

Sekuritas Jaringan

Pertemuan 4

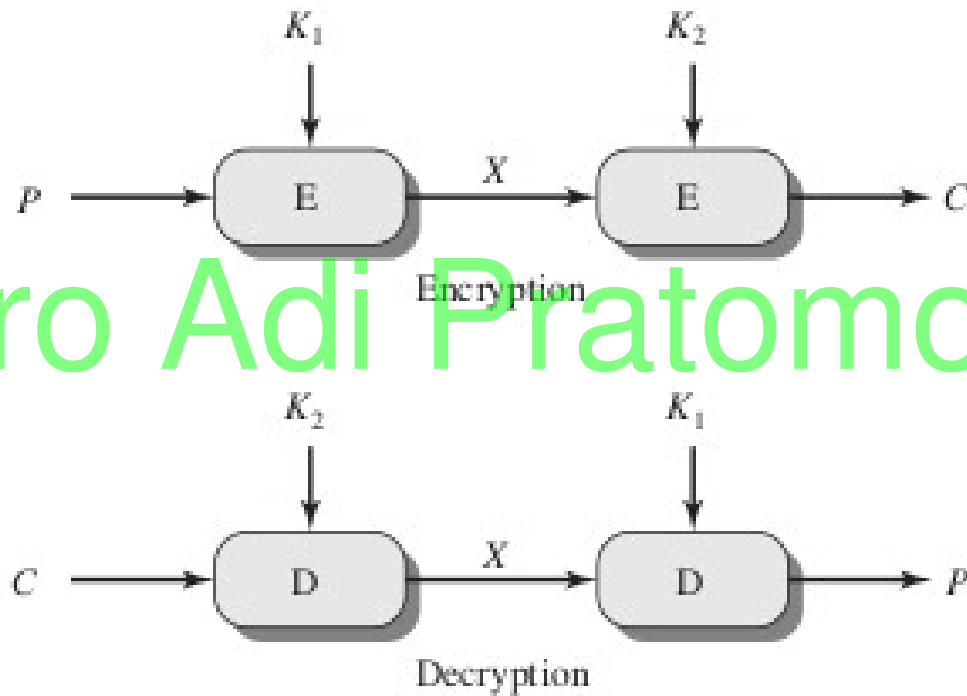
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Double DES & Triple Des

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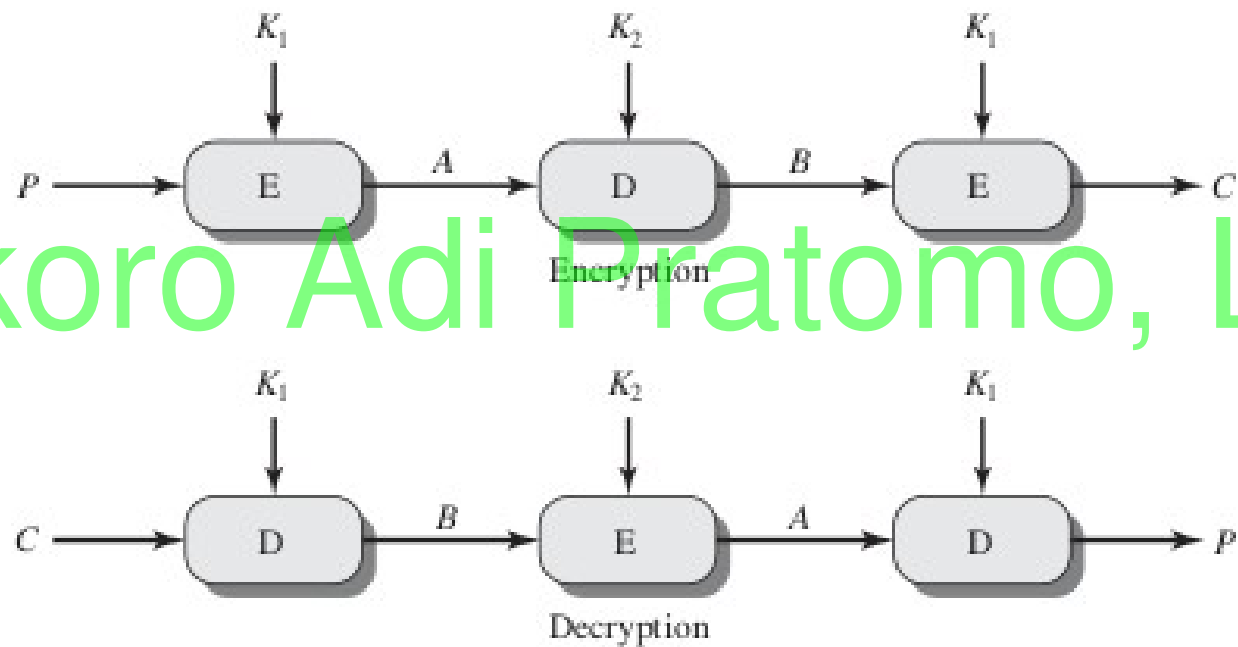
Double DES



(a) Double encryption

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Triple DES



(b) Triple encryption

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Pseudorandom Number Generator

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One Time Pad

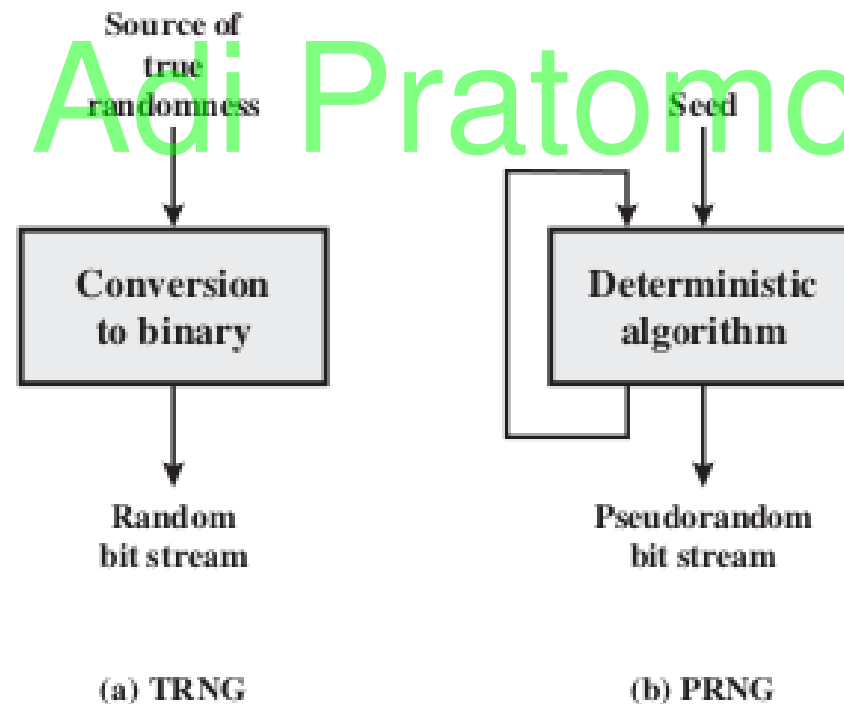
- Penggunaan kunci dengan isi yang acak dan sepanjang pesan, serta digunakan untuk sekali enkripsi

	H	E	L	L	O	message
	7 (H)	4 (E)	11 (L)	11 (L)	14 (O)	message
+	23 (X)	12 (M)	2 (C)	10 (K)	11 (L)	key
=	30	16	13	21	25	message + key
=	4 (E)	16 (Q)	13 (N)	21 (V)	25 (Z)	message + key (mod 26)
	E	Q	N	V	Z	→ ciphertext

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Generate Random Number

- True Random Number Generator (TRNG)
- Pseudorandom Number Generator (PRNG)



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PRNG Requirements

- Randomness
 - Uniformity
 - Scalability
 - Consistency
- Unpredictability
 - Forward unpredictability
 - Backward unpredictability
- Seed Requirement

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PRNG Algorithm

- Purpose Built Algorithm
 - Linear Congruential Generator
 - Blum Blum Shub Generator
- Algorithm based on existing cryptographic algorithm
 - Symmetric Cipher
 - Asymmetric Cipher
 - Hash Function & MAC

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Linear Congruential Generators

- $X_{n+1} = (aX_n + c) \bmod m$
- Dimana :
 - m : modulus : $m > 0$
 - a : multiplier : $0 < a < m$
 - c : increment : $0 \leq c < m$
 - X_0 : Seed / nilai awal : $0 \leq X_0 < m$

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Linear Congruential Generators (2)

- Jika :
 - $A = C = 1$
 - $A = 7, C = 0, M = 32, X_0 = 1$
 - $A = 5, C = 0, M = 32, X_0 = 1$
 - $A = 7^5, C = 0, M = 2^{31}, X_0 = 1$
- Bagaimana random number yang dihasilkan?

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Blum Blum Shub

- Cari dua bilangan p dan q , dimana :
 - $p \bmod 4 = q \bmod 4 = 3$
- $n = p \times q$
- Pilih random number s , dimana :
 - n dan s adalah relatively prime
 - $\text{GCD}(n, s) = 1$
- Generator :

$$X_0 = s^2 \bmod n$$

$$\text{for } i = 1 \text{ to } \infty$$

$$X_i = (X_{i-1})^2 \bmod n$$

$$B_i = X_i \bmod 2$$

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Blum Blum Shub (2)

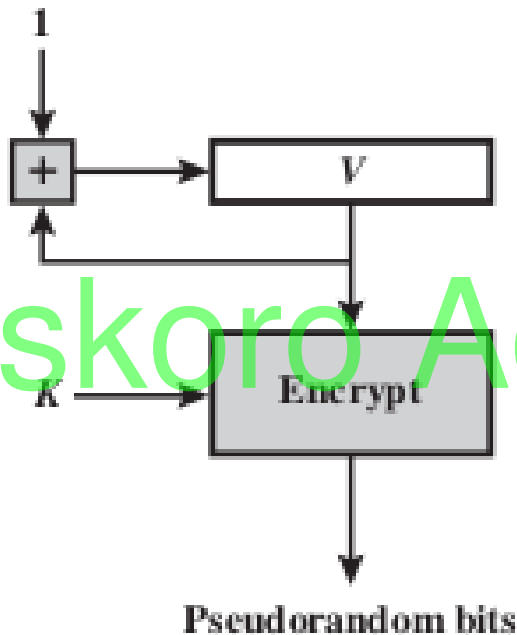
- Contoh :
 - Jika $p = 383$, $q = 503$, $s = 101355$, hitung B_1

Table 7.1 Example Operation of BBS Generator

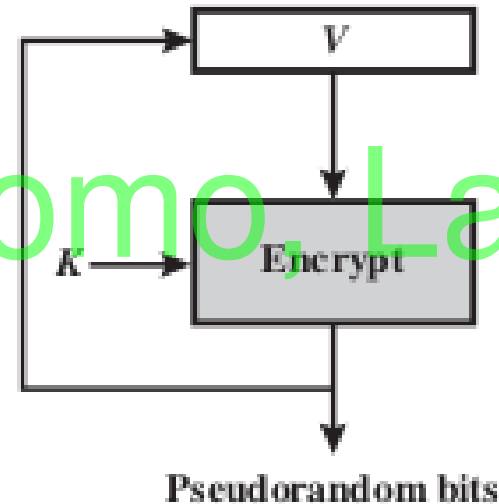
i	X_i	B_i
0	20749	
1	143135	1
2	177671	1
3	97048	0
4	89992	0
5	174051	1
6	80649	1
7	45663	1
8	69442	0
9	186894	0
10	177046	0

i	X_i	B_i
11	137922	0
12	123175	1
13	8630	0
14	114386	0
15	14863	1
16	133015	1
17	106065	1
18	45870	0
19	137171	1
20	48060	0

PRNG Using Block Cipher



(a) CTR mode



(b) OFB mode

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ANSI X9.17 PRNG

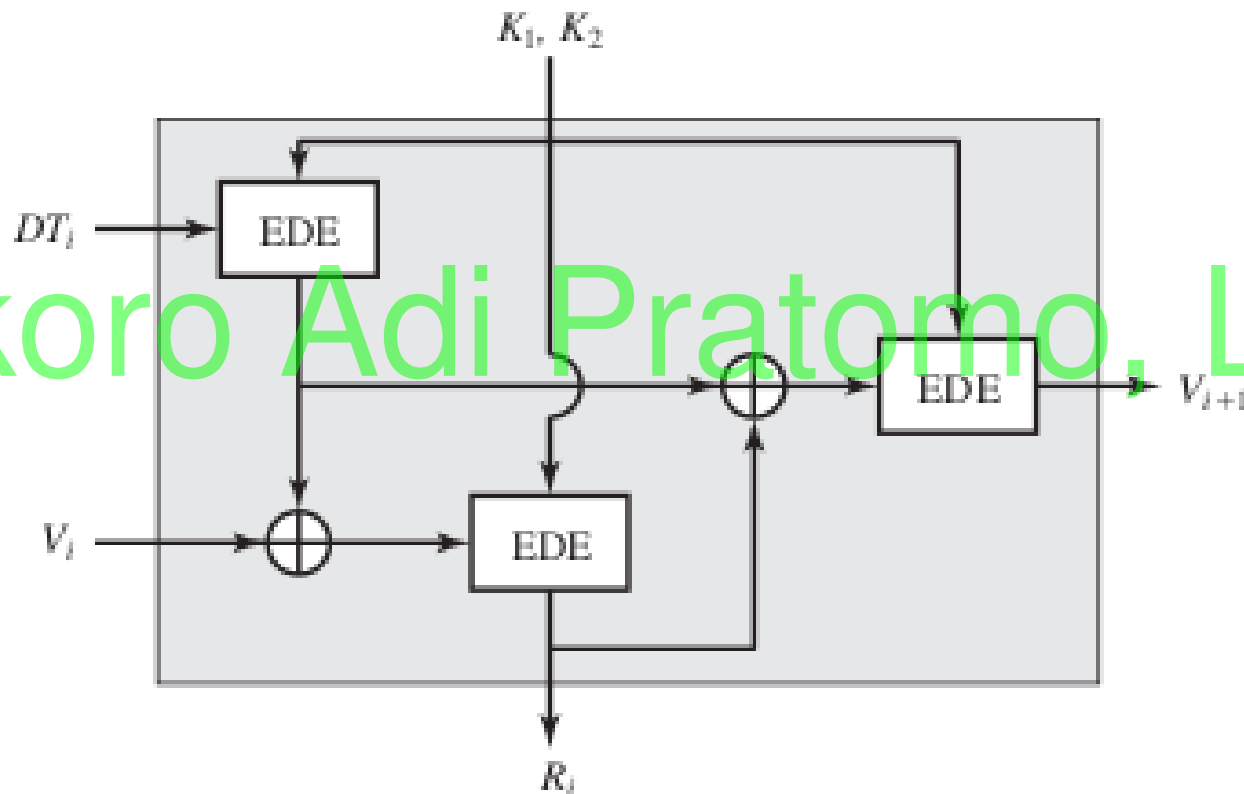


Figure 7.4 ANSI X9.17 Pseudorandom Number Generator

Stream Cipher

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Stream Cipher

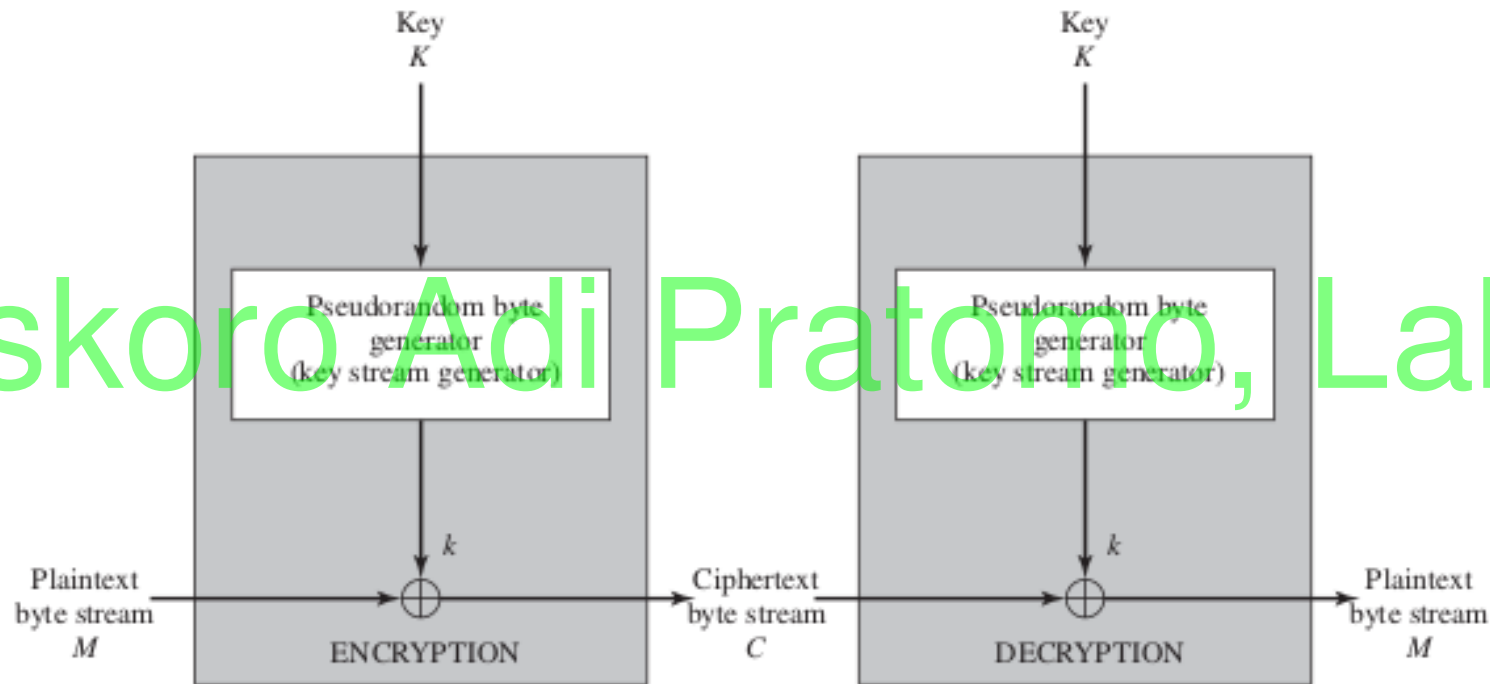


Figure 7.5 Stream Cipher Diagram

Perbandingan Kecepatan

Table 7.4 Speed Comparisons of Symmetric Ciphers on a Pentium II

Cipher	Key Length	Speed (Mbps)
DES	56	9
3DES	168	3
RC2	Variable	0.9
RC4	Variable	45

RC4

- Dibuat pada 1987
- Key size : 1-256 bytes
- Operasi : per-byte

- Digunakan di :

- Secure Socket Layer
- WEP
- WPA

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RC4 Algorithm

- Siapkan 256 byte state vector S
 - $S[0]=0, S[1]=1, S[2]=2, \dots, S[255]=255$

- Inisialisasi :

- ```
for i = 0 to 255 do
 S[i] = i;
 T[i] = K[i mod keylen];
```

- Permutasi Awal untuk S :

- ```
j = 0;
for i = 0 to 255 do
    j = (j + S[i] + T[i]) mod 256;
    Swap (S[i], S[j]);
```

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RC4 Algorithm (2)

- Keystream Generation

- $i, j = 0;$

- while (true)

- $i = (i + 1) \bmod 256;$

- $j = (j + S[i]) \bmod 256;$

- Swap ($S[i], S[j]$);

- $t = (S[i] + S[j]) \bmod 256;$

- $k = S[t];$

- Enkripsi :

- $k \text{ XOR data}$

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TUGAS PROGRAMMING

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DISKUSI