

CENG 424
Logic For Computer Science
Fall 2024-2025
Assignment 5

Lök, Yusuf Sami
e2521748@ceng.metu.edu.tr

January 19, 2025

Question 1

Relation Definitions

transition(Q_0, S, Q_1, W): This relation represents a transition from state Q_0 to Q_1 with weight W by consuming symbol S .

path(Q_e, T, S, W): This relation represents the existence of a path with a list of transitions T , ending at state Q_e , with remaining string S , and cost of weight W .

Relational Logic Formulation

1. **Initialize relational logic items:**

Create an empty list, `ListOfRelationalLogicItems`, to store all relational logic items.

2. **Initialize the starting path:**

Define the starting path with the initial state s and the input string S :

$$\text{path}(s, \text{nil}, S, 0)$$

3. **Define transitions:**

For each transition t in δ (transition relation of the NWFSM):

- (a) Find the corresponding weight w of transition t using w in the NWFSM.
- (b) Add the transition to `ListOfRelationalLogicItems`:

$$\text{transition}(Q_0, S, Q_1, w)$$

4. **Add the path extension rule:**

If there exists a transition from Q_0 to Q_1 with symbol `sym` and weight w_0 , and a path from Q_0 with remaining string `sym.r` and weight w , then the following relation holds:

$$\text{transition}(Q_0, \text{sym}, Q_1, w_0) \wedge \text{path}(Q_0, P_0, \text{sym.r}, w) \implies \text{path}(Q_1, P_0 \cup \{(Q_0, \text{sym}, Q_1)\}, \text{r}, w_0 + w)$$

5. **Add the goal statement:**

The goal is defined as the path ending at a state Q with an empty remaining string and the highest total weight W :

$$\text{goal}(P, W) \iff \text{path}(Q, P, \text{nil}, W) \wedge \forall_{Q'} \forall_{P'} (\text{path}(Q', P', \text{nil}, W') \implies W \geq W')$$

6. **Return the list:**

Return `ListOfRelationalLogicItems`.

Question 2

Finding the Highest Weighted Path for abb on NWFSM N_1

1. **Add all transitions:**

From the given NWFSM N_1 , the transitions are:

$\text{transition}(q_0, a, q_1, 4), \quad \text{transition}(q_0, a, q_3, 3),$
 $\text{transition}(q_0, b, q_2, 2), \quad \text{transition}(q_0, b, q_3, 4),$
 $\text{transition}(q_1, \epsilon, q_3, 2), \quad \text{transition}(q_1, b, q_1, 4),$
 $\text{transition}(q_2, a, q_0, 2), \quad \text{transition}(q_2, a, q_3, 4),$
 $\text{transition}(q_3, \epsilon, q_2, 7), \quad \text{transition}(q_3, b, q_1, 1)$

2. **Add initial path:**

Define the starting path:

$\text{path}(q_0, \text{nil}, \text{abb}, 0)$

3. **Add path extension rule:**

Using the transitions defined, apply the path extension rule:

$\text{transition}(Q_0, \text{sym}, Q_1, w_0) \wedge \text{path}(Q_0, P_0, \text{sym.r}, w) \implies \text{path}(Q_1, P_0 \cup \{(Q_0, \text{sym}, Q_1)\}, r, w_0 + w)$

4. **Process the string abb using resolution:**

Premises:

R1: $\text{transition}(Q_0, \text{sym}, Q_1, w_0) \wedge \text{path}(Q_0, P_0, \text{sym.r}, w) \implies \text{path}(Q_1, P_0 \cup (Q_0, \text{sym}, Q_1), r, w_0 + w)$

Goal rule:

G1: $\text{goal}(P, W) \iff \text{path}(Q, P, \text{nil}, W) \wedge$
 $\forall Q' \forall P' (\text{path}(Q', P', \text{nil}, W') \implies W \geq W')$

Derivation:

- | | |
|--|--|
| 1. $\text{transition}(q_0, a, q_1, 4)$ | [Premise] |
| 2. $\text{transition}(q_0, a, q_3, 3)$ | [Premise] |
| 3. $\text{transition}(q_1, b, q_1, 4)$ | [Premise] |
| 4. $\text{transition}(q_3, b, q_1, 1)$ | [Premise] |
| 5. $\text{transition}(q_1, \epsilon, q_3, 2)$ | [Premise] |
| 6. $\text{transition}(q_3, \epsilon, q_2, 7)$ | [Premise] |
| 7. $\text{path}(q_0, \text{nil}, \text{abb}, 0)$ | [Initial fact] |
| 8. $\text{path}(q_1, (q_0, a, q_1), \text{bb}, 4)$ | [R1: 7,1] |
| 9. $\text{path}(q_3, (q_0, a, q_3), \text{bb}, 3)$ | [R1: 7,2] |
| 10. $\text{path}(q_1, (q_0, a, q_1), (q_1, b, q_1), b, 8)$ | [R1: 8,3] |
| 11. $\text{path}(q_1, (q_0, a, q_3), (q_3, b, q_1), b, 4)$ | [R1: 9,4] |
| 12. $\text{path}(q_1, P_1, \text{nil}, 12)$ | [R1: 10,3] where $P_1 = (q_0, a, q_1), (q_1, b, q_1), (q_1, b, q_1)$ |
| 13. $\text{path}(q_1, P_2, \text{nil}, 8)$ | [R1: 11,3] where $P_2 = (q_0, a, q_3), (q_3, b, q_1), (q_1, b, q_1)$ |
| 14. $\text{path}(q_3, P_3, \text{nil}, 14)$ | [R1: 12,5] where $P_3 = P_1 \cup (q_1, \epsilon, q_3)$ |
| 15. $\text{path}(q_2, P_4, \text{nil}, 21)$ | [R1: 14,6] where $P_4 = P_3 \cup (q_3, \epsilon, q_2)$ |

Weight Derivation using Goal Rule:

Let's compare all paths ending with **nil**:

- (a) $\text{path}(q_1, P_1, \text{nil}, 12)$ from line 12
- (b) $\text{path}(q_1, P_2, \text{nil}, 8)$ from line 13

(c) `path(q3, P3, nil, 14)` from line 14

(d) `path(q2, P4, nil, 21)` from line 15

Applying G1, comparing weights: $21 > 14 > 12 > 8$

Therefore: `goal(P4, 21)`

The highest weighted path is: $(q_0, a, q_1), (q_1, b, q_1), (q_1, b, q_1), (q_1, \epsilon, q_3), (q_3, \epsilon, q_2)$ with total weight 21.

Question 3

I have chosen the Prolog programming language for this purpose.

1. Changes to the Output Format:

- The transitions should be represented as facts, such as:

`transition(q0, a, q1, 4).`

- The paths and weight relations should be defined as rules (using `:-` in Prolog) that can be recursively processed to evaluate possible paths.
- For empty transitions, a separate rule needs to be defined to change states without consuming symbols.
- The goal condition must be written as a rule to select the maximum weight path from all possible solutions.

2. How Logic Programming Tools Work:

Logic programming tools, such as Prolog, use a declarative approach, where the problem is expressed in terms of relations and logical rules. These tools automatically:

- Use facts and rules to generate all possible solutions.
- Apply backward and forward chaining to explore different combinations of transitions and paths.
- Solve the problem by searching for a path that satisfies the goal condition.
- Prolog uses a depth-first search (DFS) strategy to traverse the possible paths, and the goal condition is evaluated by unifying the paths and comparing their weights.