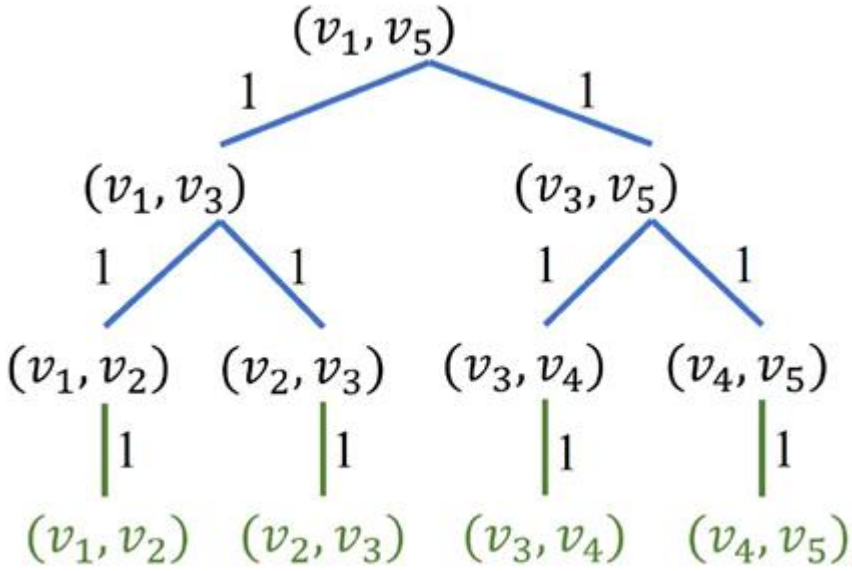




Scheduling vs Strategy Tree (Binary Tree)

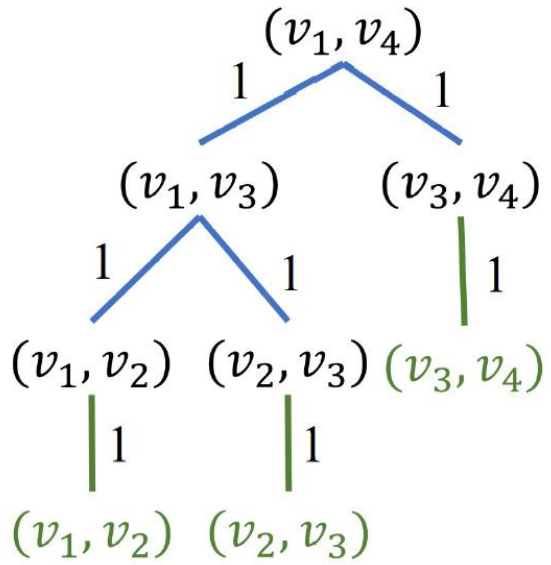
- Different schedulings lead to different binary trees

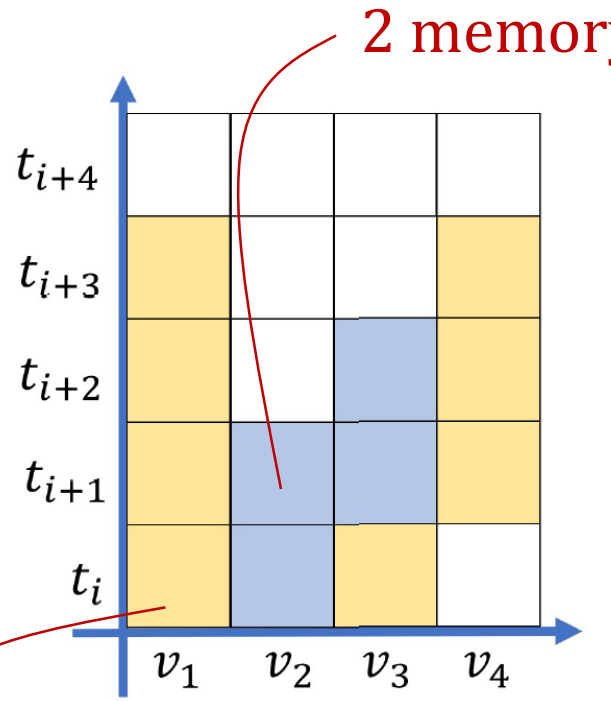




Strategy Tree vs Numerology

- Different trees lead to different numerologies





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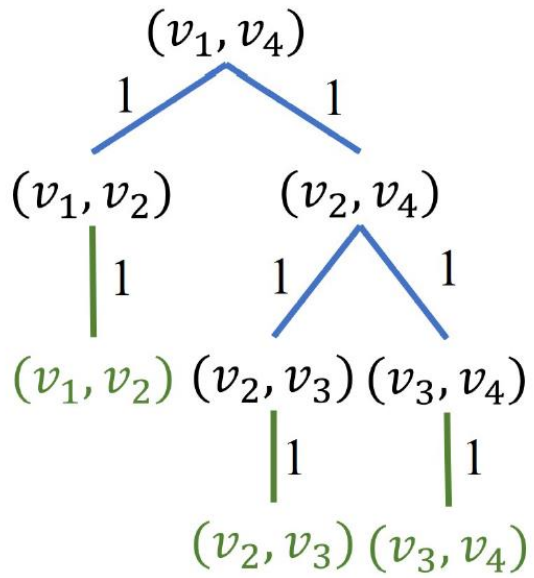

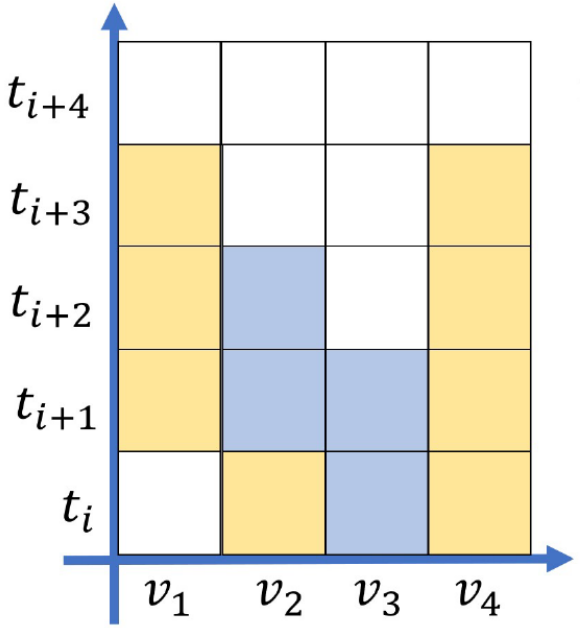

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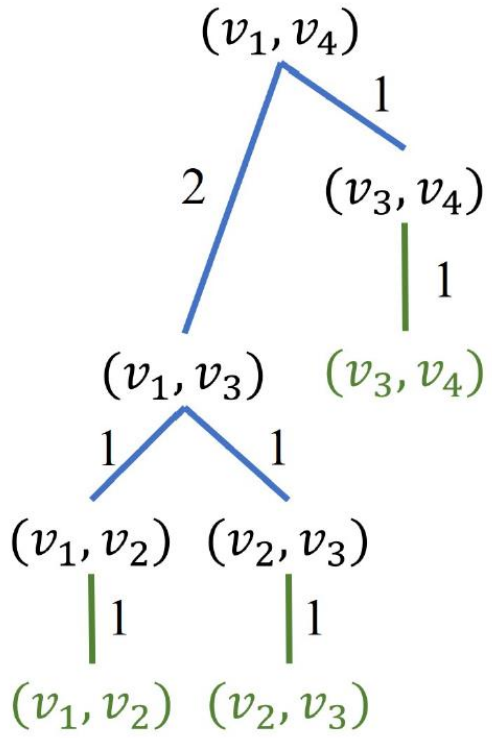
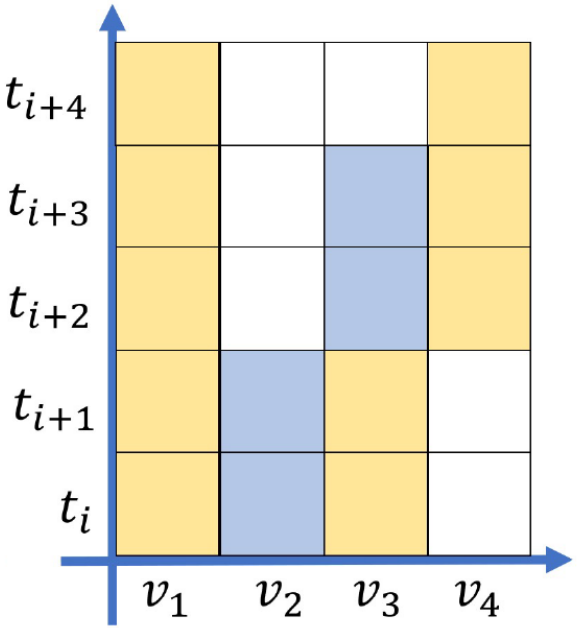
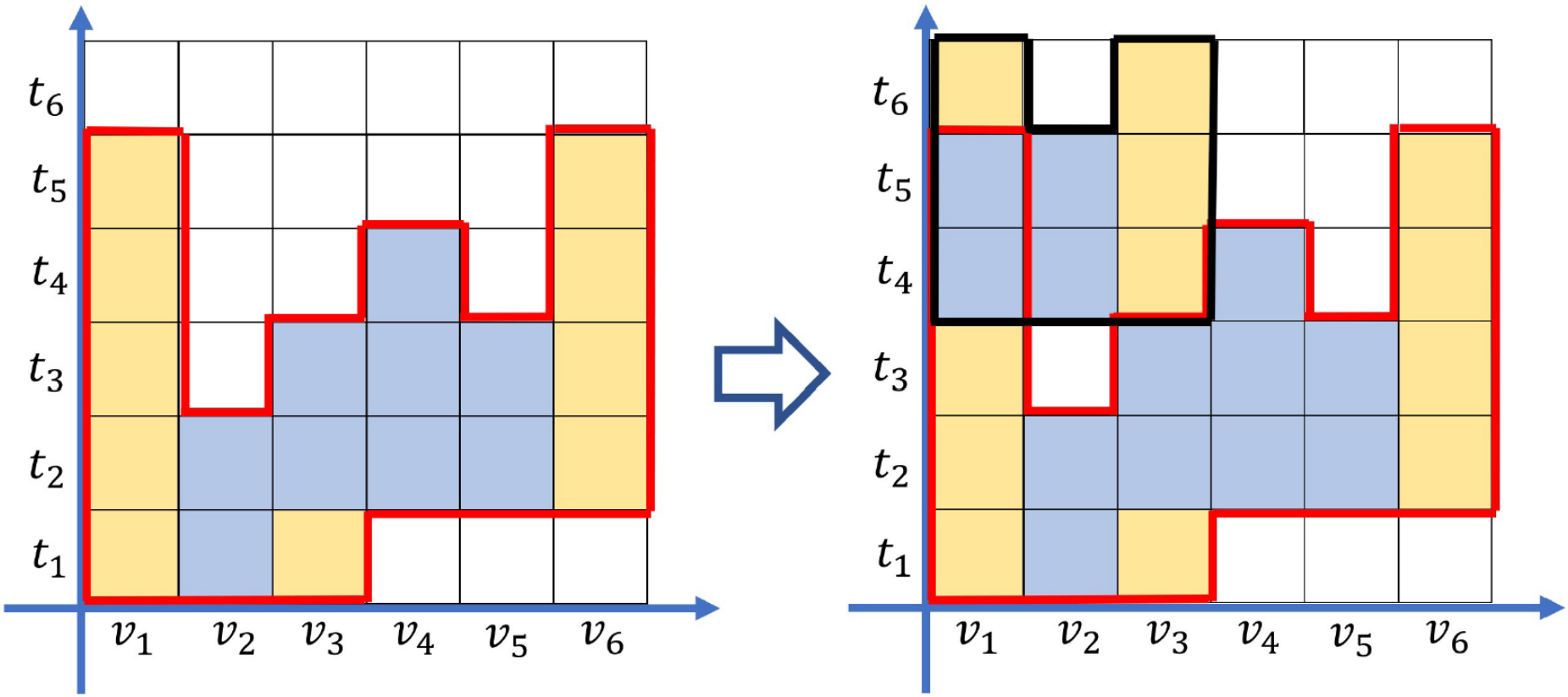




Quantum Resource Allocation Problem

- Multiple numerologies can overlap within T if the number of memory units is sufficient at that time



System Model & Problem Formulation

- Given:
- A quantum network with nodes and edges
- Each node has limited memory units
- (Assume fiber channels are always sufficient)
- Multiple source-destination (SD) requests
- Number of time slots (T: $t_1 \sim t_n$)

- Goal: maximize # accepted requests

- Constraints:
- Find a path for each accepted SD request
- Find a numerology for each accepted SD request
- Each node's memory size cannot be violated

Programming Project #3: Entanglement Routing & Numerology Selection

- Input:
  - A node-weighted network $G = (V, E)$
  - Multiple SD requests
  - A number of time slots T
- Procedure:
  - Accept or reject each SD request
  - Find a path with a numerology for accepted one
- Output:
  - The accepted SD requests
  - Their paths and numerologies
- The grade is proportional to # accepted SD requests

The Competition

- The grade is proportional to # accepted SD requests
- Basic: 60 (deadline)
  - A baseline solution (see the following pages)
  - (Meet the coding style requirements)
- Performance ranking (decided after the deadline)
  - [0%, 30%) (bottom): +0
  - [30%, 50%): + 5
  - [50%, 75%): + 10
  - [75%, 85%): + 15
  - [85%, 90%): + 20
  - [90%, 95%): + 25
  - [95%, 100%] (top): + 30
- Homework assistant (superb deadline)
  - +10

The Competition

- The grade) requests
- Basic: 60 (
  - A baselin
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- Performar
  - [0%, 30%
  - [30%, 50
  - [50%, 75
  - [75%, 85
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  - [90%, 95
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requests

eadline)

相信你們在做完作業以後

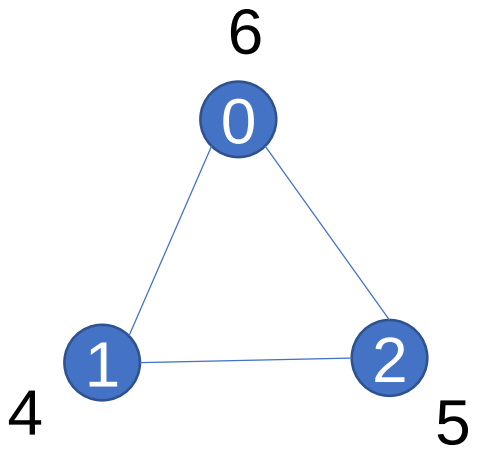
The Baseline Algorithm

- Sequentially find the minimum-hop path for each input SD request via BFS (on the original graph)
- If there is a tie, select the one that visits the node with a smaller ID first
  - 10 -> 4 -> 11 -> 15 -> 12 vs 10 -> 4 -> 9 -> 17 -> 12
  - Select 10 -> 4 -> 9 -> 17 -> 12
- Only use complete binary tree (recall that in Chap. 5) where the bottom leaf(s) is at time t_1 and side length $= 1$
- Examine whether the above path and numerology can be accommodated in the network
  - If yes, then accept the SD request
  - Otherwise, reject the SD request
- Repeat the above actions until all SD requests are examined

Input Sample: use scanf

Format:

#Nodes #Links #TimeSlots #Req
NodeID #QuantumMemories
…
LinkID LinkEnd1 LinkEnd2
…
ReqID ReqSrc ReqDst



Ex:

3 3 4 2
0 6
1 4
2 5
0 0 1
1 0 2
2 1 2
0 0 2
1 1 2

Output Sample: use printf

Format:

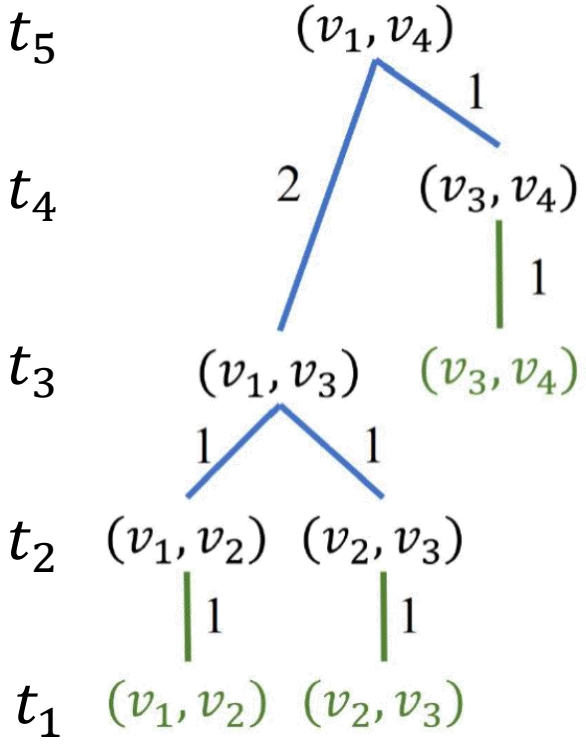
#AccReq
ReqID ReqSrc Rep1 Rep2… ReqDst
Tree Structure (detailed in the next class)
…

Output Sample: use printf

Format:

Tree Structure:

use postorder traversal




```
v1_ID v2_ID t1_time
v2_ID v3_ID t1_time
v1_ID v3_ID v1_ID v2_ID v2_ID v3_ID t3_time
v3_ID v4_ID t3_time
v1_ID v4_ID v1_ID v3_ID v3_ID v4_ID t5_time
```

Note

- Superb deadline: 12/7 Thu (adjust?)
- Deadline: 12/14 Thu (adjust?)
- Pass the test of our online judge platform
- Submit your code to E-course2
- Demonstrate your code remotely or in person with TA
- **C Source code (i.e., only .c)**
- Show a good programming style

It is worth noting that…

- The processes of entangling and swapping on the path may fail (although we don't consider it)
- In addition, strategy tree height also affects quality of entangled pair (although we don't consider it)

- For example:
- The success probability of entangling is $p = e^{-\alpha \cdot L}$, where α is a constant of the fiber material and L is the distance between the two nodes
- The success probability of swapping is q, which is related to the technique