

* Numerical Analysis

• Set-1 : Err^{ns}

$\text{Rel Err}^{\text{ns}} = \left| \frac{\text{True} - \text{Approx}}{\text{True}} \right| = \left| \frac{\text{Err}^{\text{ns}}}{\text{True}} \right|$. Other trivia from Kania's.

• Set-2 (Fin⁸ Roots of Eqs)

- Bisection Method: 1. Find $x_0 = a$, $x_1 = b$ s.t. $f(a) \cdot f(b) < 0$ 2. Next approxⁿ: $x_2 = \frac{a+b}{2}$
- Check $f(x_2)$. Based on sign, s.t. $f(x_2) \cdot f(a) < 0$ or $f(x_2) \cdot f(b) < 0$, choose next intv. Repeat it
- Always check sig of f in the mid-intv before begin → Give only 1 root in 1 go, in an intv. Chge intv more
- Regula Falsi Method (False Posn): 1. Choose x_0, x_1 s.t. $f(x_0) \cdot f(x_1) < 0$ 2. $x_2 = \frac{x_0 f(x_1) - x_1 f(x_0)}{f(x_1) - f(x_0)}$
- 3. If $f(x_0) f(x_2) < 0$, take x_2 as x_1 , else take x_2 as x_0 . (I think $x_0 < x_1$ always) $f(x_2) - f(x_0)$ → We still need to solve in calc², then choose intv. → Calc: $C = f(A)$; $D = f(B)$; $X = \frac{AD - BC}{D - C}$ to find some x_0, x_1 , i.e. A & B

• Secant Method (Mod^d RF)

- 1. Chose x_0, x_1 as any 2 pts. May not conⁿ root.
- 2. Just do $x_{n+1} = x_n - \frac{f(x_n)}{f(x_n) - f(x_{n-1})}$. Note, as you move forward, x_0, x_1, \dots are leaving beh^d. Just $f(x_n) - f(x_{n-1})$ → See this has to come fast keep moving ahead
- Calc: $C = f(A)$; $D = f(B)$; $X = \frac{AD - BC}{D - C}$! $A = B = X$. Keep note X only. Then B = X ! Recur
- Can use tab^r reprⁿ for all the above. Even if not, chill!
- Newⁿ Raphⁿ Med (Use only if $f'(x)$ & $f''(x)$ are ct^s. Also mem it) $f(x) = f(x_0) + f'(x_0)(x - x_0) + \dots$ Put A₀ Put $f(x)$ calc: $C = f(A)$; $D = f(B)$; $X = A - \frac{C}{D}$; $X = A - \frac{f(A)}{f'(A)}$
- 1. Find only 1 x_0 close enough to the root (Not x_0, x_1 like above). Get this one by calc, then choose nearby
- Check that $f'(x_0) \neq 0$. 2. $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$ [P.S.: Let actual root $x = x_0 + h$, then $f(x) = f(x_0 + h) = 0$ By Tayⁿ Ser, $f(x_0 + h) = f(x_0) + h f'(x_0) + \dots = 0$, $h \approx -\frac{f(x_0)}{f'(x_0)}$]

→ For find⁸ \sqrt{N} ($N > 0$). Consⁱ $x^2 - N = 0$. "Iter": $x_{n+1} = \frac{1}{2} [x_n + \frac{N}{x_n}]$

→ Overⁿ Itⁿ Form^a: 1. $\frac{1}{N}$: $x_{n+1} = x_n(2 - Nx_n)$ [Take $f(x) = \frac{1}{x} - N$] 2. $\frac{1}{\sqrt{N}}$: $x_{n+1} = \frac{1}{2}(x_n + \frac{1}{Nx_n})$ [$\frac{f(x)}{x^2 - \frac{1}{N}}$]

3. $\sqrt[N]{N}$: $x_{n+1} = \frac{1}{k} [(k-1)x_n + \frac{N}{x_n^{k-1}}]$ { $f(x) = x^k - N$ } 4. $\frac{1}{\sqrt[N]{N}}$: $x_{n+1} = \frac{x_n(p+1-Nx_n^p)}{p}$ { $f(x) = x^p - N$ }

→ Very fast me^d → Fails if ① $f'(x_0) = 0$ in a nbhd of the root. ② $f'(x) = 0$ at the root. This means 2 thgs:

②.1 $f(x)$ has mul^e roots at x ②.2 $f(x)$ has stat^s pt. at x ③ f, f' must be def^d & ct^s.

• Convergence Critⁿ: { x_n cgs to root x w. ord^d p, if $|x_{n+1} - x| \leq \lambda |x_n - x|^p$ f. o. $\lambda > 0$ λ : Asymp Errⁿ const

→ If $E_n = x_n - x$, we need $|E_{n+1}| \leq \lambda |E_n|^p$. For $p=1$, we require $\lambda < 1 \rightarrow$ lin^r Convergence.

→ If $|x_{n+1} - x| \leq \lambda |x_n - x| + n > 0$ or If $|x_n - x| \leq \lambda^n |x_0 - x|$, we get lin^r Convergence.

• Convergence of Bisⁿ Me^d: + n, let $c_n = \frac{a_n + b_n}{2}$, $|c_{n+1} - x| \leq \frac{|c_n - x|}{2}$, \therefore Bisⁿ Me^d is also Lin^r Convergence.

• Convergence of Sec^d Me^d: Or^r of convergence: $p = \frac{1+\sqrt{5}}{2} \approx 1.618$; $\lambda = \left[\frac{f''(x)}{2f'(x)} \right]^{\frac{1}{p-1}}$ $\rightarrow 0.618$ {P.S. by Tayⁿ Ser expanⁿ $f(x+\alpha) = \dots$ }

→ $E_{i+1} \approx \frac{E_i E_{i-1}}{\left[\frac{f''(x)}{2f'(x)} \right]}$ → Rel^b blw errⁿ vs Errⁿ Eqⁿ. Holds only if x is a simple Root.

→ But see 2019-7-12-Q5(c) for x is a mul^e root i.e. f'(x)=0, convergence is not lin^r quad^c, but lin^r for

• Convergence of Newⁿ Raphⁿ: $p=2$, $\lambda = \frac{f''(x)}{2f'(x)}$, $E_{i+1} = \left[\frac{f''(x)}{2f'(x)} \right] E_i^2$ \rightarrow ErrEq p87, we need E_0 in

• Convergence of Regula Falsiⁱ: $E_{i+1} = \left[\frac{f''(x)}{2f'(x)} \right] (x_0 - x) E_i$ {* Check p & λ !!!} $\equiv \left(\frac{E_i}{2} \right)^{\frac{1}{p}}$ → ErrEq p87,

Then: If $f(x)$ is ct^s in some $[a, b]$ that contains the root & $|f'(x)| \leq c < 1$ in $[a, b]$. Then, $\forall x_0 \in [a, b]$ t.

$\langle x_n \rangle$ given by $x_{n+1} = \phi(x_n)$, $k=0, 1, 2, \dots$ cgs to the root x of " $x = \phi(x)$ " Chebyshev's Med^d: E_i of

→ $|f'(x)| \leq c < 1$: Same as "Lipschitz Condⁿ". $C = \text{Lip}'z$ Const. 3rd or 4th

• Set-3 Diag Mat

• Dir^t Me^d: $\begin{cases} A = \text{Lower } \Delta^n: \text{ Forward Subs} \\ A = \text{Upper } \Delta^n: \text{ Back Subs} \end{cases}$ Sys^m of lin^r Eqs

That's in $\begin{cases} \text{Not w. edd} \\ \text{for } x \end{cases}$ We don't know as yet

∴ Do this much sooner

- Gauss Elim Med: 1. Remove the lead vars in successive eqns, ... chg to upper Δ form 2. Backd solve
 L $r = n$: Univ Sols, $|A| \neq 0$; $r < n$ & $b_{r+1}, b_{r+2}, \dots, b_n \neq 0$: No soln; $r > n$ & $b_{r+1}, \dots, b_n = 0$: Inf solns
 Each stage
 - Part 1 Ptg: Check 1st col in 1st stg, get the big no. on top, check 2nd col in 2nd stg, get big no. on top
 - Full Ptg: Check whole mat & get big no. on top. Never used Done when pt ele to vanish
 → Gauss-Jordan Elim Med: $[A|b] \xrightarrow{\text{Gauss Jdn}} [I|d]$. Ptg is us, Usd mostly for A^{-1} . $[A|I] \xrightarrow{\text{GI}} [I|A^{-1}]$
 - Some, Gauss Elim for $A^{-1} \equiv [A|I] \rightarrow \text{Chg A to Upp } \Delta \rightarrow \text{Chg to I} \rightarrow [I|A^{-1}]$
 - Direct Med (Iterve): Usd when coeff mat is "Sparse" or "Large."
 - Gauss-Seidel Iter Med: $x^{(k+1)} = -\frac{1}{a_{11}}(a_{12}y^{(k)} + a_{13}z^{(k)}) - b_1$, $y^{(k+1)} = -\frac{1}{a_{22}}(a_{21}x^{(k+1)} + a_{23}z^{(k)}) - b_2$, $z^{(k+1)} = -\frac{1}{a_{33}}(a_{31}x^{(k+1)} + a_{32}y^{(k+1)}) - b_3$
 - L-Calc - $A = \frac{(b_1 - B - C)}{a_{11}}$; $B = \frac{(b_2 - A - C)}{a_{22}}$; $C = \frac{(b_3 - B - A)}{a_{33}}$
 - Med of "Successive Displ": use a better val as ap
 - It's won't cgs if diag are not "heavy", i.e., Rearrange to get heavy diag & then beg. Take $x^{(0)} = y^{(0)} = z^{(0)} = 0$
- Set-4: Interpolⁿ**
 lead Δ diff's
 - Forward Diff: $\Delta y_0 = y_1 - y_0$ (1st FD); $\Delta^2 y_0 = \Delta(\Delta y_0) = \Delta y_1 - \Delta y_0 = y_2 - 2y_1 + y_0$ (2nd FD)
 - Backd - $\nabla y_1 = y_1 - y_0$ (1st BD); $\nabla^2 y_2 = \nabla y_2 - \nabla y_1$ (2nd BD)
 - Cent - $\delta y_{42} = y_2 - y_0$, i.e., $\delta y_n = y_{n+\frac{1}{2}} - y_{n-\frac{1}{2}}$, $\nabla y_{42} = y_{n+\frac{1}{2}} - y_{n-\frac{1}{2}}$
 - Shift Op^r: $E f(x) = f(x+h)$; $E^2 f(x) = f(x+2h)$; $E^n f(x) = f(x+nh)$; $E^{-n} f(x) = f(x-nh)$
 - Avrg Op^r (μ): $\mu f(x) = \frac{1}{2} [f(x+\frac{h}{2}) + f(x-\frac{h}{2})]$
 - Diff Op^r: $D f(x) = f'(x)$
 - $E = 1 + \Delta$; $E = (1 - \nabla)^{-1}$; $S = E^{\frac{1}{2}} - E^{-\frac{1}{2}}$; $\mu = \frac{1}{2}(E^{\frac{1}{2}} + E^{-\frac{1}{2}})$; $hD = \log E$ (by Tay's Thm)
- Then: n th diff's of a poly of n th deg are const & all n th deg diff's are 0, when x is equally spaced.
 [Come also True!], i.e., if x is $\frac{h}{n}$ spad & n th diff's of a fn are const \Rightarrow That fn is n th deg poly.
 • New's Binet Expan Form (For get n missed vals in mid for a fn):

$$\Delta^n y_0 = 0, \text{ i.e., } (E-1)^n y_0 = 0, \text{ i.e., } y_0 - ny_1 + \frac{n(n-1)}{2!} y_2 - \dots + (-1)^n y_0 = 0$$
 [By Binet Expan]
 For sol^r prob, expand this only & put $E^2 y_0, y_2, \dots$ etc. If n th val are known, take $\Delta^n y_0 = 0$
- If $(n+1)$ th val for a fn are give, $(n+1)$ th diff will van (not const) & for a poly of n th deg, $(n+1)$ th diff will van ($\because n$ will be const). $\therefore n+1$ val's give a poly of n th deg, but can give $\leq n$ also
 2018 Q5b-P2
- Then: New's & Forward Interpol Form, let x_0, x_1, \dots, x_n be eqly spa by dis $\approx h$. Then

$$u = \frac{x-x_0}{h}$$

$$\text{nth deg poly for } \phi(x) = y_0 + \frac{u \Delta y_0}{1!} + \frac{u(u-1)}{2!} \Delta^2 y_0 + \dots + \frac{u(u-1)\dots(u-n)}{n!} \Delta^n y_0$$
- L Usd for first val close to beg of data set L If $f(3.1)$ is reqd, Take $y_0 = f(3)$. Leave $f(1), f(2)$
 L Calc - $y_0 + A \Delta y_0 + A(A-1) \frac{\Delta^2 y_0}{2!} + \dots \rightarrow$ Press CALC, put u as A L Rem^r Cum^r Freq Type Prob
 2018 Q5b-P2
- Then: New's-Greg & Backd Diff Form: For val^r of y near the end of data set.

$$f(x) \approx \phi(x) = y_n + \frac{u \nabla y_n}{1!} + \frac{u(u+1)}{2!} \nabla^2 y_n + \dots + \frac{u(u+1)\dots(u+n-1)}{n!} \nabla^n y_n$$
 $u = \frac{x-x_n}{h}$
- Err in Diff Form: 1. For Intⁿ: $R(x) = \frac{\Delta^{n+1} y_0}{(n+1)!} u(u-1)\dots(u-n)$ $\{u = \frac{x-x_0}{h}\}$
 2. Backd Intⁿ: $R(x) = \frac{\Delta^{n+1} y_n}{(n+1)!} u(u+1)\dots(u+n)$ $\{u = \frac{x-x_n}{h}\}$ Just the next term in correct form
 form
- Intⁿ w. \neq Interpol: 1. Lag Int Form (Univ Int & Poly): $f(x) = \sum_{i=0}^n l_i(x) \cdot f_i$, wh. $l_i(x) = \frac{(x-x_0)(x-x_1)\dots(x-x_{i-1})(x-x_{i+1})\dots(x-x_n)}{(x_i-x_0)\dots(x_i-x_{i-1})(x_i-x_{i+1})\dots(x_i-x_n)}$; $\sum l_i(x) = 1$ & $l_i(x_i) = \begin{cases} 0, & i \neq j \\ 1, & i = j \end{cases}$
- Lag Int Coeffs
 $g(y) = f^{-1}(y) = \sum_{i=1}^n \frac{\text{Thres}(y)}{(y-y_i) \prod_{j \neq i} (y_j - y_i)}$ {Just x & y take oppo roles}
 rest same as Lag Form
- L General Int Form (give^r y, find x. Take it as $g(y) = f^{-1}(y)$). Can be done if f is its & more \uparrow or \downarrow)

$$g(y) = f^{-1}(y) = \sum_{i=1}^n \frac{\text{Thres}(y)}{(y-y_i) \prod_{j \neq i} (y_j - y_i)}$$
 {Just x & y take oppo roles}
 But act by it's the last t,
- coz if all Δy_0 are used in the form, Δy_0 will be 0, but Err can't be 0. We take $n+1$ as the last in only & Err = last Term
 Term

- Part ^{no} 8 Lag's Form: Take num^r of the rat^r for. Calc f(x) at x=1, 2, 3. Put in Lag Form.
 E.g. - $\frac{3x^2+x+1}{(x-1)(x-2)(x-3)}$; Take f(x) = $3x^2+x+1$, i.e. $f(x) = \frac{3x^2+x+1}{(x-1)(x-2)(x-3)}$
 Now f(x) at 1, 2, 3 we that's f(x) $\begin{array}{|c|c|c|c|} \hline x & 1 & 2 & 3 \\ \hline f(x) & 5 & 15 & 31 \\ \hline \end{array}$, i.e. $f(x) = 3x^2+x+1 = (x-2)(x-3) + \dots$
 $\therefore \frac{3x^2+x+1}{(x-1)(x-2)(x-3)} = \frac{5}{2(x-1)} - \frac{15}{2(x-2)} + \frac{31}{2(x-3)}$ → For linear poly, it's err = $\frac{1}{2}(x-x_1)(x-x_2)f''(c)$ (3)
 Fit rembd of f''(x) & get err bd & if x is inbd, use $\frac{h^2M}{8}$

E.g. - $\frac{3x^2+x+1}{(x-1)(x-2)(x-3)}$; Take $f(x) = 3x^2+x+1$, $\therefore f(x) \mid \begin{array}{cccc} 1 & 2 & 3 \\ 5 & 15 & 31 \end{array}$, $\therefore f(x) = 3x^2+x+1 = (x-2)(x-3) \cdot 5 + \dots$
 $\therefore \frac{3x^2+x+1}{(x-1)(x-2)(x-3)} = \frac{5}{2(x-1)} - \frac{15}{2(x-2)} + \frac{31}{2(x-3)}$ \rightarrow for linear poly, it's err = $\frac{1}{2} (x-x_0)(x-x_1) f''(c)$ (3)
 F'd rule bd of $f''(x)$ & \therefore get err bd $\frac{1}{2} \frac{f''(c)}{n!}$ if n is invol'd, use $\frac{n^2}{8} M$

Trunc Err Bd: $E_1(f_b, \infty) = \frac{1}{2} (x - x_0)(x - x_1) f''(z)$; **Trunc Err in Lin Int (i.e., 1st deg only)**

$|f''(x)| \leq M_2$ & $x \in [x_0, x_1]$, then $|E_2(f, x)| = |f(x) - P(x)| \leq \frac{1}{2} M_2 |(x-x_0)(x-x_1)| M_2 \rightarrow \text{Err Bd}$
 $\therefore |f(x) - P(x)| \leq \frac{1}{8} (x_1 - x_0)^2 M_2 = \frac{1}{8} h^2 M_2 \leq \varepsilon$ (given), \therefore get apt 'h' • For n^{th} deg poly L : $E_n(f, x) = \frac{T_{n+1}(x)}{(n+1)!} f^{(n+1)}(z)$
Err in Lag Int

- Lage Interpolation can be written as: $f(x) = \sum_{n=0}^{\infty} \frac{\phi(x)}{(x-x_0)(x-x_1)\dots(x-x_n)} \times f(x_n)$, where $\phi(x) = (x-x_0)(x-x_1)\dots(x-x_n)$

- Draw back of Lag intⁿ: If an ⁿ intⁿ val is inserted, the coeff's are req'd to be recalcd
 $\equiv (x-x_0)(x-x_1) \dots (x-x_n)$

$$\text{Div'd Diff's: } [x_0, x_1] = \frac{x_1 - x_0}{1^{\text{st DD}}}, \text{ where } [x_1, x_2] \text{ etc. & } [x_0, x_1, x_2] = \frac{[x_1, x_2] - [x_0, x_1]}{x_2 - x_0} \text{ etc. } \frac{2^{\text{nd DD}}}{(n+1)!}$$

$$[x_0, x_1, x_2, x_3] = [x_0, x_2, x_3] - [x_0, x_1, x_2]$$

$x_3 - x_0$

• DGS are symm, i.e. order doesn't matt;

• same as in $[x_2, x_0, x_1] = [x_0, x_1, x_2] = [x_1, x_2, x_0]$

• in the DDs of a n^{th} deg poly l are const: $[x_0, x_1] = \frac{x_1 - x_0}{n}, \dots, [x_0, x_1, \dots, x_n] = \frac{1}{n!} \Delta^n y_0$

2nd Newton's DD Form^a: $y = f(x) = y_0 + (x-x_0)[x_0, x_1] + (x-x_0)(x-x_1)[x_0, x_1, \dots, x_n]$

Unlike last time, we can do this now because $\int_a^b f(x) dx$ is the sum of areas under the curve. Not in etc. $(x-x_1)(x-x_2)$ etc. NOT ~~int by poly~~. That'll be covered later.

- Newton-Cotes Form (Gauss Quadrature Form for equidistant nodes): $I = \int_a^b f(x) dx = \sum_{x_0}^{x_n} f(x) dx$

Trapezoidal Rule ($n=1$): $\int_{x_0}^{x_1} f(x) dx = \frac{h}{2} [(f_0 + f_1) + 2(f_1 + f_2 + \dots + f_{n-1})]$

where x_0, x_1 by a start here
Not in $f''(\bar{x}) = \max f''_i$

$$\text{Error: } E = -\frac{1}{12} h^2 [y_0'' + y_1'' + \dots + y_{n-1}''] = -\frac{(b-a)}{12} h^2 y''(\bar{x}), \text{ where } y''(x) = \max \text{ of } f'' \text{ on } [a, b]$$

$\therefore \text{Err}^n \text{ in Trapezoidal rule} = \frac{2(h^2)}{\text{Even # of subint}}$

We take a 2nd deg poly to approx curve b/w x & x+2h

$$L_{F_{xx}^2}: f = -nh^5 \cdot iv(x) \leftarrow -\frac{(b-a)}{h^4} \cdot iv(\bar{x}) \quad \left\{ \begin{array}{l} \therefore nh = b-a \\ 2 \end{array} \right\} = O(h^4)$$

- Simpson's 3/8 rule ($n=3$) : $\int_a^b f(x) dx \approx \frac{8h}{3} [f(x_0) + 3(f(x_1) + f(x_2)) + 3(f(x_3) + f(x_4)) + \dots + f(x_n)]$

$$L_F = -\frac{3x^5}{30} g'(x) = O(x^4) \quad - \text{Or same as Simpson's } \frac{4}{3} \text{rd, but may be less in } \frac{4}{3} \text{nd}$$

→ Abbns: given velo's, find des; given coord, find vol. → Calc: Use TABLE for calc'd vals
 (80) → Here 80, Not 180

Weddle's Rule: Kane's rule of revs
 $\int_a^b u \, ds = \frac{(b-a)u_1 + (b+a)u_2}{2}$

- Constr Intⁿ: Allows unequal spac^s of ord. Intⁿ have n+1 subintervals
 $\int_a^b w(x) dx = w_1 f(x_1) + w_2 f(x_2) + \dots + w_n f(x_n) = \sum_{i=1}^n w_i f(x_i)$
 Abscissae → Symmetrie mit Punkt m^{it} der Int^l

$\int_{-1}^1 f(x) dx = w_1 f(x_1) + w_2 f(x_2)$

\rightarrow 2 pt form ($n=2$): $\int_{-1}^1 f(x) dx = \frac{1}{2} [f(-\frac{1}{\sqrt{3}}) + f(\frac{1}{\sqrt{3}})] \rightarrow$ Gauss Legendre Forma
 \rightarrow Exact till 3rd deg for
 $\int_{-1}^1 x^n dx = \frac{1}{2} (-\frac{1}{\sqrt{3}}) + \frac{1}{2} (\frac{1}{\sqrt{3}})$ \rightarrow Gauss Legendre Forma
 whc: ($n+1$) $P_{n+1}(x) =$
 $\rightarrow P_n(x) - n P_{n-1}(x)$.

$$\rightarrow 3 \text{ pt. form}^a: \int_{-1}^1 f(x) dx = \frac{5}{9} f\left(-\sqrt{\frac{3}{5}}\right) + \frac{8}{9} f(0) + \frac{5}{9} f\left(\sqrt{\frac{3}{5}}\right) \quad \begin{array}{l} \text{Exact till 5th} \\ \text{deg poly!} \end{array} \quad \begin{array}{l} (2n+1)\alpha P_n(x) - n P_{n-1}(x); \\ P_0(x)=1, P_1(x)=x \end{array}$$

→ The 1st & Legendre Poly⁶: $P_0(x) = 1$; $P_1(x) = x$; $P_2(x) = \frac{1}{2}(3x^2 - 1)$; $P_3(x) = \frac{1}{2}(5x^3 - 3x)$; $P_4(x) = \frac{1}{8}(35x^4 - 30x^2 + 3)$

↳ & correct we're we are giving by: $\approx x = \int_{-1}^1 \sum_{j=0}^n \frac{(x-x_j)}{(x-x_j)^2} dx$. Roots $= x_i = 0, \pm \sqrt{\frac{3}{5}} \rightarrow \text{see!}$ (4)

Q Set - 6 : Solⁿ of ODEs

• Tay's Ser^d Me^d: $y(x) = y(x_0 + h) = y_0 + hy_0' + \frac{h^2}{2!} y_0'' + \dots$, wh^e $h = x - x_0$ ($\therefore x = x_0 + h$)
 L^e y'', y''' etc. are deriv^d fr^m giving f^o only [$y'' = f_x + f_y h$; $y''' = \frac{d}{dx}(f_x + f_y h)$ etc.]
 L^e if ask^d to trunc^a af^r x⁵ term, just write x⁵ term < E → Errⁿ & get soutⁿ x.

→ Tay's Me^d for simul^s eq¹ or² ODEs: Give, $y' = f_1(x, y, z)$; $z' = f_2(x, y, z)$ L^e $y(x_0) = y_0$
 Take & Tay's Ser^d: $y_1 = y_0 + hy_0' + \frac{h^2}{2!} y_0'' + \dots$ & $z_1 = z_0 + hz_0' + \frac{h^2}{2!} z_0'' + \dots$

→ Tay's Me^d for 2nd or² ODE: Exactly same as 1st or². Get val^s for y', y'', y''', ..., put in Tay's

• Euⁿ's Me^d: Soln as a set of val^s: $y_n = y(x_n) = y_{n-1} + h f(x_{n-1}, y_{n-1}) \rightarrow$ Slow proc^s, Need sm^{ll} h to
 L^e get accⁿ, of w errⁿ get accⁿ

• Mod^d Euⁿ's Me^d: $y_r^{(n)} = y_{r-1} + \frac{h}{2} [f(x_{r-1}, y_{r-1}) + f(x_r, y_r^{(n-1)})] \rightarrow$ Also, in 1st step, it's
 L^e For it, we apply Trape^t Rule in $[x_{r-1}, x_r]$ only. This term comes only in 2nd step NOT $\frac{h}{2} f'(x_r)$

• Runge Kutta Me^d: Don't need to calc^a h² or² deriv^s, give great accⁿ, req^r for val^s only at some spec^c pts. It agrees w^r. Tay's series after the term in h^r, wh^e r is the or^r of that RK Me^d.

(1) 1st or² RK: Euⁿ's Me^d - $y_1 = y_0 + hy_0'$, $y_2 = y_0 + hy_0' + \frac{h^2}{2} y_0'' + \dots$ Match till h RK Me^d. For Euⁿ's Form^a, put $f(x_0 + h, y_1^{(0)}) = f(x_0, y_0 + h f_0)$. Then

(2) 2nd or² RK: Mod^d Euⁿ's Me^d - For Euⁿ's Form^a, put $f(x_0 + h, y_1^{(0)}) = f(x_0, y_0 + h f_0) + \frac{h}{2} f_0(\frac{\partial f}{\partial x})$ for RHS
 exp^a the LHS & use $f(x_0 + h, y_0 + h f_0) = f(x_0, y_0) + h(\frac{\partial f}{\partial x}) + h f_0(\frac{\partial f}{\partial y}) + O(h^2)$

∴ We've: $y_2 = y_0 + \frac{1}{2}(k_1 + k_2)$, wh^e $k_1 = h f_0(x_0, y_0)$, $k_2 = h f_0(x_0 + h, y_0 + k_1)$

(3) 3rd or² RK: $y_2 = y_0 + \frac{1}{6}(k_1 + 4k_2 + k_3)$, wh^e $k_1 = h f_0(x_0, y_0)$, $k_2 = h f_0(x_0 + \frac{h}{2}, y_0 + \frac{k_1}{2})$,
 $k_3 = h f_0(x_0 + h, y_0 + k_1) \rightarrow k' = h f_0(x_0 + h, y_0 + k_1)$

(4) 4th or² RK: $k_1 = h f_0(x_0, y_0)$; $k_2 = h f_0(x_0 + \frac{h}{2}, y_0 + \frac{k_1}{2})$; $k_3 = h f_0(x_0 + \frac{h}{2}, y_0 + \frac{k_2}{2})$; $k_4 = h f_0(x_0 + h, y_0 + k_3)$
 & $k = \frac{1}{6}(k_1 + 2k_2 + 2k_3 + k_4)$; ∴ $y_2 = y_0 + k$.

• Adv^r: The opⁿ is ideal, wh^e the ODE is linⁿ or non linⁿ

• RK for a sys^m of equs: $y'' = f(x, y, z)$; $y(x_0) = y_0$; $y'(x_0) = y_0'$ → Put $\frac{dy}{dx} = p$, $\frac{dp}{dx} = f_2$ & $\frac{dz}{dx} = f_3$

Given Val^s: (x_0, y_0, p_0) , Step si^s: h, k, l, t, We have:

$$\begin{aligned} k_1 &= h f_1(x_0, y_0, p_0) \\ k_2 &= h f_1(x_0 + \frac{h}{2}, y_0 + \frac{k_1}{2}, p_0 + \frac{l_1}{2}) \\ k_3 &= h f_1(x_0 + \frac{h}{2}, y_0 + \frac{k_2}{2}, p_0 + \frac{l_2}{2}) \\ k_4 &= h f_1(x_0 + h, y_0 + k_3, p_0 + l_3) \\ y_2 &= y_0 + \frac{1}{6}(k_1 + 2k_2 + 2k_3 + k_4) \end{aligned}$$

$$\begin{aligned} l_1 &= h f_2(x_0, y_0, p_0) \\ l_2 &= h f_2(x_0 + \frac{h}{2}, y_0 + \frac{k_1}{2}, p_0 + \frac{l_1}{2}) \\ l_3 &= h f_2(x_0 + \frac{h}{2}, y_0 + \frac{k_2}{2}, p_0 + \frac{l_2}{2}) \\ l_4 &= h f_2(x_0 + h, y_0 + k_3, p_0 + l_3) \\ p_2 &= p_0 + \frac{1}{6}(l_1 + 2l_2 + 2l_3 + l_4) \end{aligned}$$

Here also p_1, \dots, p_6 will also be val^s of p here, there's no x, y in f_1

- * Computer Prog 8
- Dec to Bin: Keep divⁿ by 2, Wtⁿ remⁿ b/w up. For frac^{ns}:
 - Bin to Dec: Sinceⁿ powⁿ of 2 get mult^d, start from left. For frac^{ns}, continue mult^d as $-x^{2^1} + -x^{2^2} + \dots$
 - Dec \leftrightarrow Oct, Hexa: Just remⁿ, it's done exactly as in Bin, with 8 & 16 resp.
 - Bin \leftrightarrow Oct: Groupⁿ in 3s & ch^{re} • Bin \rightarrow Hexa: Ch^{re} in 4s & ch^{re} • Hex \rightarrow Oct: Go via Bin
 - Bin +ⁿ: $\begin{array}{r} 1111 \\ + 1010 \\ \hline 11001 \end{array}$ • Bin -ⁿ: $\begin{array}{r} 1001 \\ - 0111 \\ \hline 0010 \end{array}$ • $x^n, \frac{1}{x^n}$ as in Dec. • 1's, 2's C: -ⁿ beco^s +ⁿ
 - 1C: 1 \rightarrow 0, 0 \rightarrow 1 • 2C: 1C + 1 • Both used for signed reprⁿ (i.e. msb is for sign) 0 +ve ! E.g. - For (-5) in 2C form, wr^t its bin: 101 \rightarrow 010 2C \rightarrow 011 sign 1011 2C is act^{ly} only us^f for
 - Giveⁿ any no. (+ve/-ve), if you take its 2C, you get what^r, it's the -ve of orig^r no. in 2C form
 - E.g. Take 23, take its 2C. What^r bin no. you get, is (-23) in 2C form / Not orig bin form, but 2C form.
 - Take -35, take its 2C. What^r bin no. you get, is 35 in 2C form
 - Taking 1C also does the same thing & you get the -ve in 1C form.
 - For +ⁿ, -ⁿ us^s 1C & 2C, see Book 10 (Yellow) prog 8 part. • A-B \equiv A+(B)_{2c} • 9C, 10C same.
 - Prop^s of Bool Al^a: 1. A+B=B+A, AB=BA (Commut^v) 2. A+(B+C)=(A+B)+C, A(BC)=(AB)C [Assoc^v]
 - 3. A(B+C)=AB+AC, A+BC=(A+B)(A+C) [Pf] 4. A·1+BC=A(1+B)+BC=A·1+AB+AC=A·(1+C)+AB+AC=A²+AC+AB+AC
 - 4. A+AB=A, A(A+B)=A [Absorb^v Laws] 5. A+A^TB=A+B, A(A^T+B)=AB [Absorb^v Laws]
 - 6. AB+A^TC+BC=AB+A^TC, (A+B)(A^T+C)(B+C)=(A+B)(A^T+C) [Consensus Laws] (BC=BC(A+A^T) & (B+C)=B+C+A)
 - 7. De Morgan: $\overline{AB}=\overline{A}+\overline{B}$, $\overline{A+B}=\overline{A}\overline{B}$ 8. Duals: A, A* = Ppl^e 1 by 0, 0 by 1 & (+) by (-), (-) by (+)
 - SOP & POS: As name goes • Min^m: A Prod^t here^s all vars, 0: Compl-i-mint, 1: Not
 - Max^m: A sum having all vars, 1: Compl-i-mint, 0: Not For CSOP, retain Min & mult (X+ \bar{X}) + miss X. For CPoS, retⁿ Max & add X \bar{X} + miss X.
 - Can^l SOP/POS: A SOP/POS which only has Min/Max
 - To find SOP for Truth Table: Form Min(A^TB, A^TC etc.), wherever 0/P=1, Add all to get SOP
 - To find POS for Truth Table: - Max(A+B+C etc.), wherever 0/P=0. Mul all & get POS.
 - Aid: SOP1 & POS0 (SOP: Check 1, POS: Check 0) • Each Max = Complement of corre^t Min
 - Log^c Gates: NAND: $\overline{\overline{D}} \equiv \overline{\overline{D}}$ NOR: $\overline{D} \equiv \overline{D}$, Both are Univ^c Gates.
 - ↳ For ch⁸ any for to NAND only, write it in SOP form (not CSOP). Ppl every gate by NAND NOR
 - XOR: $\overline{\overline{D}} \equiv \overline{D}$; A^TB = $\overline{A}\overline{B} + A\overline{B}$; 0/P is 1 only when odd no. of J/P are 1. 2 or 4 J/P=1 \Rightarrow 0/P=1.
 - XNOR: $\overline{\overline{D}} \equiv \overline{D}$; A^TB = $\overline{A}\overline{B} + A\overline{B}$ = AB + A^TB; 0/P 1 iff even no. of J/P are 1.
 - NOT us^s NAND: A $\overline{\overline{D}}$, us^s NOR: A \overline{D}
 - AND us^s NAND: A \overline{D} \overline{D} \overline{D} \overline{D} AB, OR us^s NOR: B \overline{D} \overline{D} \overline{D} \overline{D} A^TB } Rule^r: NAND takes AND & ~it. Only NOR. Use it & othr^r alg^c manu^s to make NAND/NOR only gates
 - Conjunction: $\wedge/\star/\cdot/\text{AND}$ • Disjunction: $\vee/\text{+}/\text{OR}$ • Cond^c Op^r: p \rightarrow q (if p then q) } Conjunction of Taut is
 - Bicond^c Op^r: p \leftrightarrow q (p iff q) • Taut^s: Idem^s True Stat • *: Idem^s False Stat also a Taut
 - Equiv^c Formu^e: (p \vee q) \wedge \neg q \Rightarrow p, [(p \vee q) p' \Rightarrow pp' + p'q' \Rightarrow p'q \Rightarrow q, Only 0th one] } Ch^{re} & & v
 - p \rightarrow q \Rightarrow \neg p \vee q } To & & + for ease
 - Most used True Arg^t = Valid, False Arg^t = Fallacy.
 - Conju^v Nor^c Form EPOS • Disjunc^v Nor^c Form = SOP • Prin^c CNF = Can^l POS, Only PDNF = CSOP
 - Use p \rightarrow q \Rightarrow \neg p \vee q, De Morgan, othr^r laws & use the meth^d to get Min/Max } E \leftrightarrow F \leftrightarrow (E \wedge F) \vee (E \neg F) \wedge (F \neg E) \vee (F \wedge E)