

The Ramsey-Theoretic Monster Mash

Boo-val Wigderson

ITS Fellows' Seminar

October 31, 2023

Dreadful was the din
Of hissing through the hall, thick swarming now
With complicated monsters head and tail

John Milton, *Paradise Lost* X.521-3

👻outline

Introduction: behemoths of Ramsey theory

Ghosts of graph Ramsey theory

Sea monsters and Ramsey multiplicity

Shapeshifters and oriented Ramsey numbers

👻outline

Introduction: behemoths of Ramsey theory

Ghosts of graph Ramsey theory

Sea monsters and Ramsey multiplicity

Shapeshifters and oriented Ramsey numbers

What is Ramsey theory?



Behemoths

Ghosts

Sea monsters

Shapeshifters



What is Ramsey theory?

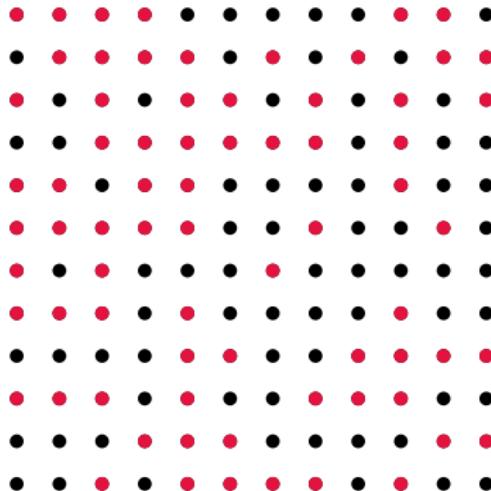
Theorem ("Folklore")

Given N points, if half are colored red, then there are $N/2$ red points.



What is Ramsey theory?

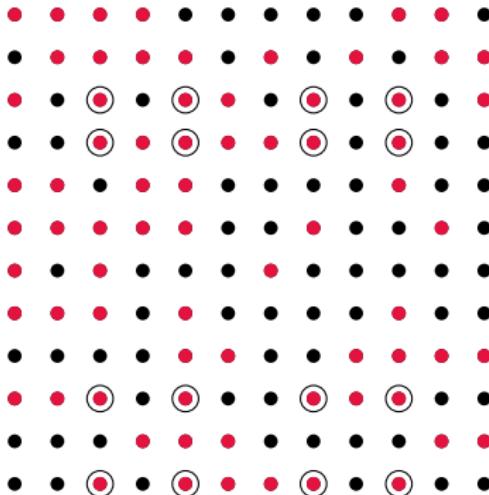
Given an $N \times N$ grid, if half the points are colored red, how large of a **red subgrid** can we find?





What is Ramsey theory?

Given an $N \times N$ grid, if half the points are colored red, how large of a **red subgrid** can we find?





What is Ramsey theory?

Theorem ("Folklore")

Given N points, if half are colored red, then there are $N/2$ red points.

Theorem (Kővári-Sós-Turán 1954)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\log N \times \log N$ red subgrid.

What is Ramsey theory?



Given N points, if half are colored red, how many **evenly spaced** red points can we find?



What is Ramsey theory?



Given N points, if half are colored red, how many **evenly spaced** red points can we find?





What is Ramsey theory?

Theorem ("Folklore")

Given N points, if half are colored red, then there are $N/2$ red points.

Theorem (Kővári-Sós-Turán 1954)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\log N \times \log N$ red subgrid.

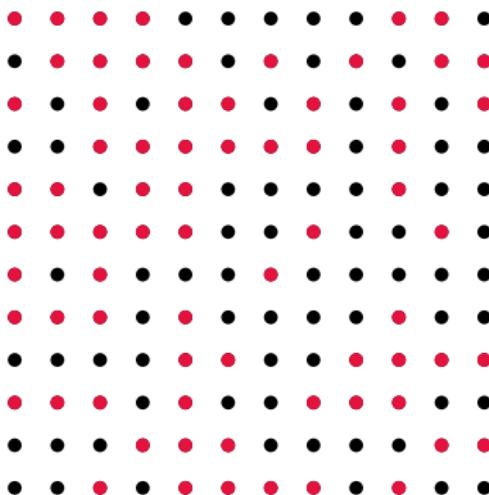
Theorem (Szemerédi 1975, Gowers 2001)

Given N points, if half the points are colored red, then there are $\log \log \log \log N$ evenly spaced red points.



What is Ramsey theory?

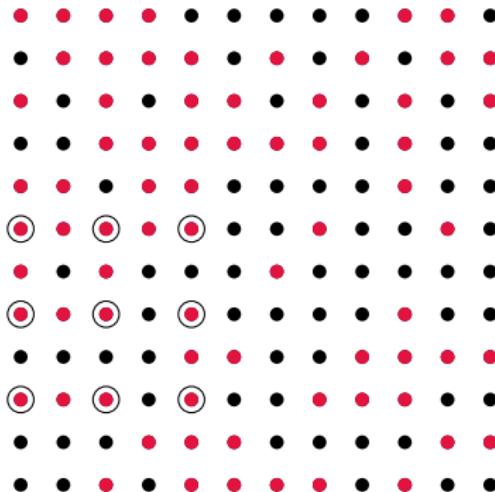
Given an $N \times N$ grid, if half the points are colored red, how large of an **evenly spaced red subgrid** can we find?





What is Ramsey theory?

Given an $N \times N$ grid, if half the points are colored red, how large of an **evenly spaced red subgrid** can we find?





What is Ramsey theory?

Theorem ("Folklore")

Given N points, if half are colored red, then there are $N/2$ red points.

Theorem (Kővári-Sós-Turán 1954)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\log N \times \log N$ red subgrid.

Theorem (Szemerédi 1975, Gowers 2001)

Given N points, if half the points are colored red, then there are $\log \log \log \log N$ evenly spaced red points.

Theorem (Furstenberg-Katznelson 1978, Nagle-Rödl-Schacht-Skokan 2006, Gowers 2007)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\sqrt{A^{-1}(N)} \times \sqrt{A^{-1}(N)}$ evenly spaced red subgrid.



What is Ramsey theory?

Theorem ("Folklore")

Given N points, if half are colored red, then there are $N/2$ red points.

Theorem (Kővári-Sós-Turán 1954)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\log N \times \log N$ red subgrid.

Theorem (Szemerédi 1975, Gowers 2001)

Given N points, if half the points are colored red, then there are $\log \log \log \log N$ evenly spaced red points.

Theorem (Furstenberg-Katznelson 1978, Nagle-Rödl-Schacht-Skokan 2006, Gowers 2007)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\sqrt{A^{-1}(N)} \times \sqrt{A^{-1}(N)}$ evenly spaced red subgrid.

Any large object contains a large structured subobject.



What is Ramsey theory?

Theorem ("Folklore")

Given N points, if half are colored red, then there are $N/2$ red points.

Theorem (Kővári-Sós-Turán 1954)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\log N \times \log N$ red subgrid.

Theorem (Szemerédi 1975, Gowers 2001)

Given N points, if half the points are colored red, then there are $\log \log \log \log N$ evenly spaced red points.

Theorem (Furstenberg-Katznelson 1978, Nagle-Rödl-Schacht-Skokan 2006, Gowers 2007)

Given an $N \times N$ grid, if half the points are colored red, then there is a $\sqrt{A^{-1}(N)} \times \sqrt{A^{-1}(N)}$ evenly spaced red subgrid.

Any large object contains a large structured subobject.

Such results exist for integers, graphs, posets, Banach spaces...

👻outline

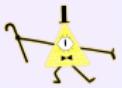
Introduction: behemoths of Ramsey theory

Ghosts of graph Ramsey theory

Sea monsters and Ramsey multiplicity

Shapeshifters and oriented Ramsey numbers

Graph Ramsey theory

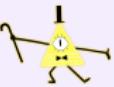


Behemoths

Ghosts

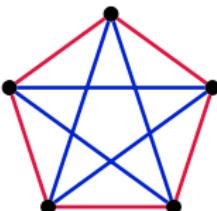
Sea monsters

Shapeshifters



Graph Ramsey theory

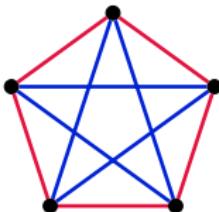
There is a 2-coloring of the edges of K_5 with no monochromatic triangle





Graph Ramsey theory

There is a 2-coloring of the edges of K_5 with no monochromatic triangle

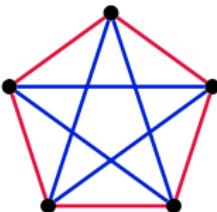


...but every 2-coloring of the edges of K_6 does have a monochromatic triangle.

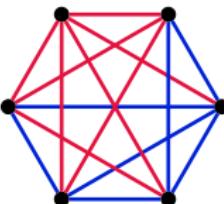


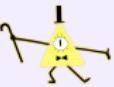
Graph Ramsey theory

There is a 2-coloring of the edges of K_5 with no monochromatic triangle



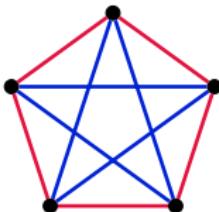
...but every 2-coloring of the edges of K_6 does have a monochromatic triangle.



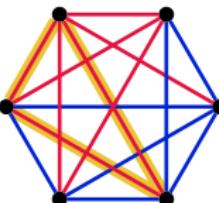


Graph Ramsey theory

There is a 2-coloring of the edges of K_5 with no monochromatic triangle



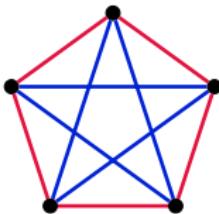
...but every 2-coloring of the edges of K_6 does have a monochromatic triangle.



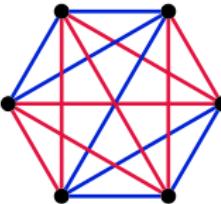


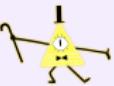
Graph Ramsey theory

There is a 2-coloring of the edges of K_5 with no monochromatic triangle



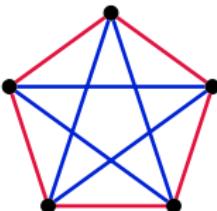
...but every 2-coloring of the edges of K_6 does have a monochromatic triangle.



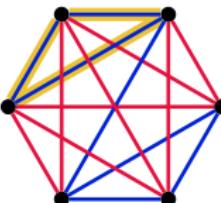


Graph Ramsey theory

There is a 2-coloring of the edges of K_5 with no monochromatic triangle



...but every 2-coloring of the edges of K_6 does have a monochromatic triangle.





Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a **construction**: a coloring of K_N with no monochromatic K_t .



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a **construction**: a coloring of K_N with no monochromatic K_t .

Theorem (Erdős 1947)

$$r(t) > 2^{t/2}.$$



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a **construction**: a coloring of K_N with no monochromatic K_t .

Theorem (Erdős 1947)

$$r(t) > 2^{t/2}.$$

Proof: Let $N = 2^{t/2}$. Consider a **random** two-coloring of $E(K_N)$.



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a construction: a coloring of K_N with no monochromatic K_t .

Theorem (Erdős 1947)

$$r(t) > 2^{t/2}.$$

Proof: Let $N = 2^{t/2}$. Consider a random two-coloring of $E(K_N)$.

$$\mathbb{E}[\#\text{monochromatic } K_t]$$



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a construction: a coloring of K_N with no monochromatic K_t .

Theorem (Erdős 1947)

$$r(t) > 2^{t/2}.$$

Proof: Let $N = 2^{t/2}$. Consider a random two-coloring of $E(K_N)$.

$$\mathbb{E}[\#\text{monochromatic } K_t] = \binom{N}{t} 2^{1 - \binom{t}{2}}$$



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a construction: a coloring of K_N with no monochromatic K_t .

Theorem (Erdős 1947)

$$r(t) > 2^{t/2}.$$

Proof: Let $N = 2^{t/2}$. Consider a random two-coloring of $E(K_N)$.

$$\mathbb{E}[\#\text{monochromatic } K_t] = \binom{N}{t} 2^{1 - \binom{t}{2}} < N^t 2^{-\frac{1}{2}t^2}$$



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a construction: a coloring of K_N with no monochromatic K_t .

Theorem (Erdős 1947)

$$r(t) > 2^{t/2}.$$

Proof: Let $N = 2^{t/2}$. Consider a random two-coloring of $E(K_N)$.

$$\mathbb{E}[\#\text{monochromatic } K_t] = \binom{N}{t} 2^{1 - \binom{t}{2}} < N^t 2^{-\frac{1}{2}t^2} = 1.$$



Ramsey numbers

$r(t)$ = minimum N so that every 2-coloring of the edges of K_N has a monochromatic K_t .

Theorem (Ramsey 1930, Erdős-Szekeres 1935)

$r(t)$ exists (i.e. is finite). In fact, $r(t) < 4^t$.

For a lower bound we need a construction: a coloring of K_N with no monochromatic K_t .

Theorem (Erdős 1947)

$$r(t) > 2^{t/2}.$$

Proof: Let $N = 2^{t/2}$. Consider a random two-coloring of $E(K_N)$.

$$\mathbb{E}[\#\text{monochromatic } K_t] = \binom{N}{t} 2^{1 - \binom{t}{2}} < N^t 2^{-\frac{1}{2}t^2} = 1.$$

So there exists a coloring of $E(K_N)$ with < 1 monochromatic K_t . □



Ghosts

Theorem (Erdős 1947)

$r(t) > 2^{t/2}$. In other words, if $N = 2^{t/2}$, then there exists a coloring of $E(K_N)$ with no monochromatic K_t .



Ghosts

Theorem (Erdős 1947)

$r(t) > 2^{t/2}$. In other words, if $N = 2^{t/2}$, then there **exists** a coloring of $E(K_N)$ with no monochromatic K_t .

The same proof shows that 99.99999% of the colorings of $E(K_N)$ have no monochromatic K_t .



Ghosts

Theorem (Erdős 1947)

$r(t) > 2^{t/2}$. In other words, if $N = 2^{t/2}$, then there **exists** a coloring of $E(K_N)$ with no monochromatic K_t .

The same proof shows that 99.99999% of the colorings of $E(K_N)$ have no monochromatic K_t . **Then where are they?**

Open problem (Erdős)

Find an **explicit** coloring on $N \geq 1.0001^t$ vertices with no monochromatic K_t .



Ghosts

Theorem (Erdős 1947)

$r(t) > 2^{t/2}$. In other words, if $N = 2^{t/2}$, then there **exists** a coloring of $E(K_N)$ with no monochromatic K_t .

The same proof shows that 99.99999% of the colorings of $E(K_N)$ have no monochromatic K_t . **Then where are they?**

Open problem (Erdős)

Find an **explicit** coloring on $N \geq 1.0001^t$ vertices with no monochromatic K_t .

Such Ramsey colorings are **ghosts**. We know they must exist, but we haven't been able to find one!



Ghosts

Theorem (Erdős 1947)

$r(t) > 2^{t/2}$. In other words, if $N = 2^{t/2}$, then there **exists** a coloring of $E(K_N)$ with no monochromatic K_t .

The same proof shows that 99.99999% of the colorings of $E(K_N)$ have no monochromatic K_t . **Then where are they?**

Open problem (Erdős)

Find an **explicit** coloring on $N \geq 1.0001^t$ vertices with no monochromatic K_t .

Such Ramsey colorings are **ghosts**. We know they must exist, but we haven't been able to find one!

Theorem (Li 2023)

*There exists an **explicit** coloring on $N \geq 2^{t^{0.00001}}$ vertices with no monochromatic K_t .*

Bell, book, and candle



Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .



Bell, book, and candle

Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .

How can we prove this? How do we **exorcise** the ghosts?



Bell, book, and candle

Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .

How can we prove this? How do we **exorcise** the ghosts?

Definition

The **book graph** $B_n^{(k)}$ consists of n copies of K_{k+1} joined along a common K_k .



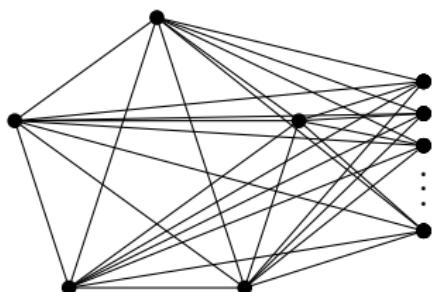
Bell, book, and candle

Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .

How can we prove this? How do we **exorcise** the ghosts?

Definition

The **book graph** $B_n^{(k)}$ consists of n copies of K_{k+1} joined along a common K_k .





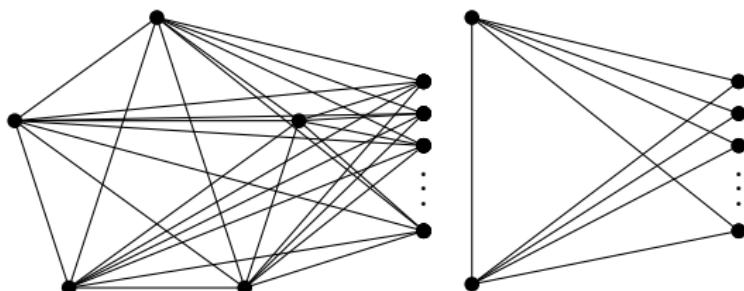
Bell, book, and candle

Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .

How can we prove this? How do we **exorcise** the ghosts?

Definition

The **book graph** $B_n^{(k)}$ consists of n copies of K_{k+1} joined along a common K_k .





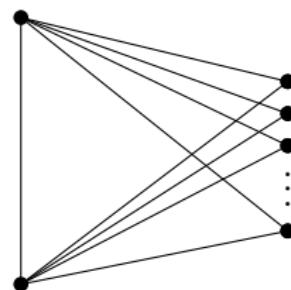
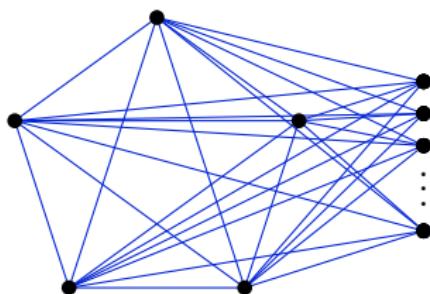
Bell, book, and candle

Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .

How can we prove this? How do we **exorcise** the ghosts?

Definition

The **book graph** $B_n^{(k)}$ consists of n copies of K_{k+1} joined along a common K_k .



Behemoths

Ghosts

Sea monsters

Shapeshifters



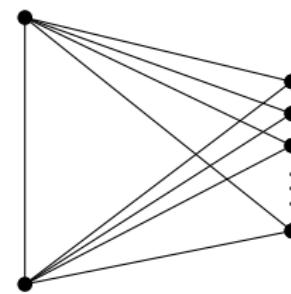
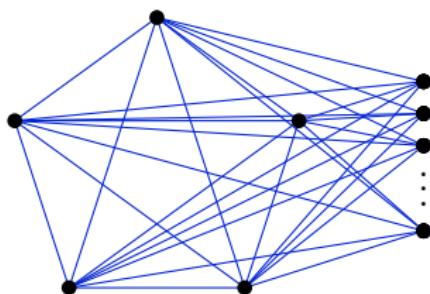
Bell, book, and candle

Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .

How can we prove this? How do we **exorcise** the ghosts?

Definition

The **book graph** $B_n^{(k)}$ consists of n copies of K_{k+1} joined along a common K_k .



Key observation: Finding a large monochromatic book in K_N helps us find a monochromatic K_t .



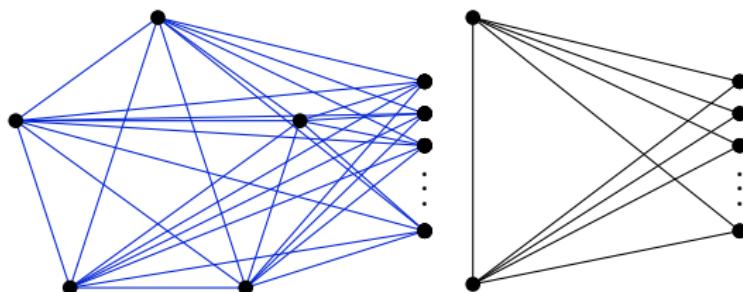
Bell, book, and candle

Theorem (Ramsey 1930): If N is sufficiently large, every coloring of $E(K_N)$ contains a monochromatic K_t .

How can we prove this? How do we **exorcise** the ghosts?

Definition

The **book graph** $B_n^{(k)}$ consists of n copies of K_{k+1} joined along a common K_k .



Key observation: Finding a large monochromatic book in K_N helps us find a monochromatic K_t .

In the n “page” vertices, it suffices to find a red K_t or a blue K_{t-k} .

Finding large books



Key observation: Finding a large monochromatic book in K_N helps us find a monochromatic K_t .



Finding large books

Key observation: Finding a large monochromatic book in K_N helps us find a monochromatic K_t .

Theorem (Conlon 2019)

Every coloring of $E(K_N)$ contains a monochromatic $B_n^{(k)}$ with

$$n \geq 2^{-k}N - o(N)$$

(and this is asymptotically tight).



Finding large books

Key observation: Finding a large monochromatic book in K_N helps us find a monochromatic K_t .

Theorem (Conlon 2019)

Every coloring of $E(K_N)$ contains a monochromatic $B_n^{(k)}$ with

$$n \geq 2^{-k}N - O_k\left(\frac{N}{\log_* N}\right)$$

(and this is asymptotically tight).



Finding large books

Key observation: Finding a large monochromatic book in K_N helps us find a monochromatic K_t .

Theorem (Conlon 2019)

Every coloring of $E(K_N)$ contains a monochromatic $B_n^{(k)}$ with

$$n \geq 2^{-k}N - O_k\left(\frac{N}{\log_* N}\right)$$

(and this is asymptotically tight).

Theorem (Conlon–Fox–W. 2022)

Every coloring of $E(K_N)$ contains a monochromatic $B_n^{(k)}$ with

$$n \geq 2^{-k}N - O_k\left(\frac{N}{(\log \log \log N)^{1/25}}\right).$$



Finding large books

Key observation: Finding a large monochromatic book in K_N helps us find a monochromatic K_t .

Theorem (Conlon 2019)

Every coloring of $E(K_N)$ contains a monochromatic $B_n^{(k)}$ with

$$n \geq 2^{-k}N - O_k\left(\frac{N}{\log_* N}\right)$$

(and this is asymptotically tight).

Theorem (Conlon–Fox–W. 2022)

Every coloring of $E(K_N)$ contains a monochromatic $B_n^{(k)}$ with

$$n \geq 2^{-k}N - O_k\left(\frac{N}{(\log \log \log N)^{1/25}}\right).$$

This result is still **far too weak** to improve the bound $r(t) < 4^t$.

The book algorithm

Theorem (Erdős-Szekeres 1935, Erdős 1947)

$$2^{t/2} < r(t) < 4^t.$$

The book algorithm

Theorem (Erdős-Szekeres 1935, Erdős 1947)

$$2^{t/2} < r(t) < 4^t.$$

Theorem (Campos-Griffiths-Morris-Sahasrabudhe 2023)

$$r(t) < 3.993^t.$$

The book algorithm

Theorem (Erdős-Szekeres 1935, Erdős 1947)

$$2^{t/2} < r(t) < 4^t.$$

Theorem (Campos-Griffiths-Morris-Sahasrabudhe 2023)

$$r(t) < 3.993^t.$$

They introduced a “book algorithm” which can find **some** appropriate monochromatic book.

The book algorithm

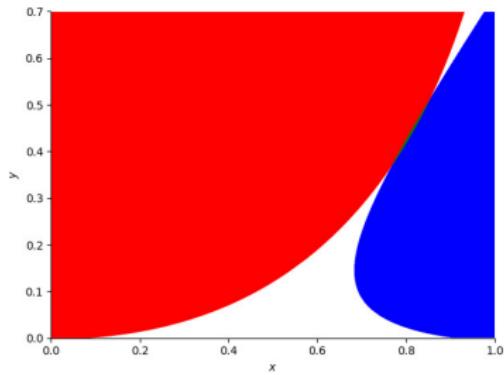
Theorem (Erdős-Szekeres 1935, Erdős 1947)

$$2^{t/2} < r(t) < 4^t.$$

Theorem (Campos-Griffiths-Morris-Sahasrabudhe 2023)

$$r(t) < 3.993^t.$$

They introduced a “book algorithm” which can find **some** appropriate monochromatic book.



The book algorithm

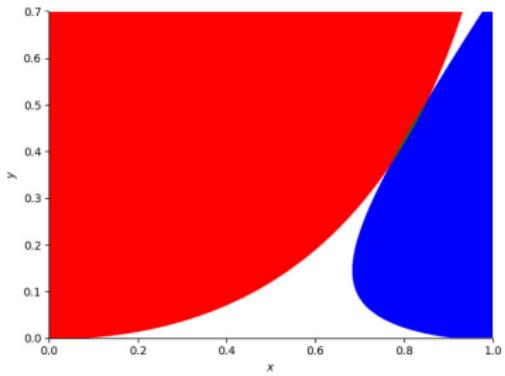
Theorem (Erdős-Szekeres 1935, Erdős 1947)

$$2^{t/2} < r(t) < 4^t.$$

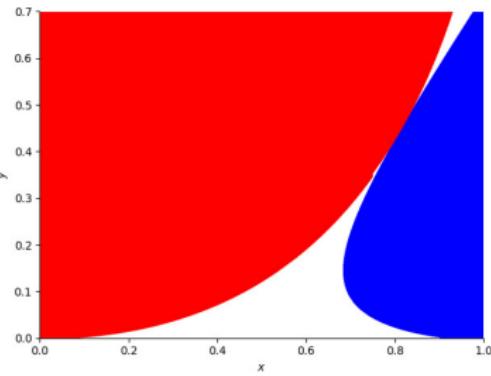
Theorem (Campos-Griffiths-Morris-Sahasrabudhe 2023)

$$r(t) < 3.993^t.$$

They introduced a “book algorithm” which can find **some** appropriate monochromatic book.



Behemoths



Ghosts

Sea monsters

Shapeshifters

👻outline

Introduction: behemoths of Ramsey theory

Ghosts of graph Ramsey theory

Sea monsters and Ramsey multiplicity

Shapeshifters and oriented Ramsey numbers

Edges vs. triangles



Behemoths

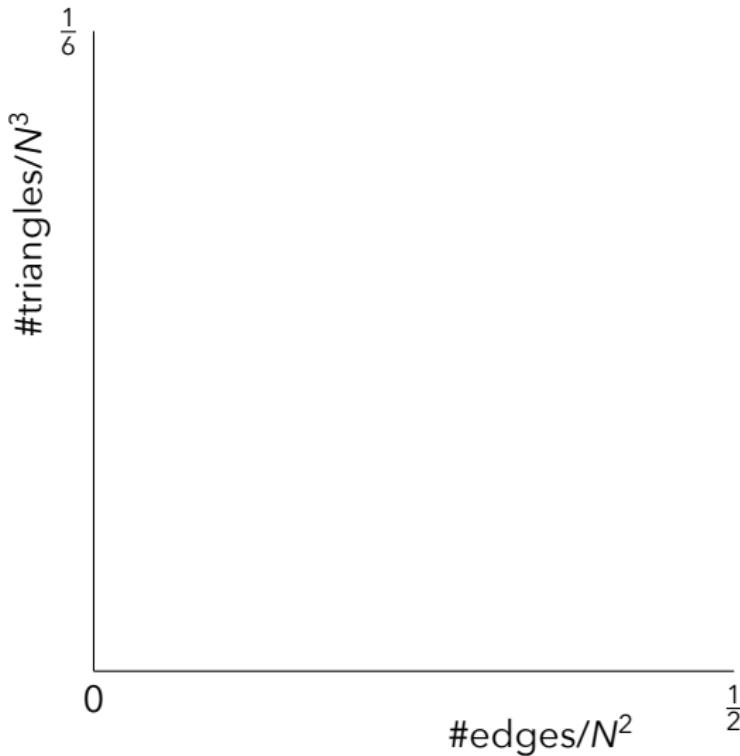
Ghosts

Sea monsters

Shapeshifters



Edges vs. triangles



Behemoths

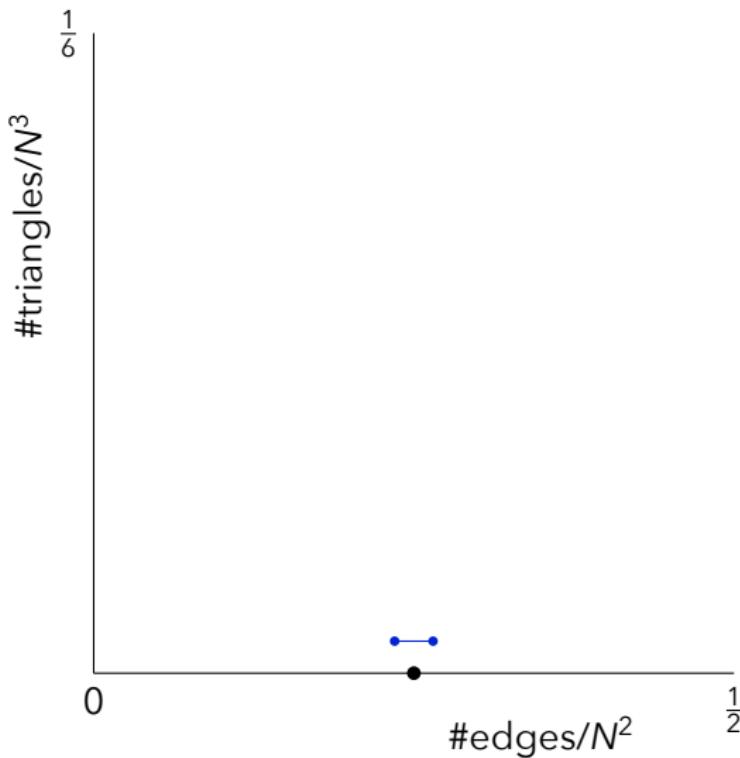
Ghosts

Sea monsters

Shapeshifters



Edges vs. triangles



Behemoths

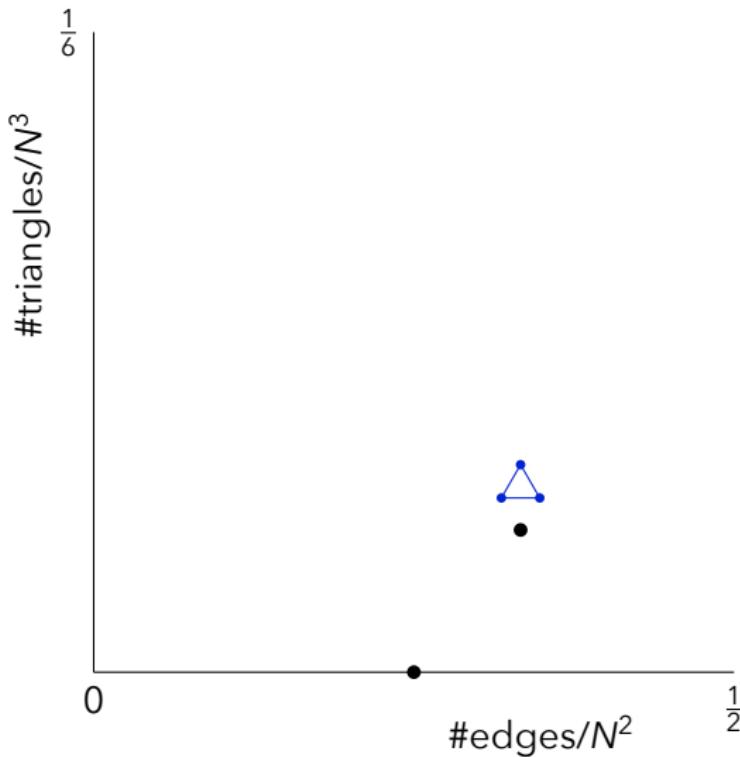
Ghosts

Sea monsters

Shapeshifters



Edges vs. triangles



Behemoths

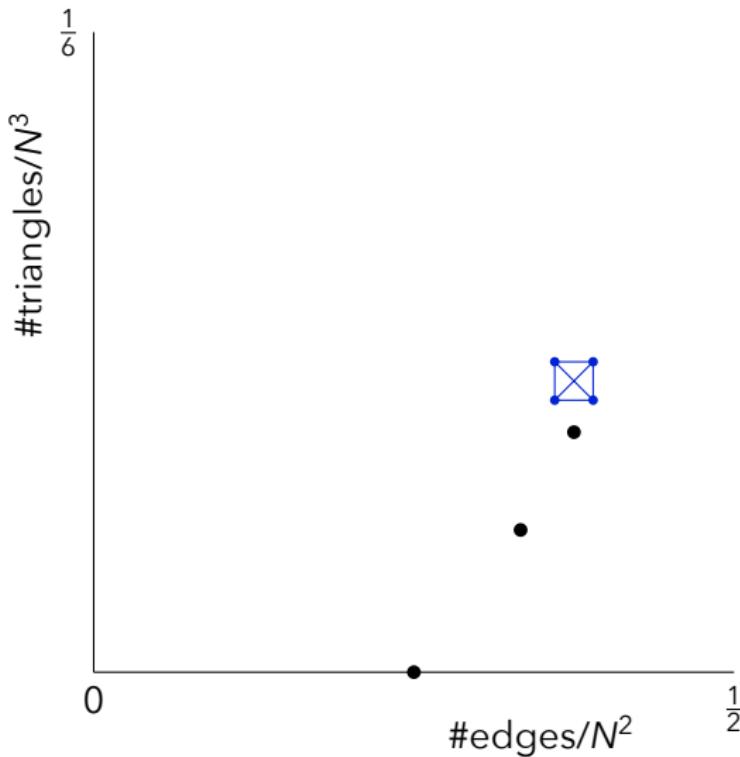
Ghosts

Sea monsters

Shapeshifters



Edges vs. triangles



Behemoths

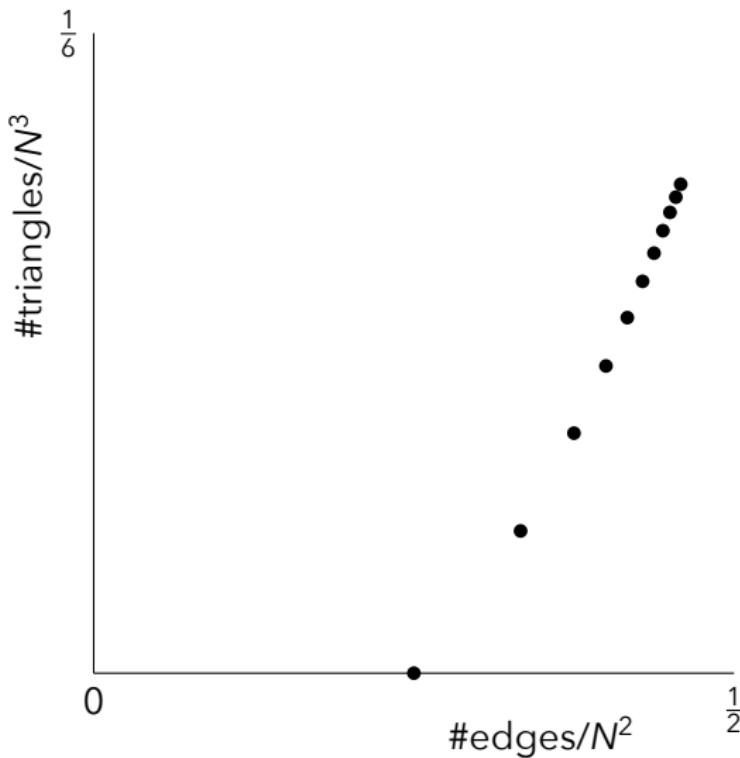
Ghosts

Sea monsters

Shapeshifters



Edges vs. triangles



Behemoths

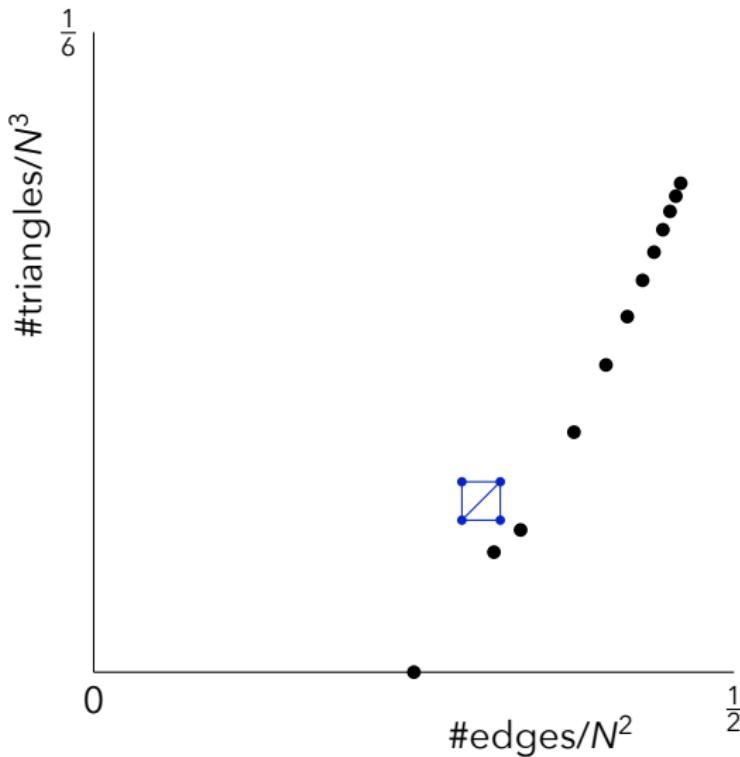
Ghosts

Sea monsters

Shapeshifters



Edges vs. triangles



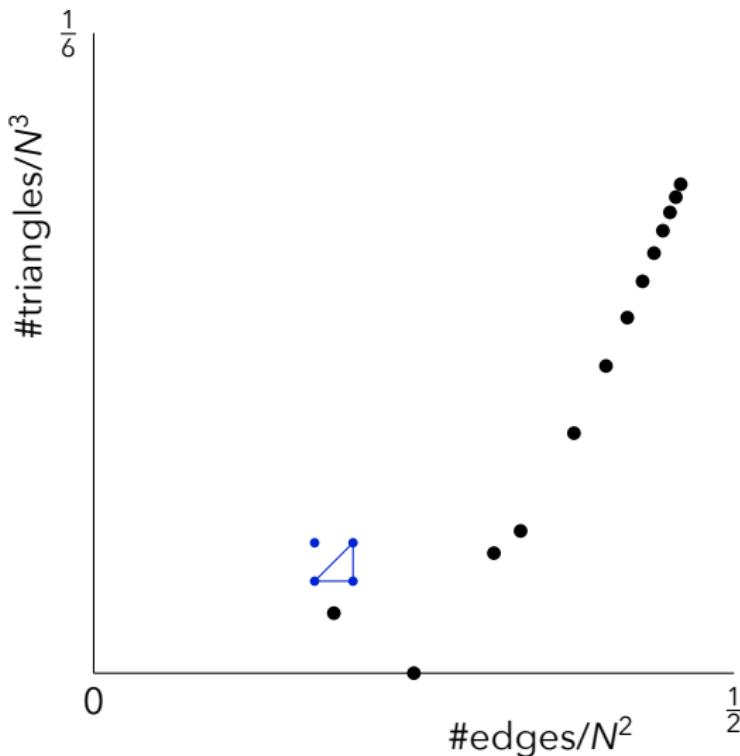
Behemoths

Ghosts

Sea monsters

Shapeshifters

Edges vs. triangles



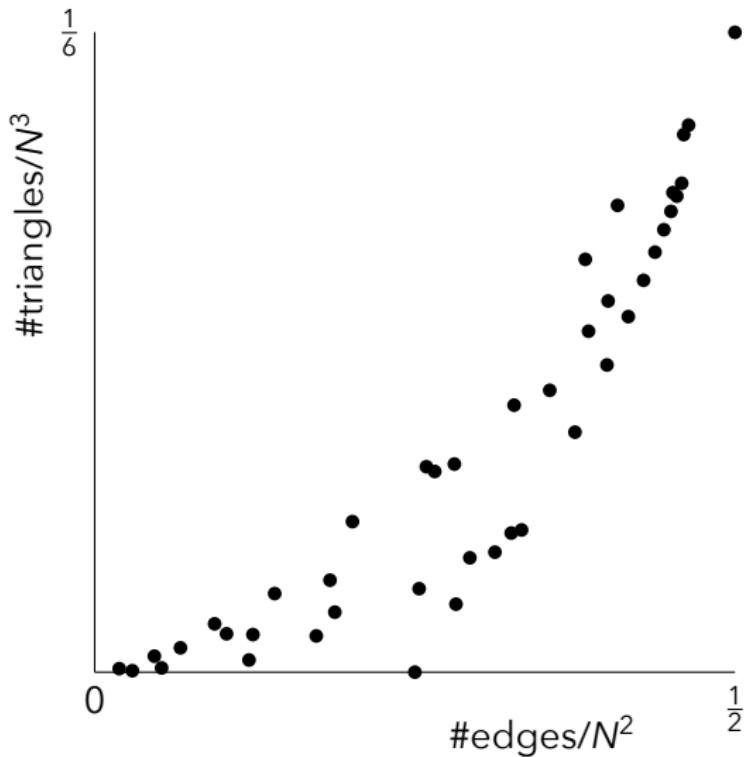
Behemoths

Ghosts

Sea monsters

Shapeshifters

Edges vs. triangles



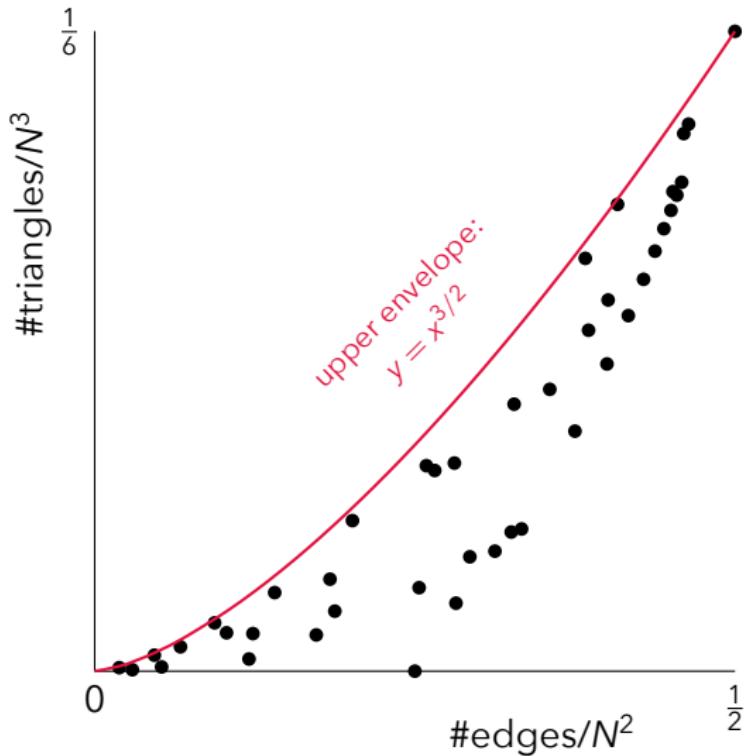
Behemoths

Ghosts

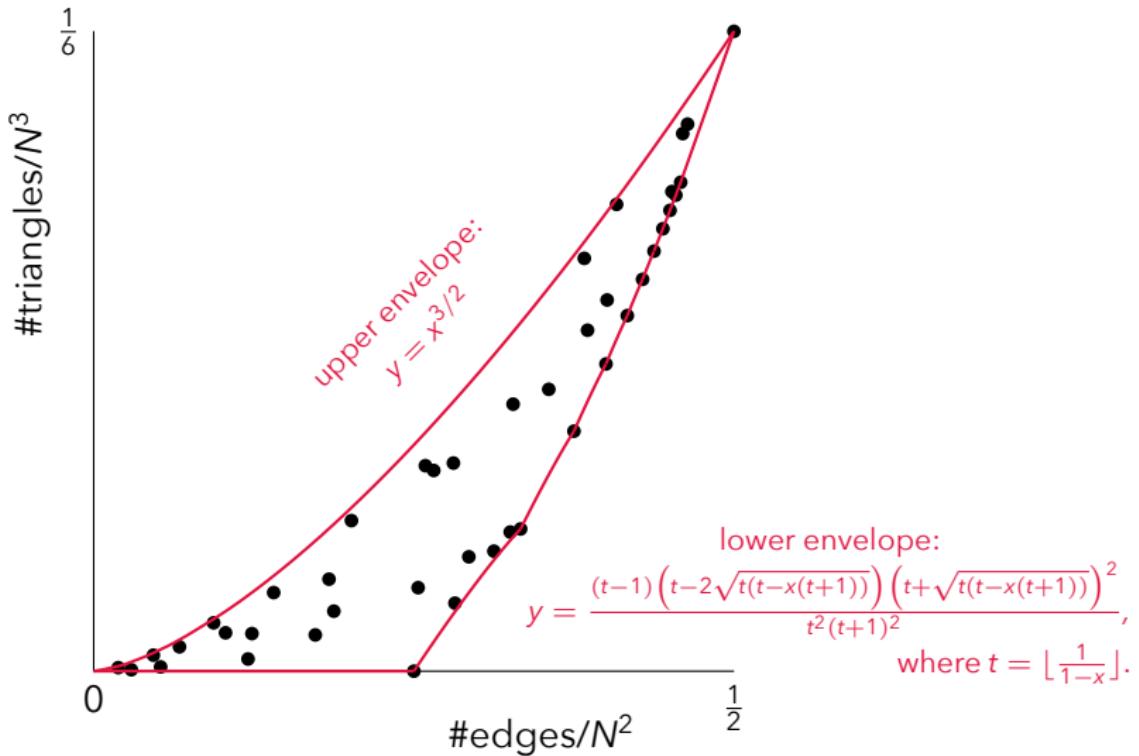
Sea monsters

Shapeshifters

Edges vs. triangles



Edges vs. triangles





Densities

Theorem (Razborov 2008)

The lower envelope for edges vs. K_3 is given by the function

$$y = \frac{(t-1) \left(t - 2\sqrt{t(t-x(t+1))} \right) \left(t + \sqrt{t(t-x(t+1))} \right)^2}{(t(t+1))^2},$$

where $t = \lfloor \frac{1}{1-x} \rfloor$.



Densities

Theorem (Razborov 2008, Nikiforov 2011, Reiher 2016)

The lower envelope for edges vs. K_r is given by the function

$$y = \frac{(t-1)! \left(t - 2\sqrt{t(t-x(t+1))} \right) \left(t + \sqrt{t(t-x(t+1))} \right)^{r-1}}{(t(t+1))^{r-1}(t-r+1)!},$$

where $t = \lfloor \frac{1}{1-x} \rfloor$.



Densities

Theorem (Razborov 2008, Nikiforov 2011, Reiher 2016)

The lower envelope for edges vs. K_r is given by the function

$$y = \frac{(t-1)! \left(t - 2\sqrt{t(t-x(t+1))} \right) \left(t + \sqrt{t(t-x(t+1))} \right)^{r-1}}{(t(t+1))^{r-1}(t-r+1)!},$$

where $t = \lfloor \frac{1}{1-x} \rfloor$.

Conjecture (Sidorenko 1993)

If H is bipartite, the lower envelope for edges vs. H is given by

$$y = x^m,$$

where $m = e(H)$.



Densities

Theorem (Razborov 2008, Nikiforov 2011, Reiher 2016)

The lower envelope for edges vs. K_r is given by the function

$$y = \frac{(t-1)! \left(t - 2\sqrt{t(t-x(t+1))} \right) \left(t + \sqrt{t(t-x(t+1))} \right)^{r-1}}{(t(t+1))^{r-1}(t-r+1)!},$$

where $t = \lfloor \frac{1}{1-x} \rfloor$.

Conjecture (Sidorenko 1993)

If H is bipartite, the lower envelope for edges vs. H is given by

$$y = x^m,$$

where $m = e(H)$.

"A random graph minimizes the number of copies of H , among all graphs with the same number of edges."

Ramsey multiplicity



Conjecture (Sidorenko 1993)

For **bipartite** H , a random graph minimizes the number of H copies.

Ramsey multiplicity



Conjecture (Sidorenko 1993)

For **bipartite** H , a random graph minimizes the number of H copies.

Can such a statement be true for **general** H ?



Ramsey multiplicity

Conjecture (Sidorenko 1993)

For **bipartite** H , a random graph minimizes the number of H copies.

Can such a statement be true for **general** H ?

Conjecture (Erdős 1962, Burr-Rosta 1980)

For **any** H , a random **coloring** minimizes the number of
monochromatic copies of H .



Ramsey multiplicity

Conjecture (Sidorenko 1993)

For *bipartite* H , a random graph minimizes the number of H copies.

Can such a statement be true for *general* H ?

Conjecture (Erdős 1962, Burr-Rosta 1980)

For *any* H , a random *coloring* minimizes the number of *monochromatic* copies of H .

Theorem (Goodman 1959)

This is true for $H = K_3$.



Ramsey multiplicity

Conjecture (Sidorenko 1993)

For **bipartite** H , a random graph minimizes the number of H copies.

Can such a statement be true for **general** H ?

Conjecture (Erdős 1962, Burr-Rosta 1980)

For **any** H , a random **coloring** minimizes the number of **monochromatic** copies of H .

Theorem (Goodman 1959)

This is true for $H = K_3$.

Theorem (Thomason 1989)

This is **false** for $H = K_4$!



Ramsey multiplicity

Conjecture (Sidorenko 1993)

For **bipartite** H , a random graph minimizes the number of H copies.

Can such a statement be true for **general** H ?

Conjecture (Erdős 1962, Burr-Rosta 1980)

For **any** H , a random **coloring** minimizes the number of **monochromatic** copies of H .

Theorem (Goodman 1959)

This is true for $H = K_3$.

Theorem (Thomason 1989)

This is **false** for $H = K_4$!

There exists a coloring of $E(K_N)$ with $< \frac{1}{33} \binom{N}{4}$ monochromatic K_4 (vs. $\frac{1}{32} \binom{N}{4}$ in a random coloring).



Ramsey multiplicity

Conjecture (Sidorenko 1993)

For **bipartite** H , a random graph minimizes the number of H copies.

Can such a statement be true for **general** H ?

Conjecture (Erdős 1962, Burr-Rosta 1980)

For **any** H , a random **coloring** minimizes the number of **monochromatic** copies of H .

Theorem (Goodman 1959)

This is true for $H = K_3$.

Theorem (Thomason 1989)

This is **false** for $H = K_4$!

There **exists** a coloring of $E(K_N)$ with $< \frac{1}{33} \binom{N}{4}$ monochromatic K_4 (vs. $\frac{1}{32} \binom{N}{4}$ in a random coloring).

m-fold cover of an orthogonal tower with maximal Witt index.



A simpler monster

Conjecture (Erdős 1962, Burr-Rosta 1980)

For any H , a random coloring minimizes the number of monochromatic copies of H .



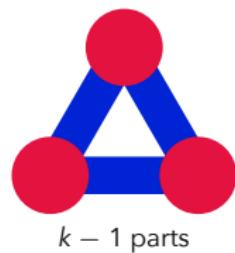
A simpler monster

Conjecture (Erdős 1962, Burr-Rosta 1980)

For any H , a random coloring minimizes the number of monochromatic copies of H .

Theorem (Fox 2008)

If H has chromatic number k and $\gg k^2$ edges, the [Turán coloring](#) beats the random coloring.





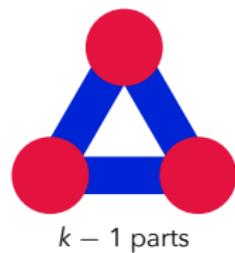
A simpler monster

Conjecture (Erdős 1962, Burr-Rosta 1980)

For any H , a random coloring minimizes the number of monochromatic copies of H .

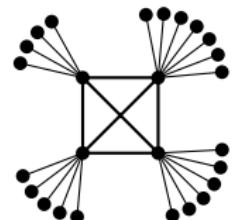
Theorem (Fox 2008)

If H has chromatic number k and $\gg k^2$ edges, the [Turán coloring](#) beats the random coloring.



Theorem (Fox-W. 2023)

If $H = K_k + \text{many}$ pendant edges, the Turán coloring minimizes the number of monochromatic copies of H .





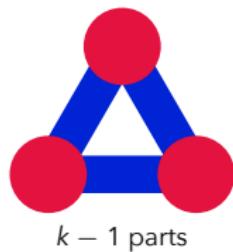
A simpler monster

Conjecture (Erdős 1962, Burr-Rosta 1980)

For any H , a random coloring minimizes the number of monochromatic copies of H .

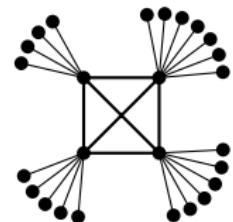
Theorem (Fox 2008)

If H has chromatic number k and $\gg k^2$ edges, the [Turán coloring](#) beats the random coloring.



Theorem (Fox-W. 2023)

If $H = K_k + \text{many}$ pendant edges, the Turán coloring minimizes the number of monochromatic copies of H .



Open problem: Which coloring minimizes the number of K_4 ?

👻outline

Introduction: behemoths of Ramsey theory

Ghosts of graph Ramsey theory

Sea monsters and Ramsey multiplicity

Shapeshifters and oriented Ramsey numbers

Ramsey numbers of graphs and digraphs





Ramsey numbers of graphs and digraphs

The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

Ramsey numbers of graphs and digraphs



The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every edge orientation of K_N contains a **transitive** K_t .

Ramsey numbers of graphs and digraphs



The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every **N -vertex tournament** contains a **transitive** K_t .

Ramsey numbers of graphs and digraphs



The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every **N -vertex tournament** contains a **transitive** K_t .

$$2^{t/2} < \vec{r}(t) < 2^t.$$



Ramsey numbers of graphs and digraphs

The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *Ramsey number* $r(H)$ of a graph H is the minimum N such that every 2-coloring of $E(K_N)$ contains a monochromatic copy of H .

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every **N -vertex tournament** contains a **transitive** K_t .

$$2^{t/2} < \vec{r}(t) < 2^t.$$

Ramsey numbers of graphs and digraphs



The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *Ramsey number* $r(H)$ of a graph H is the minimum N such that every 2-coloring of $E(K_N)$ contains a monochromatic copy of H .

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every **N -vertex tournament** contains a **transitive** K_t .

$$2^{t/2} < \vec{r}(t) < 2^t.$$

The *oriented Ramsey number* $\vec{r}(H)$ of an **acyclic digraph** H is the minimum N such that every N -vertex tournament contains a copy of H .

Ramsey numbers of graphs and digraphs



The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *Ramsey number* $r(H)$ of a graph H is the minimum N such that every 2-coloring of $E(K_N)$ contains a monochromatic copy of H .

Chvátal-Rödl-Szemerédi-Trotter (1983): If H has t vertices and maximum degree Δ , then $r(H) = O_\Delta(t)$.

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every **N -vertex tournament** contains a **transitive** K_t .

$$2^{t/2} < \vec{r}(t) < 2^t.$$

The *oriented Ramsey number* $\vec{r}(H)$ of an **acyclic digraph** H is the minimum N such that every N -vertex tournament contains a copy of H .

Ramsey numbers of graphs and digraphs



The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *Ramsey number* $r(H)$ of a graph H is the minimum N such that every 2-coloring of $E(K_N)$ contains a monochromatic copy of H .

Chvátal-Rödl-Szemerédi-Trotter (1983): If H has t vertices and maximum degree Δ , then $r(H) = O_\Delta(t)$.

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every **N -vertex tournament** contains a **transitive** K_t .

$$2^{t/2} < \vec{r}(t) < 2^t.$$

The *oriented Ramsey number* $\vec{r}(H)$ of an **acyclic digraph** H is the minimum N such that every N -vertex tournament contains a copy of H .

Bucić-Letzter-Sudakov (2019): If H has t vertices and maximum degree Δ , is it true that $\vec{r}(H) = O_\Delta(t)$?



Ramsey numbers of graphs and digraphs

The *Ramsey number* $r(t)$ is the minimum N such that every 2-edge-coloring of K_N contains a **monochromatic** K_t .

$$2^{t/2} < r(t) < 3.993^t.$$

The *Ramsey number* $r(H)$ of a graph H is the minimum N such that every 2-coloring of $E(K_N)$ contains a monochromatic copy of H .

Chvátal-Rödl-Szemerédi-Trotter (1983): If H has t vertices and maximum degree Δ , then $r(H) = O_\Delta(t)$.

Theorem (Fox-He-W. 2022)

No! For any $C > 0$, there exist bounded-degree H with $\vec{r}(H) > t^C$.

The *oriented Ramsey number* $\vec{r}(t)$ is the minimum N such that every **N -vertex tournament** contains a **transitive** K_t .

$$2^{t/2} < \vec{r}(t) < 2^t.$$

The *oriented Ramsey number* $\vec{r}(H)$ of an **acyclic digraph** H is the minimum N such that every N -vertex tournament contains a copy of H .

Bucić-Letzter-Sudakov (2019): If H has t vertices and maximum degree Δ , is it true that $\vec{r}(H) = O_\Delta(t)$?



Shapeshifters

Theorem (Fox-He-W. 2022)

For any $C > 0$, there exist bounded-degree H with $\vec{r}(H) > t^C$.



Shapeshifters

Theorem (Fox-He-W. 2022)

For any $C > 0$, there exist bounded-degree H with $\vec{r}(H) > t^C$.

Theorem (Fox-He-W. 2022)

$\vec{r}(H)$ is “small” if H has “few edge length scales”.



Shapeshifters

Theorem (Fox-He-W. 2022)

For any $C > 0$, there exist bounded-degree H with $\vec{r}(H) > t^C$.

Theorem (Fox-He-W. 2022)

$\vec{r}(H)$ is “small” if H has “few edge length scales”.





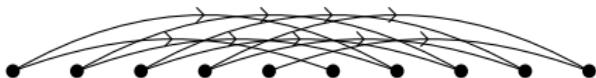
Shapeshifters

Theorem (Fox-He-W. 2022)

For any $C > 0$, there exist bounded-degree H with $\vec{r}(H) > t^C$.

Theorem (Fox-He-W. 2022)

$\vec{r}(H)$ is “small” if H has “few edge length scales”.





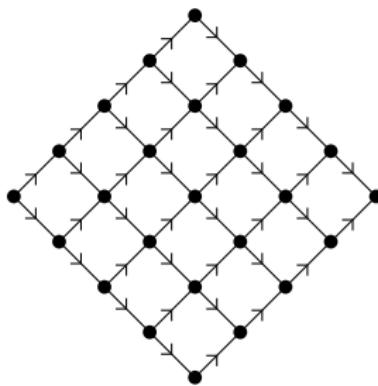
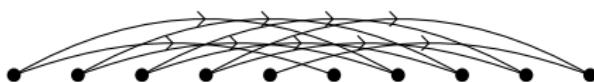
Shapeshifters

Theorem (Fox-He-W. 2022)

For any $C > 0$, there exist bounded-degree H with $\vec{r}(H) > t^C$.

Theorem (Fox-He-W. 2022)

$\vec{r}(H)$ is “small” if H has “few edge length scales”.





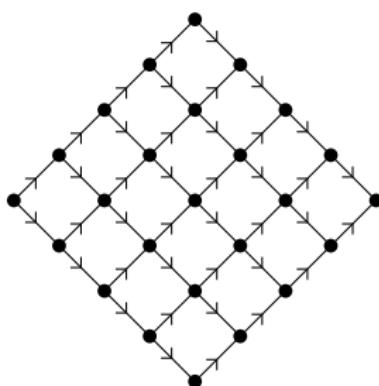
Shapeshifters

Theorem (Fox-He-W. 2022)

For any $C > 0$, there exist bounded-degree H with $\vec{r}(H) > t^C$.

Theorem (Fox-He-W. 2022)

$\vec{r}(H)$ is “small” if H has “few edge length scales”.



The digraphs for which $\vec{r}(H)$ is “large” are **shapeshifters**: they have many edges at every length scale, despite having **bounded degree**.

Thank you!