

→ Which bindings we can access → Scope?

explicit bindings:

`val x = 1`

`var x = 2`

`:`

`x = 3`

implicit/indirect:

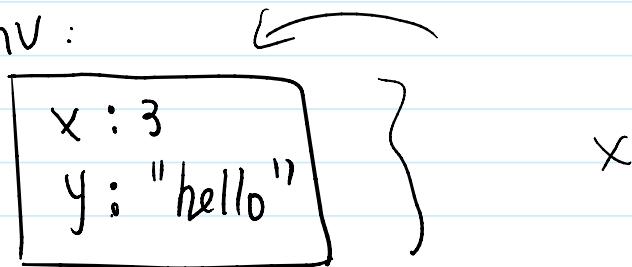
`def foo(a:Int) = { ... }`

`foo(3)`

* assumption: `a:3`

Sets of bindings → environment

env:



`x = y`

`x: 3` `x: 10`

RHS → evaluate it
→ look it up in the env.

LHS
↳ Look up the binding

Lots of choices:

`val y = 2`
`def foo(x:Int) = {`
 `x + y`
`}`
`foo(3)`

`foo(3)`
`val y = 2`
`def foo(x:Int) = {`
 `x + y`
`}`

In Scala, this can work!

```
val y = z  
val z = 2|
```

In scala, it will complain.

Scope → area in which a set of bindings is active.

2 main 'styles' of scoping

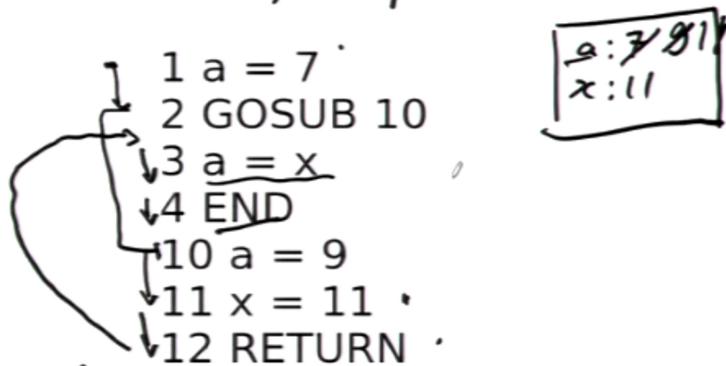
- Static Scoping
 - modern languages : Java, Scala
- Dynamic Scoping
 - older languages : early LISP

Static Scope

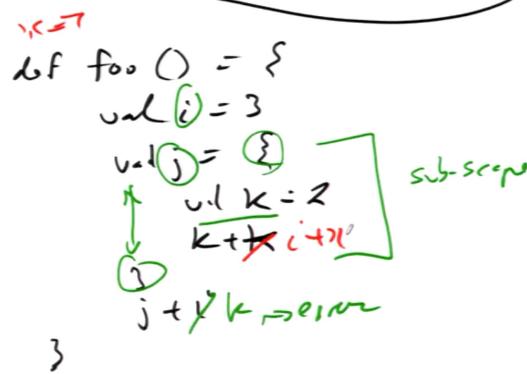
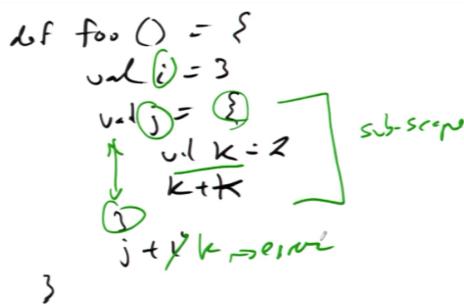
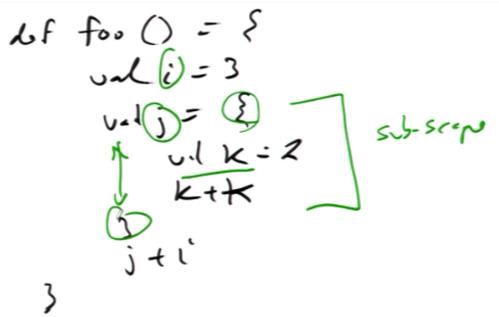
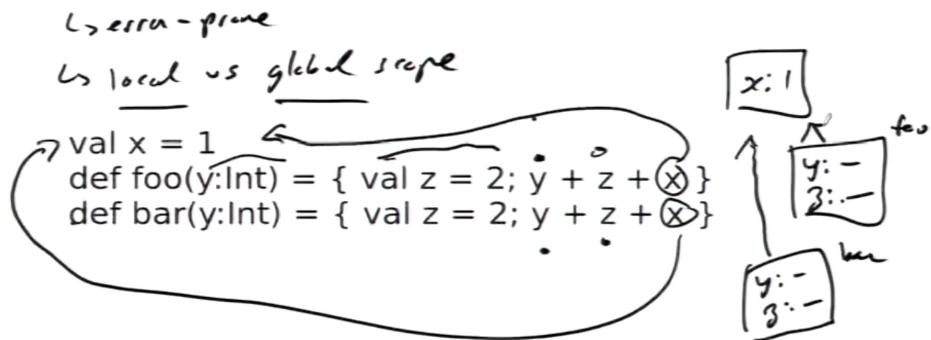
- lexical scoping
- find which vars you can access by looking at in the code. (no dynamic/runtime info needed)

Simplest model → early Basic

- one scope
- no procs/funcs



→ error-prone
→ local vs. global scope



First, my scope →
parent's scope → ...

Algol-style

→ A name declared in a scope is known to its scope & all inner scopes down to the point at which an identically named variable is declared.

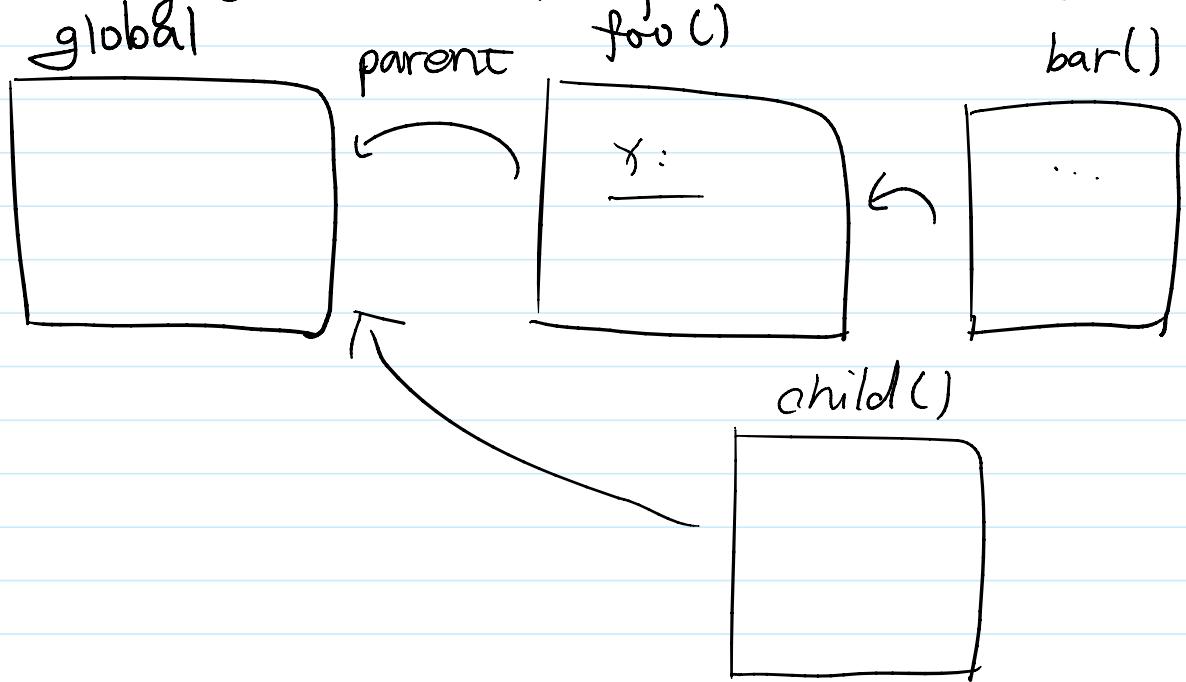
foo () {
 val i = 3 {
 val i = 4 {
 :
 }
}

shadowing

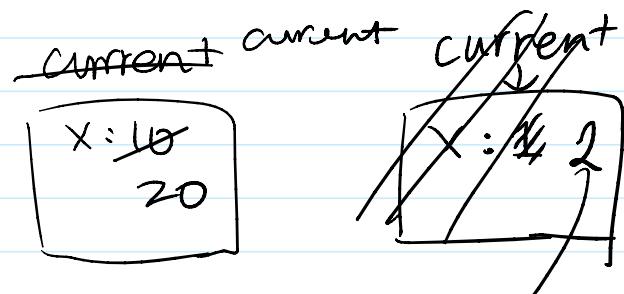
\vdots
 \vdots
 F } shadowing

How does it work?

→ nesting (or chain) of environments



```
{
  var x = 10
  {
    var x = 1
    x = x + x
    x = x + x
    x
  }
}
```



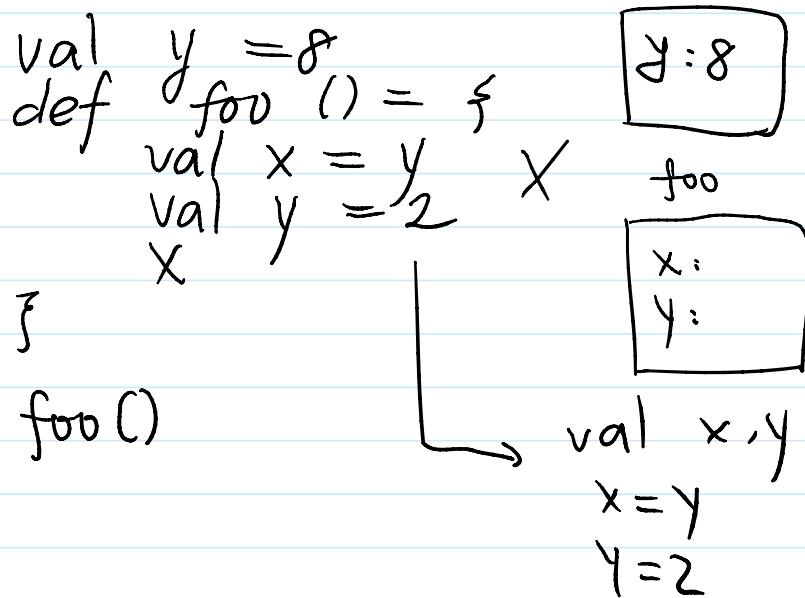
aside: JavaScript

```
function foo () {
  var x = 1
  { var x = 2 }
  return x
} // return 2.
// In Scala,
// return 1.
```

{ var x = 2 } || return 1;
return x
}

Declaration Order

→ Modula-3 style (A language)
→ declaration order doesn't matter



Why unordered declarations?

→ mainly recursive

def foo() = { ; def bar = {
 ; ; ;
 bar() ; ; ;
} ; ; ;
}; ; ;

→ No solution unless treated unordered.

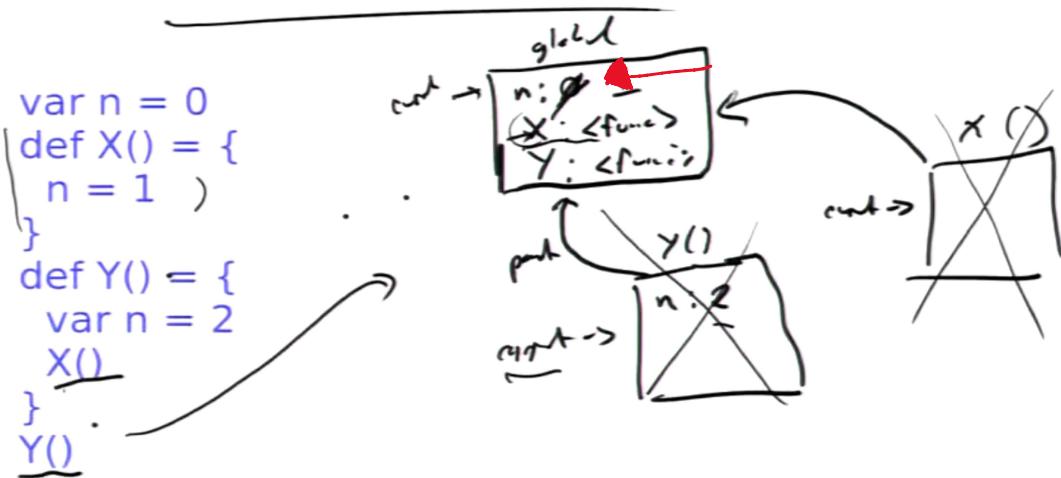
Would it make sense?

val $x = y$ (wouldn't work!)

$$\begin{aligned} \text{val } y &= x \\ \rightarrow x &= y \cup \{a\} \\ y &= x \cup \{b\} \end{aligned}$$

→ That is solution

$$\begin{array}{ll} x = \{a, b\} & ; \{a, b, c\} \\ y = \{a, b\} & ; \{a, b, c\} \end{array}$$



- Current env → static following the func calls.
- Parent env → env in which the func was defined.

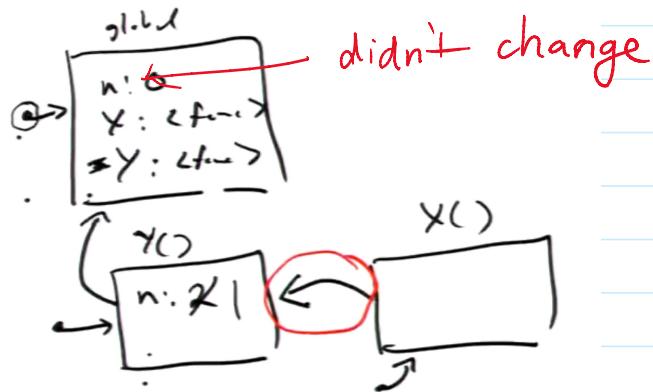
→ look up symbols

→ start at current → follow the parent ptrs.

Dynamic Scoping

→ parent ptrs are based on the current env.

```
var n = 0  
def X() = {  
    n = 1  
}  
def Y() = {  
    var n = 2  
    X()  
}  
Y()
```



→ This seems good?
→ errors

```
def foo() = {}  
def bar() {  
    var x = ...  
    foo():  
}
```

// In scala, it is broken.
// In dynamic, it is okay.