Searching Arrays

Linear Search

- Find where x occurs in A
 - return some index where x appears
 - \circ for x=5, return 3
- Linear search algorithm:
 - look for it in each place until we find it
- First attempt:

```
int search(int x, int[] A, int n)
{
  for (int i = 0; i < n; i++)
    {
     if (A[i] == x) return i;
    }
}</pre>
```

x: 5

A: $\begin{bmatrix} 0 & 1 & 2 & 3 & 4 \\ 7 & 3 & 12 & 5 & 8 \end{bmatrix}$

- Remember safety!
 - A[i]: i should be *provably* in bounds
 - on is the length of A

x: 5

```
A: 7 3 12 5 8
```

Contracts!

```
int search(int x, int[] A, int n)

//@requires n == \length(A);

{
   for (int i = 0; i < n; i++)

//@loop_invariant 0 <= i;

{
    if (A[i] == x) return i;
   }
}</pre>
```

- What if x does not occur in A?
 - return something that cannot possibly be an index
 - \circ -1



```
A: 7 3 12 5 8
```

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
{
  for (int i = 0; i < n; i++)
  //@loop_invariant 0 <= i;
  {
    if (A[i] == x) return i;
  }
  return -1;
}</pre>
```

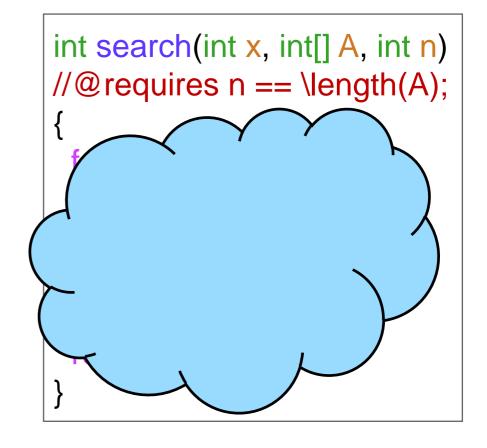
- How will a caller use search?
 - check if element was in A
 - > if the returned value is not -1
 - o if so, do something with that position
 - > e.g., update the value

x: 12

	0	1	2	3	4
A:	7	3	12	5	8

Caller

```
int k = search(12, A, 5);
if (k!= -1) {
    A[k] = 13; // changes 12 to 13
}
...
```



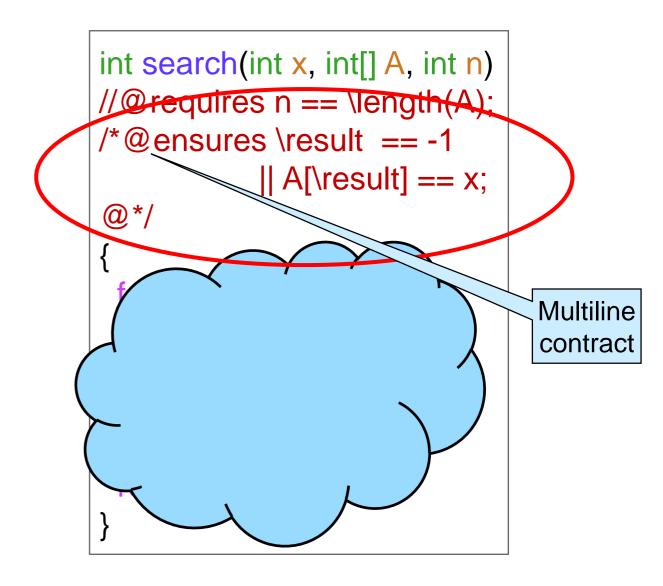
- How does the caller know how search behaves?
 - that -1 is a valid returned value
 - that A[k] contains 12
- Add postconditions!

Caller

```
int k = search(12, A, 5);
if (k!= -1) {
    A[k] = 13; // changes 12 to 13
    }
...
```

x: 12

```
A: 7 3 12 5 8
```



- Can we be sure that A[\result] is safe?
 - Extend the postcondition

x: 12

```
A: 7 3 12 5 8
```

```
int search(int x, int[] A, int n)
//@ requires n == \length(A);
/*@ensures \result == -1
            (0 \le \text{result && result < n && A[result] == x);}
@*/
 for (int i = 0; i < n; i++)
 //@loop_invariant 0 <= i;
  if (A[i] == x) return i;
 return -1;
```

- A[\result] == x won't be called if \result is out of bounds
 - && short-circuits evaluation

- Is search correct?
 - The postconditions are met when the preconditions hold
 - We'll have to prove that

- x: 12
- A: 7 3 12 5 8

- Does it do what we expect?
 - o find x in A
 - Looks plausible

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures \result == -1
             || (0 \le \text{result \&\& result < n \&\& A[\result] == x)};
@*/
 for (int i = 0; i < n; i++)
 //@loop_invariant 0 <= i;
  if (A[i] == x) return i;
 return -1;
```

Contract Exploits

Is this version of search correct?
 Postconditions are met when preconditions hold
 Definitely!

```
A: 7 3 12 5 8
```

- Does it do what we expect?
 - o find x in A
 - No!!!!

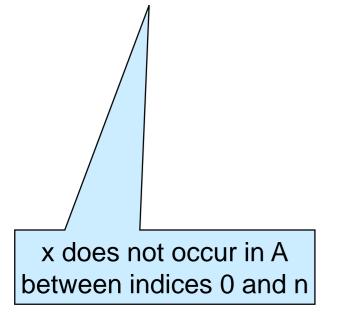
➤ always returns -1

x: 12

- This is a contract exploit
 - Postconditions are met when preconditions hold
 - > the function is correct
 - but it does not do what we expect

Fixing this Contract Exploit

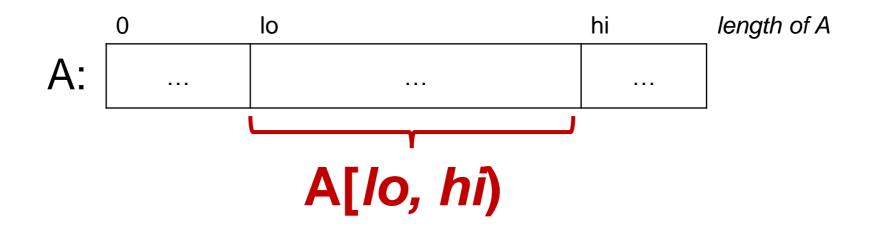
- We want search to return -1 only if x does not occur in A
 - Strengthen the postcondition to say just that
 - o !is_in(x, A, 0, n)



```
A: 0 1 2 3 4
A: 7 3 12 5 8
```

x: 12

Array Segments, in Math



A:

12

8

Segment of array A between index *lo* included and index *hi* excluded

- O Examples:
 - > A[1, 4) contains 3, 12, 5
 - ➤ A[2, 3) contains 12
 - \rightarrow A[0,5) is the entire array A
 - > A[3, 3) does not contain any element: it is an empty segment
 - > A[4, 2) does not make sense
- o we want

 $0 \le lo \le hi \le length of A$

Fixing this Contract Exploit

Let's define x ∈ A[lo, hi), in math

```
x \in A[lo, hi) = \begin{cases} false & \text{if } lo = hi \\ true & \text{if } lo \neq hi \text{ and } A[lo] = x \\ x \in A[lo+1, hi) & \text{if } lo \neq hi \text{ and } A[lo] \neq x \end{cases}
```

- Let's implement it as is_in(x, A, Io, hi)
 - This is a specification function
 - > transcription of math
 - obviously correct
 - □ used interchangeably in proofs
 - > meant to be used in contracts
 - > often recursive
 - > often no postconditions

```
bool is_in(int x, int[] A, int lo, int hi)
//@requires 0 <= lo && lo <= hi && hi <= \length(A);
{
   if (lo == hi) return false;
   return A[lo] == x || is_in(x, A, lo+1, hi);
}</pre>
```

- then, is_in(x, A, 0, n) implements $x \in A[0, n)$
 - o is x in the array segment A[0, n)? i.e., is x in A?

Fixing this Contract Exploit

Fixed code for search

```
int search(int x, int[] A, int n)
//@ requires n == \length(A);
/* @ensures (\result == -1 && !is_in(x, A, 0, n))
             || (0 \le \text{result \&\& result < n \&\& A[\result] == x)};
@*/
 for (int i = 0; i < n; i++)
 //@loop_invariant 0 <= i;</pre>
  if (A[i] == x) return i;
 return -1;
```

- Is it correct?
 - The postconditions are met when the preconditions hold

Correctness

Correctness

- search has two return statements
 - both must satisfy the postcondition
- the postcondition is a disjunction (||)
 - satisfying one branch is enough

return i on line 10

```
○ To show: if n = \lambda(A), then either \beta(A), then  either \ result = -1 & x \not\in A[0, n) or  0 \le result < n & A[result] = x
```

Looks promising

```
A. \ by line 10
```

B.
$$0 \le i$$
 by line 8

C.
$$i < n$$
 by line 7

D.
$$A[i] = x$$
 by line 10



return i on line 10

○ **To show**: if n = Vength(A), then either \result = -1 && $x \notin A[0, n)$ or $0 \le \text{Vresult} < n && A[\text{Vresult}] = x$

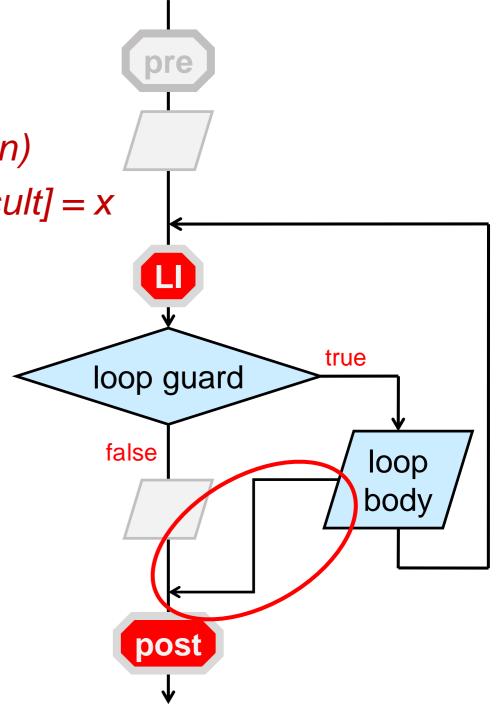
A. |result = i| by line 10

B. $0 \le i$ by line 8

C. i < n by line 7

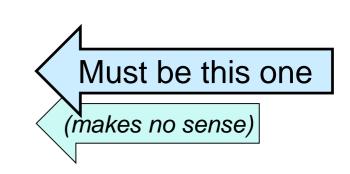
D. A[i] = x by line 10

- We did not use EXIT
 - when we return inside the loop,
 the loop invariant is not checked again



return -1 on line 12

```
○ To show: if n = \operatorname{length}(A), then either \operatorname{result} = -1 \&\& x \not\in A[0, n) or 0 \le \operatorname{result} < n \&\& A[\operatorname{result}] = x
```



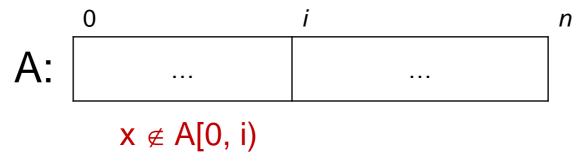
We must prove

```
x ∉ A[0, n)

No point-to argument to do so!
```

!is_in(x, A, 0, n)

• What do we know as we start iteration i of the loop?



- \circ x \notin A[0, i)
- o why?
 - > Because we looked there and didn't find x
- This is something we believe to be true at every iteration of the loop
 - O A loop invariant!
 - Well, a candidate loop invariant
 - We need to prove it is valid

return -1 on line 13

```
○ To show: if n = \operatorname{length}(A), then either \operatorname{result} = -1 \&\& x \not\in A[0, n) or 0 \le \operatorname{result} < n \&\& A[\operatorname{result}] = x
```

```
Must be this one
```

- We still need to prove
 - x ∉ A[0, i) is a valid loop invariant
 - 2. $x \notin A[0, n)$

```
bool is_in(int x, int[] A, int lo, int hi)
//@requires 0 <= lo <= hi <= \length(A);
{
  if (lo == hi) return false;
  return A[lo] == x || is_in(x, A, lo+1, hi);
}</pre>
```

$x \notin A[0, i)$ is a valid loop invariant

INIT:

```
➤ To show: x ∉ A[0, i) initially
```

A. i = 0 by line 7

B. $x \in A[0, 0) ==$ false by definition of is_in

C. $x \notin A[0, i) == true$ by math



- A[0,0) is the empty array segment
 - > Nothing is in it

```
bool is_in(int x, int[] A, int lo, int hi)
//@requires 0 <= lo <= hi <= \length(A);
{
  if (lo == hi) return false;
  return A[lo] == x || is_in(x, A, lo+1, hi);
}</pre>
```

$x \notin A[0, i)$ is a valid loop invariant

PRES:

```
➤ To show: if x \notin A[0, i), then x \notin A[0, i')
A. i' = i+1 by line 7
B. x \notin A[0, i+1) iff x \notin A[0, i) and A[i] \neq x by def. of is_in
C. x \notin A[0, i) by assumption
D. A[i] = x??
```

- - a) If true: we return on line 11
 - We exit the function
 - We won't check the loop invariant again
 - b) If false: we continue with the loop
 - We will check the loop invariant again
 - \Box $x \notin A[0, i+1)$

by B, C, D(b)

When returning from inside a loop, we don't need to show preservation

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
3. /*@ensures (... !is_in(x, A, 0, n))
               || ...;
  @*/
   for (int i = 0; i < n; i++)
    //@loop_invariant 0
  //@loop_invariant !is_in(x, A, 0, i);
10.
     if (A[i] == x) return i;
12.
13. return -1;
14.
```

return -1 on line 13

- We must prove
 - 1. $x \notin A[0, i)$ is a valid loop invariant \checkmark
 - $2. x \notin A[0, n)$

return -1 on line 13

- We must still prove $x \notin A[0, n)$
- When the loop terminates, we know that

```
\bigcirc x \notin A[0, i) by line 9
```

 $\circ i \ge n$ by line 7

To conclude x ∉ A[0, n)
 we need i = n



 Add i ≤ n as another loop invariant

○ Is it valid?

Left as exercise

return -1 on line 13

- We must still prove x ∉ A[0, n)
- When the loop terminates, we know that

```
A. x \notin A[0, i) by line 9
B. i \ge n by line 7
C. i \le n by line 8
D. i = n by B, C
E. x \notin A[0, n) by A, D
```



Scope

- When the loop terminates, we know that
 D. i = n
 by B, C
- We cannot record this with an //@assert
 - the variable i is not defined outside of the for loop
 - this mention of i would be out of scope

Compilation error

Final Code for search

```
int search(int x, int[] A, int n)
//@ requires n == \length(A);
/*@ensures (\result == -1 && !is_in(x, A, 0, n))
            || (0 \le \text{result \&\& result < n \&\& A[result] == x)};
@*/
 for (int i = 0; i < n; i++)
 //@loop_invariant 0 <= i && i <= n;
 //@loop_invariant !is_in(x, A, 0, i);
  if (A[i] == x) return i;
 return -1;
```

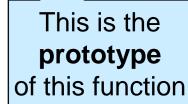
- We proved it safe and correct
- Does it do what we expect?
 - O Yes!

Testing

Client View

A caller of search can only rely on its contracts

- We may not be able to see the source code
 - > it may have been written by someone else
 - > it may be part of a library



- Can there be an implementation that satisfies these contracts but does not do what we expect?
 - > An implementation that is correct, but wrong
 - Ocan there be contract exploits?

More Contract Exploits

```
int search(int x, int[] A, int n)
//@requires n == \length(A);
/*@ensures (\result == -1 && !is_in(x, A, 0, n))
            || (0 \le \text{result \&\& result < n \&\& A[result] == x)};
@*/
 for (int i = 0; i < n; i++)
 //@loop_invariant 0 <= i && i <= n;
 //@loop_invariant !is_in(x, A, 0, i);
  A[i] = x;
                           // puts x in A[0]
  if (A[i] == x) return i; // and returns
return -1;
```

Even More Contract Exploits

```
int search(int x, int[] A, int n)
//@ requires n == \length(A);
/*@ensures (\result == -1 && !is_in(x, A, 0, n))
            || (0 \le \text{result \&\& result < n \&\& A[\result] == x)};
@*/
 for (int i = 0; i < n; i++)
 //@loop_invariant 0 <= i && i <= n;
 //@loop_invariant !is_in(x, A, 0, i);
  A[i] = x + 1;
                      // puts x+1 everywhere
  if (A[i] == x) return i; // will never return here
return -1;
```

Protecting against Contract Exploits

- The function changes the array
 - The caller has no way to know based on contracts
- What to do?
 - O Even stronger contracts?
 - ➤ Check that the array doesn't change
 - > Cannot be done in C0
 - But other languages support this

Unit testing

- ➤ Call search with a variety of inputs and check that it returns the expected value
- Usually impractical to test with all possible inputs
 - Look for inputs where errors are likely

In practice:

- write strong contracts
 - use them to reason about your code
- do thorough unit testing
 - with contracts on for smaller tests

Create a test file and write tests in its main function

- For each test
 - define input values
 - use assert
 to check that
 the function returns
 the expected result

```
int main() {
           Creates test array
               A = [3, -7]
                                      // Test #1
                                       int[] A = alloc_array(int, 2);
                                      A[0] = 3;
    3 is at index 0 of A:
                                      A[1] = -7;
search(3, A, 2) should return 0
                                       assert(search(3, A, 2) == 0);
                                       assert(search(-7, A, 2) == 1);
        -7 is at index 1 of A
                                       assert(search(42, A, 2) == -1);
           42 is not in A
                                       assert(A[0] == 3);
                                       assert (A[1] == -7);
        A wasn't changed
                                      return 0;
```

use assert to check that the function returns the expected result

assert

- aborts execution if its argument evaluates to false
- continues with the next line if it evaluates to true
- assert is not a contract
 - it is executed even when compiling without -d
 - o we cannot use \length in it
- //@assert is a contract
 - o it is executed only when compiling with -d
 - we <u>can</u> use \length in it

```
int main() {
 // Test #1
 int[] A = alloc_array(int, 2);
 A[0] = 3;
 A[1] = -7;
 assert(search(3, A, 2) == 0);
 assert(search(-7, A, 2) == 1);
 assert(search(42, A, 2) == -1);
 assert(A[0] == 3);
 assert (A[1] == -7);
 return 0;
```

- Edge cases are inputs at the edge of the input range
 - first element of an array
 - last element of an array
 - o empty array
 - 1-element array
- Test as many edge cases as possible

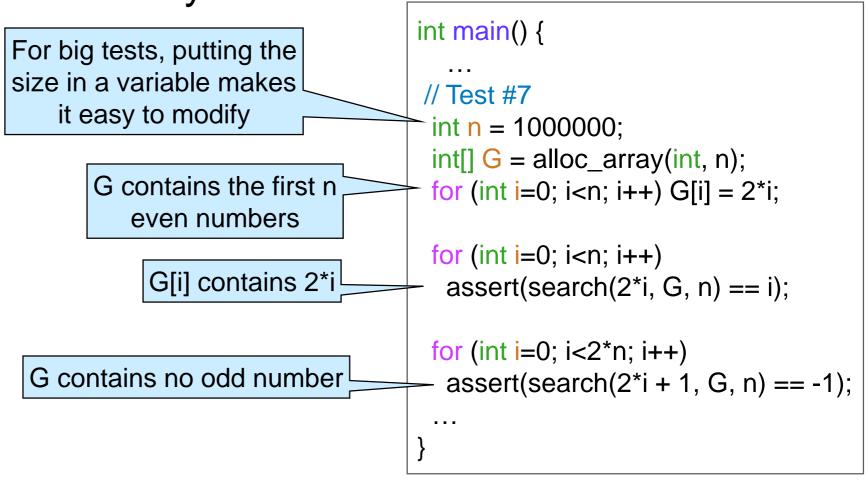
```
Creates test array
                                 int main() {
        B = [10, 11, 12, 13]
                                  // Test #2
                                  int[] B = alloc_array(int, 4);
10 is the first element of B
                                  for (int i=0; i<4; i++) B[i] = i+10;
 and 13 the last element
                                  assert(search(10, B, 4) == 0);
                                  assert(search(13, B, 4) == 3);
                                 // Test #3
        Nothing is in the
                                  int[] C = alloc_array(int, 0);
                                  assert(search(8, C, 0) == -1);
          empty array
                                 // Test #4
                                  int[] D = alloc_array(int, 1);
            Testing a
                                  D[0] = 122;
         1-element array
                                  assert(search(122, D, 1) == 0);
```

- Test inputs that are easily mishandled
 - sorted arrays
 - > with values that are
 - □ too small
 - □ too big
 - □ just right

```
int main() {
                              // Test #5
                              int[] E = alloc\_array(int, 6);
                              for (int i=0; i<6; i++) E[i] = i+1;
E is the sorted array
                              assert(search(-3, E, 6) == -1);
E = [1, 2, 3, 4, 5, 6]
                              assert(search(4, E, 6) == 3);
                              assert(search(9, E, 6) == -1);
                              // Test #6
                              int[] F = alloc_array(int, 6);
                              for (int i=0; i<6; i++) F[i] = 6-i;
         F is E in
                              assert(search(-3, F, 6) == -1);
      reverse order
                              assert(search(4, F, 6) == 2);
                              assert(search(9, F, 6) == -1);
```

 For good measure, include some big inputs and test them systematically

these are called stress tests



- best would be to use random inputs
 - > we will see later how to do that

Do not test implementation details

 anything that the function description leaves open-ended

```
H is initialized with the default int H = [0, 0, 0, 0, 0]
```

// Test #8
int[] H = alloc_array(int, 5);
assert(search(0, H, 5) == 0);
return 0;

Example: array with duplicate elements

- nothing tells us the index of which occurrence search will return
 - our implementation returns the first
 - □ but other implementations may return
 - the last
 - the middle occurrence
 - a random occurrence
 - ...

int main() {