

1. We saw that  $r_o(k; 2) = r(k; 2)$  and that  $r_o(k; q) \leq r(k; q)$  for all  $k, q$ . Prove an inequality in the other direction, namely that

$$r(k; q) \leq r_o(r_o(\cdots r_o(r_o(k; q); q-1) \cdots; 3); 2).$$

for any  $q \geq 3$ .

2. (a) Prove that

$$r_o(k; 2k-1) = 2k$$

for all  $k \geq 2$ .

★(b) Determine  $r_o(k; q)$  exactly for all  $q \geq 2k$ .

★(c) For any fixed  $\alpha \in (0, 2)$ , determine

$$\lim_{k \rightarrow \infty} \frac{r_o(k; \alpha k)}{k}.$$

3. Let  $1 \leq \ell \leq q-1$  be integers, and let  $\binom{[q]}{\ell}$  denote the collection of all  $\ell$ -element subsets of  $[q]$ . A  $(q, \ell)$ -set coloring is a function  $\chi : E(K_N) \rightarrow \binom{[q]}{\ell}$ ; in other words, rather than assigning every edge of  $K_N$  a single color out of  $q$  options, we assign every edge a list of  $\ell$  colors from a palette of size  $q$ . We say that  $v_1, \dots, v_k \in V(K_N)$  form a *color-intersecting clique* if there is a color that appears in all of the  $\binom{k}{2}$  lists associated to the edges they span, that is, if  $\bigcap_{1 \leq i < j \leq k} \chi(v_i v_j) \neq \emptyset$ . The *set coloring Ramsey number*  $r_s(k; (q, \ell))$  is the least  $N$  such that every  $(q, \ell)$ -set coloring of  $E(K_N)$  contains a color-intersecting clique of order  $k$ .

- (a) Prove that  $r_s(k; (q, 1)) = r(k; q)$ .  
 (b) Prove that  $r_s(k; (q, \ell)) \leq r_s(k; (q, \ell-1))$  for any  $2 \leq \ell \leq q-1$ . Conclude that  $r_s(k; (q, \ell)) \leq r(k; q)$  for all  $1 \leq \ell \leq q-1$ .  
 (c) Prove that  $r_s(k; (q, q-1)) = r_o(k; q)$ .  
 (d) Combining parts (a) and (c) with our known bounds on  $r(k; q)$  and  $r_o(k; q)$ , conclude the following. There exist absolute constants  $c, C$  such that for any  $k \geq q \geq 2$ , we have

$$2^{ckq} \leq r_s(k; (q, 1)) \leq 2^{Ckq \log q} \quad \text{and} \quad 2^{\frac{ck}{q}} \leq r_s(k; (q, q-1)) \leq 2^{\frac{Ck}{q} \log q}.$$

In other words, at both extremes  $\ell = 1$  and  $\ell = q-1$ , we have a  $\Theta(\log q)$  gap between the upper and lower bounds.

- (e) Prove that, for every  $\varepsilon > 0$  there exists some  $B > 0$  such that the following holds. If  $\ell \geq \varepsilon q$ , then  $r_s(k; (q, \ell)) \leq 2^{Bkq}$ .  
 (f) Using Theorem 8.1.4, prove the following. For every  $x \geq 1$ , there exists  $D > 0$  such that

$$r_s(k; (q, q-x)) \leq 2^{\frac{Dk}{q} \log q}.$$

Note that this bound is much stronger than that given in (e).

- ★(g) Prove that, for every  $\varepsilon > 0$ , there exists some  $\delta > 0$  such that the following holds. If  $\varepsilon q \leq \ell \leq (1 - \varepsilon)q$ , then

$$r_s(k; (q, \ell)) \geq 2^{\delta k q}.$$

This shows that the bound in (e) is tight up to the value of  $B$  when  $\varepsilon q \leq \ell \leq (1 - \varepsilon)q$ . On the other hand, (f) shows that the upper bound  $\ell \leq (1 - \varepsilon)q$  cannot be entirely removed.

4. (a) Prove that Theorem 8.2.4 is equivalent to the following statement. For every  $C > 0, k \in \mathbb{N}$ , the following holds for sufficiently large  $N$ . Consider a coloring  $\chi : E(K_N) \rightarrow \{\text{red, blue}\}$ , and suppose that  $\chi$  contains no monochromatic clique of order  $C \log N$ . Then for every coloring  $\psi : E(K_k) \rightarrow \{\text{red, blue}\}$ , there is a  $k$ -vertex subset  $S$  of  $K_N$  such that the restriction of  $\chi$  to  $S$  equals  $\psi$  (up to permutations of the vertices).

(b) State and prove a generalization of (a) to colorings with more than two colors.

- ⋄ 5. Prove that if  $G$  is a  $k$ -universal graph, then  $G$  has at least  $2^{(k-1)/2}$  vertices.