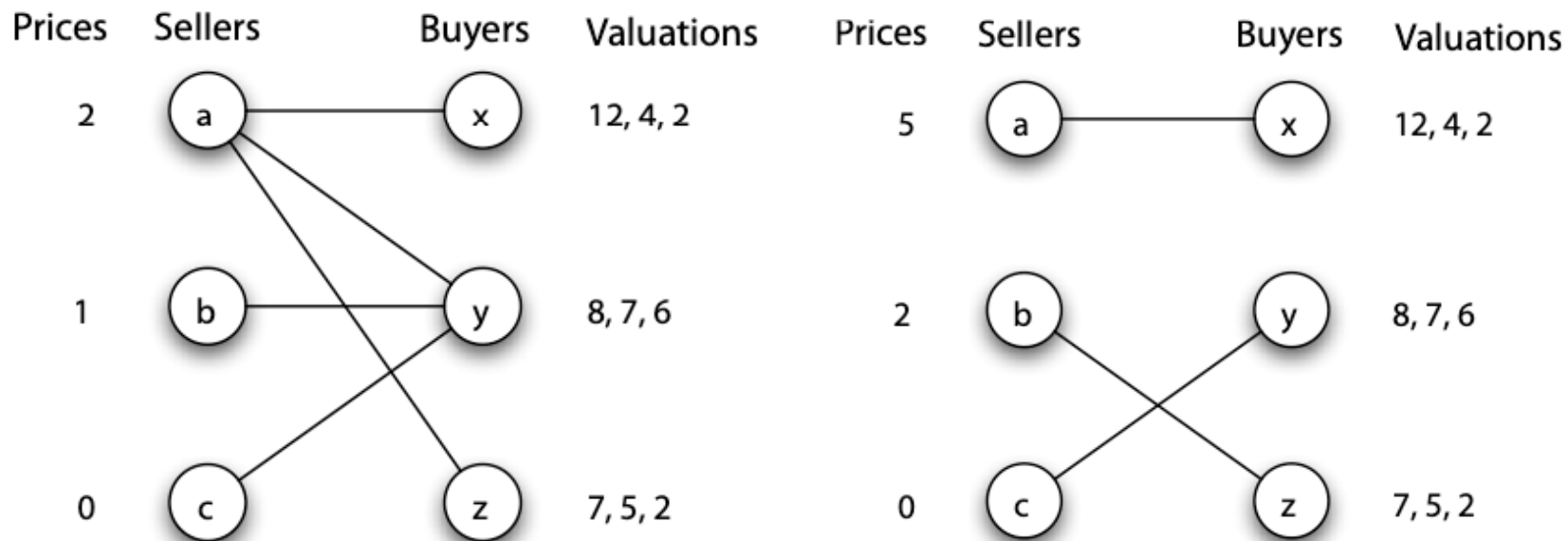


VE444: Networks

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Matching Markets

Market clearing price



A set of price is **market clearing price** if the resulting preferred-seller graph has a **perfect matching**.

- Coordination over tie-breaking allowed
- Multiple sets of market clearing price could exist

Optimality of the Market-clearing Prices

Optimality of the Market-clearing Prices: For any set of **market-clearing prices**, a perfect matching in the resulting preferred-seller graph has the **maximum total valuation** of any assignment of sellers to buyers

Reason:

M: perfect matching

Total Payoff of buyers in M = Total Valuation of buyers in M – Sum of all prices

Optimality of the Market-clearing Prices (v2): For any set of market-clearing prices, a perfect matching in the resulting preferred-seller graph has the maximum sum of payoffs of all sellers to buyers

Existence of Market-clearing prices

Existence of Market-clearing prices: For any set of buyer valuations, there exists a set of market-clearing price.

- Why it always exists?
 - Construct a procedure that stops only when market-clearing prices are found
 - This procedure has limited rounds.

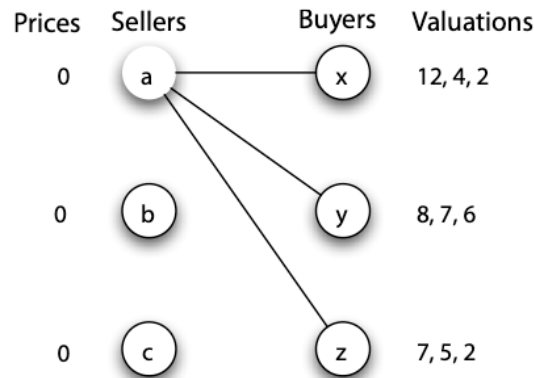
How to find a set of Market-Clearing Prices?

- Two parts: price probing(increase) and price reduction
- Procedure:
 1. At the start of each round, a set of prices, with the smallest one equal to 0
 2. Construct the preferred-seller graph and check whether there is a perfect matching
 3. Stops when perfect matching exists, output current prices
 4. If not, identify a constricted set of buyers, S and their neighbors $N(S)$
 5. Each seller in $N(S)$ simultaneously raises his price by one unit
 6. Reduce price to guarantee the smallest price equal to 0.
 7. Using the updated price to start a new round

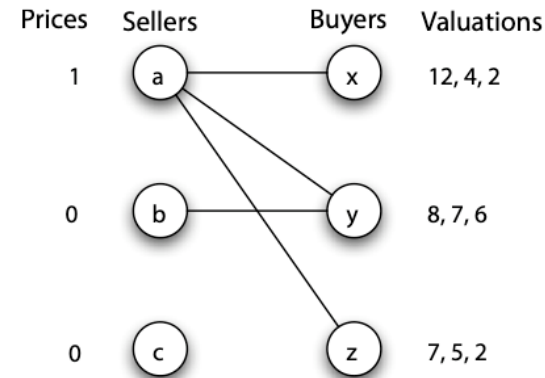
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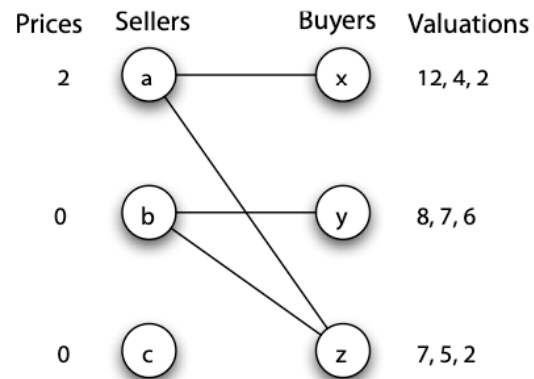
How to find a set of Market-Clearing Prices?



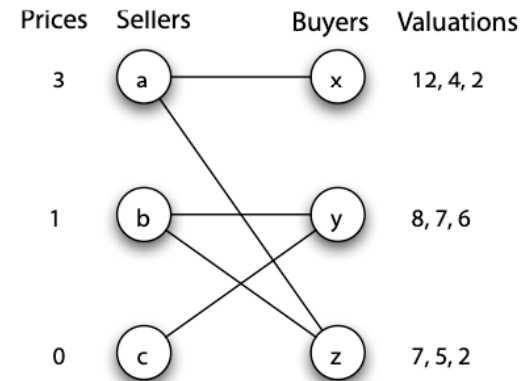
(a) Start of first round



(b) Start of second round



(c) Start of third round



(d) Start of fourth round

Existence of the market clearing price

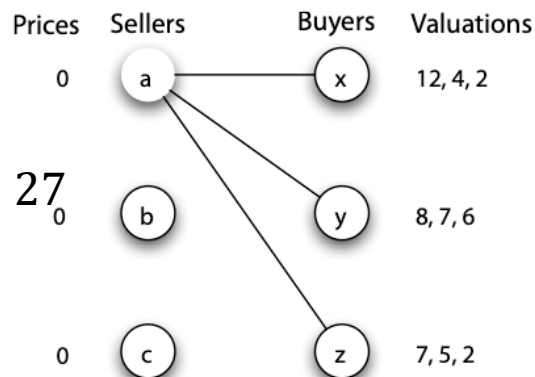
- This procedure stops after finite rounds.
- **Potential energy:**
 - Potential energy of a buyer: maximum payoff under current prices
 - Potential energy of a seller: current price
 - Potential energy of the auction: sum of potentials of all buyers and sellers
- Potential energy at the begin: $P \geq 0$
- Potential energy at the start of each round at least zero
- Potential energy only changes when the prices change
 - Price reduction: no change
- Price probing (S : constricted buyer set):
 - Each seller potential goes up by one unit
 - Each buyer potential goes down by one unit
 - But $S \geq N(S)$

Existence of the market clearing price

$$P_s: 0$$

$$P_b: 12 + 8 + 7 = 27$$

$$P_t: 27$$

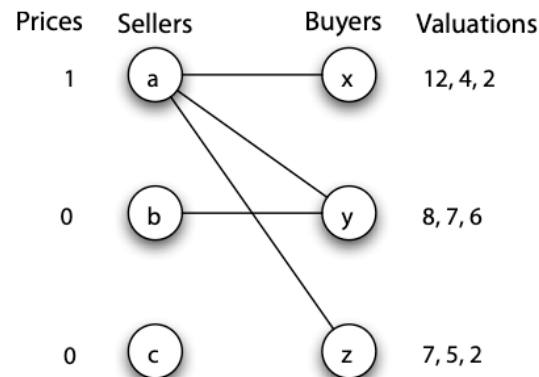


(a) Start of first round

$$P_s: 1$$

$$P_b: 11 + 7 + 6 = 24$$

$$P_t: 25$$

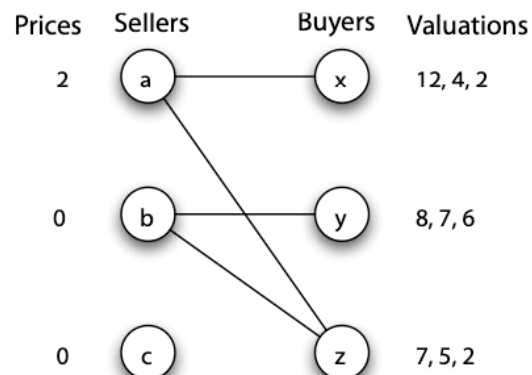


(b) Start of second round

$$P_s: 2$$

$$P_b: 10 + 7 + 5 = 22$$

$$P_t: 24$$

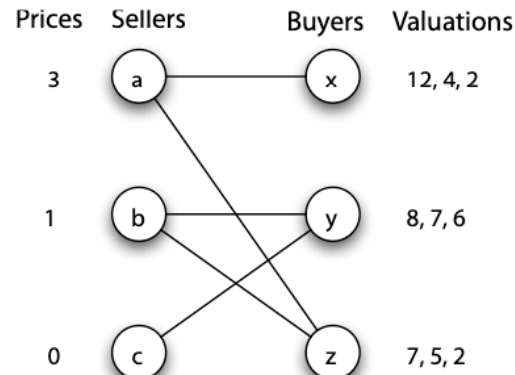


(c) Start of third round

$$P_s: 4$$

$$P_b: 9 + 6 + 4 = 19$$

$$P_t: 23$$

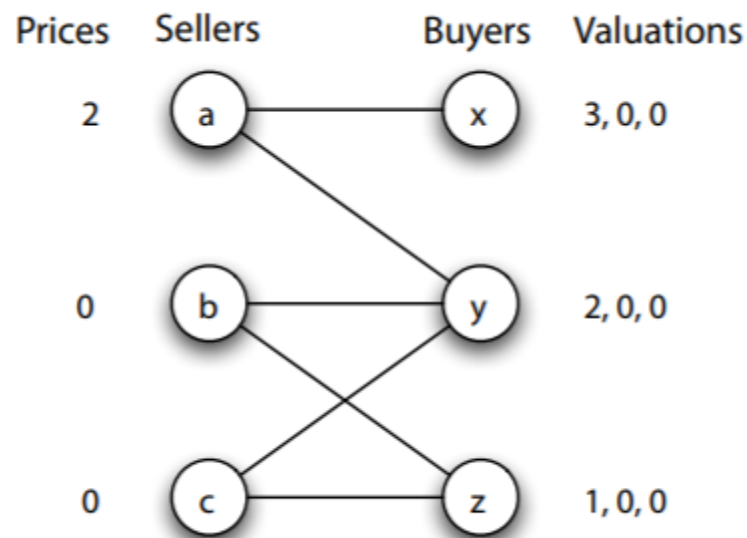
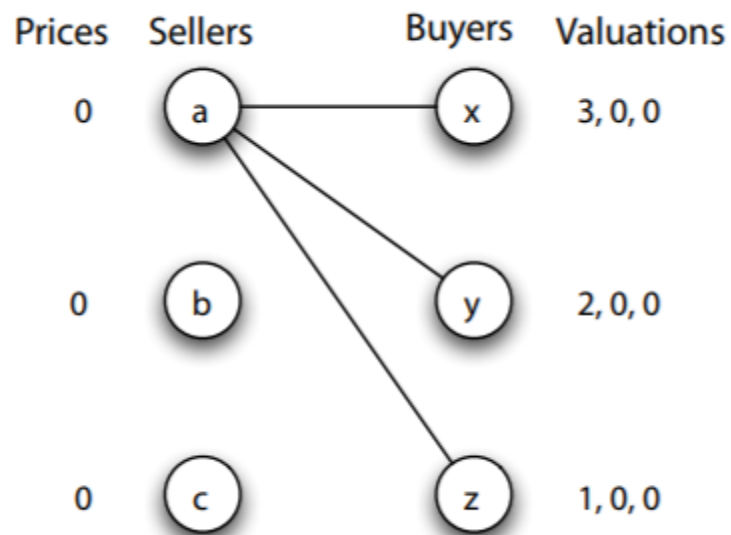


(d) Start of fourth round

Relationship with auction?

Relationship with single-item auctions

From one seller, n buyers to perfect matching model



Homework

During the winter break, n students would love to choose a foreign university to do the exchange problem. Each student has a list of universities that they want to go. Naturally, each university has some capacities. For simplicity, let's assume each university can accommodate 3 students at most. *Describe an algorithm that finds if all the students can go to their desired universities for exchange.* (Hint: Embedding the problem into the appropriate graph is the key. The rest is just class content)

Summary

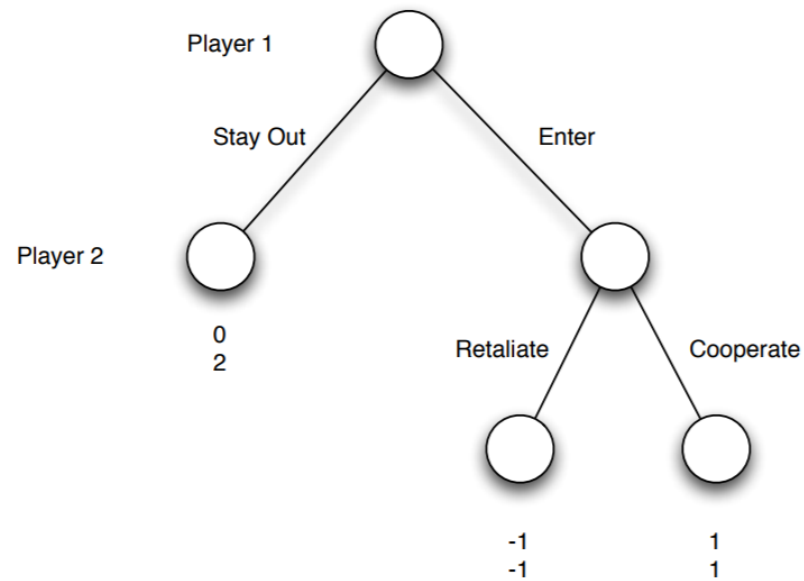
- Perfect matching v.s. constricted set
- Alternating BFS -> perfect matching
 - Augmented path || constricted set
- Existence of the market clearing price
- Optimality of the market clearing price

Dynamic game

- Dynamic game: some player or set of players moves first, others players respond.
 - Example: Board games, negotiations
- Market Entry Game: Consider a region where Firm 2 is currently the only serious participant in a given line of business, and Firm 1 is considering whether to enter the market.
 - The first move in this game is made by Firm 1, who must decide whether to stay out of the market or enter it.
 - If Firm 1 chooses to stay out, then the game ends, with Firm 1 getting a payoff of 0 and Firm 2 keeping the payoff from the entire market
 - If Firm 1 chooses to enter, then the game continues to a second move by Firm 2, who must choose whether to cooperate and divide the market evenly with Firm 1, or retaliate and engage in a price war.
 - If Firm 2 cooperates, then each firm gets a payoff corresponding to half the market
 - If Firm 2 retaliates, then each firm gets a negative payoff.
- **Extensive form** representation of a game:

Dynamic game

- Extensive form representation of a game:
- Backward induction



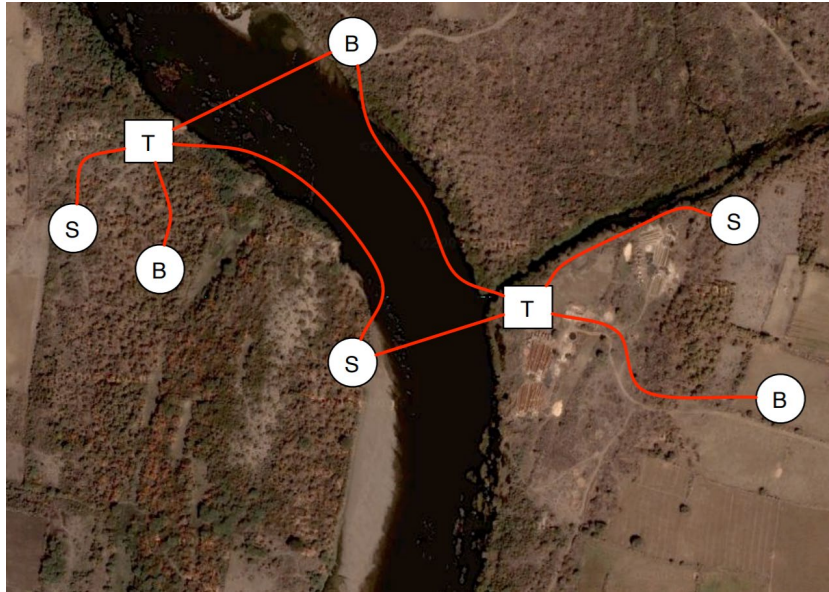
Markets with intermediaries

- Single seller, multiple buyers, private valuations -> auction
- Multiple seller, multiple buyers, direct interaction -> matching
- **Trade with Intermediaries:** individual buyers and sellers trade through intermediaries
 - Example: Agricultural goods trading, financial assets markets

Network model

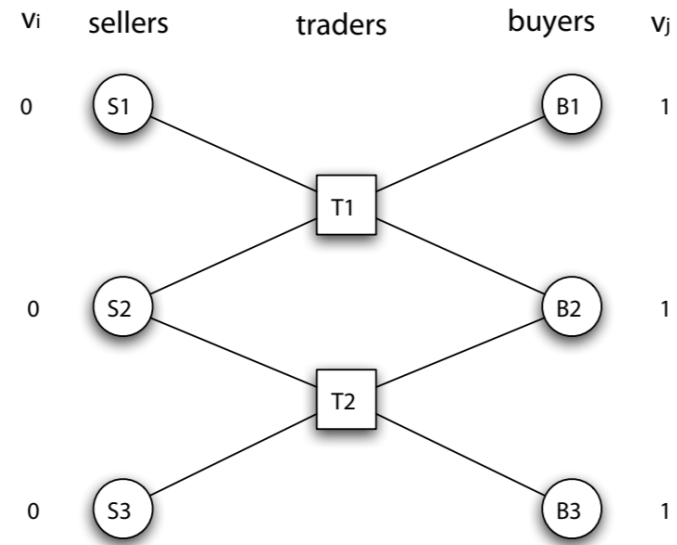
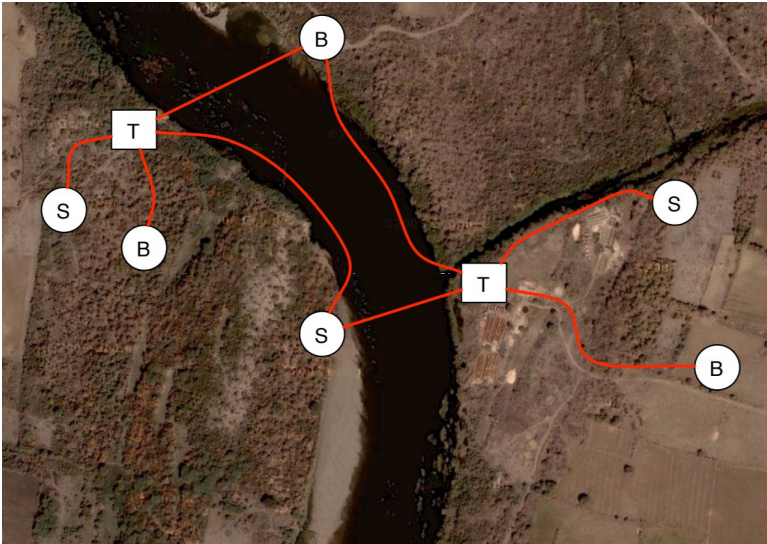
- Individual buyers and sellers often trade through intermediaries
- Not all buyers and sellers have access to the same intermediaries
- Not all buyers and sellers trade at the same price

Example: Agricultural trading in a developing country



Poor transportation networks
Perishability of the products

Network model



Single type

Indivisible units

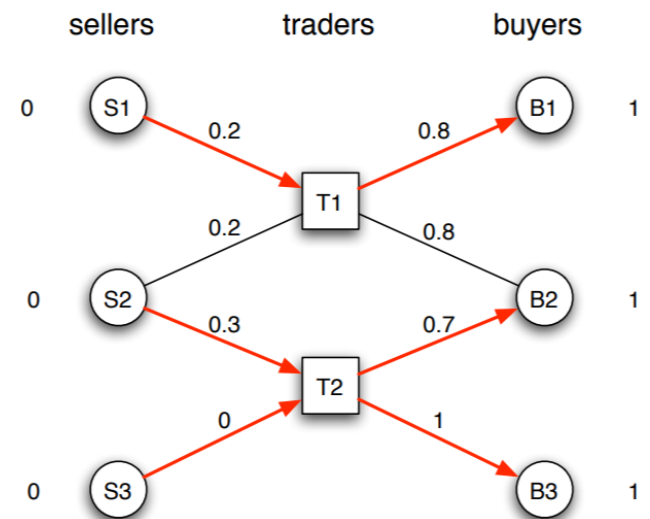
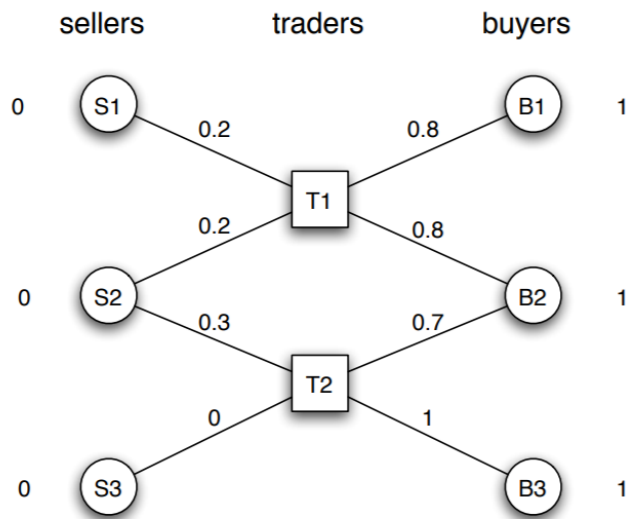
Valuation known to all buyers, sellers, and traders

Difference between matching:

- Intermediaries
- Same valuation
- Network is fixed

Prices in the Network model

Dynamic game: traders first set prices, sellers and buyers react to these prices



Trader's strategy: choice of bid b_{ti} and ask prices a_{tj}

Buyer/Seller's strategy: which trader to deal with or stay out of the transaction

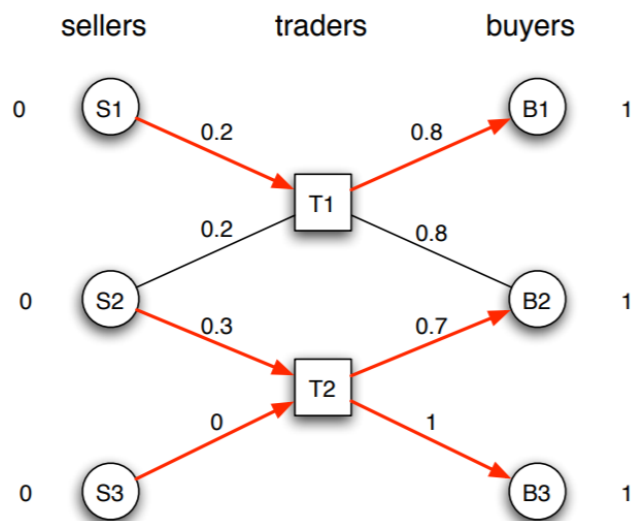
Indifference exist: tie-breaking

Payoff for players

Seller's payoff: Trade b_{ti} , no trade v_i (=0 currently);

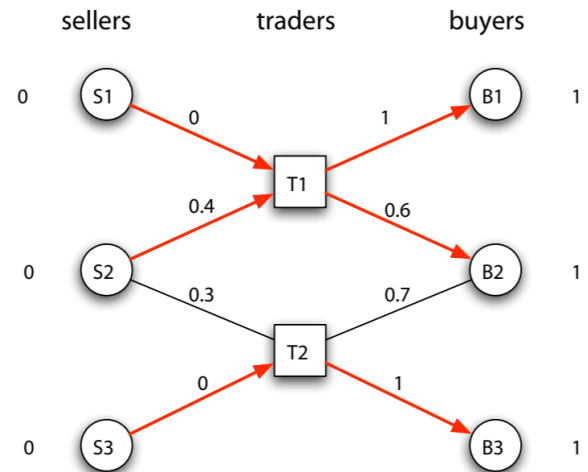
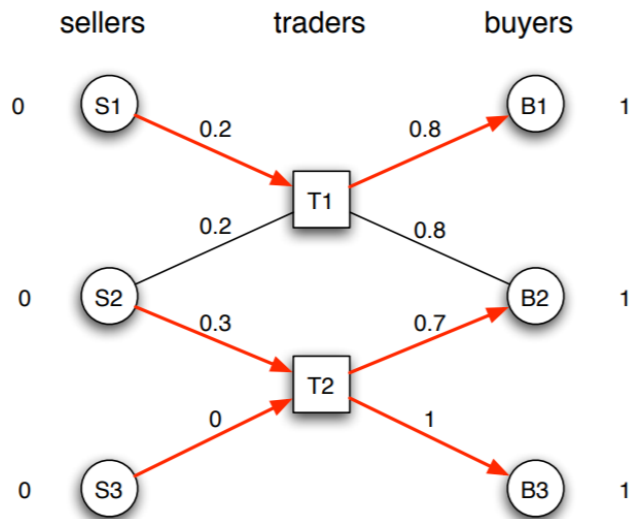
Buyer's payoff: Trade: $v_j - a_{tj}$; No Trade: 0

Trader's payoff: sum of asks – sum of bids



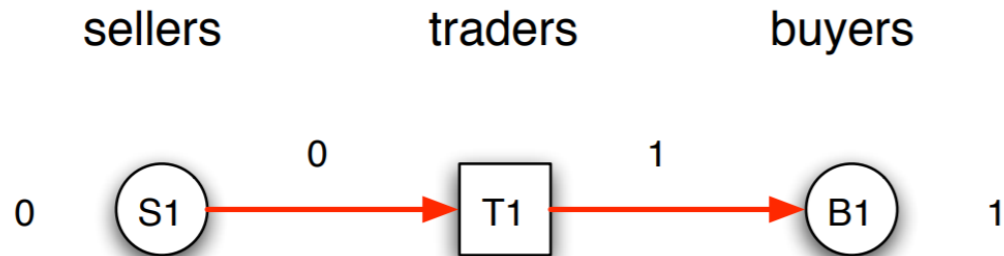
Equilibrium?

T1 to S2 and B2
T1 to S1 and B1

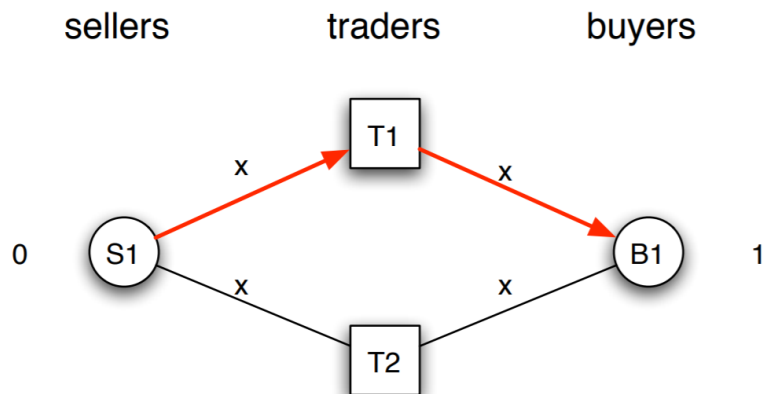


Simple building blocks

Monopoly case:

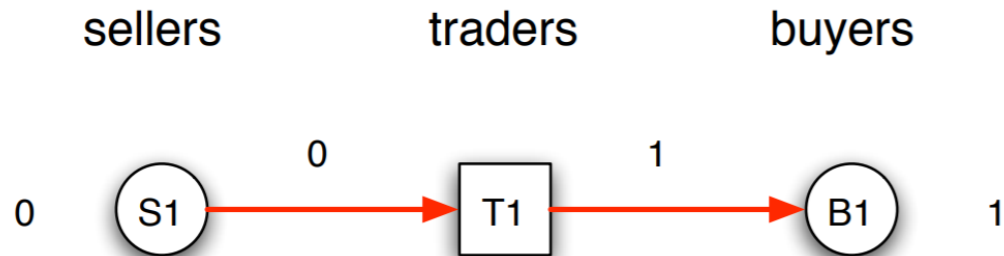


Perfect competition:

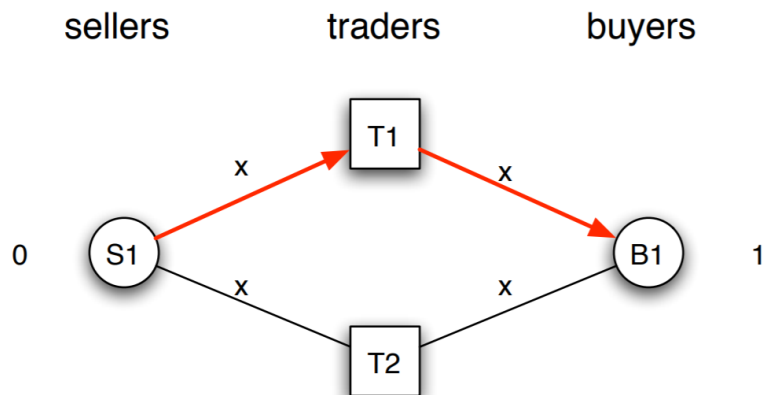


Reasoning

Monopoly case:



Perfect competition:



1. Trader performing the trade at the equilibrium must have a payoff of 0
2. Trader not performing the trade must also have bid and ask values of x

The equilibria for the trading network

