

# Lecture 1

Yutong Yan

## 1 The Central Paradigm of Computer Science

- The central *paradigm* in computer science is that an algorithm **A** is good if:
  - **A** runs in *polynomial time* in the input size  $n$ .
  - That is, **A** runs in time  $T(n) = O(n^k)$  for some constant number  $k$ .
    - $T(n) = 100n + 55$
    - $T(n) = \frac{1}{2}n^2 + 999 \log n$
    - $T(n) = 6n^7 + 900000n^2 - \sqrt{n}$
  - An algorithm is *bad* if it runs in exponential time.
    - $T(n) = 2^n + 100n^5$
    - $T(n) = 1.0000000001^n - n^3 - n$
  - An algorithm is *good* if it runs in *polynomial time* in the input size  $n$ .

e.g.

		Input Size $n$		
		10	100	1000
Runtime of Algorithm	$n$	10	100	1000
	$n^2$	100	10000	1000000
	$2^n$	$10^3$	$10^{30}$	$10^{300}$

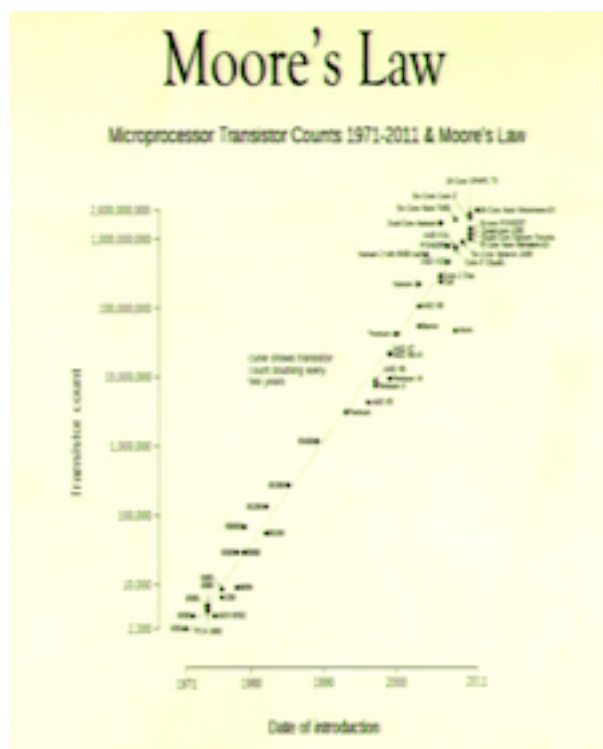
## 2 Good versus Bad (Algorithms)

- For example, consider the problem of sorting  $n$  numbers.
  - A Good Algorithm: **MergeSort** runs in time  $O(n \cdot \log n)$
  - A Bad Algorithm: **BruteForce Search** runs in time  $O(n \cdot n!) \gg 2^n$

### 3 An Equivalent Characterization

- This central *paradigm* has an equivalent formulation
  - $A$  runs in *polynomial time* in the input size  $n$ .
  - The input sizes that  $A$  can solve, in a fixed amount  $T$  of time, *scales multiplicatively* with increasing computational power.

Input Sizes solved in Time T		
	Power = 1	Power = 2
$n$	$T$	$2T$
Runtime of Algorithm $n^2$	$\sqrt{T}$	$\sqrt{2} \cdot \sqrt{T}$
$2^n$	$\log T$	$1 + \log T$

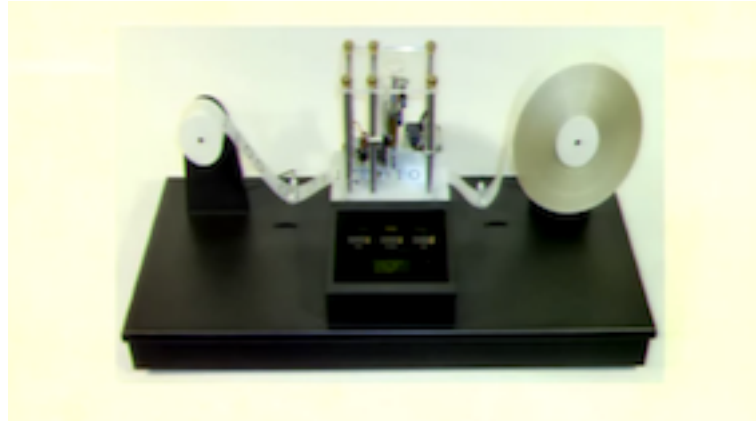


- Moore's Law: Computational power *doubles* roughly every two years.
  - Functional time algorithms will never be able to solve large problems.

- The practical implications are perhaps simpler to understand with this latter formulation.
- Thus, improvements in hardware will *never* overcome ***bad algorithm design***.
- Indeed, the current dramatic breakthroughs in computer science are based upon better (faster and higher performance) algorithmic techniques.

## 4 Robustness

- This measure of quality or "goodness" is ***robust***



- All reasonable models of algorithms are polynomial time equivalent.
  - Otherwise one model could perform, say, an exponential number of operations in the time another model took to perform just one.
- The standard formal model is the **Turing Machine**.

## 5 Cryptography