Quantum state testing beyond the polarizing regime and quantum triangular discrimination

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- 1 Quantum state testing and the class QSZK
- 2 Main result: quantum state testing beyond the polarizing regime
- 3 Which parameter regime is easy for the class QSZK?
- Open problems

Statistical Difference Problem meets statistical zero-knowledge

Definition 1.1 (Statistical zero-knowledge, informal) An interactive proof protocol admits the (statistical) zero-knowledge property if verifier's view (\mathcal{P}_0) is (statistically) indistinguishable from "verifier's view" (\mathcal{P}_1) generated by a poly-time simulator.



- ▶ Public-coin is sufficient. All SZK protocols can be transformed into a form that all messages from verifier (V) to prover (P) are public coins [Okamoto'00].
- Intuitively, the views \mathcal{P}_0 and \mathcal{P}_1 can be treated as distributions p_0 and p_1 , respectively. Then statistical indistinguishablity is on $\mathrm{SD}(p_0,p_1):=\frac{1}{2}\|p_0-p_1\|_1$.

Definition 1.2 (Statistical Difference Problem, SDP) [SV03]. Given efficiently sampable (namely, using polynomial-size Boolean circuits) distributions p_0 and p_1 , decide whether $\mathrm{SD}(p_0,p_1)\geq \alpha$ or $\mathrm{SD}(p_0,p_1)\leq \beta$.

Theorem 1.3 [Sahai-Vadhan'03, Goldreich-Sahai-Vadhan'98]. Statistical Difference Problem is SZK-complete. Specifically, (α, β) -SDP is in SZK if $\alpha^2 - \beta > 0$ for constant α and β .

From quantum ℓ_1 norm to Quantum State Distinguishability Problem

An n-qubit quantum state ρ is a $2^n \times 2^n$ matrix such that $\mathrm{Tr}(\rho) = 1$ and $\rho \succeq 0$.

Classical and quantum ℓ_1 norms

- ▶ Classical: statistical distance $SD(p_0, p_1) = \frac{1}{2} \sum_{x \in \mathcal{S}} |p_0(x) p_1(x)|$.
- ▶ Q (option 1): trace distance $td(\rho_0, \rho_1) := \frac{1}{2} Tr |\rho_0 \rho_1|$.
- $\qquad \qquad \mathsf{Q} \text{ (option 2): "measured ℓ_1 distance" } \mathrm{td}^{\mathrm{meas}}(\rho_0,\rho_1) := \sup_{\mathcal{E}} \left\{ \mathrm{SD} \left(p_0^{(\mathcal{E})}, p_1^{(\mathcal{E})} \right) \right\} :$
 - \diamond measurement $\mathcal{E}=(E_1,\cdots,E_{2^n})$ such that $\sum_i E_i=I$ and $E_i\succeq 0 \ \forall i;$
 - $\diamond \ \ \text{induced distribution} \ p_k^{(\mathcal{E})} = (\operatorname{Tr}(\rho_k E_1), \cdots, \operatorname{Tr}(\rho_k E_{2^n})) \ \text{for} \ k \in \{0,1\}.$

Theorem 1.4 [Helstrom'76]. For any ρ_0 and ρ_1 , $td(\rho_0, \rho_1) = td^{meas}(\rho_0, \rho_1)$.

Definition 1.5 (Quantum State Distinguishability Problem, QSDP) [Wat02]. Given efficiently preparable (namely, using polynomial-size quantum circuits) quantum states ρ_0 and ρ_1 , decide whether $\operatorname{td}(\rho_0,\rho_1) \geq \alpha$ or $\operatorname{td}(\rho_0,\rho_1) \leq \beta$.

Theorem 1.6 [Watrous'02, Watrous'09]. QUANTUM STATE DISTINGUISHABILITY PROBLEM is QSZK-complete. Specifically, (α, β) -QSDP is in QSZK if $\alpha^2 - \beta > 0$ for const α and β .

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Polarization lemma: SZK and QSZK containments

Polarization lemma [Sahai-Vadhan'03]

There is an efficient algorithm to construct p_k' for $k \in \{0,1\}$ such that

$$\diamond$$
 Yes: $SD(p_0, p_1) \ge \alpha$;

$$\diamond \ \ \mathsf{Yes} \colon \operatorname{SD}(p_0',p_1') \geq 1 - \epsilon;$$

$$\diamond$$
 No: $SD(p_0, p_1) \leq \beta$;

$$\diamond$$
 No: $\mathrm{SD}(p_0',p_1') \leq \epsilon$;

where the dimension of p_k is 2^n .

where the dimension of p_k' is $2^n \cdot \ln(1/\epsilon).$

The construction of (p_0', p_1') is based on an appropriate composition of:

- ▶ Direct product lemma (yes instances): $(p_0^{\otimes l}, p_1^{\otimes l})$;
- $\qquad \qquad \mathsf{XOR} \ \mathsf{lemma} \ \big(no \ \mathsf{instances} \big) \colon \left(2^{-l} \sum_{i_1 \oplus \cdots \oplus i_l = 0} p_{i_1} \otimes \cdots \otimes p_{i_l}, 2^{-l} \sum_{i_1 \oplus \cdots \oplus i_l = 1} p_{i_1} \otimes \cdots \otimes p_{i_l} \right).$

By inspection, the proof techniques in [SV03] can achieve:

Theorem 2.1 [SV03,GSV98].
$$(\alpha,\beta)$$
-SDP is in SZK if $\alpha^2(n) - \beta(n) \ge 1/O(\log n)$.

- ▶ **Q1**: What about the parameter regime $\alpha^2 < \beta < \alpha$?
- ▶ **Q2:** Could we make the promise gap $\alpha^2(n) \beta(n) \ge 1/\text{poly}(n)$?

Main result on quantum state testing and QSZK

Theorem 2.3 [Berman-Degwekar-Rothblum-Vasudevan'19]. (α, β) -SDP is in SZK if $\alpha^2(n) - \beta(n) \ge 1/\mathrm{poly}(n)$. Additionally, there are two new SZK-complete problems:

- \diamond Jensen-Shannon Divergence Problem: (α, β) -JSP is in SZK if $\alpha \beta \ge 1/\text{poly}$.
- \diamond Triangular Discrimination Problem: (α, β) -TDP is in SZK if $\alpha \beta \ge 1/\text{poly}$;
- Examining existing approaches to polarization: $TDP \leftrightarrow original polarization$ lemma [SV03] and $JSP \leftrightarrow original polarization [GV99];$
- ▶ **Proof** by reductions: Entropy Difference \rightarrow JSP \rightarrow TDP \rightarrow 1/poly-SDP.

<u>Theorem 2.4.</u> (α,β) -QSDP is in QSZK if $\alpha^2 - \sqrt{2\ln 2}\beta \ge 1/\mathrm{poly}$. In addition, there are two new QSZK-complete problems:

- \diamond Quantum Jensen-Shannon Divergence Problem: $(\alpha,\beta)\text{-QJSP is in QSZK if }\alpha(n)-\beta(n)\geq 1/\mathrm{poly}(n).$
- \diamond Measured Quantum Triangular Discrimination Problem: (α,β) -measQTDP is in QSZK if $\alpha(n)-\beta(n)\geq 1/O(\log n)$;

What we need: Quantum counterparts of classical distances...

Quantum analogues of the triangular discrimination

Triangular discrimination: $TD(p_0, p_1) := \frac{1}{2} \sum_{x \in \mathcal{S}} \frac{(p_0(x) - p_1(x))^2}{p_0(x) + p_1(x)}$

Quantum analogues of the triangular discrimination

- $\ \, \underline{\mathsf{Option}\; \mathbf{1}\! :}\; \mathrm{QTD}(\rho_0,\rho_1) := \tfrac{1}{2} \mathrm{Tr} \Big((\rho_0-\rho_1)(\rho_0+\rho_1)^{-1/2} (\rho_0-\rho_1)(\rho_0+\rho_1)^{-1/2} \Big);$
- $\diamond \ \underline{\mathsf{Option}\ 2\text{::}}\ \mathrm{QTD}^{\mathrm{meas}}(\rho_0,\rho_1) := \sup\nolimits_{\mathsf{measurement}}\ {}_{\mathcal{E}}\left\{\mathrm{TD}\Big(p_0^{(\mathcal{E})},p_1^{(\mathcal{E})}\Big)\right\}.$

<u>Theorem 2.5.</u> Inequalities on quantum analogues of triangular discrimination:

	Classical	Quantum
SD vs. TD	$\mathrm{SD}^2 \leq \mathrm{TD} \leq \mathrm{SD}$ [Topsøe'00]	$td^2 \leq QTD^{meas} \leq QTD \leq td$
JS vs. TD	$rac{1}{2} \mathrm{TD} \leq \mathrm{JS} \leq \ln 2 \cdot \mathrm{TD}$ [Topsøe'00]	$\frac{1}{2}$ QTD 2 \leq QJS \leq QTD
H ² vs. TD	$ m H^2 \leq TD \leq 2H^2 \; [Le \; Cam'86]$	$\begin{aligned} \frac{1}{2}B^2 &\leq QTD^{meas} \leq B^2 \\ \frac{1}{2}B^2 &\leq QTD \leq B \end{aligned}$

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Easy regimes for the class QSZK

Theorem 3.1 (Easy regimes for the class QSZK).

Parameter regimes	$(1-\epsilon,\epsilon)$ - $\overline{\mathrm{SDP}}$	$(1 - \epsilon, \epsilon)$ - $\overline{\text{QSDP}}$
$\epsilon = 0$	in NP	in BQP
$\epsilon = 0$	Folklore	This work
$\epsilon \le 2^{-n/2 - 1}$	in PP	in PP
$\epsilon \leq z$	[Bouland-Chen-Holden-Thaler-Vasudevan'19]	This work
$\epsilon \ge 2^{-n^{1/2 - \gamma}}$	SZK-hard	QSZK-hard
for $\gamma \in (0,1/2)$	Implicitly in [Sahai-Vadhan'03]	Implicitly in [Watrous'02]

 $The proof is mainly based on different usages of the SWAP test \cite{test [Buhrman-Cleve-Watrous-de Wolf'01]}.$

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Conclusions and open problems

Take-home messages

- A classical distance may have several quantum counterpart, and a unique counterpart will make our life much easier.
- We define two quantum counterparts of the triangular discrimination, and demonstrate inequalities between these distances and other common distances.
- By employing these inequalities, we improve the QSZK containment of QSDP to non-polarizing regimes via two new QSZK-complete problems.
- Easy regimes for QSZK indicates that length-preserving polarization seems unlikely unless QSZK ⊆ PP.

Open problems

- Is there any other applications of these quantum analogues of the triangular discrimination? For instance, triangular discrimination can be used to improve the communication compleixty lower bound of the point chasing problem [Yehudayoff'20].
- Better upper bound for (α,β)-QSDP with α − β ≥ 1/poly?
 The best known bound is PSPACE which is implicitly shown in [Watrous'02].
- Quantum analogue of the set lower bound protocol and better bounds for QSZK?

