

Problem 5.1: Syntax of first-order logic

Recall the formal syntax of first-order logic, especially the definition of terms and sentences.

Problem 5.1.1: (Taken from [2] Exercise 2.1)

Let $\mathcal{F} = \{d, f, g\}$ be a set of symbols, where d is a constant, f a function symbol with two arguments, and g a function symbol with three arguments. Which of the following expressions are terms over \mathcal{F} ? (You may assume that x, y, z are symbols for variables.)

1. $g(d, d)$
2. $f(x, g(y, z), d)$
3. $g(x, f(y, z), d)$
4. $g(x, h(y, z), d)$

Problem 5.1.2: (Taken from [2] Exercise 2.2)

Let m be a constant, f a function symbol with one argument, and S and B two predicate symbols, each with two arguments. Which of the following expressions are sentences in first-order logic? Specify a reason for failure for expressions which are not. (You may assume that x, y, z are symbols for variables.)

1. $S(m, x)$
2. $B(m, f(m))$
3. $B(B(m, x), y)$
4. $B(x, y) \Rightarrow [\exists z S(z, y)]$
5. $S(x, y) \Rightarrow S(y, f(f(x)))$

Problem 5.2: Universal and existential quantifiers

Let us abbreviate the predicate “ x is taking the Bus” by $B(x)$, and “ x has a Ticket” by $T(x)$. Suppose also that we only consider the universe of discourse where there are only three people: Alice, Bob, and Charlie.

Universal quantifiers

Problem 5.2.1: For each of our protagonists, **Table 1** lists whether they have a ticket, and whether they take the bus. Suppose that we want to formalize the sentence “All people who take the bus have a ticket”. Given the characteristics in the table below, do you think the sentence evaluates to true? Or false?

Table 1

x	$B(x)$	$T(x)$
Alice	True	True
Bob	False	True
Charlie	False	False

Problem 5.2.2: Now, consider the formula

$$\forall x \quad B(x) \wedge T(x).$$

By using the extended interpretation for universal quantifiers (in other words, the transformation to propositional logic), decide whether this formalization is true or false. Do you think this formula is a faithful formalization of “All people who take the bus have a ticket”? If not, what would be a better formula?

Existential quantifiers

Problem 5.2.3: We now look instead at the truth assignments given in **Table 2**. Suppose that we now want to formalize the sentence “Some people who take the bus have a ticket”¹. Given the characteristics in the table below, do you think the sentence evaluates to true? Or false?

Table 2

x	$B(x)$	$T(x)$
Alice	False	True
Bob	False	True
Charlie	False	False

Problem 5.2.4: Now, consider the formula

$$\exists x \quad B(x) \Rightarrow T(x).$$

By using the extended interpretation for existential quantifiers, decide whether this formalization is true or false. Do you think this formula is a faithful formalization of “Some people who take the bus have a ticket”? If not, what would be a better formula?

Problem 5.3: Formalization to First Order Logic

Problem 5.3.1: (*Taken from [2] Exercise 2.1*)

Use the predicates

$$\begin{aligned} A(x, y) : & \quad x \text{ admires } y \\ P(x) : & \quad x \text{ is a professor} \end{aligned}$$

and the constant

$$d : \quad \text{Dan}$$

to translate the following sentences into first-order logic:

1. Dan admires every professor.
2. Some professor admires Dan.
3. Dan admires himself.

¹This has to be understood as “There is **at least** one person that takes the bus, who has a ticket”.

Problem 5.3.2: (Taken from [2] Exercise 2.1)

Use the predicates

$$\begin{aligned} A(x, y) : & \quad x \text{ attended } y \\ S(x) : & \quad x \text{ is a student} \\ L(x) : & \quad x \text{ is a lecture} \end{aligned}$$

to translate the following sentences into first-order logic:

1. No student attended every lecture.
2. No lecture was attended by every student.
3. No lecture was attended by any student.

Problem 5.3.3: (Taken from [1])

Let h stand for Holmes (Sherlock Holmes) and m for Moriarty (Professor Moriarty). Let us abbreviate “ x can trap y ” by $T(x, y)$. Give the symbolic rendition of the following:

1. Holmes can trap anyone who can trap Moriarty.
2. Holmes can trap anyone whom Moriarty can trap.
3. Holmes can trap anyone who can be trapped by Moriarty.
4. If anyone can trap Moriarty, then Holmes can.
5. If everyone can trap Moriarty, then Holmes can.
6. Anyone who can trap Holmes can trap Moriarty.
7. No one can trap Holmes unless he can trap Moriarty.
8. Everyone can trap someone who cannot trap Moriarty.
9. Anyone who can trap Holmes can trap anyone whom Holmes can trap.

References

- [1] D. Gries and F. B. Schneider. *A Logical Approach to Discrete Math*. Springer-Verlag New York, Inc., New York, NY, USA, 1993.
- [2] M. Huth and M. D. Ryan. *Logic in computer science - modelling and reasoning about systems (2. ed.)*. Cambridge University Press, 2004.