Extending OCaml with copatterns and first-class labels for Free

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1 SYNTAX

Here we describe the syntax of our extension. This one is conflictless with the actual syntax used in OCaml 4.04. In the following section:

- Programs with a yellow background stand for the original code, while blue ones represent generated
- Data constructors indiced with a small cross (such as Stream^x) are inaccessible for the programmer.
- Infix symbol (\triangleright) corresponds to the reverse-application operator such that ($x \triangleright f$) is equivalent to f(x).
- Type constructors codata and query are abstracts and automatically imported from module Pervasives.

1.1 Codata types

Codata types are introduced with the **type** keyword.

```
type \alpha !stream = {
   Head : \alpha;
   Tail : \alpha !stream;
```

```
type (\sigma, \alpha) stream =
    | Stream<sup>x</sup> : { dispatch : \sigma.(\sigma query, \alpha) stream \rightarrow \sigma} \rightarrow (codata, \alpha) stream
                  : (\alpha \text{ query}, \alpha) \text{ stream}
                  : (((codata, \alpha) stream) query, \alpha) stream
let head = function Stream<sup>x</sup> { dispatch } → dispatch Head
let tail = function Stream<sup>x</sup> {dispatch} → dispatch Tail
```

```
type (\alpha, \beta)!product = {
   Fst : \alpha;
   Snd : \beta;
```

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```
type (\sigma, \alpha, \beta) product =

| Product* : {dispatch : \sigma.(\sigma query, \alpha, \beta) product \rightarrow \sigma} \rightarrow (codata, \alpha, \beta) product

| Fst : (\alpha query, \alpha, \beta) product

| Snd : (\beta query, \alpha, \beta) product

let fst = function Product* {dispatch} \rightarrow dispatch Fst

let snd = function Product* {dispatch} \rightarrow dispatch Snd
```

1.2 Copattern matching

```
let zeros = comatch zs: int !stream with

| zs#Head → 0

| zs#Tail → zs
```

```
let zeros =
let rec zs: (codata, int) stream =
let dispatch: type \sigma.(\sigma \text{ query, int}) stream \rightarrow \sigma = \text{function}
| \text{Head} \rightarrow 0
| \text{Tail} \rightarrow \text{zs}
in Stream {dispatch}
```

1.3 Unnesting copatterns

Here, we process more or less as in [?].

```
let fibonacci = comatch fib : int !stream with
  | fib#Head → 0
  | fib#Tail#Head → 1
  | fib#Tail#Tail → zipWith (+) fib fib#Tail
```

```
let fibonacci =
let rec fib: (codata, int) stream =
let f1 =
let dispatch: type \sigma.(\sigma query, int) stream → \sigma = function
| Head → 1
| Tail → zipWith (+) fib (fib ▷ tail)
in Stream<sup>x</sup> {dispatch}
in
let dispatch: type \sigma.(\sigma query, int) stream → \sigma = function
| Head → 0
| Tail → f1
in Stream<sup>x</sup> {dispatch}
in fib
```

1.4 BNF

REFERENCES

[] Anton Setzer, Andreas Abel, Brigitte Pientka, and David Thibodeau. 2014. Unnesting of copatterns. In *International Conference on Rewriting Techniques and Applications*. Springer, 31–45.

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