

Q. 20th Jan 2026

$$\Sigma = \{a, b, c\}$$

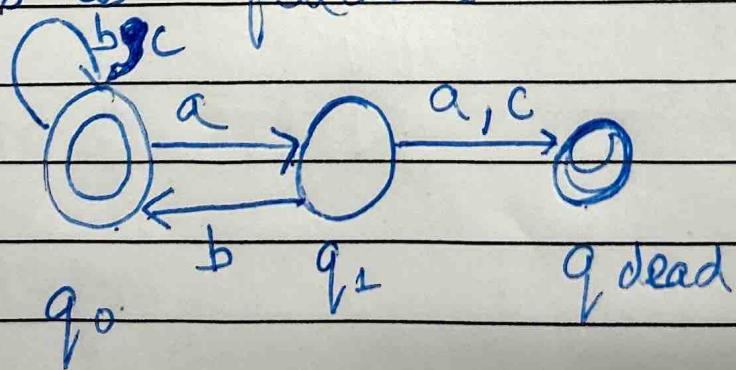
$A(n), B(n), C(n)$: n position contains a, b, c respectively.

$S(n, y) : y$ is successor of n .

Now. $\forall n (A(n) \rightarrow \exists y (S(n, y) \wedge B(y)))$

So every 'a' is followed by a 'b'

Is this regular? Yes. Finite automata is as follow:



So it is regular.

Yes, every language defined by such FOL over formula regular.

Because :

- 1) FOL over position can only talk about local relationship (adjacency, ordering, existence, etc)
- 2) They can't count unboundedly

Yes, every regular language can be expressed by such an FOL formula?

Sketch of proof:

- Let DFA have states q_0, \dots, q_n
- We introduce many predicates $Q_i(n)$ s.t. at position n , DFA state is q_i
- Now A) Initial state at position 1

2) Transition using successor:

$\forall n \forall y (S(n, y) \rightarrow$ correct transition
holds)

3) Last position must be accepting state.