

# CENG 280

## Formal Languages and Abstract Machines

Spring 2022-2023

### Homework 3

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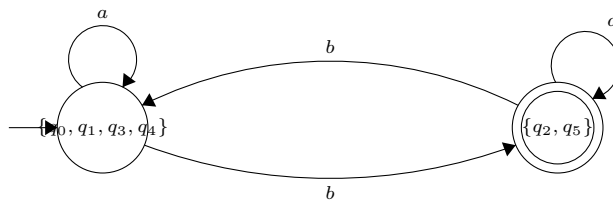
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## Answer for Q1

1. We have first two classes which are the sets that has final states and non-final states. After verifying that the states goes to the same equivalence classes with letters 'a' and 'b', we reached the end of the algorithm. We stopped since we were able to reach the same equivalence class in  $\equiv_1$ .

$$\equiv_0 = \{q_2, q_5\}, \{q_0, q_1, q_3, q_4\}$$

$$\equiv_1 = \{q_2, q_5\}, \{q_0, q_1, q_3, q_4\}$$



2. For  $\{q_2, q_5\}$ , the regular expression is  $L$ .

For  $\{q_0, q_1, q_3, q_4\}$ , the regular expression is  $(Lba^* \cup a^*)$ . We can loop with a or empty string or we can reach this state by passing through the  $\{q_2, q_5\}$ .  $L$  includes all strings accepts the language, then we can reach the  $\{q_0, q_1, q_3, q_4\}$  with a "b" from  $\{q_2, q_5\}$ . Then we have a "a" loop. Therefore, it is  $(Lba^* \cup a^*)$

3. In order to show that  $L'$  is not regular by using Myhill-Nerode theorem we need to prove that  $L'$  has infinite number of equivalence classes. For all distinct  $m+n$  values, there should be 1 distinct equivalence class because each distinct  $m+n$  value should have a distinct  $k+2u$  value. Therefore, there are infinite number of equivalence classes.

$$m+n = 0 \rightarrow k+2u=0$$

$$m+n = 1 \rightarrow k+2u=1$$

$$m+n = 2 \rightarrow k+2u=2$$

$$m+n = 3 \rightarrow k+2u=3$$

.....

.....

Therefore, for each  $m+n$  value there exist a unique distinct  $k+2u$  value. Since there are infinitely many equivalence classes,  $L'$  is not regular. To illustrate that:

$$[e] \not\approx [a] \not\approx [aa] \not\approx [aaa] \not\approx [aaaa] \dots$$

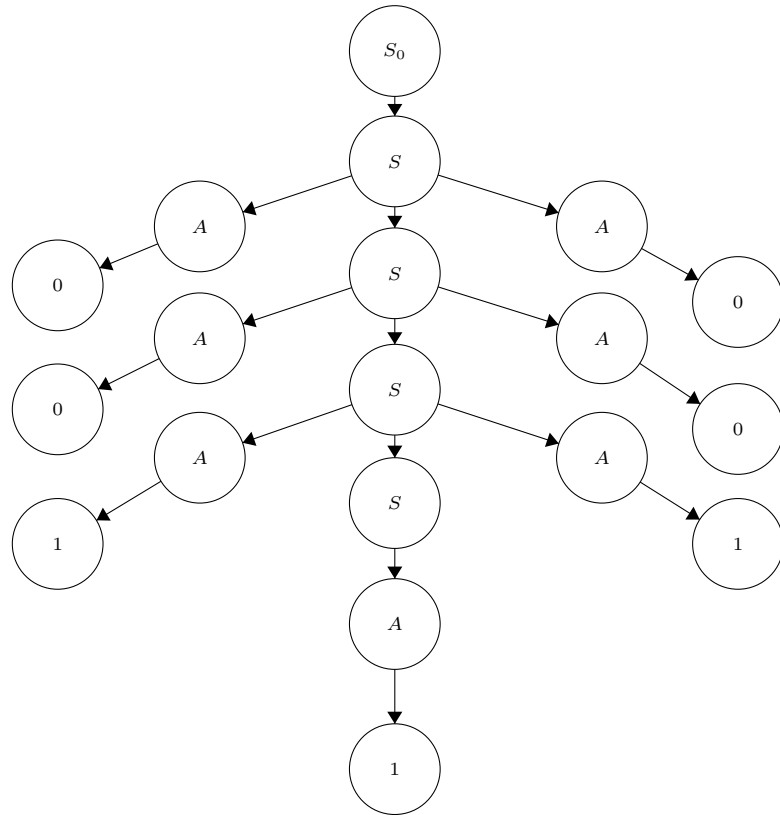
## Answer for Q2

1.  $G = (V_1, \Sigma, R_1, S_0)$  where  $V_1 = \{S_0, B, D\} \cup \Sigma$   
 $\Sigma = \{a, b\}$  and  
 $R_1 = \{S_0 \rightarrow Bb|DS_0|S_0DB,$   
 $D \rightarrow bDa|aDb|DD|e,$   
 $B \rightarrow e|Bb\}$

2.  $G = (V_1, \Sigma, R_1, S_0)$  where  $V_1 = \{S_0, S_1, S_2\} \cup \Sigma$   
 $\Sigma = \{0, 1, 2\}$  and  
 $R_1 = \{S_0 \rightarrow S_1S_2,$   
 $S_1 \rightarrow 0S_11|e,$   
 $S_2 \rightarrow 1S_22|e\}$

3.  $G = (V_1, \Sigma, R_1, S_0)$  where  $V_1 = \{S_0, S, A\} \cup \Sigma$   
 $\Sigma = \{0, 1\}$  and  
 $R_1 = \{S \rightarrow A \mid ASA,$   
 $S_0 \rightarrow S$   
 $A \rightarrow 0|1\}$

Parse tree for string 0011100 is:



## Answer for Q3

1. The (context-free) language is that accepts strings with the same first and last characters with at least two characters, as well as the empty string.  $\{w \mid \text{the first and last characters of } w \text{ are the same and } |w| \geq 2 \text{ or } w = \epsilon, \text{ and } w \in \{0, 1\}^*\}$
2. The (context-free) language is that accepts strings that contain at least two ones.  $\{w \mid w \text{ contains at least two ones, and } w \in \{0, 1\}^*\}$