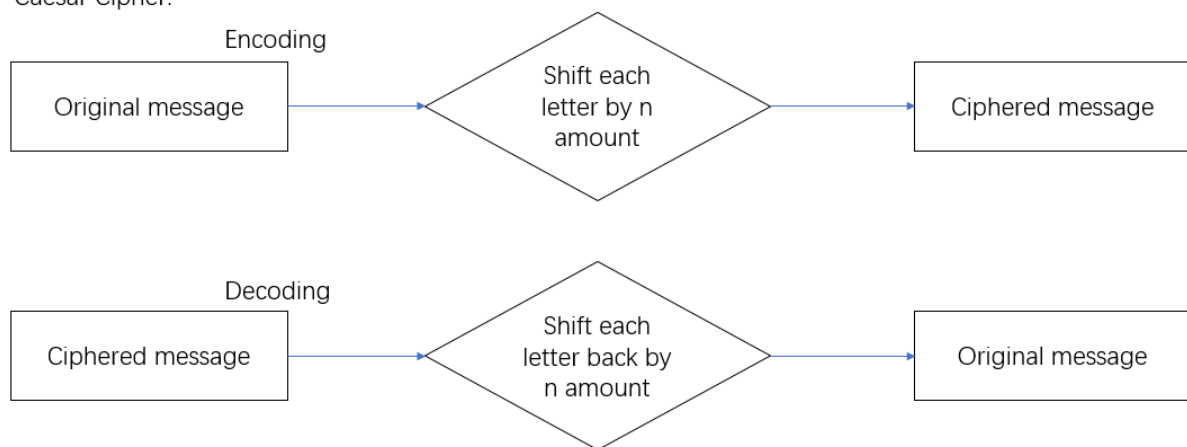


# File name: README.pdf  
# Created By: Yuxiao Wu  
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# No collaborators, no late days  
# Source: Textbook, wiki

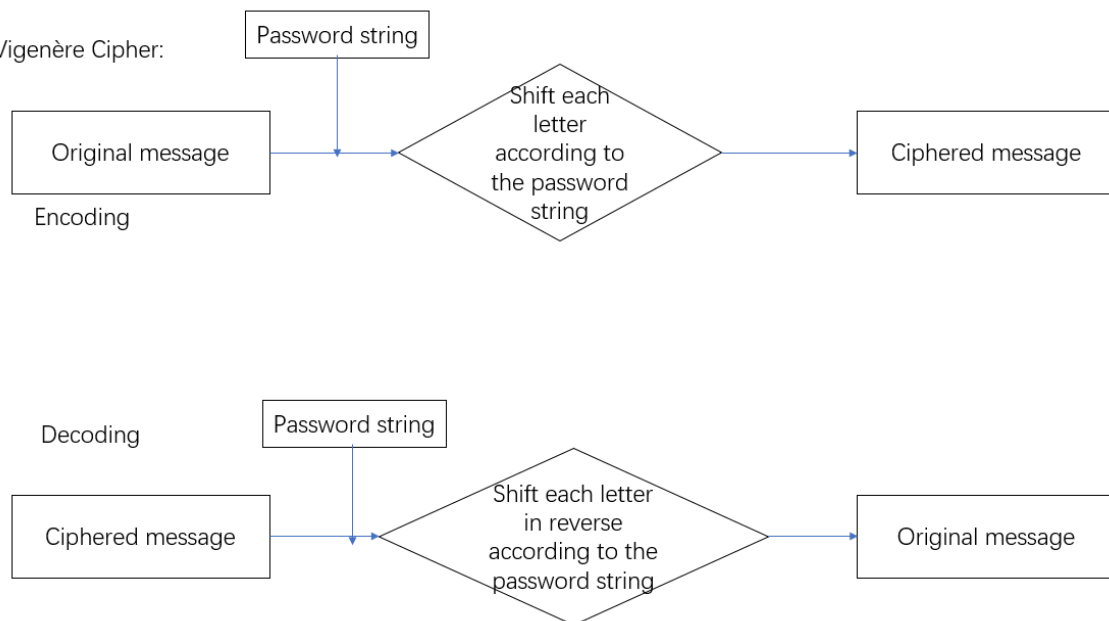
## Part 0: Design

In this lab we are going to build three different encrypting methods which are Caesar cipher, Vigenère cipher and one time pad method.

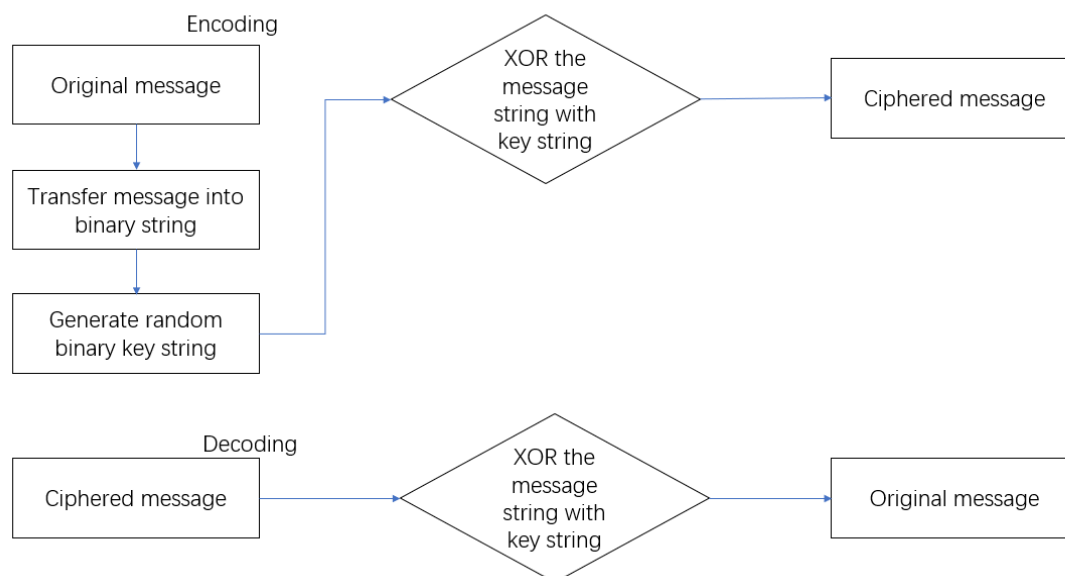
Caesar Cipher:



Vigenère Cipher:



One Time Pad:



### Extra Credit

#### 1. Break Vigenère

See in cipher.py

#### 2. Do some external research and write a paragraph explaining why and how OTP works and briefly illustrate (5-10 sentences) why it's secure (3 points)

The one-time pad (OTP) is an encryption technique, in this technique, the original message is paired with a random generated key. Then each bit of the original message is encrypted by combining it with the key bit. one-time pad technique cannot be cracked if the four conditions are met:

1. The key must be truly random.
2. The key must be at least as long as the plaintext.
3. The key must never be reused in whole or in part.
4. The key must be kept completely secret by the communicating parties.

Different from the Caesar cipher and Vigenère cipher, the rotation for every letter in ciphered message using One time pad is total random. So, it is extremely hard to break the one time pad using the same way breaking Caesar cipher and Vigenère cipher. For example, in order to decrypting a five-letter word, attacker has to try  $26^5$  which is more than 10 million times in the worst case to find the original message.

#### 3. Do some external research and write a paragraph explaining why reusing keys makes OTP insecure (2 points)

For example, let  $p_1$  and  $p_2$  be two original message to send,  $c_1$  and  $c_2$  be the ciphered message, and  $k$  be the common key used for these two messages. Then we have the

$$c_1 = p_1 \oplus k$$

math equations:  $c_2 = p_2 \oplus k$  we can derive that  $p_1 \text{ xor } p_2 = c_1 \text{ xor } c_2$ . Because the attacker can see both  $c_1$  and  $c_2$ , so the attacker can know the result of  $p_1 \text{ xor } p_2$ . After that, the attacker can use heuristic cryptanalysis with possibly a few ambiguities to recover the original message by finding the overlaps of the original message.

4. Write a function to attack OPT  
See in cipher.py.