

Homework 1

1) a) $\lim_{n \rightarrow \infty} \frac{\log_2(n^2+1)}{n} \xrightarrow{\text{l'Hopital's rule}} \lim_{n \rightarrow \infty} \frac{2n}{n^2+1} \xrightarrow{\text{l'Hopital's rule}} \lim_{n \rightarrow \infty} \frac{2}{2n} = 0$

so $\log_2(n^2+1) \in O(n)$: True.

b) $\lim_{n \rightarrow \infty} \frac{\sqrt{n^2+n}}{n} = \lim_{n \rightarrow \infty} \sqrt{\frac{n^2+n}{n^2}} = \lim_{n \rightarrow \infty} \sqrt{1 + \frac{1}{n}} = 1$

so $\sqrt{n^2+n} \in \Omega(n)$: True.

c) $\lim_{n \rightarrow \infty} \frac{n^{n-1}}{n^n} = \lim_{n \rightarrow \infty} \frac{\overbrace{n \cdot n \cdot n \dots n}^{n-1}}{\underbrace{n \cdot n \cdot n \dots n}_n} = \lim_{n \rightarrow \infty} \frac{1}{n} = 0$

so $n^{n-1} \notin \Omega(n^n)$. Statement is false.

d) if $f(n) \in O(g(n))$ then we can say $O(f(n)) \subset O(g(n))$

$\lim_{n \rightarrow \infty} \frac{(2^n + n^2)}{4^n} = \lim_{n \rightarrow \infty} \left(\frac{1}{2} \right)^n + \frac{n^2}{4^n} = \underbrace{\lim_{n \rightarrow \infty} \left(\frac{1}{2} \right)^n}_0 + \underbrace{\lim_{n \rightarrow \infty} \frac{n^2}{4^n}}_0 = 0 + 0 = 0$

so $2^n + n^2 \in O(4^n)$ and $O(2^n + n^2) \subset O(4^n)$. True.

e) $\lim_{n \rightarrow \infty} \frac{2 \log_2 \sqrt{n}}{3 \log_2 n^2} = \lim_{n \rightarrow \infty} \frac{\frac{2}{3} \log_2 n}{6 \log_2 n} \xrightarrow{\text{l'Hopital's rule}} \lim_{n \rightarrow \infty} \frac{\frac{1}{n \ln(2)}}{\frac{1}{n \ln(2)}} = \lim_{n \rightarrow \infty} \frac{\ln 2}{\ln 2} = \frac{\ln 2}{\ln 2}$
 (we can ignore the constants)

Result is smaller than 0. So statement is false.

f) $\lim_{n \rightarrow \infty} \frac{\log_2 \sqrt{n}}{(\log_2 n)^2} = \lim_{n \rightarrow \infty} \frac{\frac{1}{2} \cdot \log_2 n}{\log_2 n \cdot \log_2 n} = \lim_{n \rightarrow \infty} \frac{1}{2 \log_2 n} = 0$

$\log_2 \sqrt{n} \in O(\log_2 n)^2$
 $(\log_2 n)^2 \notin O(\log_2 \sqrt{n})$) Statement is false.

$$2) 10^n > 2^n > n^2 = 8^{\lg n} > n^2 \lg n > n^2 > \sqrt{n} > \lg n$$

Just for fun
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$$\Rightarrow 10^n - 2^n$$

Let a^x, b^x if $a^x > b^x$, $a > b$.

$$\text{so } \underline{10^n > 2^n}$$

$$\Rightarrow 2^n - n^3$$

$$\lim_{n \rightarrow \infty} \frac{n^3}{2^n} = 0 \quad \left(\begin{array}{l} \text{from growth} \\ \text{function definition} \\ (x^a < a^x) \end{array} \right) \quad \text{so } \underline{2^n > n^3}$$

$$\Rightarrow n^3 - 8^{\lg n}$$

$$8^{\lg n} = n^{\lg 8} = n^3 \quad \text{so } \underline{n^3 = 8^{\lg n}}$$

$$\Rightarrow 8^{\lg n} - n^2 \lg n$$

$$8^{\lg n} = n^3$$

$$\lim_{n \rightarrow \infty} \frac{n^2 \lg n}{n^3} = \lim_{n \rightarrow \infty} \frac{\lg n}{n} \xrightarrow{\text{L'Hopital's rule}} \lim_{n \rightarrow \infty} \frac{1}{n} = 0 //$$

$$\text{so } \underline{8^{\lg n} > n^2 \lg n}$$

$$\Rightarrow n^2 \lg n - n^2$$

$$\lim_{n \rightarrow \infty} \frac{n^2}{n^2 \lg n} = \lim_{n \rightarrow \infty} \frac{1}{\lg n} = 0 //$$

$$\underline{n^2 \lg n > n^2}$$

$$\Rightarrow n^2 - \sqrt{n}$$

$$n^2 - n^{1/2}$$

Let $x^a - x^b$ $x^a > x^b$ iff $a > b$

$$\text{so } \underline{n^2 > \sqrt{n}}$$

$$\Rightarrow \sqrt{n} - \lg n$$

$$\lim_{n \rightarrow \infty} \frac{\lg n}{\sqrt{n}} \Rightarrow \text{assume } n = x^2$$

$$\lim_{x \rightarrow \infty} \frac{2 \lg x}{x} = \lim_{x \rightarrow \infty} \frac{2}{x} = 0 //$$

$$\text{so } \underline{\sqrt{n} > \lg n}$$

3) a) In this function the basic operation is comparison operators. Since there is no nested loops or recursive calls in the function complexity only depends on loop. Loops iterates only size of array times. So complexity is $O(n) = \Omega(n) = \Theta(n)$

b) In this function, if we examine the pattern of this loop, whenever even number assigned to i , count is increased 1 in the next iteration. After that loop iterates one more time and calculates $i = (i-1)^*$; statement and generates even i and repeats same pattern. So we assume this as every else statement it increases count. The growth of i is depends on else statement.

Since $i = i^2 - i$ statement increases i exponentially the complexity is $O(\log(\log n))$.

4) a) $\int_0^1 f(x) dx \leq f(x) \leq \int_1^{n+1} f(x) dx$

$\int x^2 \log x dx \Rightarrow \frac{x^3 \log x}{3} - \int \frac{x^2}{3} dx = \frac{x^3 \log x}{3} - \frac{x^3}{9} = \frac{x^3 (3 \log x - 1)}{9} \Rightarrow$ let's put this in a statement

$u v - \int v du$
 $u = \log x \quad v = x^2$
 $u' = \frac{1}{x} \quad v' = \frac{x^2}{3}$

$\frac{x^3 (3 \log x - 1)}{9} \Big|_0^n \leq f(x) \leq \frac{x^3 (3 \log x - 1)}{9} \Big|_1^{n+1}$

$\frac{n^3 (3 \log n - 1)}{9} - \frac{0 (3 \log 0 - 1)}{9} \leq f(x) \leq \frac{(n+1)^3 (3 \log(n+1) - 1)}{9} + \frac{1}{9}$

formula.
 Alternation
 difference.

$$\frac{n^3(\log n - 1)}{9} - \underbrace{\frac{0 \cdot (\log 0 - 1)}{9}}_{\log 0 \text{ is undefined}} \leq f(x) \leq \frac{(n+1)^3(\log(n+1) - 1)}{9} + \frac{1}{9}$$

for this part, let's try this;

$$1 + \int_1^n x^2 \log x \, dx = 1 + \left(\frac{x^3(\log x - 1)}{9} \right) \Big|_1^n = 1 + \frac{n^3(\log n - 1)}{9} + \frac{1}{9}$$

so statement becomes;

$$\frac{n^3(\log n - 1)}{9} + \frac{10}{9} \leq f(x) \leq \frac{(n+1)^3(\log(n+1) - 1)}{9} + \frac{1}{9}$$

so $f(x) \in O(n^3 \log n)$

b) $\int_0^n f(x) \, dx \leq f(x) \leq \int_1^{n+1} f(x) \, dx$

$$\int_0^n x^3 \, dx \leq f(x) \leq \int_1^{n+1} x^3 \, dx$$

$$\frac{x^4}{4} \Big|_0^n \leq f(x) \leq \frac{x^4}{4} \Big|_1^{n+1} \Rightarrow \frac{n^4}{4} \leq f(x) \leq \frac{(n+1)^4}{4} \text{ so } \underline{\underline{f(x) \in O(n^4)}}$$

c) $\int_0^n \frac{1}{2\sqrt{x}} \, dx \leq f(x) \leq \int_1^{n+1} \frac{1}{2\sqrt{x}} \, dx$

$$\sqrt{x} \Big|_0^n \leq f(x) \leq \sqrt{x} \Big|_1^{n+1}$$

$$\sqrt{n} \leq f(x) \leq \sqrt{n+1} - 1 \text{ so } \underline{\underline{f(x) \in O(\sqrt{x})}}$$

$$d) \int_0^1 \frac{1}{x} dx \leq f(x) \leq \int_1^{n+1} \frac{1}{x} dx$$

$$\ln x \Big|_0^1 \leq f(x) \leq \ln(x) \Big|_1^{n+1}$$

$$\underbrace{\ln 0 - \ln 1}_{\text{undefined?}} \leq f(x) \leq \ln(n+1)$$

alternatively for

$$1 + \int_1^x \frac{1}{x} dx \leq f(x)$$

$$1 + \ln x \Big|_1^x \leq f(x)$$

$$\ln(n)+1 \leq f(x) \leq \ln(n+1) \quad \text{so } \underline{\underline{f(x) \in O(\ln(n))}}$$

5) Best case: first element is searching element.

$$L[1] = x \quad B(n) \in O(1)$$

Worst case: Element does not in the list or $L[n] = x$ (last element)

$W(n) \in O(n)$ which n is length of list.

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