# VE444: Networks

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# PageRank: The "Flow" Model

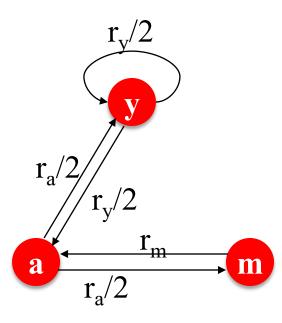
- A page is important if it is pointed to by other important pages
- Define a "rank"  $r_j$  for node j

$$r_j = \sum_{i \to j} \frac{r_i}{d_i}$$

 $d_i$  ... out-degree of node i

You might wonder: Let's just use Gaussian elimination to solve this system of linear equations. Bad idea!

The web in 1839



#### "Flow" equations:

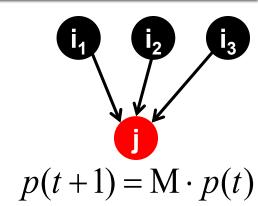
$$r_y = r_y/2 + r_a/2$$

$$r_a = r_y/2 + r_m$$

$$r_m = r_a/2$$

### Random walk

- Where is the surfer at time *t*+1?
  - Follow a link uniformly at random  $p(t+1) = M \cdot p(t)$



Suppose the random walk reaches a state

$$p(t+1) = M \cdot p(t) = p(t)$$

then p(t) is stationary distribution of a random walk

- Our original rank vector r satisfies  $r = M \cdot r$ 
  - So, r is a stationary distribution for the random walk

# Solution: Random Teleports

- Google's solution that does it all:
  - At each step, random surfer has two options:
  - With probability  $\beta$ , follow a link at random
  - With probability  $1-\beta$ , jump to some random page
- PageRank equation [Brin-Page, 98]

$$r_j = \sum_{i \to i} \beta \frac{r_i}{d_i} + (1 - \beta) \frac{1}{N}$$

of node i

This formulation assumes that *M* has no dead ends. We can either preprocess matrix *M* to remove all dead ends or explicitly follow random teleport links with probability 1.0 from dead-ends.

# Link analysis in Modern Web search

- Combining Links, text, and usage data
  - "I am a student at <u>SJTI-JI</u>". (Anchor text)
  - "For more information, click here"

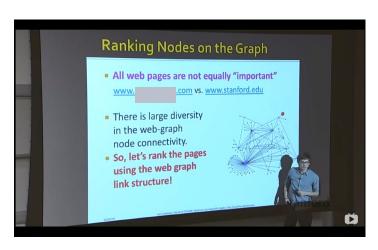


# Link analysis in Modern Web search

- Current ranking functions are constantly evolving
- Always expect the world to react to what you do (Game theoretic view)
  - Search Engine Optimization/Marketing(SEO, SEM)
  - An example last week

#### Links as Votes

- Idea: Links as votes
  - Page is more important if it has more links
    - In-coming links? Out-going links?
- Think of in-links as votes:
  - www.stanford.edu has 23,400 in-links
  - www.nobodyhomepg.com has 1 in-link
- Are all in-links equal?
  - Links from important pages count more
  - Recursive question!



https://www.bilibili.com/video/BV1jE41177A4?p=13

# **Sponsored Search Markets**

## Google revenue breakdown

- In 2019, Google made \$161.85 billion in revenue
- Out of this amount, \$134.81 billion from advertising!

70.9%

#### The search advertising market



#### Welcome to The **Keuka Lake** Wine Trail

Information about seven wineries on **Keuka Lake** in the Finger Lakes district. Offers a trail map, event calendar, winery descriptions, tourist services, ...

www.keukawinetrail.com/ - 13k - Cached - Similar pages - Note this

#### A complete guide to the Keuka Lake Wine Country

your own, follow the **Keuka Lake** Wine Trail, or book a wine tour and leave the driving to a pro. From casual to gourmet, hotdogs to haute cuisine, ...

www.keukalake.com/ - 24k - Cached - Similar pages - Note this

#### Keuka Lake - Wikipedia, the free encyclopedia

**Keuka Lake** is an unusual member of the Finger Lakes because it is Y-shaped instead of long and narrow. Because of its shape, it was referred to in the past ...

en.wikipedia.org/wiki/Keuka\_Lake - 26k - Cached - Similar pages - Note this

#### Seneca Lake (New York) - Wikipedia, the free encyclopedia

The two main inlets are Catharine Creek at the southern end and the Keuka Lake Outlet.

Sponsored Links

#### Keuka Lake Lodging

Lakeside vacation rentals on the Finger Lakes in upstate New York. FingerLakesPremierProperties.com

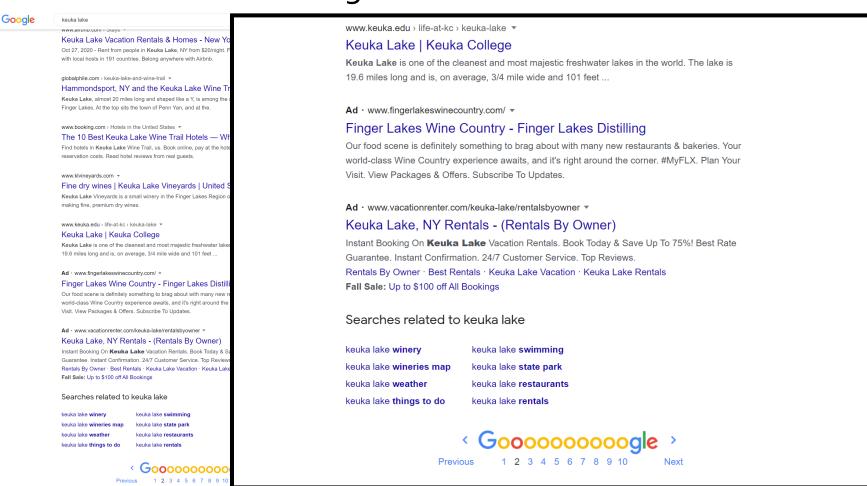
#### Keuka Lake Real Estate

Looking for Information about Keuka Lake Real Estate? www.MarkMalcolm.com New York

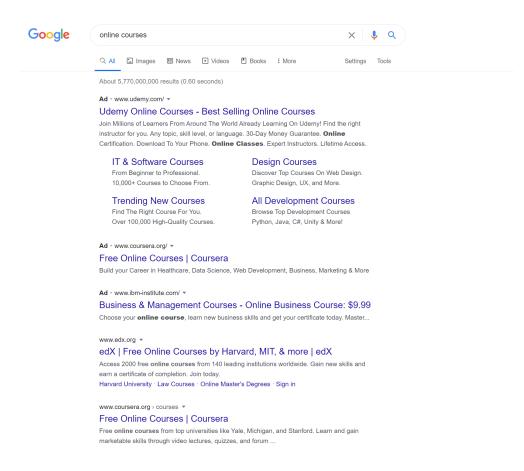
#### Finger Lakes Real Estate

Find your dream home; Lakefront, Lakeview, Cottage, Land or Farm! www.winetrailproperties.com New York

#### The search advertising market



#### The search advertising market



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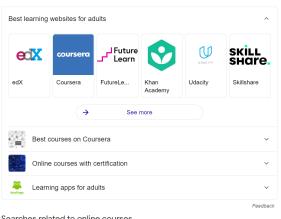
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- We have studied auction
- We have studied matching
- We study them in separation!

Sponsored search market is a combination of matching and auction!

### Three problems

- Question 1: how should advertisers pay for the ads?
- Question 2: how should websites allocate the ads slots to the advertisers?
- Question 3: how should websites set the prices on the ad slots

Sponsored search market is a combination of matching and auction!

## How to sale ads in search engines

### Keyword-base advertisement

- Targeted advertisement, as compared to traditional forced advertisement
- Low cost, especially for less frequently searched keywords
- Easy to evaluate the effectiveness, i.e., the number of users clicked

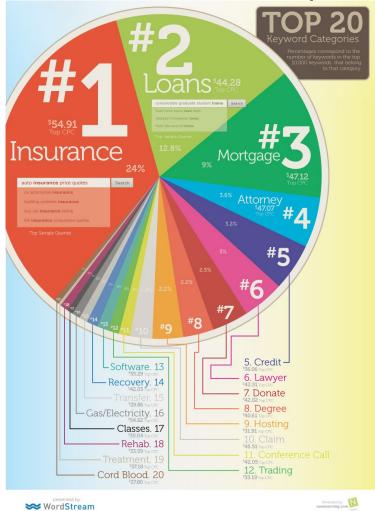
Question 1: how should advertisers pay for the ads?

### Payment methods by the websites

- By showing:
  - CPM (cost per mille/cost per thousand impressions)
- By action:
  - CPA (cost per action)
  - CPC (cost per click)
- By sale:
  - CPO (cost per order)

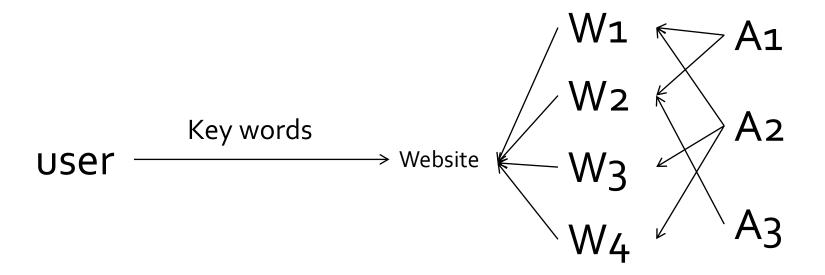
## Payment methods by the websites

Different key words have different prices



### Keyword-base advertisement

Question 2: how should websites allocate the ads slots to the advertiser?



Ad slots Advertisers

### Modeling

Can we model it into the matching problem we have studied before?

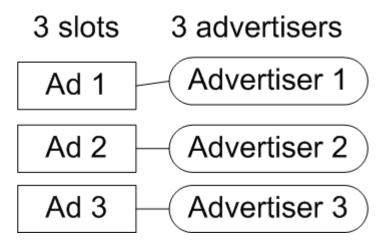
From a real world scenario → problem formulation (recall the matching market that we have learned before)

### Some basic concepts

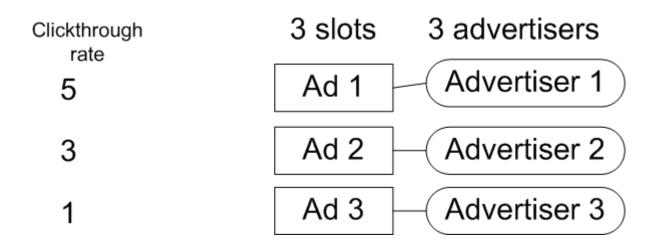
- Advertising slots
- Clickthrough rate: the expected clicks per hour on an advertising slot
- Advertiser's revenue per click: the expected revenue of every click

### Some basic concepts

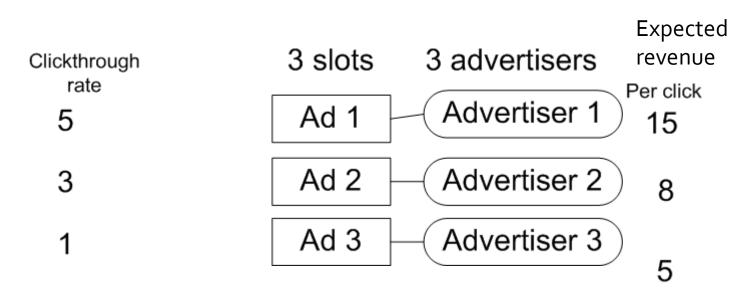
- Advertising slots
- Clickthrough rate: the expected clicks per hour on an advertising slot
- Advertiser's revenue per click: the expected revenue of every click
- Advertiser's valuation (per unit time): clickthrough rate \* revenue per click
- Advertiser's payoff: expected revenue price



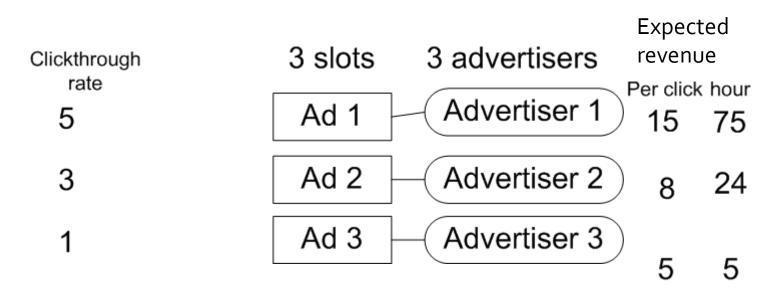
3 slots and 3 advertisers



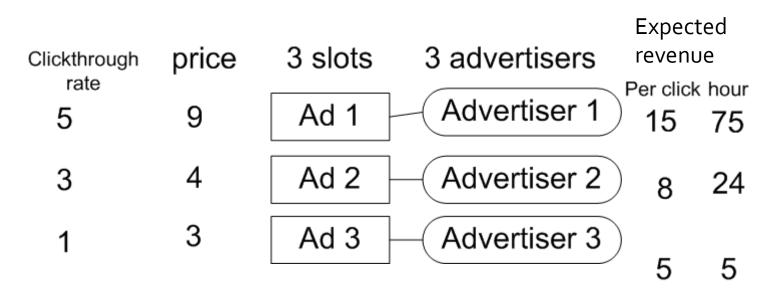
- 3 slots and 3 advertisers
- Clickthrough rate



- 3 slots and 3 advertisers
- Clickthrough rate, expected revenue per click



- 3 slots and 3 advertisers
- Clickthrough rate, expected revenue per click
- Valuation = Expected revenue can be calculated = clickthrough rate \* expected revenue per click



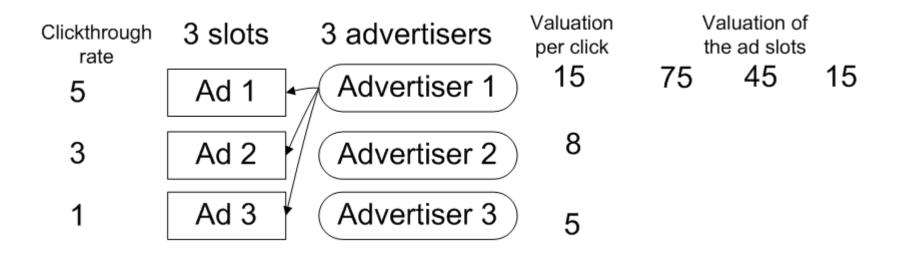
- 3 slots and 3 advertisers
- Clickthrough rate, expected revenue per click
- Valuation = Expected revenue can be calculated = clickthrough rate \* expected revenue per click
- Assume we know the price

Clickthrough	price	3 slots	3 advertisers	Expect revenu		Expected payoff
rate 5	9	Ad 1	Advertiser 1	Per click	hour 75	Per click
3	4	Ad 2	Advertiser 2	8	24	4
1	3	Ad 3	Advertiser 3	5	5	2

- 3 slots and 3 advertisers
- Clickthrough rate, expected revenue per click
- Valuation = Expected revenue can be calculated = clickthrough rate \* expected revenue per click
- Assume we know the price, expected payoff per click

Ol' el Herenad		0 -1-1-		Expected revenue		Expected payoff	
Clickthrough rate	price	3 slots	3 advertisers	Per click	hour	Per click	hour
5	9	Ad 1	Advertiser 1	15	75	6	30
3	4	Ad 2	Advertiser 2	8	24	4	12
1	3	Ad 3	Advertiser 3	5	5	2	2

- Advertiser's valuation (per unit time): clickthrough rate \* expected revenue per click
- Advertiser's payoff: (expected revenue price) \* clickthrough rate

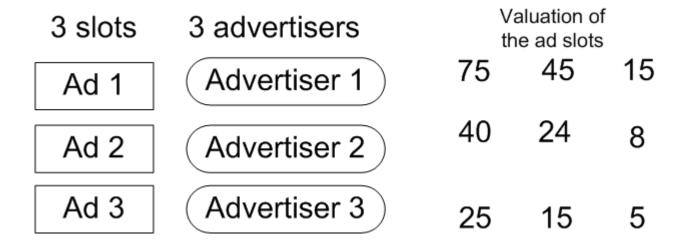


 Different slots have different clickthrough rates, so their values are different

Clickthrough rate	3 slots	3 advertisers	Valuation per click			
5	Ad 1	Advertiser 1	15	75	45	15
3	Ad 2	Advertiser 2	8	40	24	8
1	Ad 3	Advertiser 3	5	25	15	5

- Different slots have different clickthrough rates, so their values are different
- Different advertiser's valuation on per click is different, so the valuation of each ad slot is also different

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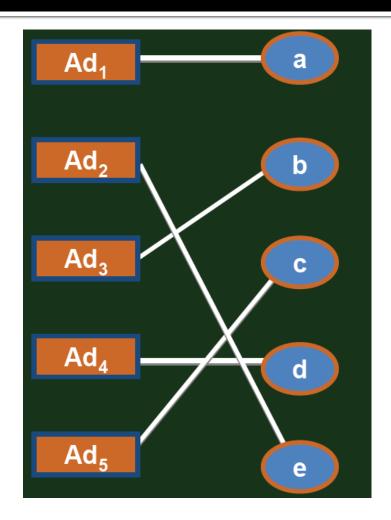
- Seems familiar?
- We have formulated the problem as a matching market

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### Pricing the ad slots or the clicks

# Question 3: how should websites set the prices on the ad slots

- Matching market: market clearing prices
- Vickrey-Clarke-Groves (VCG) mechanism
- Generalized Second-Price Auction (GSP)



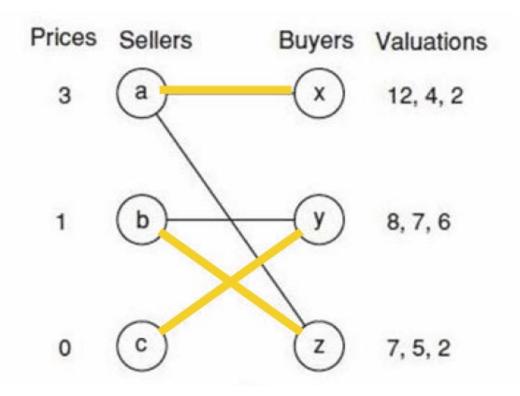
### A summary

- General backgrounds on online ads
- Some general concepts
  - Ad slots, clickthrough rate, advertiser's revenue, advertiser's valuation, advertiser's payoff, etc
- The advertisement market
  - Payment methods
  - Ad slots allocation
  - Pricing the ad slots

# Price setting: Matching of advertiser and ad slots

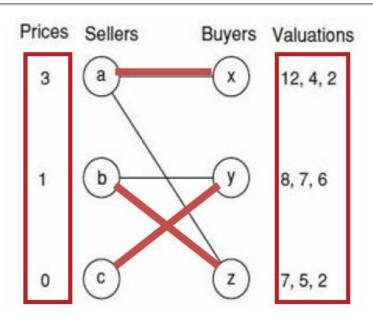
# **Matching market**

Matching market: market clearing prices



#### Recall some basic elements

- A set of sellers and buyers
- Every buyer has a valuation on every object from the seller
- Every seller has a price for each object
- Every buyer try to maximize his payoff
- If the preferred sellers graph has a perfect match, then this is an optimal matching, the price in optimal matching is the market-clearing price



Preferred sellers graph

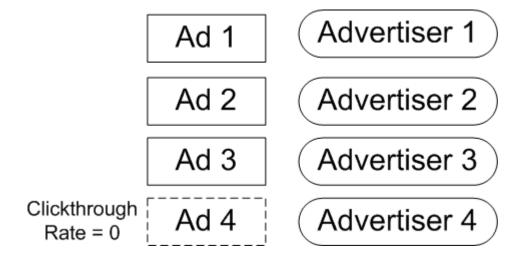
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## The good property of market-clearing prices

- Market-clearing prices satisfy: optimal matching, maximize the total valuations, maximize the total payoff
- If we know the valuations of the advertisers, every seller uses the market-clearing prices for each slot

## The good property of market-clearing prices

- If there are more advertisers than ad slots, we can simply add one slot with clickthrough rate = o
- 4 advertisers competing for 3 ad slots



## The limitation of market matching

 We have an assumption: every advertiser needs to announce their valuations on the ad slots

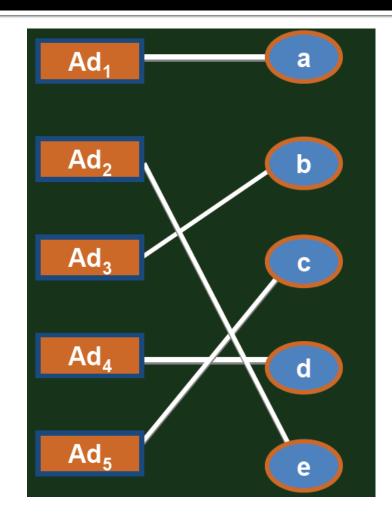
This is not realistic (this can be commercial secrets)

Is there anything else we can do?

# Review: pricing the ad slots or the clicks

# Question 3: how should websites set the prices on the ad slots

- Matching market: market clearing prices
- Vickrey-Clarke-Groves (VCG) mechanism
- Generalized Second-Price Auction (GSP)



# Vickrey-Clarke-Groves (VCG) mechanism

### **VCG** mechanism

- When the advertiser cannot articulate his valuations on the ad slots, e.g., because of secrecy
- Encouraging Truthful Bidding in Matching Markets: The VCG Principle
- This is primarily an auction,
- i.e., if you don't want to say, let's bid

## Which auction format to choose?

- First price auction?
  - There are many disadvantages for first price auction
- Can we do second-price auction?
  - We know that for a single item, second-price auction has very good property: truthful bid
- VCG is an extension of single item second-price auction to multiple items

## Review: Second-price sealed-bid auctions

#### An object

- Different people may have different values for it, v1, v2, ..., vk. These are true value/intrinsic value, i.e., player i will pay at most vi.
- Players don't know other's true values
- Every one has a bid, assume b1 > b2 > ... > bk
- By the rule of second-price auction, the payoff of the highest bidder is v1 – b2 and others are 0; here, b1 is the highest bid, so v1 is the true value
- Truthful bid is a dominant strategy
- Harm due to the existence of the winner

## An example

slots	advertisers	valuations	prices	slots	advertisers	valuations
a	x	30, 15, 6	13	а	×	30, 15, 6
b	У	20, 10, 4	3	b	У	20, 10, 4
С	Z	10, 5, 2	0	c	Z	10, 5, 2

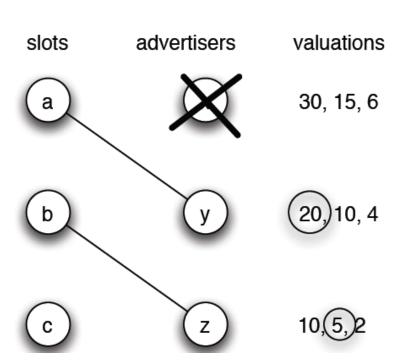
3 slots a, b, c, 3 advertisers x, y z, the valuations are as shown. Using market matching, we can see that the optimal matching is a-x, b-y, c-z, with a total valuation of 42

## An example

If x weren't there, y would do better by 20-10=10, and z would do better by 5-2=3

The total harm x caused is 13

This is the price x should pay



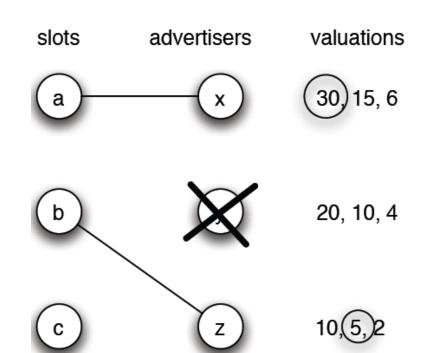
## An example

If y weren't there, x would be unaffected, and z would do better by 5-2=3

The total harm y caused is 3

This is the price y should pay

If z weren't there, z cause no harm to x and y



# Formulating the "harm"

Let  $V_B^S$  be the maximum total valuation over all possible perfect matchings

 $V_{B-j}^{S-i}$  be the value of the optimal matching when we take out buyer j and seller i

 $V_{B-j}^{S}$  be the value of the optimal matching when we take out buyer j

Let 
$$p_{ij} = V_{B-j}^{S} - V_{B-j}^{S-i}$$

This is how everyone do without j (where i is there for sell), and how everyone do with j (where i is taken by j)

# The VCG price setting procedure

- Ask buyers to announce valuations for the items (these announcement need not be truthful)
- Choose a socially optimal assignment of items to buyers, i.e., a perfect matching that maximizes the total valuation of each buyer for what they get.
- Charge each buyer the appropriate VCG price: if buyer j receives item i under optimal matching, then charge j a price pij (see pij in the last slide)

# Analyze the VCG procedure

Claim: If items are assigned and prices computed according to the VCG procedure, then truthfully announcing valuations is a dominant strategy for each buyer, and the resulting assignment maximizes the total valuation of any perfect matching of slots and advertisers.

- Two good properties:
  - Maximize total valuation
  - Truthful-telling is a dominant strategy

- Claim: If items are assigned and prices computed according to the VCG procedure, then truthfully announcing valuations is a dominant strategy for each buyer, and the resulting assignment maximizes the total valuation of any perfect matching of slots and advertisers.
- There are two parts. The second part on maximizing total valuation is easier: as long as they report their valuation truthfully, the total valuation is maximized as the matching is perfect matching

- For the first part, truthful-telling, we want to show that no one has the incentive to deviate
- That is:
- If buyer j bids true valuation and gets slot i, his payoff is

$$v_{ij} - p_{ij}$$

Assume he changes his bid and gets slot h, his payoff is

$$v_{hj} - p_{hj}$$

We need to show

$$v_{ij} - p_{ij} \ge v_{hj} - p_{hj}$$

We can translate

$$v_{ij} - p_{ij} \ge v_{hj} - p_{hj}$$

to

$$v_{ij} - [V_{B-j}^S - V_{B-j}^{S-i}] \ge v_{hj} - [V_{B-j}^S - V_{B-j}^{S-h}]$$

which is

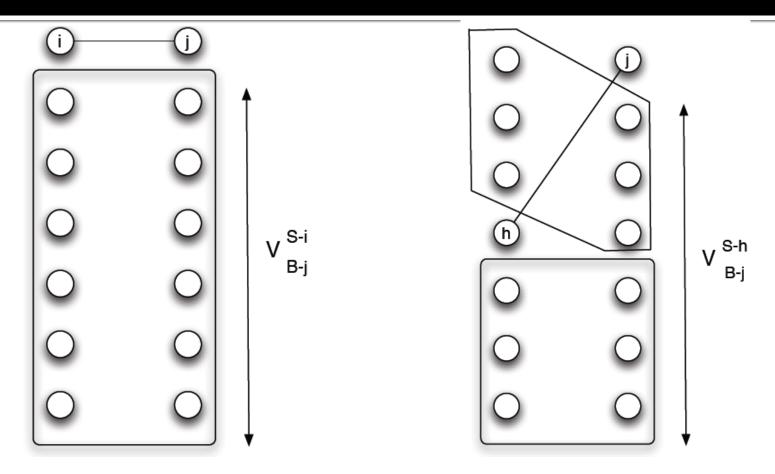
$$v_{ij} + V_{B-j}^{S-i} \ge v_{hj} + V_{B-j}^{S-h}$$

We know that

$$v_{ij} + V_{B-j}^{S-i} = V_B^S$$

and

$$v_{hj} + V_{B-j}^{S-h} \le V_B^S$$



(a)  $v_{ij} + V_{B-j}^{S-i}$  is the maximum valuation of any matching.

(b)  $v_{hj} + V_{B-j}^{S-h}$  is the maximum valuation only over matchings constrained to assign h to j.

# Price setting: Generalized Second-Price Auction (GSP)

# The GSP price setting procedure

- Each advertiser j announces a bid consisting of a single number bj, i.e., the price he is willing to pay per click
- After each advertiser submits a bid, GSP awards each slot i to the ith highest bidder, at a price per click equal to the (i+1)st highest bid
- In other words, each advertiser who is shown on the results page is paying a price per click equal to the bid of the advertiser just below them

# The GSP price setting procedure

- n slots: click rates r1, r2, ..., rn, let's assume in descending order
- n advertisers: bids, b1, b2, ..., bn, let's assume in descending order
- Assign r1 to the 1<sup>st</sup> advertiser, pay by b2, assign r2 to the 2<sup>nd</sup> advertiser, pay by b3, ...
- We can create an artificial advertiser to make this balance
- Set the pay as o or a "club joining fee"

## Some basic elements of GSP

- GSP's CPC (cost per click): b<sub>i+1</sub>
- Payoff = vi \* ri b<sub>i+1</sub> \* ri
- vi is the valuation of advertiser i

# **Analysis of GSP**

- Truth-telling may not be the dominant strategy for GSP
- In VCG, truth-telling is the dominant strategy

# An example for not truth-telling

clickthrough rates 10	slots	advertisers	revenues per click 7
4	b	У	6
0	С	Z	1

- 3 slots, a, b, c, clickthrough rate 10, 4, o
- 3 advertisers, x, y, z, every click produces a revenue of 7,
   6, 1 for them

# An example for not truth-telling

clickthrough rates 10	slots	advertisers	revenues per click 7
4	b	У	6
0	С	Z	1

- If truth-telling, x will get slot a, pay at a price of 6, his payoff will be 7\*10 - 6\*10 = 10
- If x bid 5, what will happen?

# An example for not truth-telling

clickthrough rates 10	slots	advertisers	revenues per click 7
4	b	У	6
0	C	Z	1

- If truth-telling, x will get slot a, pay at a price of 6, his payoff will be 7\*10 - 6\*10 = 10
- If x bid 5, what will happen?
- x will get slot b, pay 1, payoff at 7\*4-1\*4 = 24

clickthrough rates 10	slots	advertisers	revenues per click 7	5
4	b	У	6	4
0	С	Z	1	2

- The bids of 5, 4, 2 form a Nash Equilibrium
- Now, x gets a, with payoff 7\*10 4\*10 = 30. If x bids below 4, he will get b, with payoff 7\*4 2\*4 = 20
- Now, y gets b, with payoff 6\*4 2\*4 = 16. If y bids above 5, he will get a, with payoff 6\*10 5\*10 = 10. If y bids below 2, he will get c, with payoff 0

clickthrough rates 10	slots	advertisers	revenues per click 7	5
4	b	У	6	4
0	С	Z	1	2

- The bids of 5, 4, 2 form a Nash Equilibrium
- This is also social optimal 70+24+0=94

clickthrough rates 10	slots	advertisers	revenues per click 7	3
4	b	У	6	5
0	С	Z	1	1

- The bids of 3, 5, 1 also form a Nash Equilibrium
- Now, x gets b, with payoff 7\*4 1\*4 = 24. If x bids above 5, he will get a, with payoff 7\*10 5\*10 = 20. If x bids below 1, he gets c, with payoff o
- Now, y gets a, with payoff 6\*10 3\*10 = 30. If y bids below 3, he will get b, with payoff 6\*4 1\*4 = 20.

clickthrough rates 10	slots	advertisers	revenues per click 7	3
4	b	У	6	5
0	С	Z	1	1

- The bids of 3, 5, 1 also form a Nash Equilibrium
- The social welfare is 7\*4 + 6\*10 + 1\*0 = 88
- This is not social optimal

## Some characteristics of GSP

- Truthful-telling is not a dominant strategy
- There can be multiple equilibria, including social optimal and also non-social optimal

# The revenue of the search engine of GSP

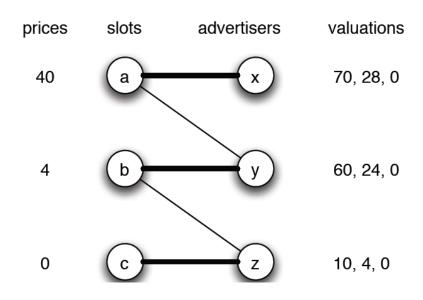
clickthrough rates	slots	advertisers	revenues per click		
10	a	x	7	5	3
4	b	У	6	4	5
0	С	Z	1	2	1

- For the equilibrium of 5, 4, 2: 4\*10 + 2\*4 = 48
- For the equilibrium of 3, 5, 1: 3\*10 + 1\*4 = 34

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# The revenue of the search engine of VCG



- The harm is (60 24) + (4 0) = 44
- VCG may not provide a higher revenue to the search engine (of course, this depends on which equilibrium it ends up to in GSP)

# Why GSP

- Most search engine currently adopts GSP
- VCG targets on social optimal
- VCG is difficult to compute, difficult to understand by many
- Search engine companies care more on revenue, not social optimal
- VCG has assumptions that, no collusion, the price is only related to clickthrough rate, etc

# A summary

- How to auction the advertisement slots
- This is a combination of matching and auction
- VCG has excellent theoretical properties
  - Social optimal
  - Truthful-telling
- GSP (and its variants) has been used widely in practice
  - Easy to understand
  - Search engine may get more