## 5. Longest Palindromic Substring

Given a string s, return the longest palindromic substring in s.

```
from typing import List
class Solution:
    def longestPalindrome(self, s: str) -> str:
       result = ""
        for i in range(len(s)):
           1, r = i, i
            while l \ge 0 and r < len(s) and s[l] == s[r]:
                if (r - l + 1) > len(result):
                    result = s[l:r + 1]
                r += 1
           1, r = i, i + 1
            while l \ge 0 and r < len(s) and s[l] == s[r]:
                if (r - l + 1) > len(result):
                   result = s[1:r + 1]
                1 -= 1
               r += 1
        return result
```

Example 1:

Input: s = "babad" Output: "bab"

Explanation: "aba" is also a valid answer.

Example 2:

Input: s = "cbbd" Output: "bb"



```
from typing import List
class Solution:
   def longestPalindrome(self, s: str) -> str:
       result = ""
       for i in range(len(s)):
            for l, r in ((i,i), (i,i+1)):
                while l \ge 0 and r < len(s) and s[l] == s[r]:
                    if (r - l + 1) > len(result):
                       result = s[l:r + 1]
                    1 -= 1
                    r += 1
       return result
```

"abbad" 5 for i in range(len(s)):

```
first, check when I and r point to same index
1, r = i, i
                                       while 1 and r pointers are in bound,
                                        and s[l] and s[r] are the same
while l \ge 0 and r < len(s) and s[l] == s[r]:
                           (r-1+1) is the size of the substring,
                           if (r-1+1) is greater than length of result,
                           then save the new result
          if (r - l + 1) > len(result):
                 result = s[1:r + 1]
          1 -= 1
          r += 1
```

```
abbad
                                                   abbad
next, check when 1 and r point to adjacent index
1, r = i, i + 1
```

```
1 and r pointers are in bound, and s[l] and s[r] are palindromes
while l \ge 0 and r < len(s) and s[l] == s[r]:
                            (r-1+1) is the size of the substring,
                            if (r-1+1) is greater than length of result,
                            then save the new result
       if (r - l + 1) > len(result):
              result = s[1:r + 1]
      1 -= 1
      r += 1
```

```
"abbad" 5
               checks for situation where
               two adjacent letters are the same
for i in range(len(s)):
                                                                           ((4,4),(4,5))
                                                               ((3,3),(3,4))
                                  ((1,1),(1,2))
                                                ((2,2),(2,3))
        for 1, r in ((i,i), (i,i+1)):
                                          while 1 and r pointers are in bound,
                                          and s[l] and s[r] are the same
                while l \ge 0 and r < len(s) and s[l] == s[r]:
                                           (r-1+1) is the size of the substring,
                                           if (r-1+1) is greater than length of result,
                                           then save the new result
                        if (r - 1 + 1) > len(result):
                               result = s[1:r + 1]
                        1 -= 1
                        r += 1
      1. Loop through the string, s.
```

a. set left (1) and right (r) pointers the same as

check to see if the s[l] and s[r] are the same.

that it accounts for situations where there are

expand left(l=1) and right (r+=1), and check to

two adjacent characters that are the same, and

i, expand left (1-=1) and right (r+=1), and

b. set left (l) as i, and right (r) as i + 1, so

see if the s[1] and s[r] are the same.

2. There are two parts: