

VIETNAM NATIONAL UNIVERSITY,
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UNIVERSITY OF SCIENCE
FACULTY OF INFORMATION TECHNOLOGY

Project 01: Ares's Adventure Report

CS14003 – INTRODUCTION TO ARTIFICIAL INTELLIGENCE

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1 Group Information

- **Subject:** Introduction to Artificial Intelligence.
- **Class:** 23CLC09.
- **Lecturer:** Bui Duy Dang, Le Nhut Nam.
- **Team members:**

No.	Fullname	Student ID	Email
1	Ngo Nguyen The Khoa	23127065	nntkhoa23@clc.fitus.edu.vn
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- **Tools:**
 - **Git**, **GitHub**: Source code version control.
 - **CapCut**: Video editing.
 - **ChatGPT**, **Gemini** and **DeepSeek**.
 - **Visual Studio Code**: Code editor (Python, Latex).

2 Project Information

- **Name:** Ares's Adventure.
- **Developing Environment:** Visual Studio Code (Windows).
- **Programming Language:** Python.
- **Libraries and Tools:**
 - **rye:** A comprehensive project and package management solution for Python.
 - **PyGame:** A Python game maker library.
 - **numpy:** used for large, multi-dimensional arrays and matrices, along with a large collection of high-level mathematical functions to operate on these arrays.
 - **scipy:** pre-implementation of the Hungarian Min-matching algorithm.

3 Work assignment table

No.	Task Description	Assigned to	Rate
1	Implement BFS, DFS	Minh Duy	100%
2	Implement UCS, Dijkstra	Anh Khoa	100%
3	Implement A*, GBFS	The Khoa	100%
4	Implement Swarm	Anh Khoa	100%
5	Optimize heuristic function using Hungarian for min-matching	The Khoa, Minh Duy	100%
6	Optimize the number of expanded nodes by using deadlock detection	The Khoa, Anh Khoa	100%
7	Video Editing	Minh Duy	100%
8	Report	All members	100%
<p style="text-align: center;"><i>All requirements are completed!</i></p> <p style="text-align: center;"><i>No errors occur while running the program</i></p>			

4 Self-evaluation

No.	Criteria	Score
1	Implement BFS correctly.	100%
2	Implement DFS correctly.	100%
3	Implement UCS correctly.	100%
4	Implement A* correctly.	100%
5	Implement GBFS correctly.	100%
6	Generate at least 10 test cases for each level with different attributes.	100%
7	Result (output file and GUI).	100%
8	Video to demonstrate all algorithms for some test cases.	100%
9	Report.	100%
10	Implement, written report for Dijkstra's Algorithm and Swarm Algorithm.	100%

5 Algorithms' Implementations

5.1 Breadth First Search

Breadth First Search (BFS), as the name says, explores the search space in the increasing order of the depth and the costs of traveling from one state to another is assumed to be a positive number. Typically, this algorithm is often associated with the concept of stack and queue and pushing and popping from the stack.

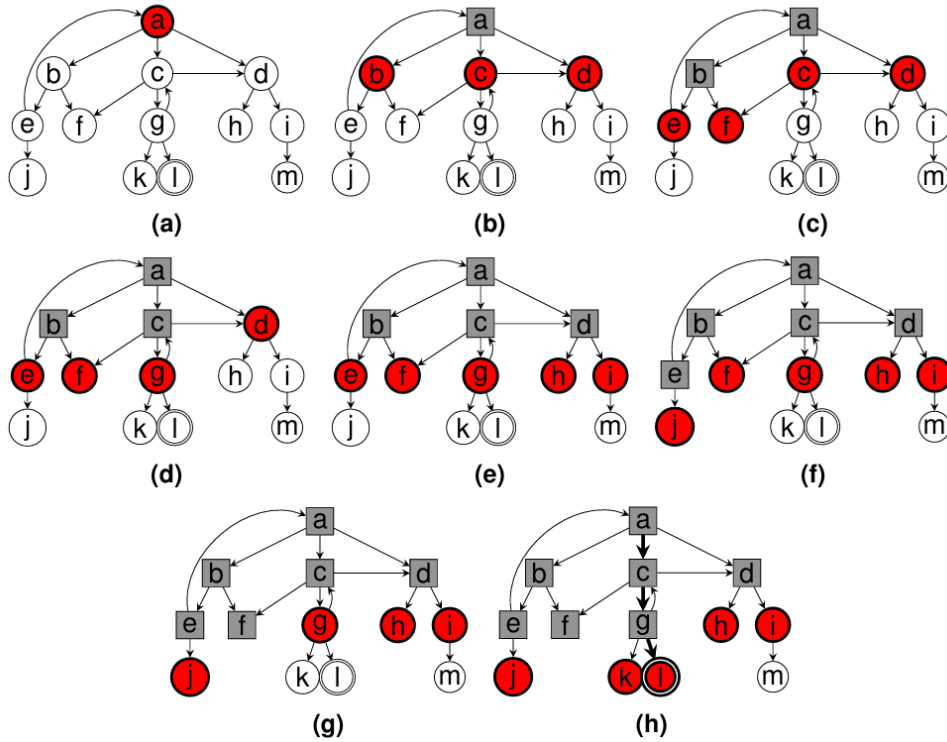


Figure 1: Breadth First Search (BFS)

Pseudocode

Algorithm 1 Breadth First Search (*start*, *goal*)

```
1: queue gets [start]
2: while queue is not empty do
3:   node gets dequeue(queue)
4:   if node = goal then
5:     return path
6:   end if
7:   for all neighbor in valid moves do
8:     if neighbor not visited then
9:       mark neighbor as visited
10:      enqueue(queue, neighbor)
11:    end if
12:  end for
13: end while
14: return failure
```

Implementation

Time and Space Complexity

Time Complexity: $O(V + E)$, where V is the number of vertices and E is the number of edges. Each node and edge is processed once.

Space Complexity: $O(V)$ in the worst case, as the queue stores all nodes at the widest level of the graph.

5.2 Depth First Search

Depth First Search (DFS) is a special case of backtracking search algorithm. The search starts from the root and proceeds to the farthest node before backtracking. The difference between this and the backtracking is that this stops the search once a goal is reached and does not care if it is not minimum.

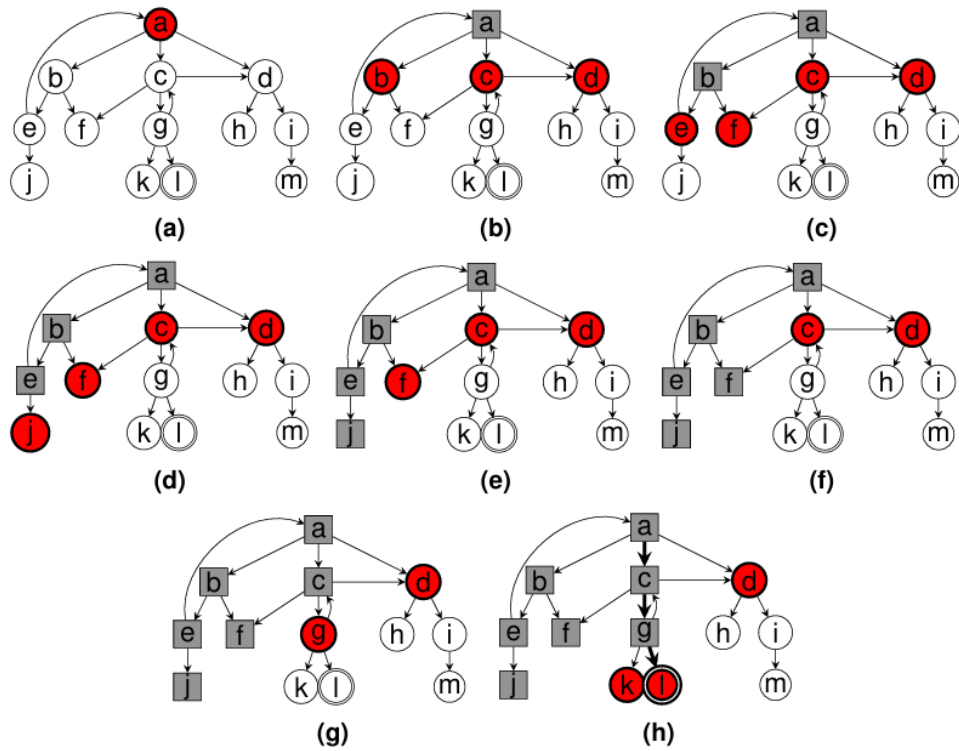


Figure 2: Depth First Search

Pseudocode

Algorithm 2 Depth First Search (*start*, *goal*)

```

1: stack  $\leftarrow$  [start]
2: while stack is not empty do
3:   node  $\leftarrow$  pop(stack)
4:   if node = goal then
5:     return path
6:   end if
7:   for all neighbor in valid moves do
8:     if neighbor not visited then
9:       mark neighbor as visited
10:      push(stack, neighbor)
11:    end if
12:  end for
13: end while
14: return failure

```

Implementation

Time and Space Complexity

5.3 Uniform Cost Search

For any search problem, Uniform Cost Search (UCS) is the better algorithm than the previous ones. The search algorithm explores in branches with more or less same cost. This consist of a priority queue where the path from the root to the node is the stored element and the depth to a particular node acts as the priority. UCS assumes all the costs to be non negative. While the DFS algorithm gives maximum priority to maximum depth, this gives maximum priority to the minimum cumulative cost.

Pseudocode

Algorithm 3 Uniform Cost Search (*start*, *goal*)

```
1: priority queue  $\leftarrow$  [(start, cost = 0)]
2: while priority queue is not empty do
3:   (node, cost)  $\leftarrow$  dequeue(priority queue)
4:   if node = goal then
5:     return path
6:   end if
7:   for all neighbor in valid moves do
8:     new cost  $\leftarrow$  cost + move cost
9:     if neighbor not visited or new cost < previous cost then
10:      mark neighbor as visited
11:      enqueue(priority queue, (neighbor, new cost))
12:    end if
13:  end for
14: end while
15: return failure
```

Implementation**Time and Space Complexity****5.4 Dijkstra's Algorithm****Pseudocode**

Algorithm 4 Dijkstra's Algorithm (*start*, *goal*)

```

1: priority queue  $\leftarrow$  [(start, cost = 0)]
2: distances[start]  $\leftarrow$  0
3: while priority queue is not empty do
4:   (node, cost)  $\leftarrow$  dequeue(priority queue)
5:   if node = goal then
6:     return distances
7:   end if
8:   for all neighbor in valid moves do
9:     new cost  $\leftarrow$  cost + move cost
10:    if new cost < distances[neighbor] then
11:      distances[neighbor]  $\leftarrow$  new cost
12:      enqueue(priority queue, (neighbor, new cost))
13:    end if
14:  end for
15: end while
16: return distances

```

Implementation**Time and Space Complexity****5.5 Greedy Best-First Search**

Greedy Best-First Search is similar to A* algorithm. The difference is that the vertices in the priority queue are ordered only by the estimated remaining distance to the solution. It has to be noted that complete best first search is not optimal.

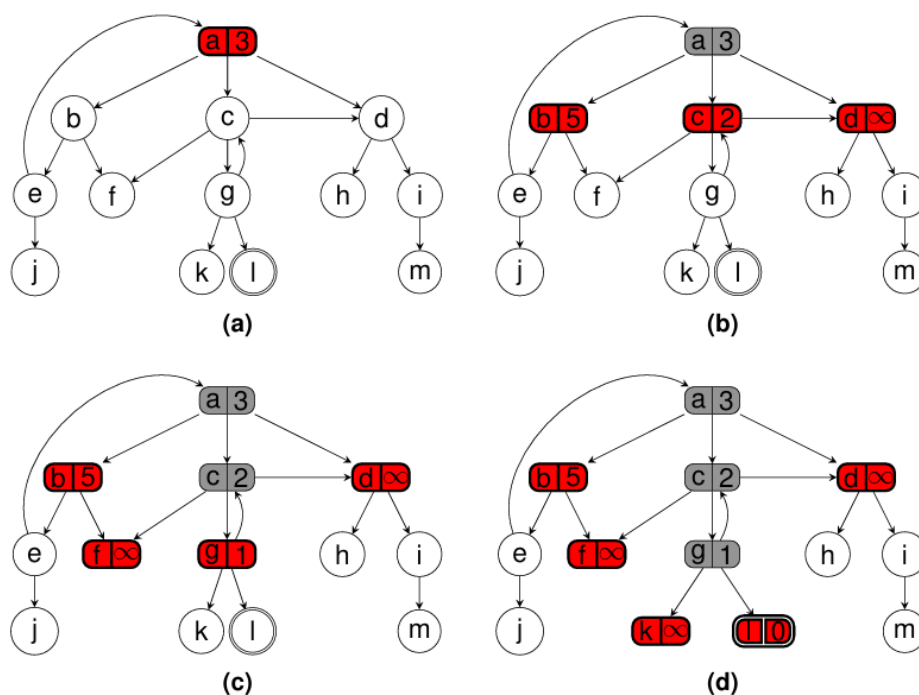


Figure 3: Greedy Best-First Search

Pseudocode

Algorithm 5 Greedy Best-First Search (*start*, *goal*, *heuristic*)

```

1: priority queue  $\leftarrow [(start, heuristic(start))]$ 
2: while priority queue is not empty do
3:   node  $\leftarrow$  dequeue(priority queue)
4:   if node = goal then
5:     return path
6:   end if
7:   for all neighbor in valid moves do
8:     if neighbor not visited then
9:       mark neighbor as visited
10:      enqueue(priority queue, (neighbor, heuristic(neighbor)))
11:    end if
12:  end for
13: end while
14: return failure

```

Implementation

Time and Space Complexity

5.6 A* Search with heuristic

A* algorithm is one of the popular technique used in path finding and graph traversals. This algorithm completely relies on heuristics for computing the future cost of a problem. This

algorithm is equivalent to the uniform cost search with modified edge cost. This heuristic is chosen according to the case where the algorithm is implemented, thus emphasizing the importance of domain knowledge. This algorithm is consistent if the modified cost is greater than zero.

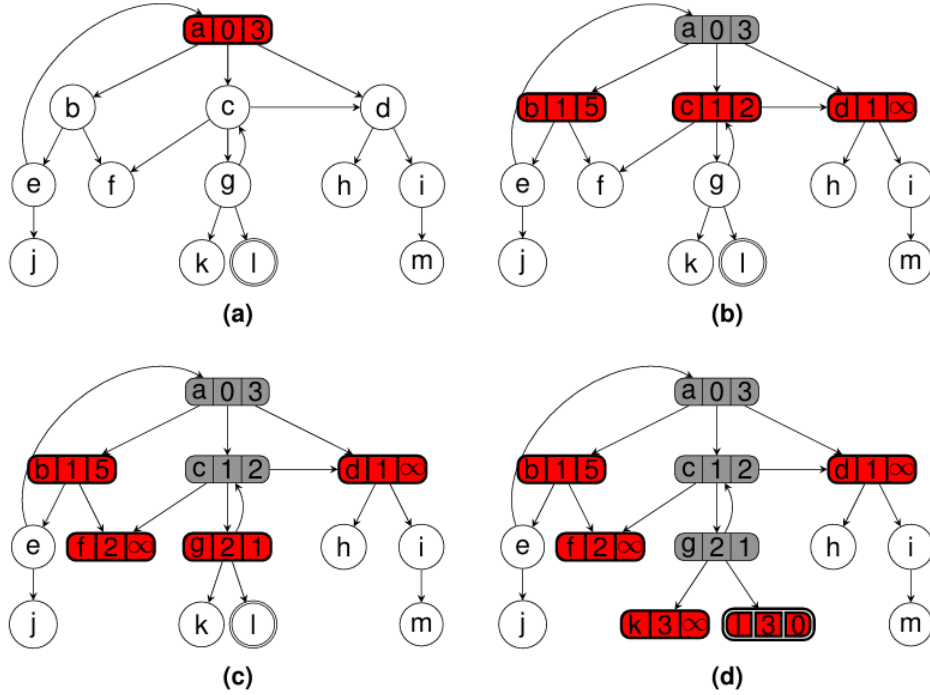


Figure 4: A* Algorithm

Pseudocode

Algorithm 6 A* Search (*start*, *goal*, *heuristic*)

```

1: priority queue  $\leftarrow [(start, cost = 0, estimated\ total\ cost = heuristic(start))]$ 
2: while priority queue is not empty do
3:   (node, cost)  $\leftarrow$  dequeue(priority queue)
4:   if node = goal then
5:     return path
6:   end if
7:   for all neighbor in valid moves do
8:     new cost  $\leftarrow$  cost + move cost
9:     estimated total cost  $\leftarrow$  new cost + heuristic(neighbor)
10:    if neighbor not visited or new cost < previous cost then
11:      mark neighbor as visited
12:      enqueue(priority queue, (neighbor, new cost, estimated total cost))
13:    end if
14:  end for
15: end while
16: return failure

```

Implementation

Time and Space Complexity

5.7 Swarm Algorithm

Pseudocode

Implementation

Time and Space Complexity

6 App Screenshots

7 References

1. [Rye documentation](#)