

*ISIT 2013, Istanbul*

*July 9, 2013*

# Minimax Universal Sampling for Compound Multiband Channels

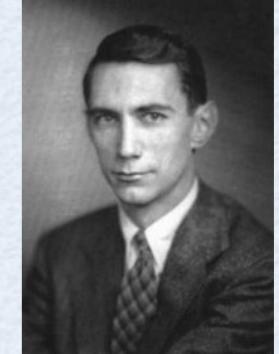
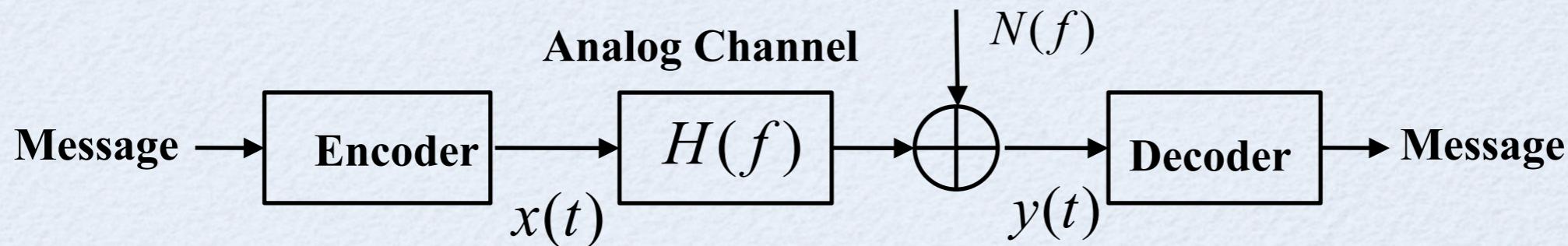
**Yuxin Chen, Andrea Goldsmith, Yonina Eldar**

*Stanford University*

*Technion*

# Capacity of Undersampled Channels

- Point-to-point channels

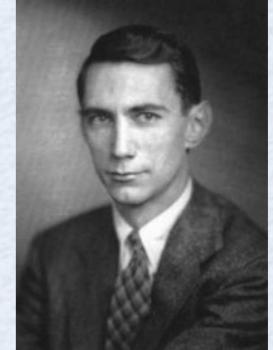
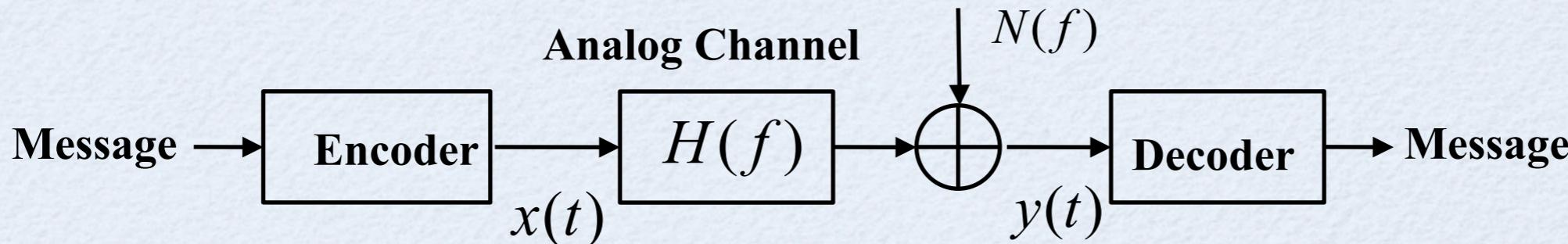


C. E. Shannon

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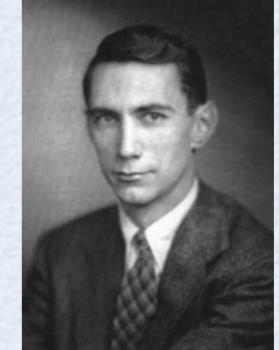
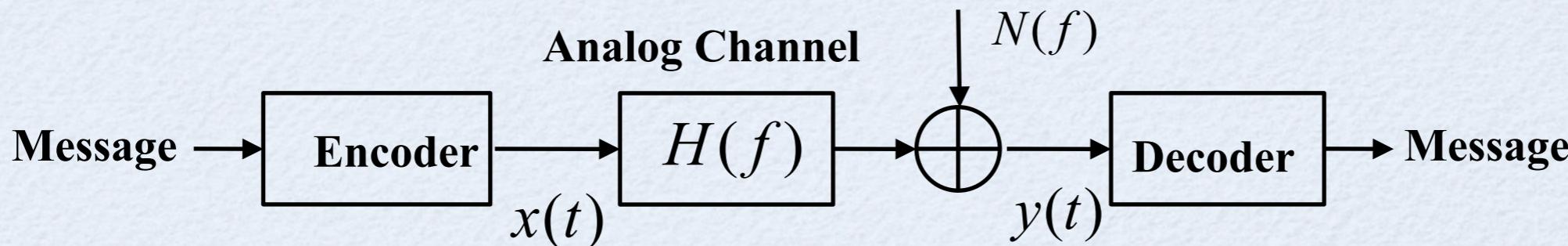
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H. Nyquist

- Sub-Nyquist sampling well explored in Signal Processing
  - *Landau-rate sampling, compressed sensing, etc.*
  - *Objective metric: MSE*

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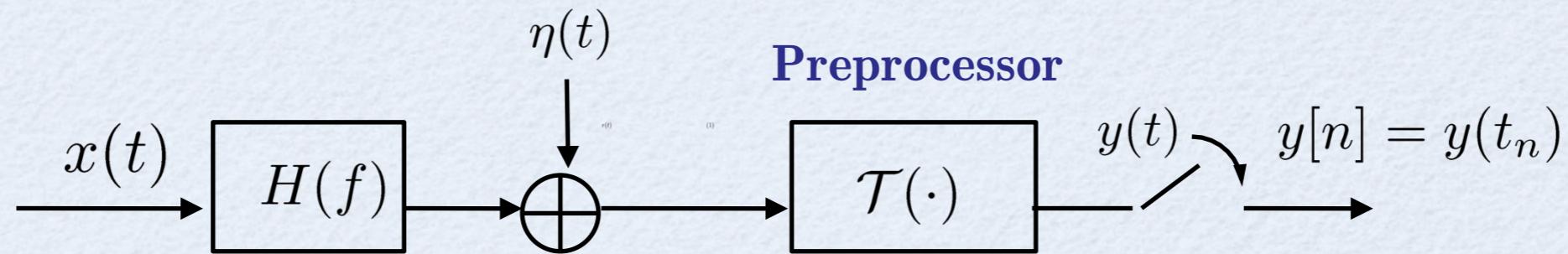
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H. Nyquist

- Sub-Nyquist sampling well explored in Signal Processing
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  - *Objective metric: MSE*
- *Question: which sub-Nyquist samplers are optimal in terms of CAPACITY?*

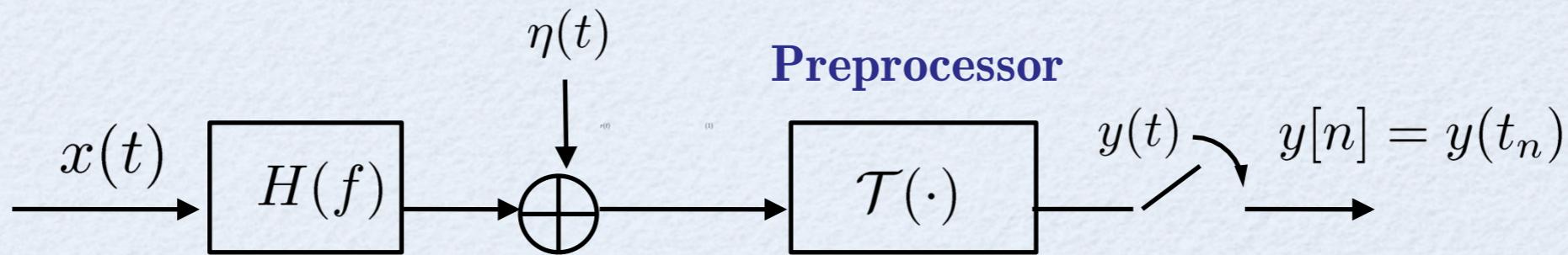
# Prior work: Channel-specific Samplers

- Consider linear time-invariant sub-sampled channels

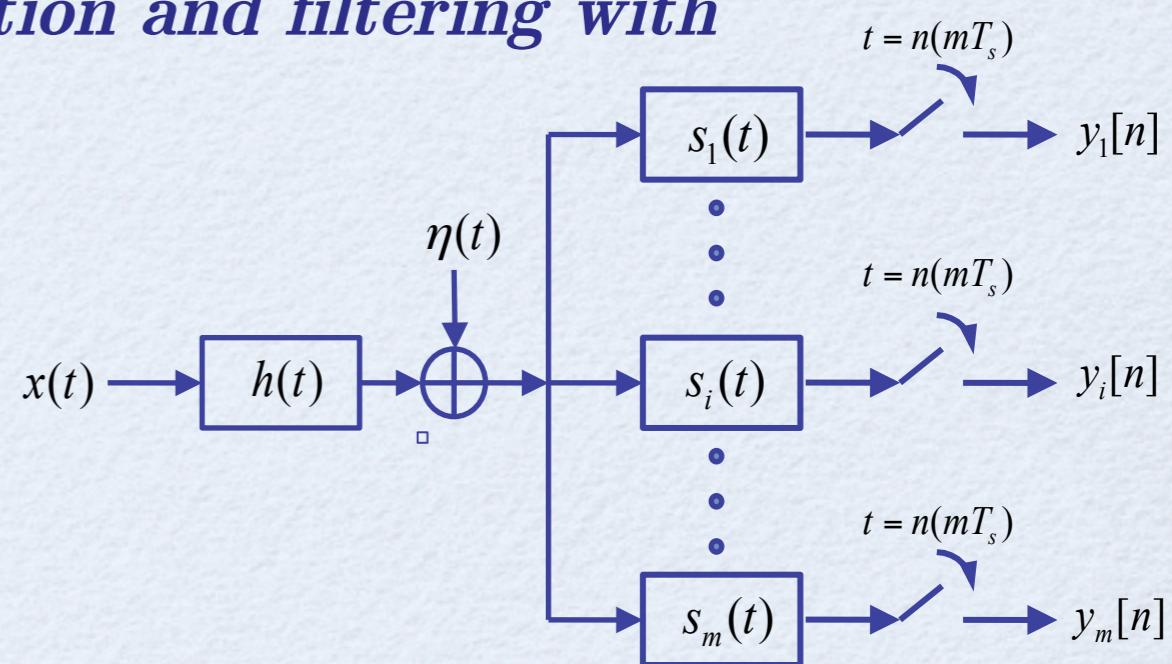


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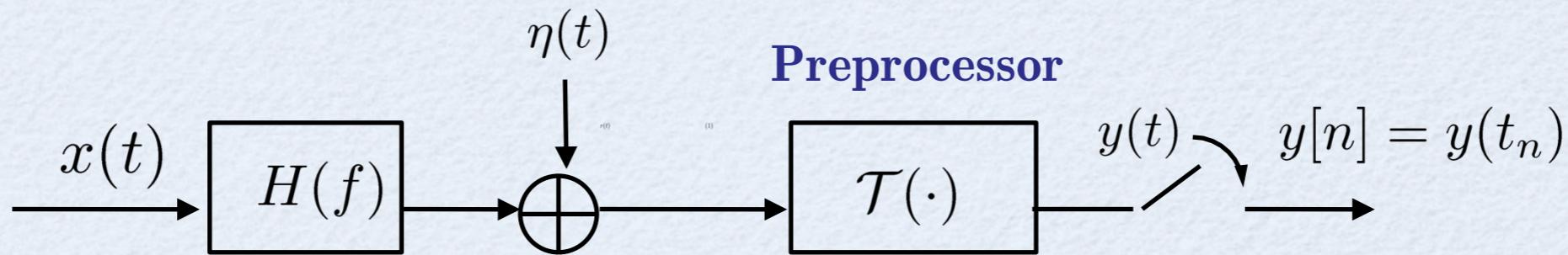


- The channel-optimized sampler (**optimized for a single channel**)
  - (1) a filter bank followed by uniform sampling
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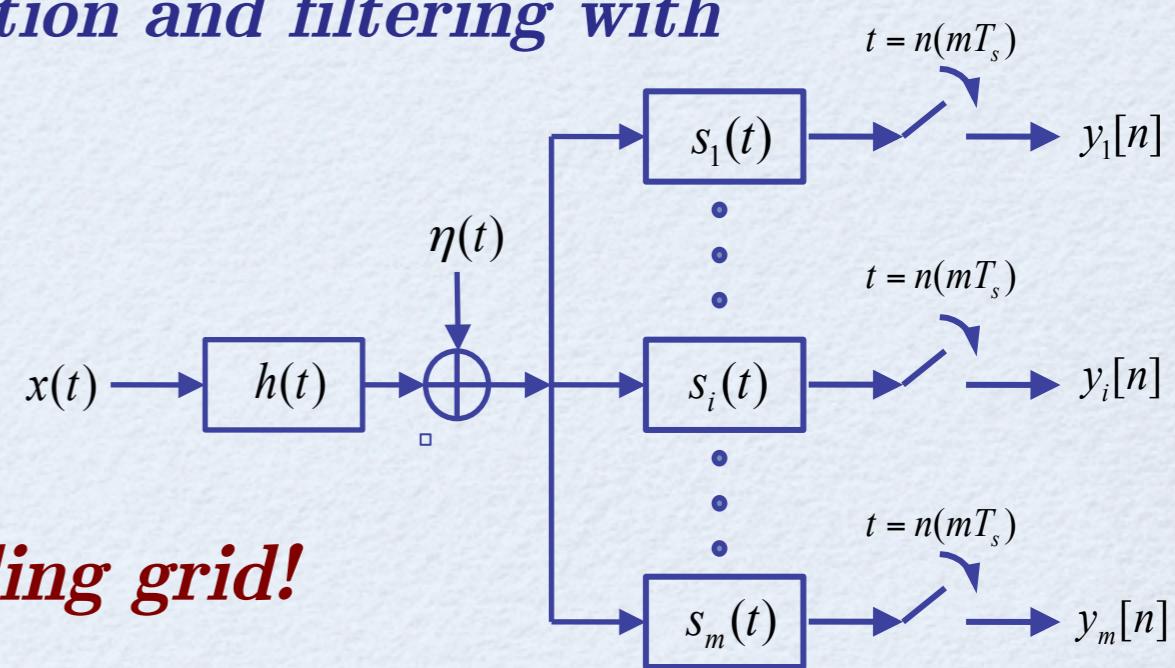


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- Suppresses Aliasing**

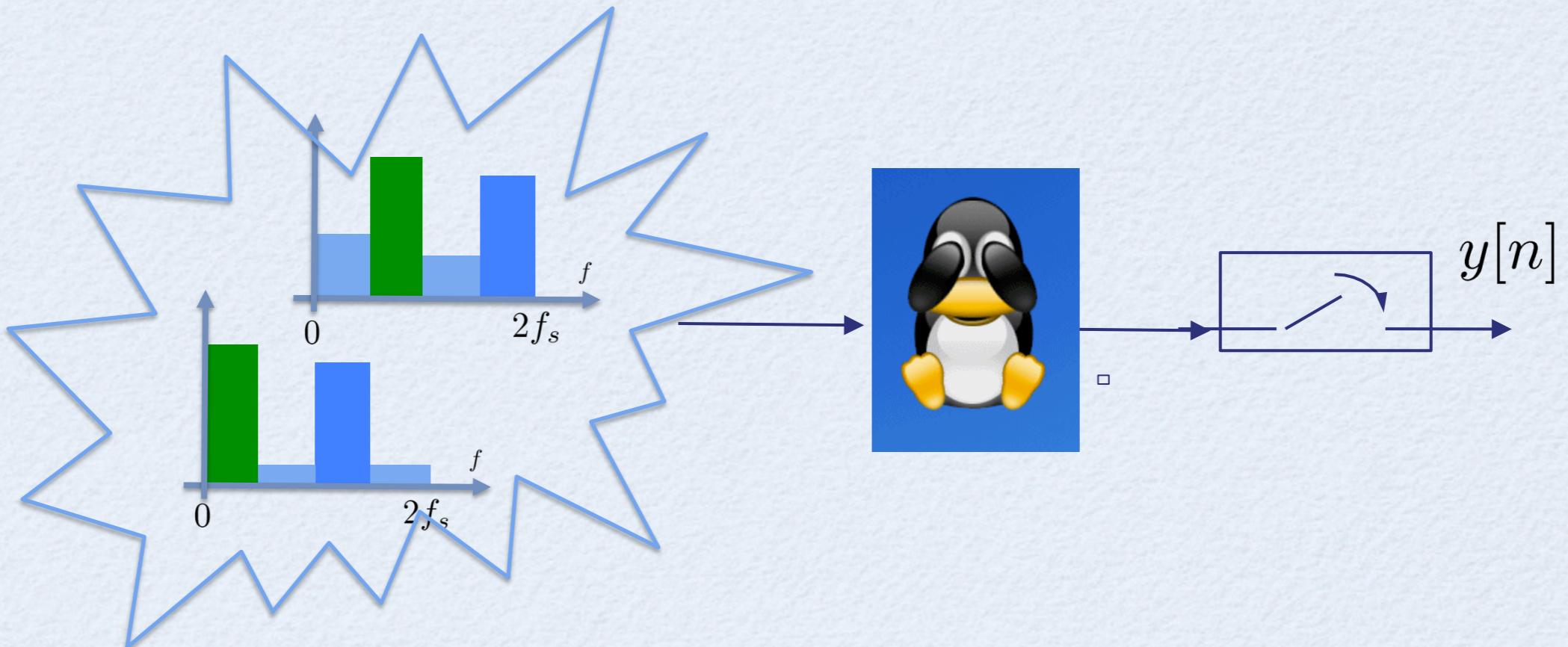
- No need to use non-uniform sampling grid!**



# Universal Sampling for Compound Channels

*The channel-optimized sampler suppresses aliasing*

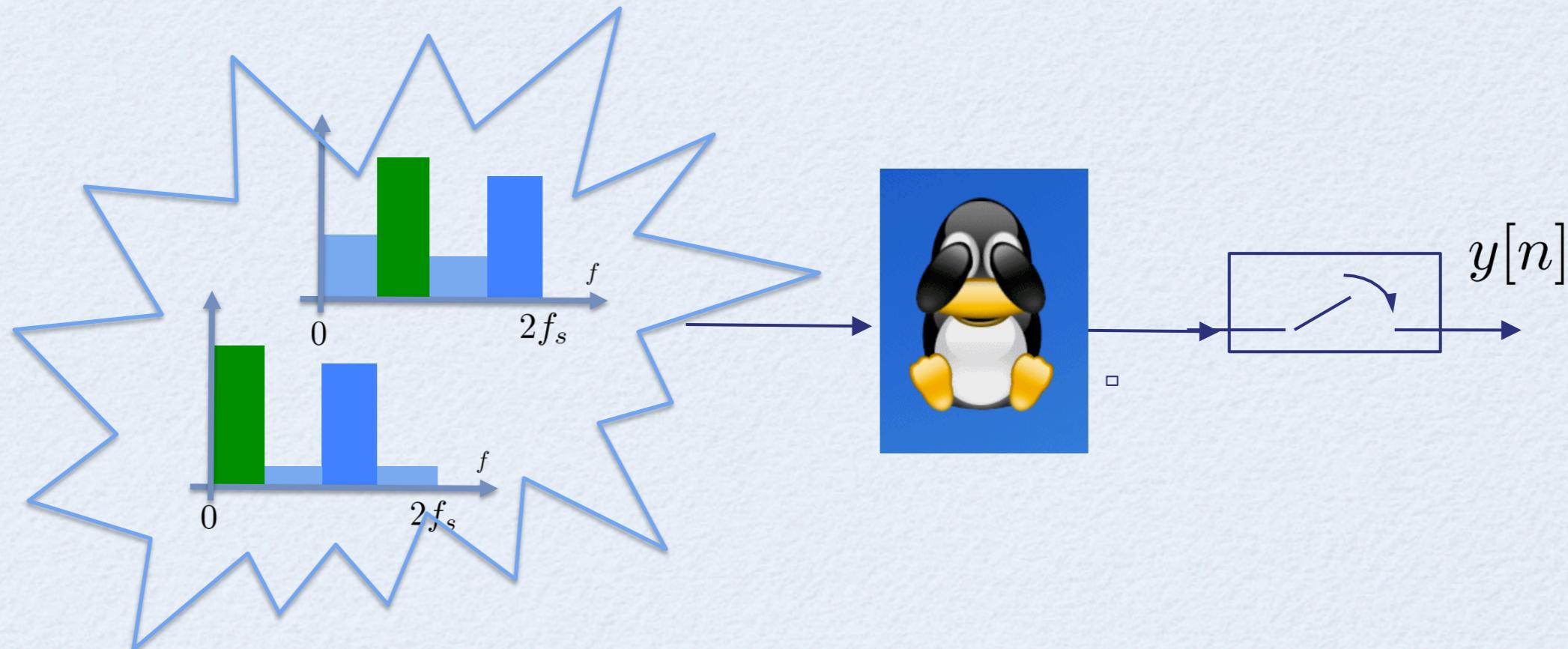
- What if there are a collection of channel realizations?



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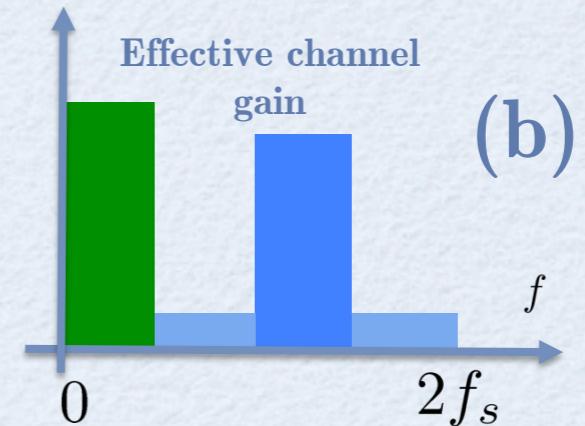
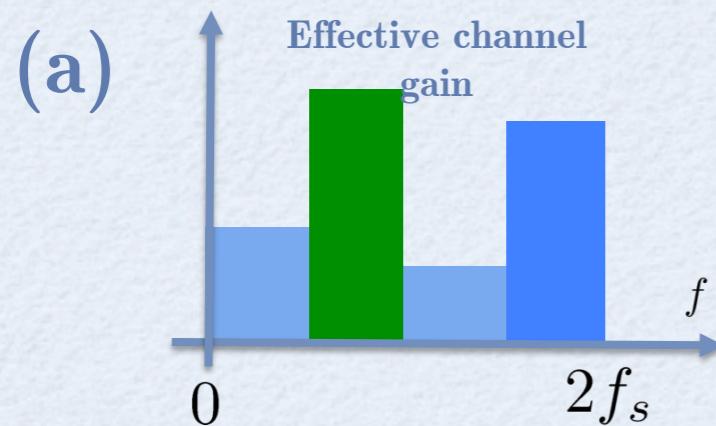
- What if there are a collection of channel realizations?



- Universal (*channel-blind*) Sampling
  - A sampler is typically integrated into the hardware
  - Need to operate *independently* of instantaneous realization

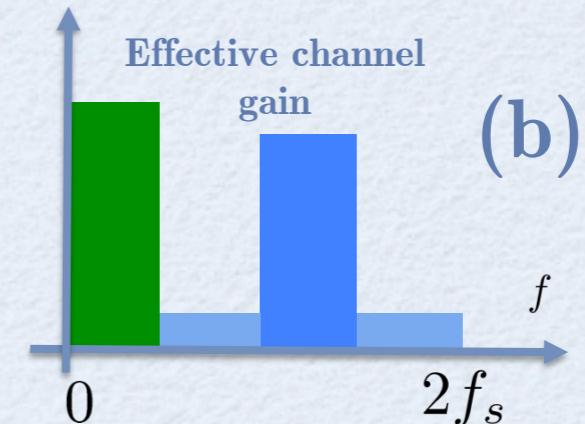
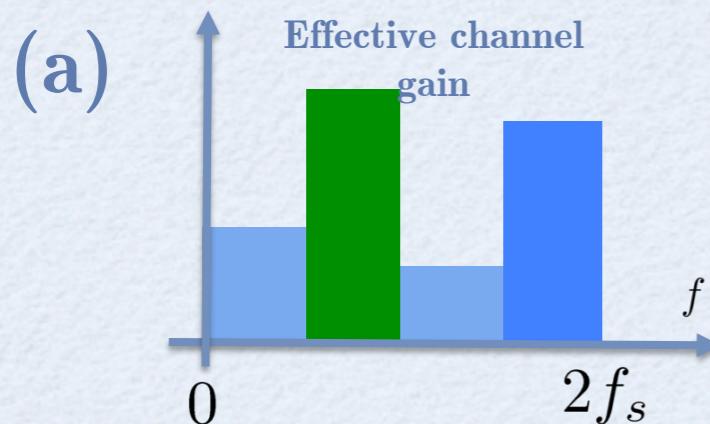
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*Consider 2 possible channel realizations .....*



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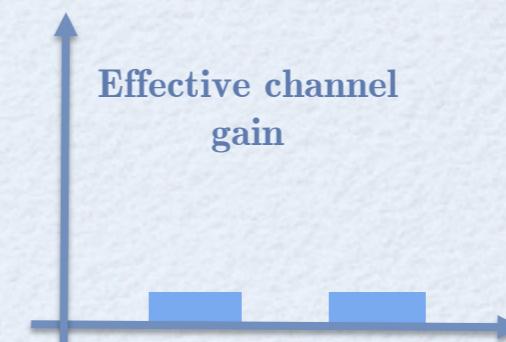
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*optimal sampler for (a)*

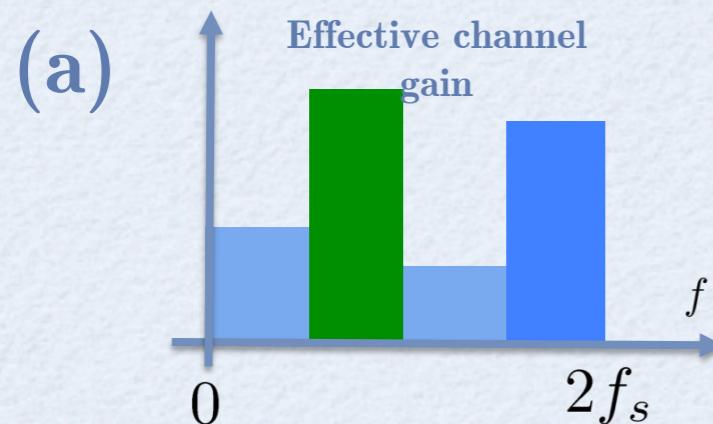


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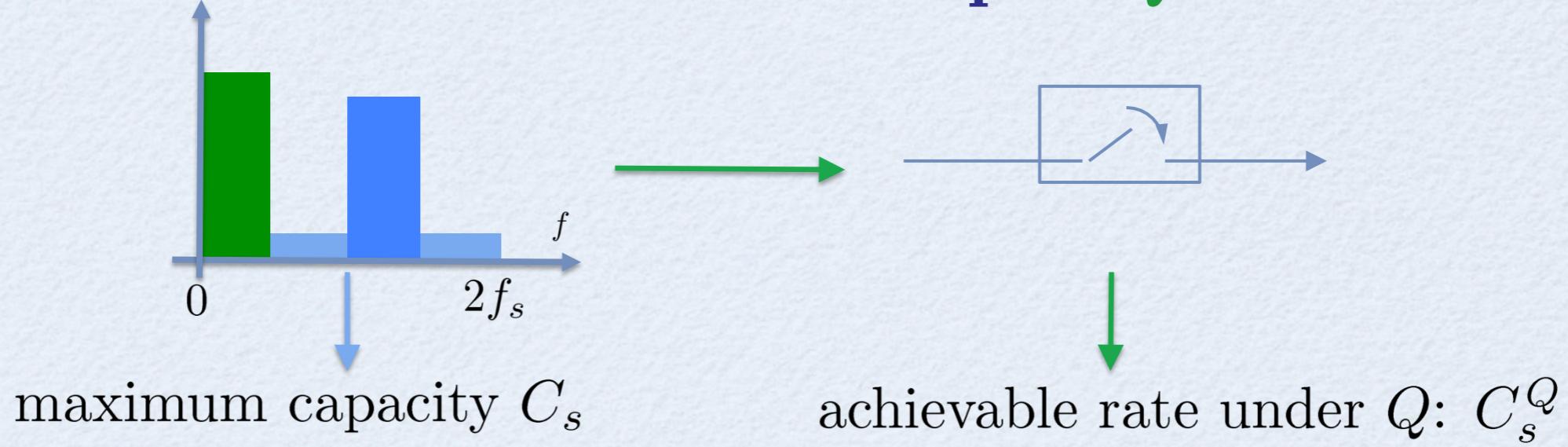
*Far from optimal!*



- No single linear sampler can maximize capacity for all realizations!
- Question: how to design a universal sampler *robust to different channel realizations*

# Robustness Measure: Minimax Capacity Loss

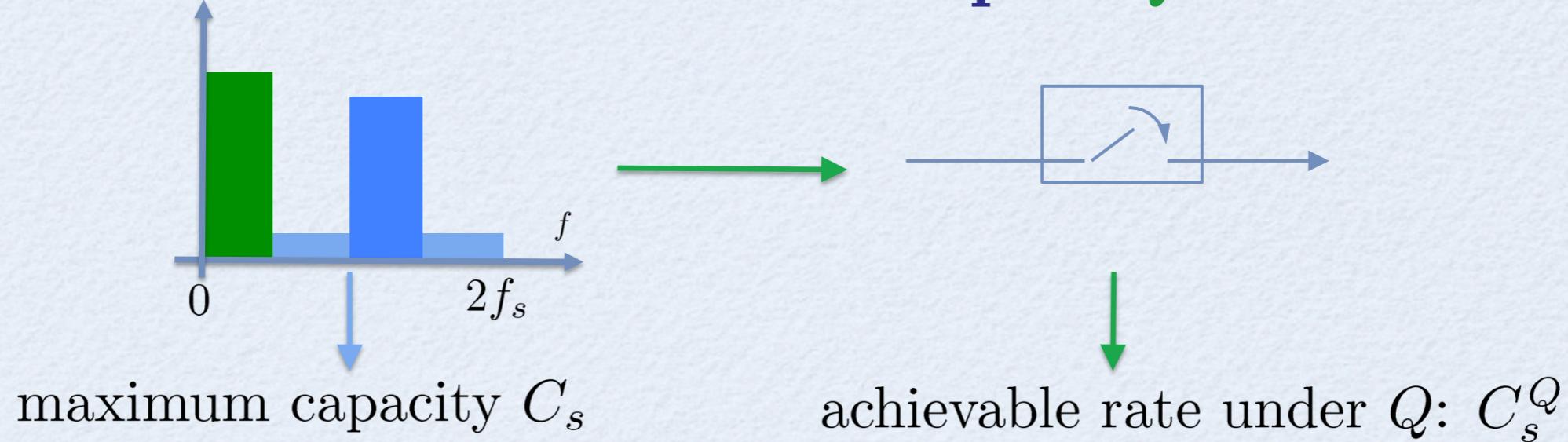
- Consider a channel state  $s$  and a sampler  $Q$ :



**Capacity Loss:**  $L_s^Q := C_s - C_s^Q$

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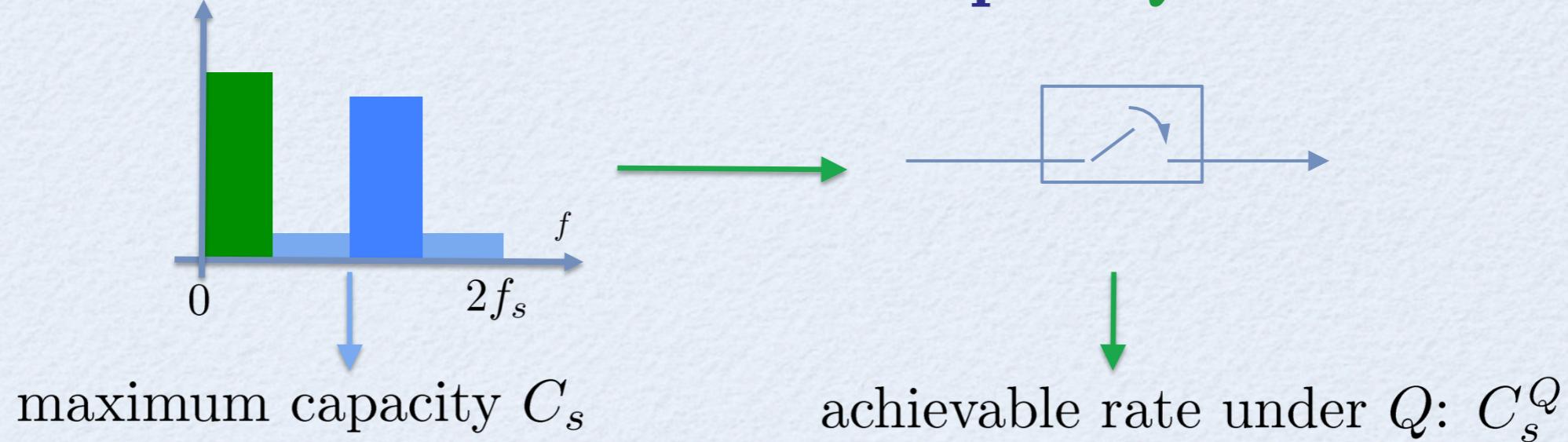
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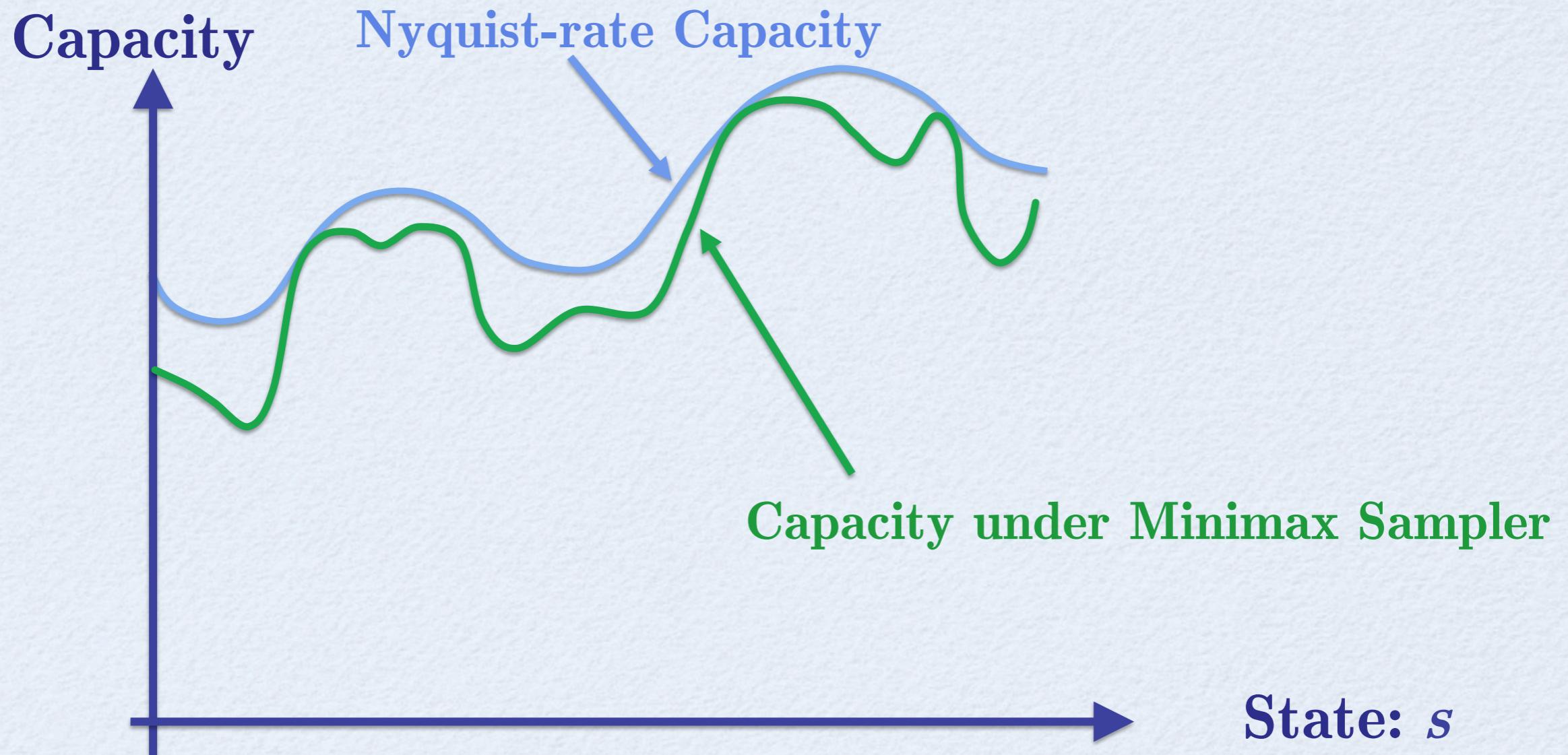
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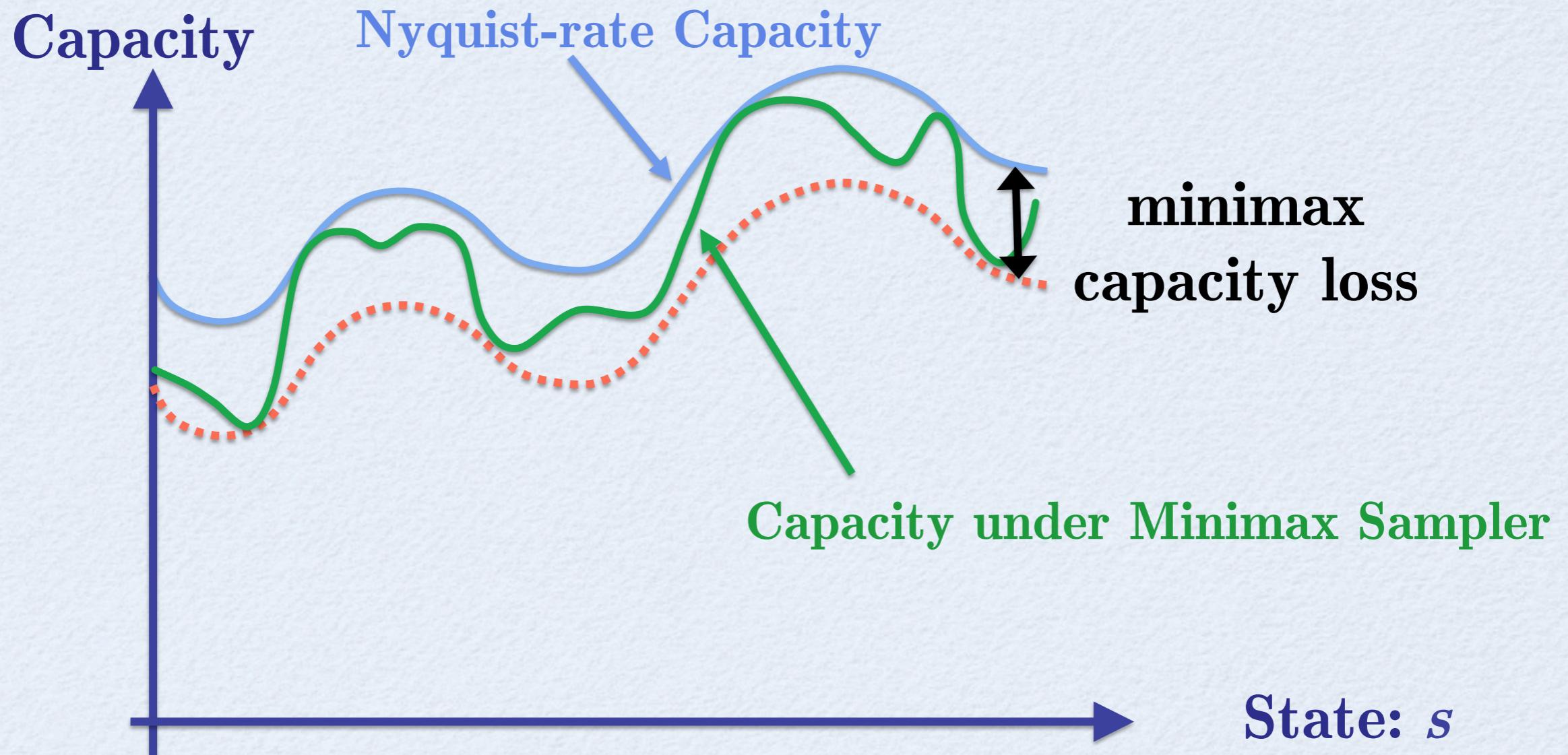
optimize over a large class of samplers  
accounting for all channel states  $s$

$Q^* = \arg \min_Q \max_s L_s^Q$  -- **Minimax Sampler**

# Minimax Universal Sampling



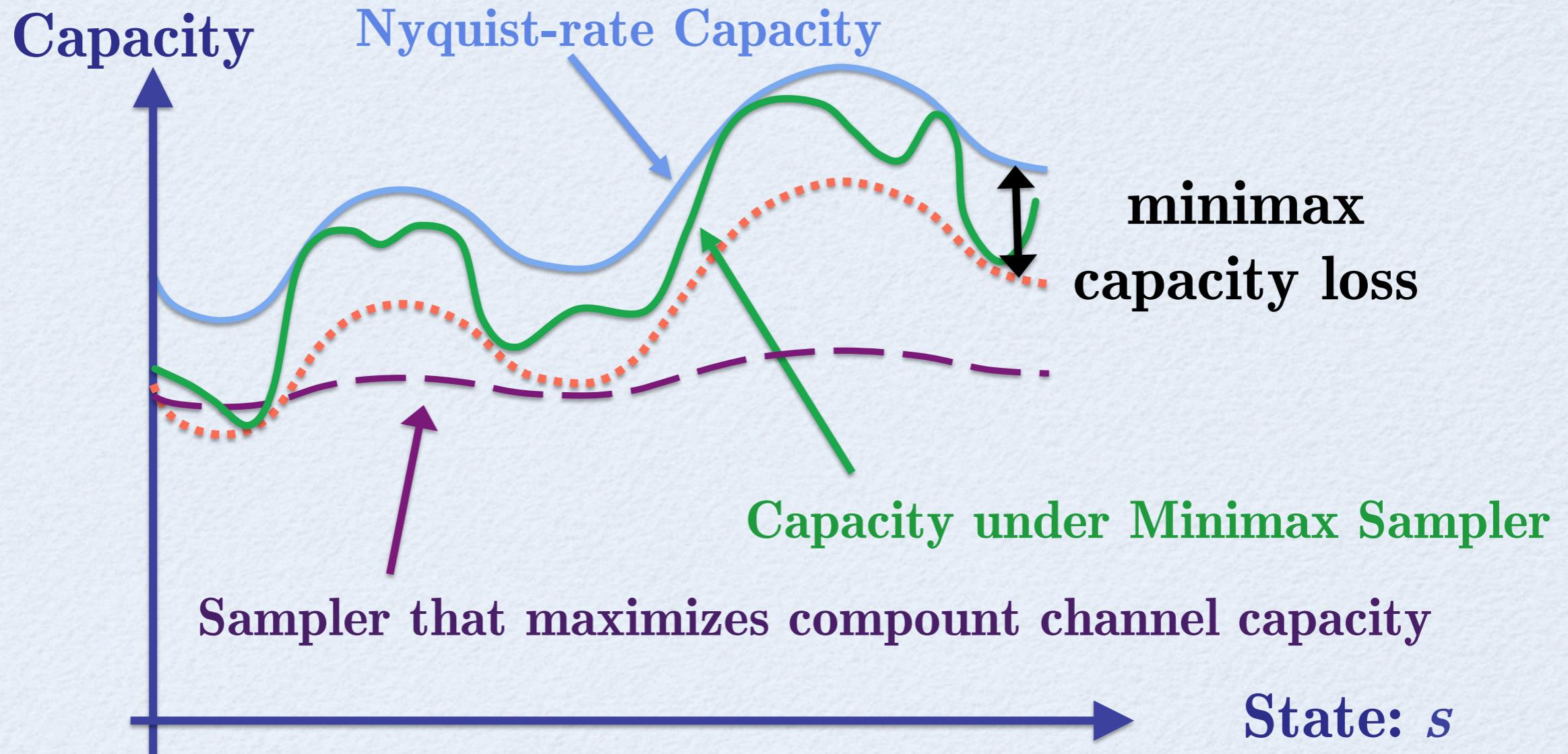
# Minimax Universal Sampling



- A sampler that minimizes the worse-case capacity loss due to universal sampling

$$Q^* = \arg \min_Q \max_s C_s - C_s^Q$$

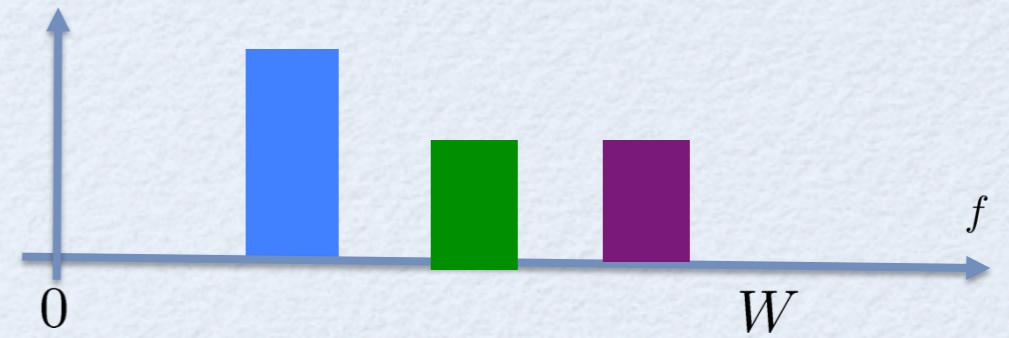
# Minimax Universal Sampling



- A sampler that maximizes compound channel capacity  $\hat{Q} = \arg \max_Q \min_s C_s^Q$
- A sampler that minimizes the worse-case capacity loss due to universal sampling  $Q^* = \arg \min_Q \max_s C_s - C_s^Q$

# Focus on Multiband Channel Model

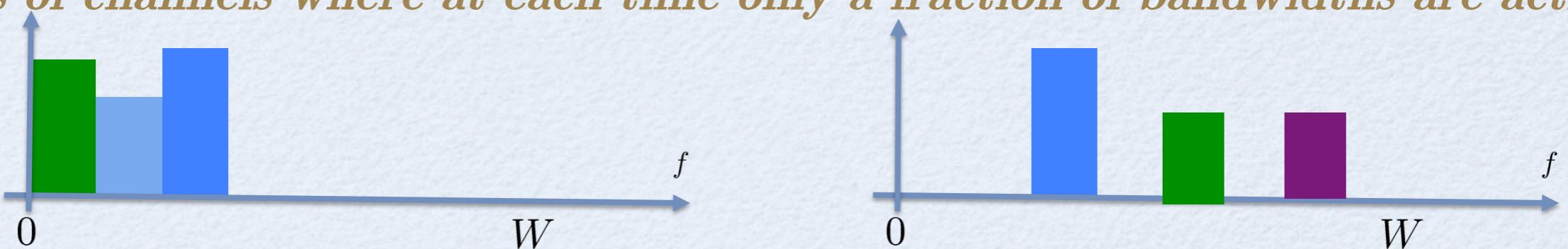
*A class of channels where at each time only a fraction of bandwidths are active.*



**$k$  out of  $n$  subbands are active.**

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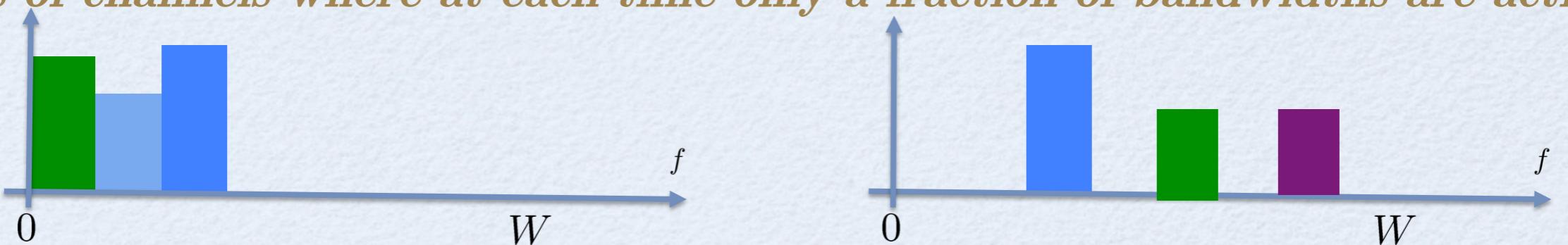
**$k$  out of  $n$  subbands are active.**

Sparsity ratio:  $\beta := k/n$

Undersampling ratio:  
 $\alpha := f_s/W$

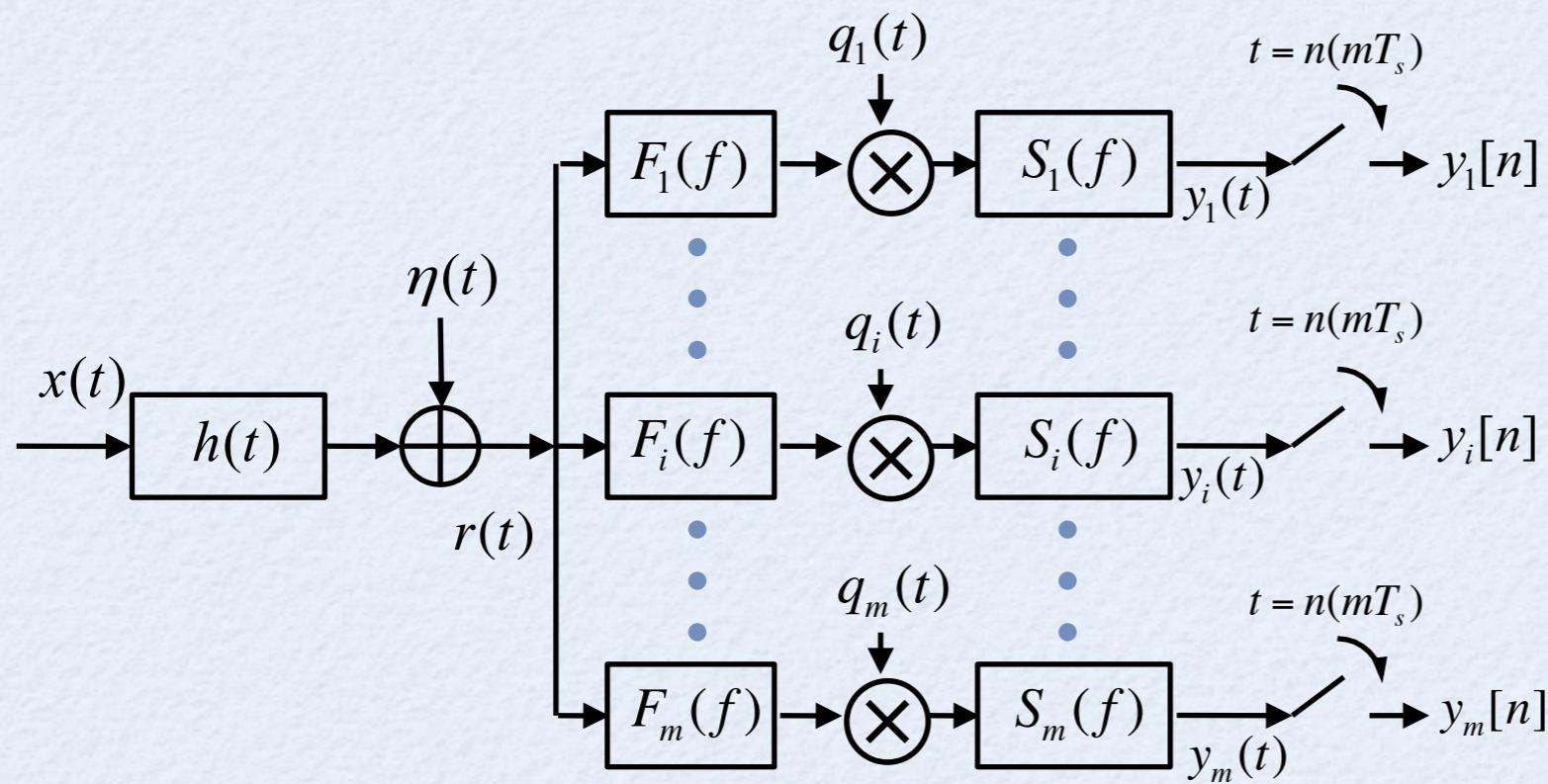
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**$m$ -branch sampling with modulation and filtering:**



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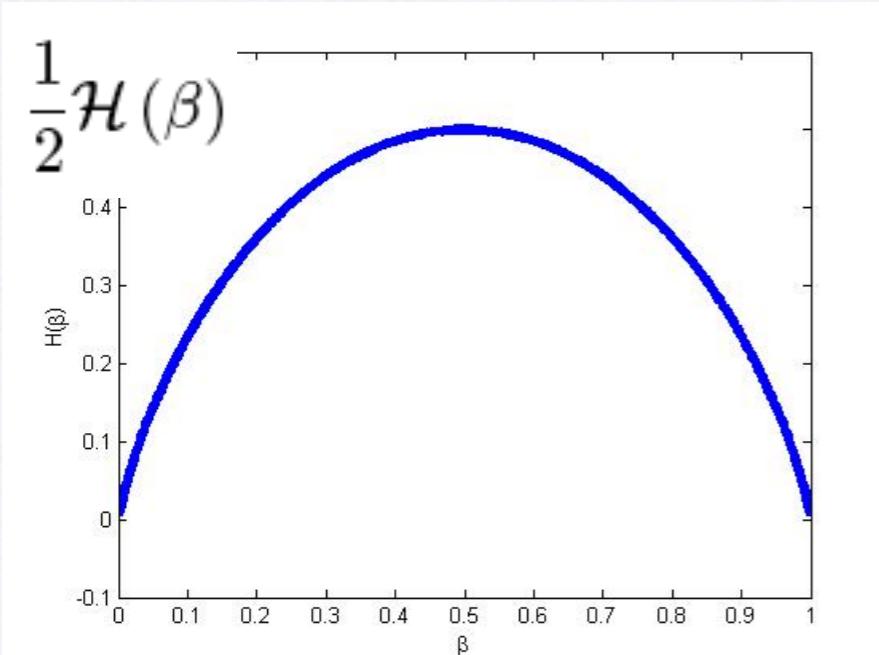
# Converse: Landau-rate Sampling ( $\alpha=\beta$ )

Sparsity ratio:  $\beta := k/n$

Undersampling ratio:  $\alpha := m/n = f_s/W$

**Theorem (Converse):** The minimax capacity loss *per Hertz* obeys:

$$\inf_Q \max_{s \in \binom{[n]}{k}} L_s^Q \geq \frac{1}{2} \left\{ \mathcal{H}(\beta) - \frac{2}{\sqrt{\text{SNR}_{\min}}} - \frac{\log n}{n} \right\}$$



Sparsity ratio:  $\beta := k/n$



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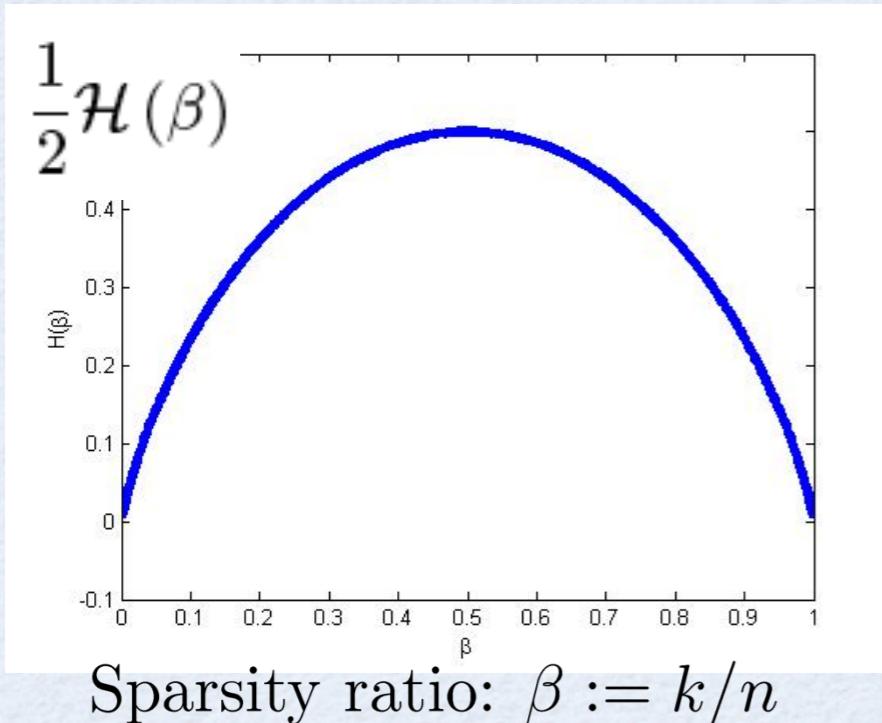
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*At high SNR and large n,*

*minimax capacity loss determined by subband uncertainty*

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**Key observation for the proof :**

$$\sum_s \exp(L_s^Q) \approx \text{constant}$$

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*The minimax sampler achieves equivalent loss across all channel states*

# Achievability: Landau-rate Sampling ( $\alpha=\beta$ )

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- ***Deterministic optimization is NP-hard (non-convex).***

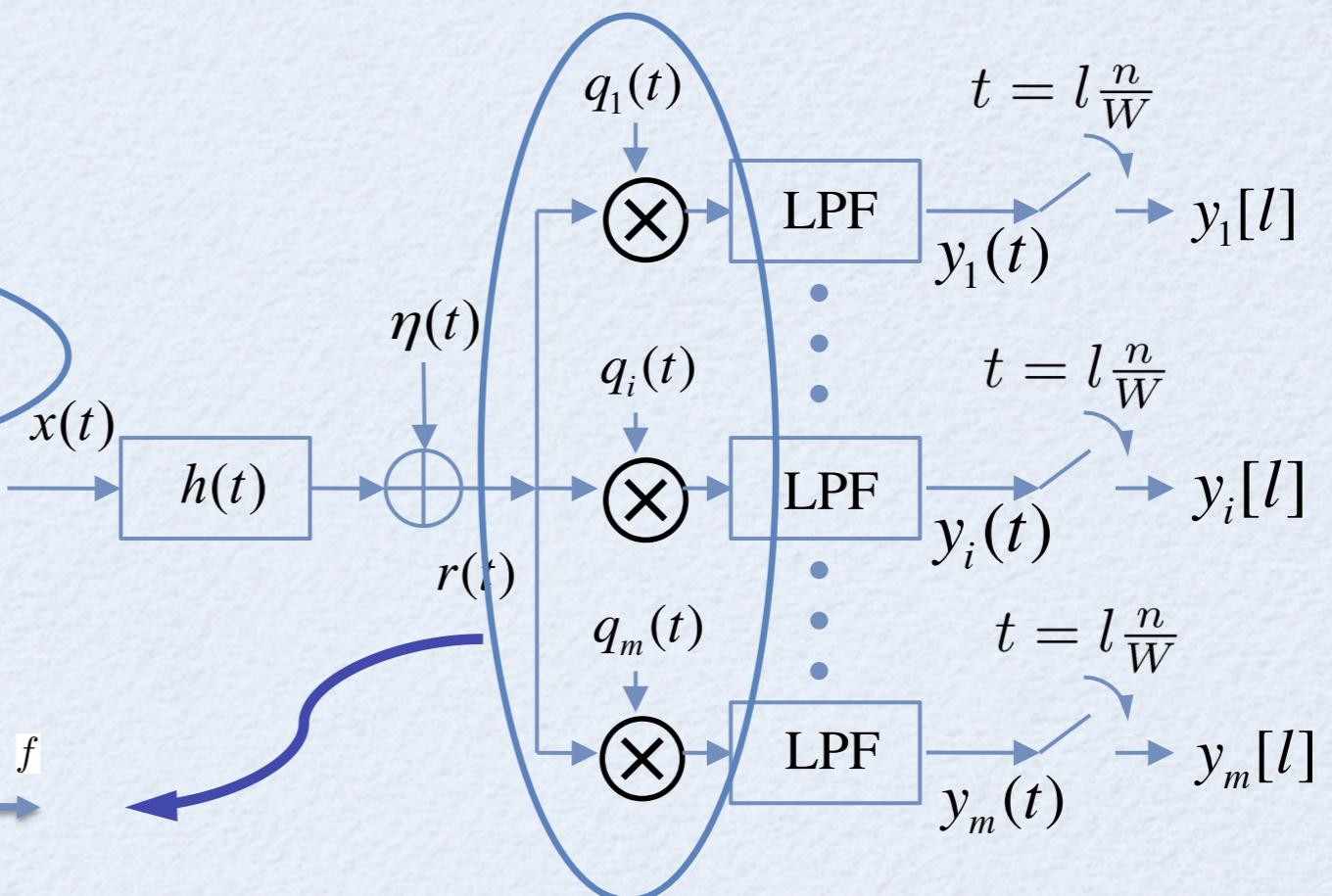
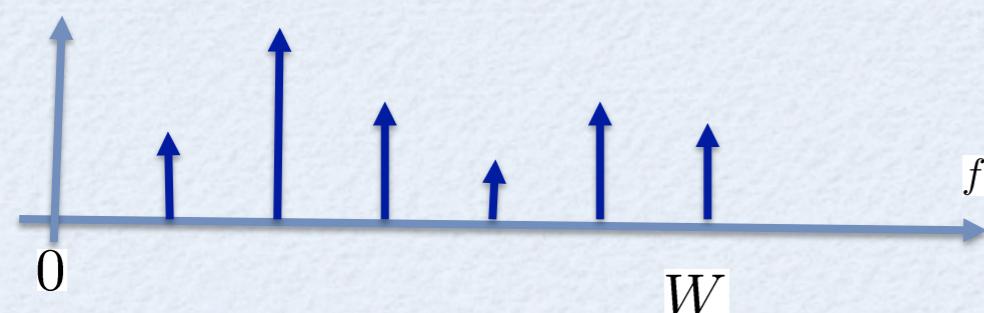
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- Hope: **random sampling**

**Fourier transform of periodic sequence is a spike-train**



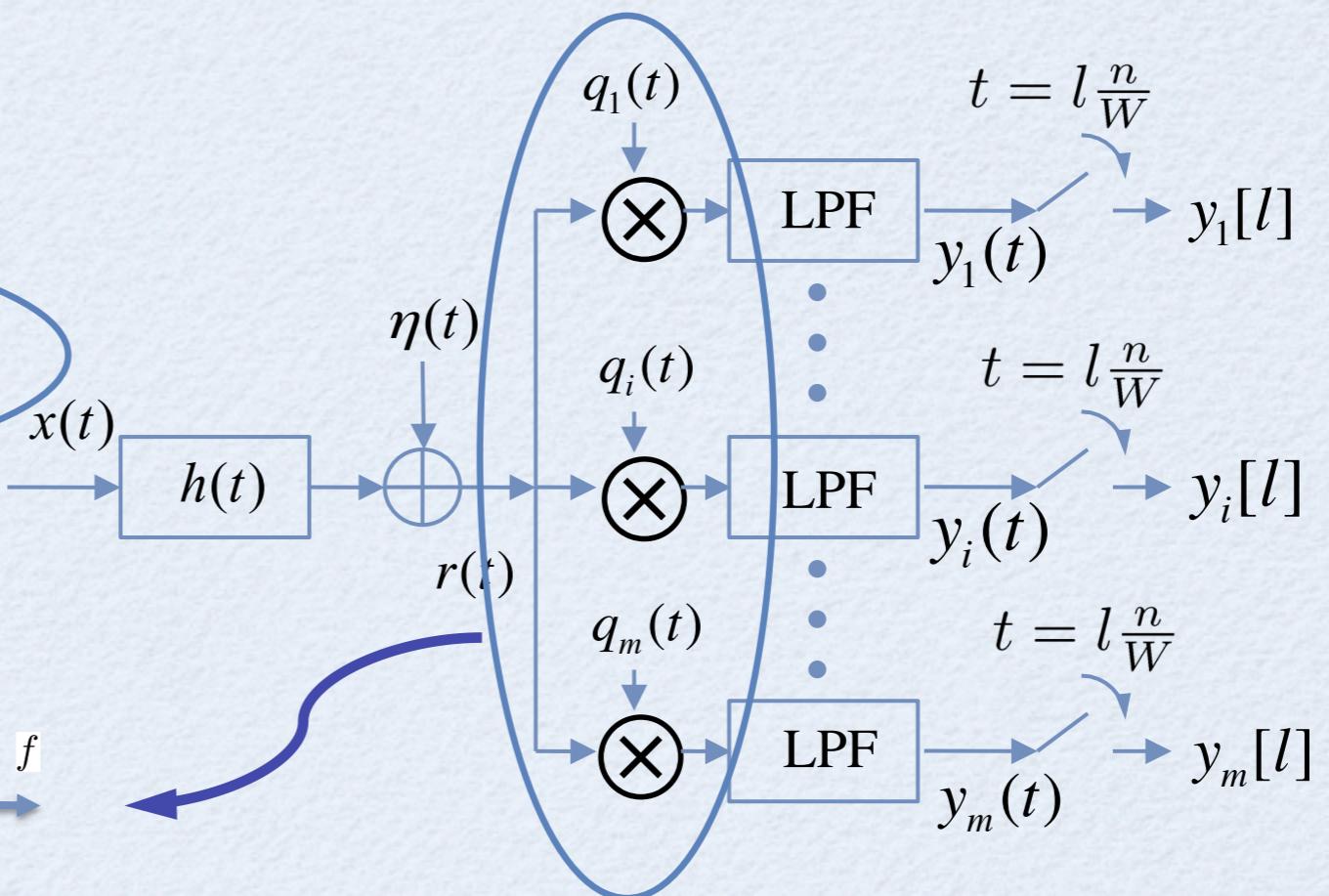
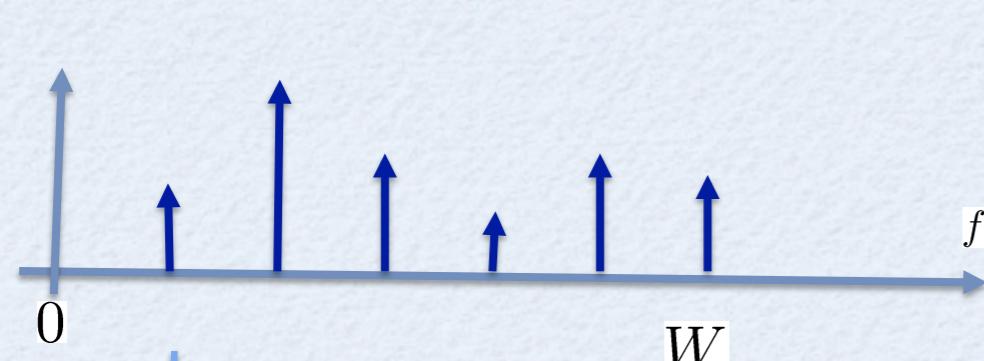
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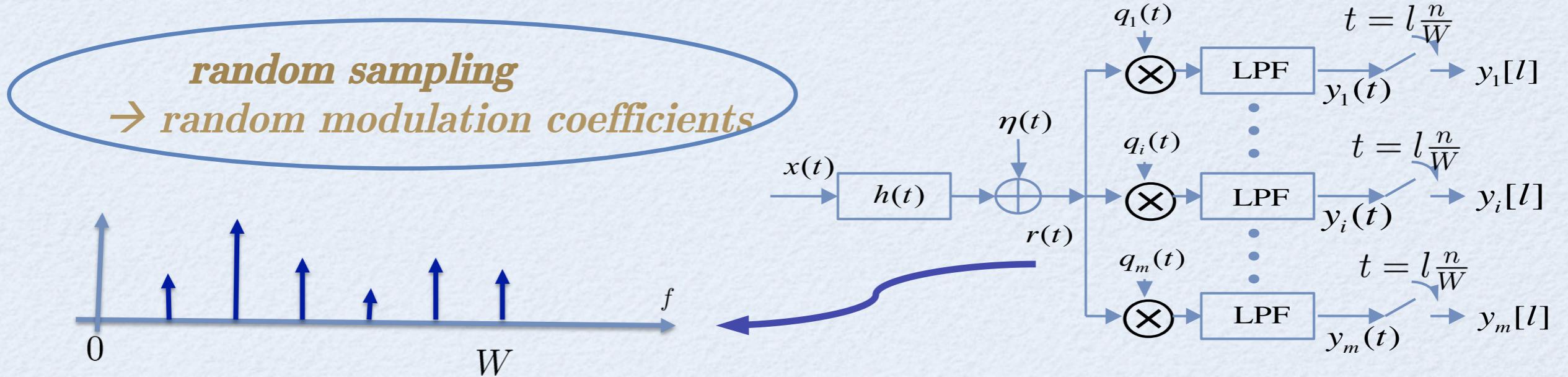


A sampling system is called *independent random sampling* if the coefficients of the spike-train are independently and randomly generated.

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**Theorem (Achievability):** The capacity loss *per Hertz* under **independent random sampling** is

$$\forall s \in \binom{[n]}{k} : \quad L_s^Q \leq \frac{1}{2} \left\{ \mathcal{H}(\beta) + \frac{5 \log k}{n} + \frac{\beta}{\text{SNR}_{\min}} \right\}$$

with probability exceeding  $1 - e^{-\Omega(n)}$ .

# Implications: Landau-rate Sampling ( $\alpha=\beta$ )

Theorem (*Converse*):

$$\inf_Q \max_{s \in \binom{[n]}{k}} L_s^Q \geq \frac{1}{2} \left\{ \mathcal{H}(\beta) - \frac{2}{\sqrt{\text{SNR}_{\min}}} - \frac{\log n}{n} \right\}$$

Theorem (*Achievability*): Under **independent random sampling** (with zero mean and unit variance), with **exponentially** high probability,

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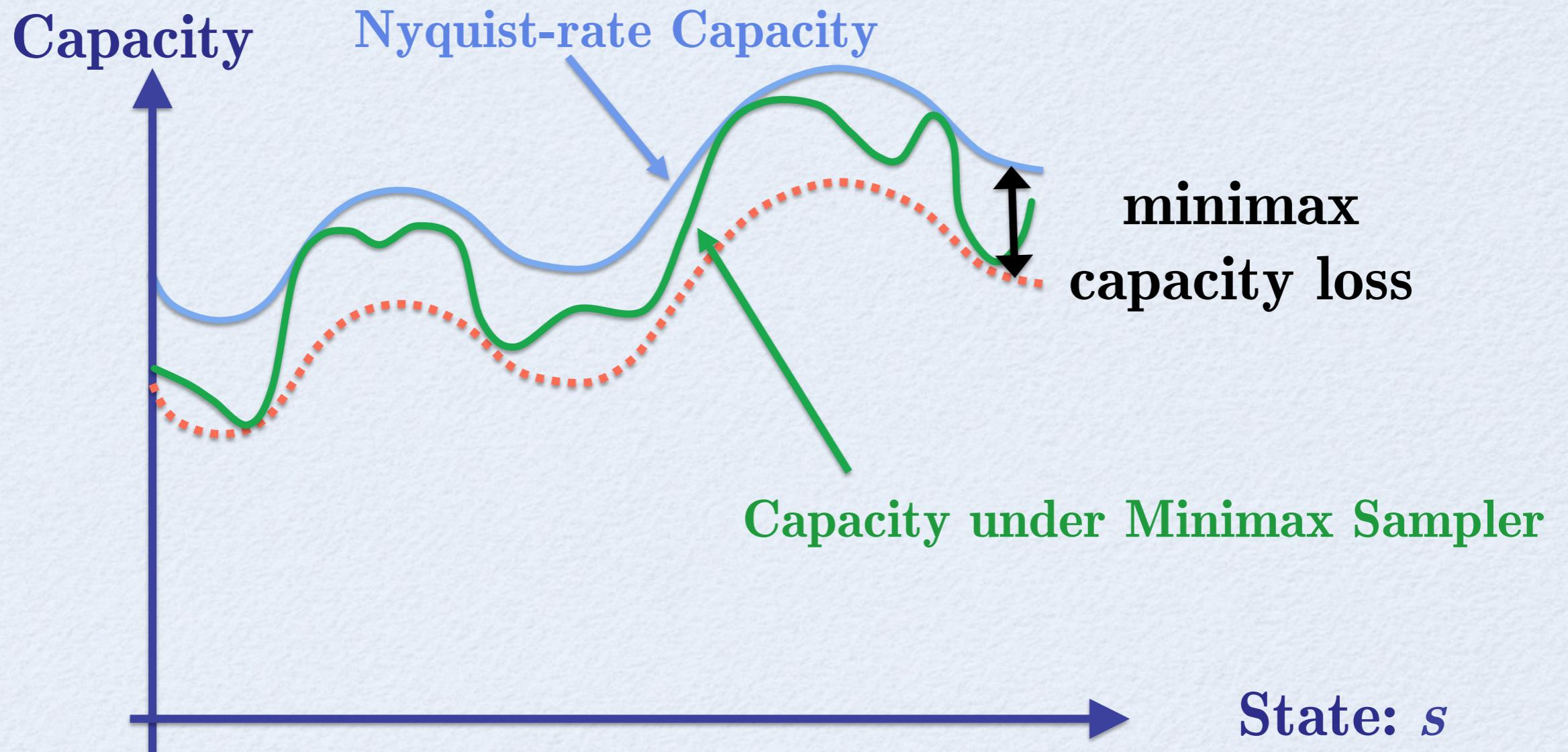
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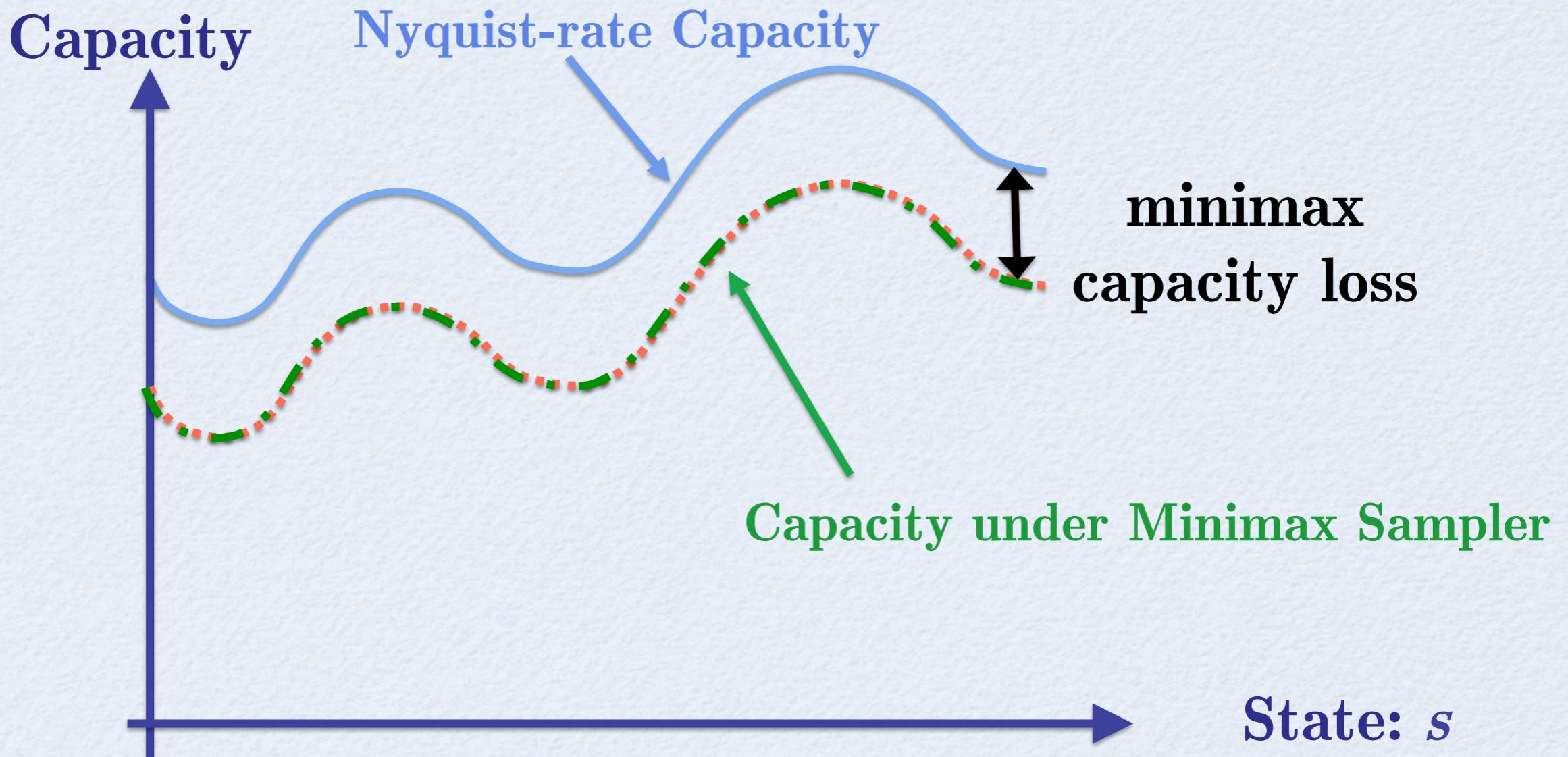
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- Random sampling is **Minimax**
- Sharp concentration – exponentially high probability
- **Universality phenomena:**
  - *A large class of distributions can work!*
    - *Gaussian, Bernoulli, uniform...*
  - *No need for i.i.d. randomness*
    - *can be a mixture of Gaussian, Bernoulli, uniform...*

# Capacity Loss for Multiband Channels



# Capacity Loss for Multiband Channels



Minimax sampling yields *equivalent capacity loss* over all possible channel realizations when SNR and  $n$  are large!

# Converse: Super-Landau Sampling ( $\alpha > \beta$ )

Sparsity ratio:  $\beta := k/n$

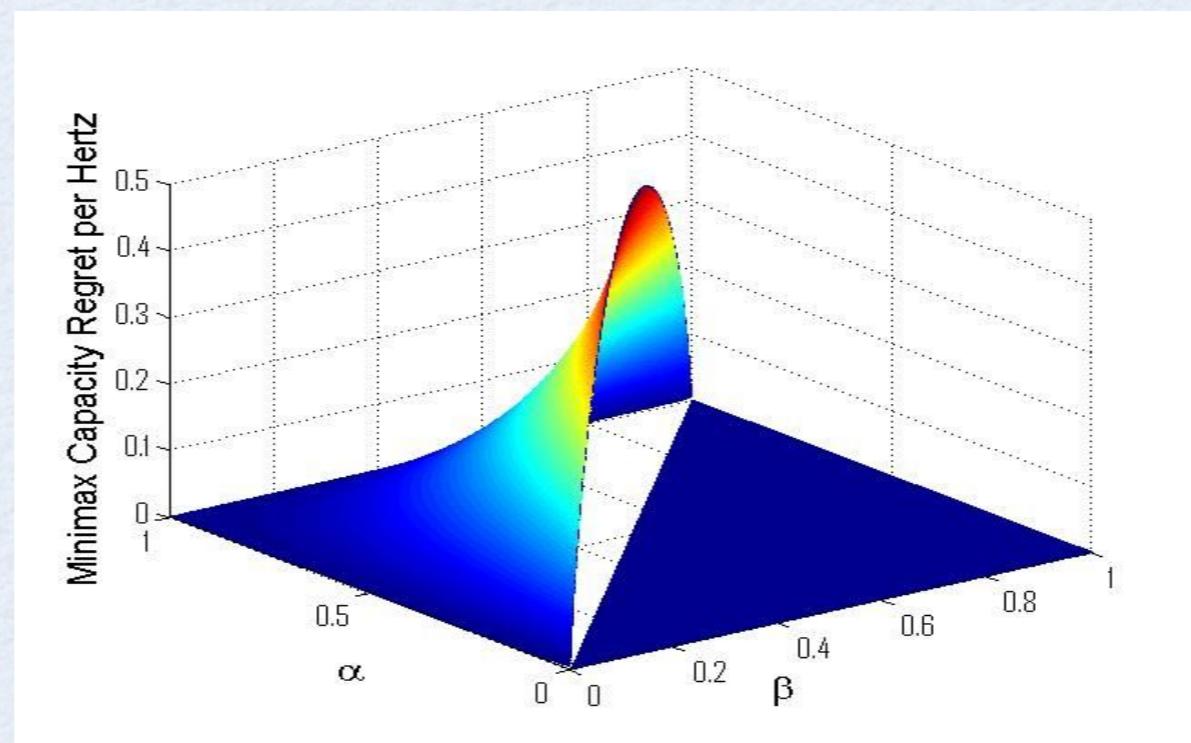
Undersampling ratio:  $\alpha := m/n = f_s/W$

**Theorem (Converse):** The minimax capacity loss *per Hertz* obeys:

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- Capacity gain due to oversampling is

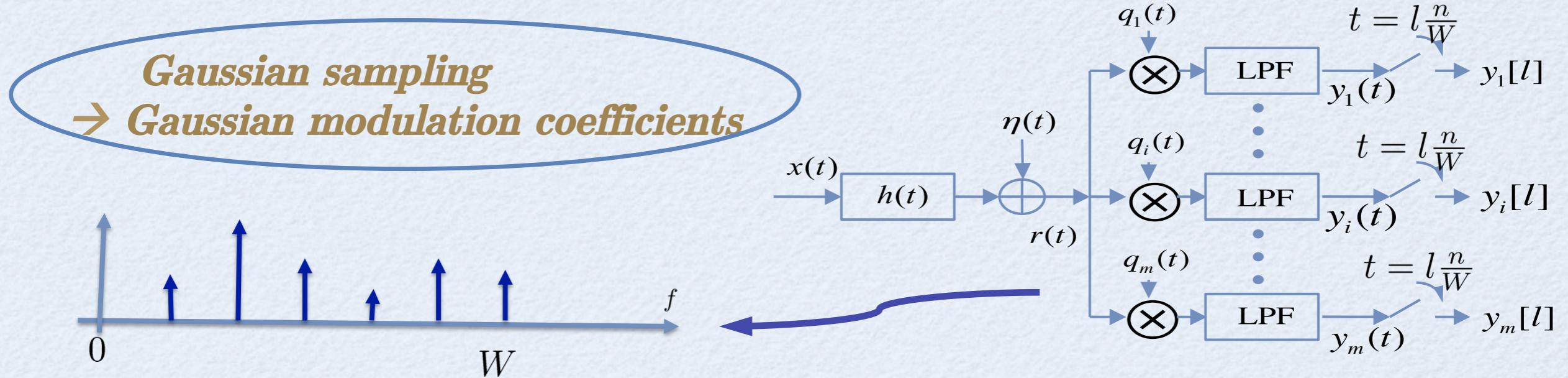
$$\frac{1}{2} \alpha \mathcal{H} \left( \frac{\beta}{\alpha} \right)$$



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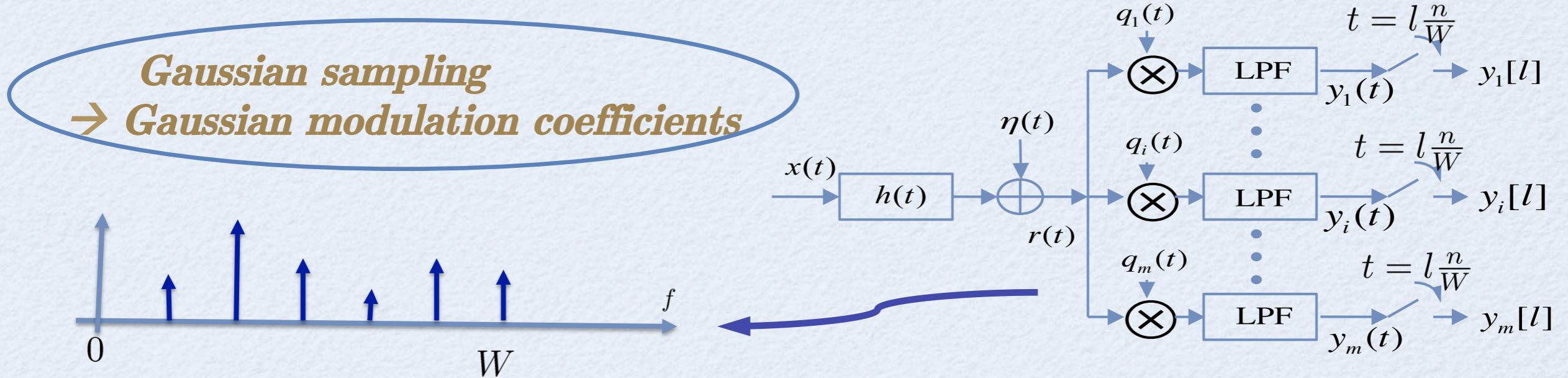
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**Theorem (Achievability):** If  $\alpha + \beta < 1$ , then the capacity loss *per Hertz* under **i.i.d. Gaussian random sampling** is

$$\forall s \in \binom{[n]}{k} : L_s^Q \leq \frac{1}{2} \left\{ \mathcal{H}(\beta) - \alpha \mathcal{H}\left(\frac{\beta}{\alpha}\right) + O\left(\frac{\log^2 n}{\sqrt{n}}\right) + \frac{\beta}{\text{SNR}_{\min}} \right\}$$

with probability exceeding  $1 - e^{-\Omega(n)}$ .

# Implications: super-Landau sampling ( $\alpha=\beta$ , $\alpha+\beta<1$ )

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**Theorem (Achievability):** Under **i.i.d. Gaussian random sampling**, with **exponentially** high probability

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- Gaussian sampling is **Minimax !**
- Sharp concentration: **exponentially** high probability
- ***Universality phenomena not shown...***
  - We have only shown the results for i.i.d. Gaussian sampling

# Concluding Remarks

- *Minimax Capacity Loss*
  - A new *metric* to characterize *robustness against different channel realizations*
  - For multiband channels, it depends only on undersampling factor and sparsity ratio
- *The power of random sampling*
  - Near-optimal in an overall sense (minimax)
  - Large random samplers behave *in deterministic ways*  
(sharp concentration + universality)
- *A Non-Asymptotic analysis of random channels*

# Full-Length Paper

- Y. Chen, A. J. Goldsmith, and Y. C. Eldar,  
“**Minimax Capacity Loss under Sub-Nyquist Universal Sampling**”, *submitted to IEEE Trans Info Theory*, [arxiv.org/abs/1304.7751](https://arxiv.org/abs/1304.7751), April 2013,

Thank You!