1. Prove that, for any  $r \ge 1$ ,

$$r(\underbrace{3,3,\ldots,3}_{r \text{ times}},k) = \Omega\left(\frac{k^{r+1}}{(\log k)^{C_r}}\right),$$

for some constant  $C_r$  depending only on r. Note that this matches the upper bound of Theorem 2.1.5 up to the logarithmic factor.

- 2. Prove that the graph  $\Lambda_q$  has no O'Nan configuration.
- 3. Let q be a prime power. Construct a graph  $\Pi_q$  with vertex set  $V(\Pi_q) = \mathbb{F}_q^2$ , in which two vertices  $(x_1, y_1), (x_2, y_2)$  are adjacent if and only if  $x_1x_2 + y_1y_2 = 1$ .
  - (a) Prove that  $\Pi_q$  is  $C_4$ -free.
  - $\star$  (b) Prove that  $\Pi_q$  satisfies the assumptions of Lemma 4.3.1 with  $\beta = \Theta(1/q)$  and  $R = \Theta(q^{3/2})$ .

*Remark:* You should feel free to assume this result without proof, or to prove it with logarithmic losses in the value of R. The only way I know how to prove this involves techniques (from spectral graph theory) that we will not cover in this class, but I believe there should be an "elementary" proof. Please let me know if you find one!

(c) Use these results, plus a natural generalization of Lemma 3.1.1, to deduce a lower bound for the off-diagonal graph Ramsey number  $r(C_4, K_k)$ , namely

$$r(C_4, K_k) = \Omega\left(\frac{k^{\frac{3}{2}}}{(\log k)^C}\right)$$

for an absolute constant C > 0. What value of C can you get?

- ? (d) Improve the lower bound to  $r(C_4, K_k) \ge k^{\frac{3}{2} + \varepsilon o(1)}$  for any absolute constant  $\varepsilon > 0$ .
- 4. Let q be a prime power. We define a graph  $\Gamma_q^{(5)}$  to be the natural five-dimensional analogue of  $\Gamma_q$ . Namely,  $\Gamma_q^{(5)}$  is a bipartite graph with parts  $P \cup L$ , where P is identified with  $\mathbb{F}_q^5$ , and L comprises all lines in  $\mathbb{F}_q^5$  whose direction is of the form  $(1, z, z^2, z^3, z^4)$  for some  $z \in \mathbb{F}_q$ .
  - (a) Prove that  $\Gamma_q^{(5)}$  is  $C_4$ -free and  $C_{10}$ -free.
  - (b) Define a natural analogue  $G_q^{(5)}$  of  $G_q$ . It is natural to suppose that  $G_q^{(5)}$  is  $C_5$ -free with probability 1; show that this is *not* the case.

*Hint:* Show that  $\Gamma_q^{(5)}$  is not  $C_8$ -free. Use this to find a  $C_5$  in  $G_q^{(5)}$ .

<sup>\*</sup> means that a problem is hard.

<sup>?</sup> means that a problem is open.

 $<sup>\</sup>Leftrightarrow$  means that a problem is on a topic beyond the scope of the course.

- \*\*(c) Without further modifications, this approach cannot be used to lower-bound  $r(C_5, K_k)$ , since  $G_q^{(5)}$  has copies of  $C_5$ . Can you nonetheless salvage a similar approach, and obtain a lower bound on  $r(C_5, K_k)$ ?
- 5. (a) Prove that r(T;q) = O(qn) for every  $q \ge 2$  and every n-vertex tree T.
  - $\star$  (b) Prove that  $r(T;q) = \Theta(qn)$  for every  $q \ge 2$  and every n-vertex tree T.
- $\oplus$  6. A *subdivision* of a graph H is obtained from H by replacing every edge of H by a path of some length (not necessarily the same length for all edges, and paths of length 1 are allowed, so that H is a subdivision of itself). A famous conjecture of Hajós asserts that if  $\chi(G) \ge k$ , then G contains a subdivision of  $K_k$  as a subgraph.
  - (a) Prove that Hajós' conjecture is true for  $k \leq 3$ .
  - $\star$  (b) Prove that Hajós' conjecture is true for k=4.
    - (c) Prove that Hajós' conjecture for k = 5 implies the four-color theorem. Conclude that it is probably pretty hard to prove the k = 5 case.
    - (d) Prove that if Hajós' conjecture is true, then  $r(k) \leq 3k^3$ . Conclude that Hajós' conjecture is false.
    - (e) Prove that if Hajós' conjecture is true, then  $r(3, k) \leq 12k$ . Conclude that Hajós' conjecture is false.
- ↑7. A classical fact in graph theory is that there exist triangle-free graphs of arbitrarily high chromatic number. A standard proof, taught in most introductory graph theory courses, uses the Mycielski construction. In this exercise, you will see two alternative Ramsey-theoretic proofs.
  - (a) For an integer N, let  $S_N$  be a graph with vertex set  $\binom{[\![N]\!]}{2}$ , where we think of the vertices of  $S_N$  as ordered pairs (a,b) with  $1 \leqslant a < b \leqslant N$ . The edges of  $S_N$  consist of all pairs of the form ((a,b),(b,c)) for a < b < c. Prove that  $S_N$  is triangle-free, and that  $\chi(S_N) \to \infty$  as  $N \to \infty$ .
  - (b) The graph  $G_q$  constructed in class is triangle-free; prove that  $\chi(G_q) \to \infty$  as  $q \to \infty$ .
- $\oplus$ 8. (a) Let  $K_{\mathbb{N}}$  denote the complete graph whose vertex set is  $\mathbb{N}$ . Prove the "infinite Ramsey theorem": for any positive integer q, and any q-coloring of  $K_{\mathbb{N}}$ , there is an infinite monochromatic clique.
  - $\star$  (b) Prove that the finite and infinite Ramsey theorems are equivalent. Hint: This fact is often called "compactness", and you may want to use something else called compactness in the proof.
- $\oplus$  9. Prove that there is an infinite set  $S \subseteq \mathbb{N}$  such that for every  $a, b \in S$ , the number a + b has an even number of prime factors (counted without multiplicity).