

# Abstract Binding Trees

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## 1 Preliminaries

Fix a set  $\mathcal{S}$  of *sorts*. We will say  $s$  *sort* when  $s \in \mathcal{S}$ . A valence  $\vec{p} \parallel \vec{q}.s$  specifies an expression of sort  $s$  which binds symbols in  $\vec{p}$  and variables in  $\vec{q}$ .

$$\frac{s \text{ sort} \quad p_i \text{ sort } (i \leq m) \quad q_i \text{ sort } (i \leq n)}{p_0, \dots, p_m \parallel q_0, \dots, q_n.s \text{ valence}}$$

An arity  $(\vec{v})_s$  specifies an operator of sort  $s$  with arguments of valences  $\vec{v}$ . We will call the set of valences  $\mathcal{V}$ , and the set of arities  $\mathcal{A}$ .

$$\frac{s \text{ sort} \quad v_i \text{ valence } (i \leq n)}{(v_0, \dots, v_n)_s \text{ arity}}$$

Let  $\mathbb{I}$  be the infinite set of symbols; let  $\mathbb{F}$  be the free cocartesian category over  $\mathbb{I}$ . Then, fix a covariant presheaf of operators  $\mathcal{O} : \mathbb{F} \times \mathcal{A} \rightarrow \mathbf{Set}$  such that the arrows in  $\mathbb{F}$  lift to renamings of operators' parameters; via the Grothendieck construction, we can also consider the set  $\int \mathcal{O}$  of operators  $(U, \rho, \vartheta)$  for  $\vartheta \in \mathcal{O}(U, \rho)$ .

$$\frac{\vartheta \in \mathcal{O}(U, \rho)}{U \vdash \vartheta : \rho}$$

The judgment  $U \vdash \vartheta : \rho$  supports the structural principles of weakening and exchange, because of the functoriality of  $\mathcal{O}$ .

**Examples** Operators are defined by specifying the fibres of  $\int \mathcal{O}$  in which they reside. For instance, consider the lambda calculus with a single sort, **exp**, about whose operators we may assert the following:

$$\begin{aligned} U \vdash \lambda : (\cdot \parallel \text{exp}.\text{exp})\text{exp} \\ U \vdash \text{ap} : (\cdot \parallel \cdot.\text{exp}, \cdot \parallel \cdot.\text{exp})\text{exp} \end{aligned}$$

However, consider the extension of the calculus with assignables. Then, we shall have the following:

$$\begin{aligned}
U &\vdash \text{decl} : (\cdot \parallel \cdot . \text{exp}, \text{exp} \parallel \cdot . \text{exp}) \text{exp} \\
U, u &\vdash \text{get}\{u\} : (\cdot) \text{exp} \\
U, u &\vdash \text{set}\{u\} : (\cdot \parallel \cdot . \text{exp}) \text{exp}
\end{aligned}$$

Weakening can be seen as inducing a “degeneracy map” on operators, whereas a renaming  $u \mapsto v$  will take  $\text{get}\{u\}$  to  $\text{get}\{v\}$ .

## 2 Contexts

In general, we will have three kinds of context: metacontexts, variable contexts, and parameter contexts. A metacontext  $\Omega$  consists of bindings of valences to metavariables; a variable context  $\Gamma$  is a collection of bindings of sorts to variables, and a parameter context  $\Upsilon$  is a collection of bindings of sorts to symbols.

$$\frac{}{\cdot \text{ mctx}} \quad \frac{\Omega \text{ mctx} \quad v \text{ valence} \quad \mathsf{M} \notin |\Omega|}{\Omega, \mathsf{M} : v \text{ mctx}}$$

$$\frac{}{\cdot \text{ vctx}} \quad \frac{\Gamma \text{ vctx} \quad s \text{ sort} \quad x \notin |\Gamma|}{\Gamma, x : s \text{ vctx}}$$

$$\frac{}{\cdot \text{ sctx}} \quad \frac{\Upsilon \text{ vctx} \quad s \text{ sort} \quad u \notin |\Upsilon|}{\Upsilon, u : s \text{ sctx}}$$

## 3 Abstract Binding Trees

Let the judgment  $\Omega \triangleright \Upsilon \parallel \Gamma \vdash M : s$  presuppose  $\Omega \text{ mctx}$ ,  $\Upsilon \text{ sctx}$ ,  $\Gamma \text{ vctx}$  and  $s \text{ sort}$ , meaning that  $M$  is an abstract binding tree of sort  $s$ , with metavariables in  $\Omega$ , parameters in  $\Upsilon$ , and variables in  $\Gamma$ .

$$\begin{array}{c}
\frac{\Gamma \ni x : \textcolor{brown}{s}}{\Omega \triangleright \Upsilon \parallel \Gamma \vdash x : \textcolor{brown}{s}} \textit{var} \\
\\
\frac{\begin{array}{l} \Omega \ni \mathsf{M} : \textcolor{brown}{p}_0, \dots, \textcolor{brown}{p}_m \parallel \textcolor{brown}{q}_0, \dots, \textcolor{brown}{q}_n . \textcolor{brown}{s} \\ \Upsilon \ni u_i : \textcolor{brown}{p}_i \quad (i \leq m) \\ \Omega \triangleright \Upsilon \parallel \Gamma \vdash M_i : \textcolor{brown}{q}_i \quad (i \leq n) \end{array}}{\Omega \triangleright \Upsilon \parallel \Gamma \vdash_{\mathsf{M}} \{u_0, \dots, u_m\} (u_0, \dots, u_m; M_0, \dots, M_n) : \textcolor{brown}{s}} \textit{mvar} \\
\\
\frac{\begin{array}{l} |\Upsilon| \vdash \mathfrak{S} : \textcolor{brown}{p}_0, \dots, \textcolor{brown}{p}_m \parallel \textcolor{brown}{q}_0, \dots, \textcolor{brown}{q}_n . \textcolor{brown}{s} \\ \Upsilon \ni u_i : \textcolor{brown}{p}_i \quad (i \leq m) \\ \Omega \triangleright \Upsilon \parallel \Gamma \vdash M_i : \textcolor{brown}{q}_i \quad (i \leq n) \end{array}}{\Omega \triangleright \Upsilon \parallel \Gamma \vdash_{\mathfrak{S}} \{u_0, \dots, u_m\} (u_0, \dots, u_m; M_0, \dots, M_n) : \textcolor{brown}{s}} \textit{app}
\end{array}$$