

ISIT 2014

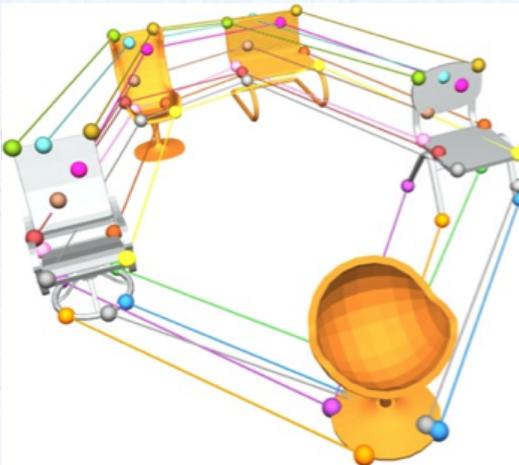
Information Recovery from Pairwise Measurements

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Joint Data Analysis

- **Data-Intensive Applications**



- many **relevant** but **heterogeneous** data
- exploit **data correlation** to improve individual task

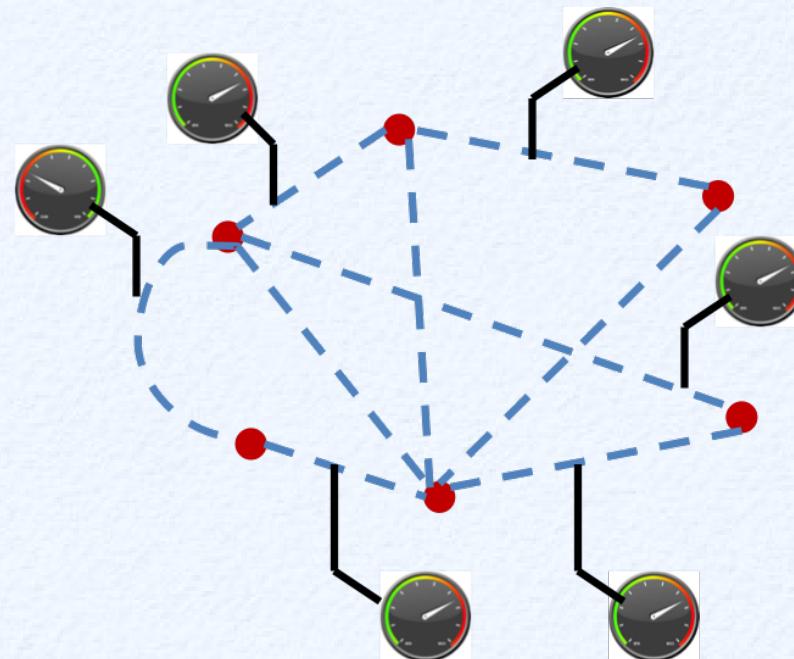
This Talk: Info Recovery from Graph-Based Measurements

- **Information Network**

- n vertices

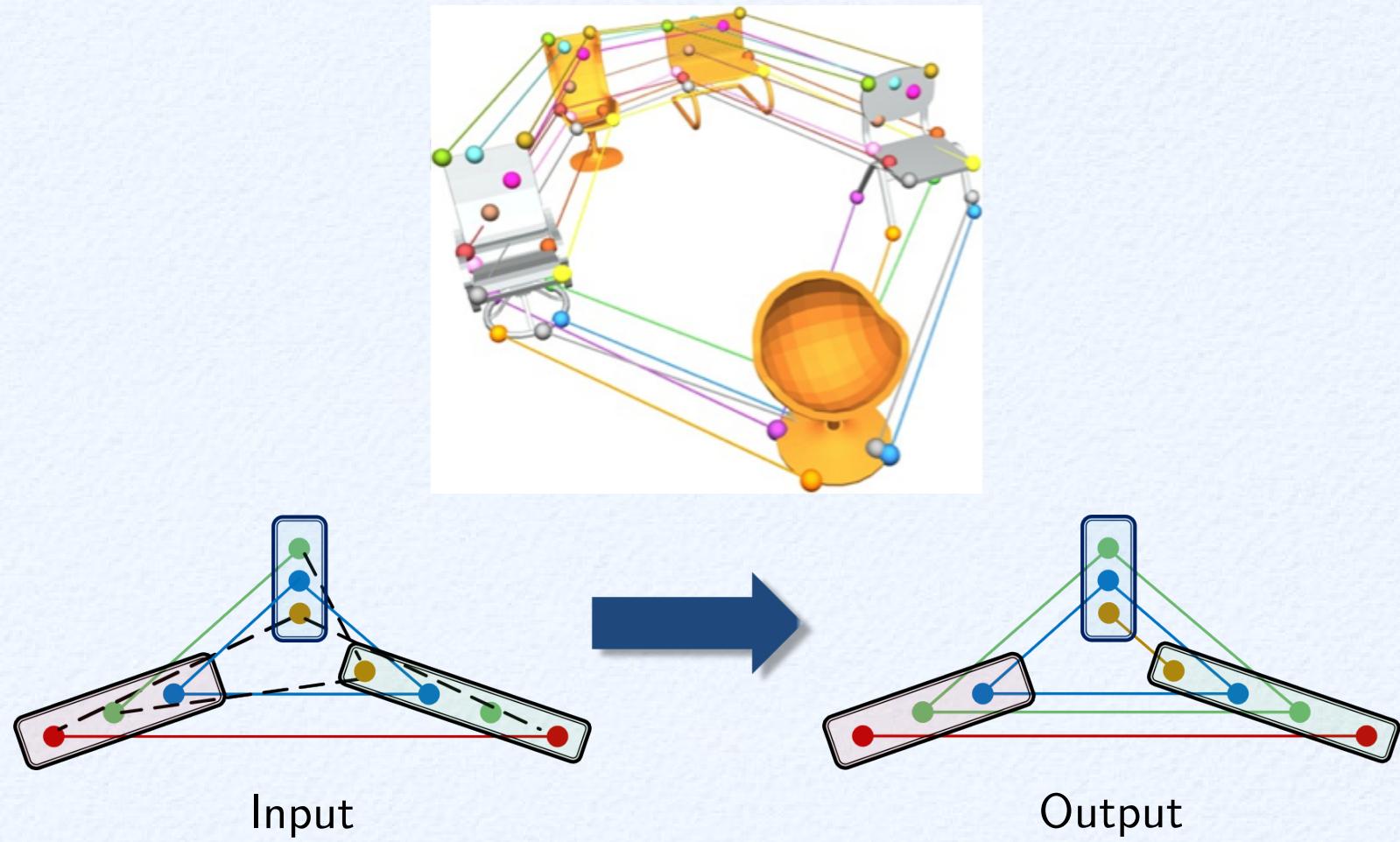
- **Measurements**

- *taken on a few edges*



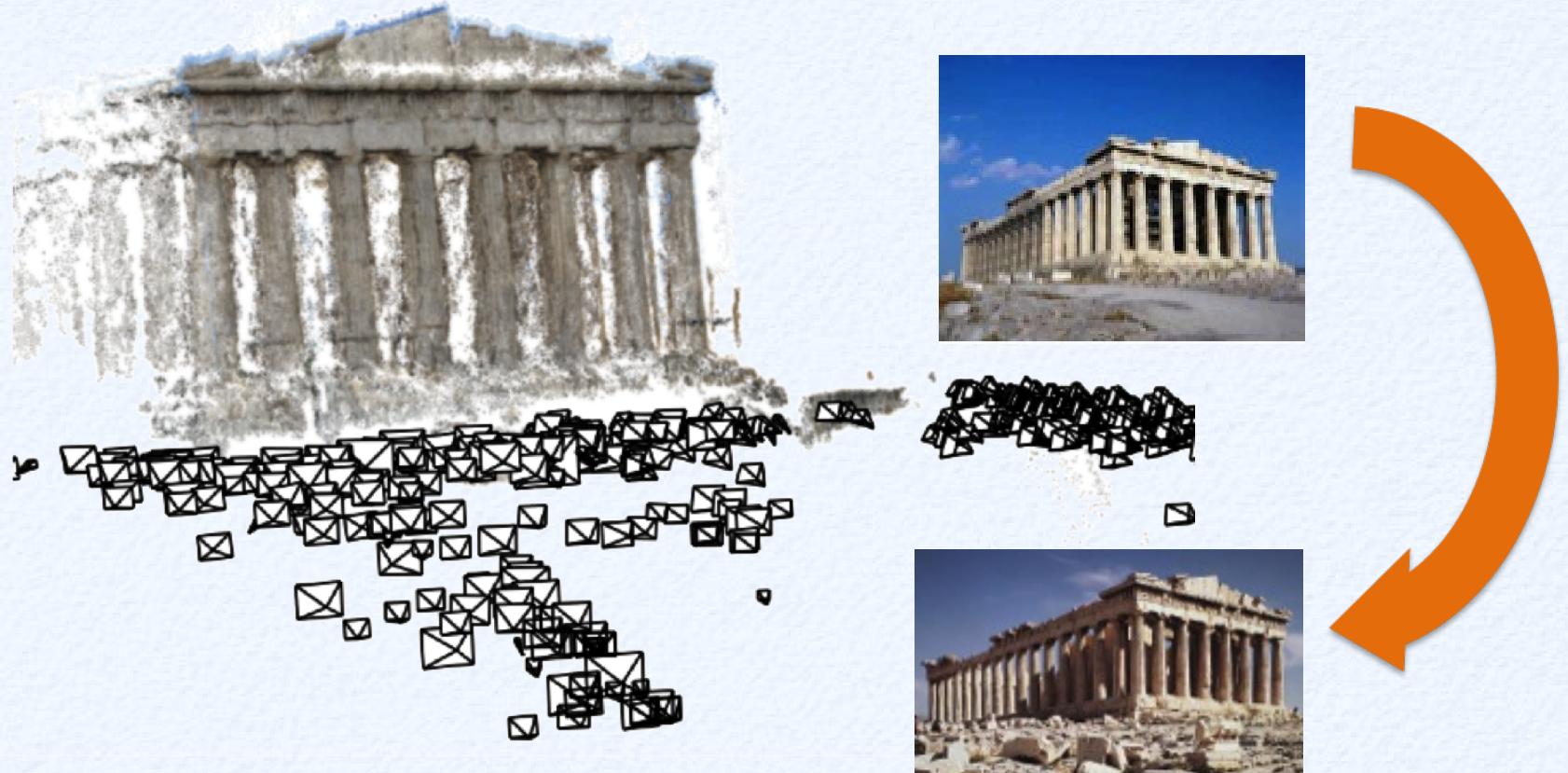
Motivating Application: Joint Graph Matching

- **Joint Graph Matching:** identify feature correspondences across instances



Motivating Application: Multi-Image Registration

- **Structure from Motion:** estimate 3D structures from 2D image sequences

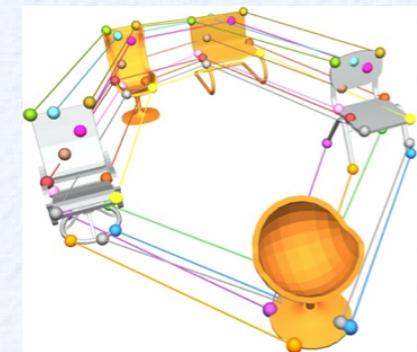


- jointly align all images given relative camera poses between each pair

Lack of Benchmark

- **A Flurry of Activity in Algorithm Design**

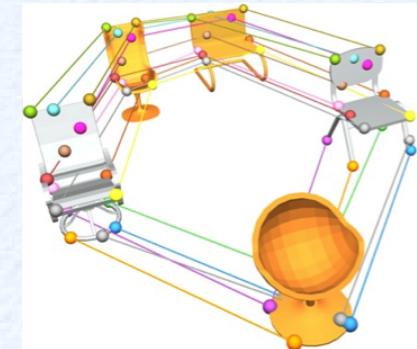
- Joint Matching (HuangGuibas'2013)
 - works if input error rate $\leq 50\%$
- Rotation Registration (WangSinger'2012)
 - works if input error rate $\leq 50\%$



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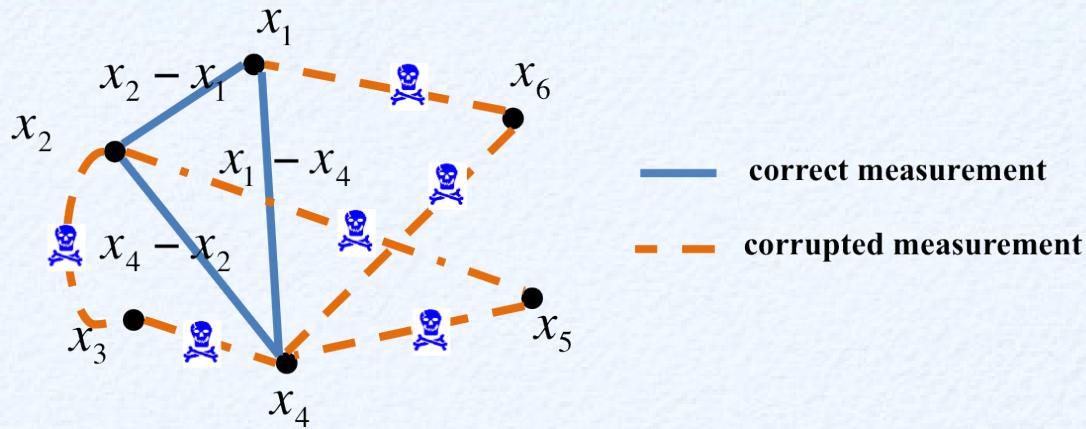


- **Lack of Benchmark:**

- What are the fundamental performance limits?



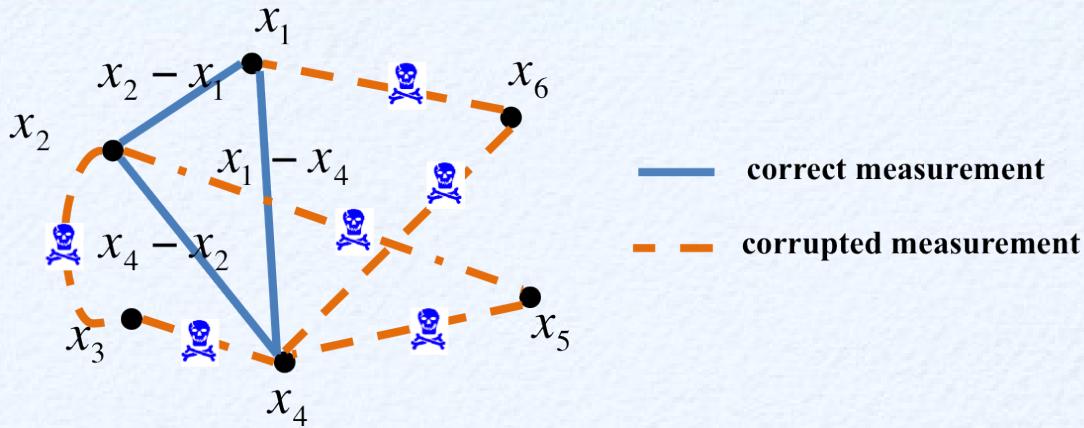
Problem Setup: A Unified Formulation



- **Information Network**

- n vertices
- vertex i takes value $x_i \in \{1, \dots, M\}$

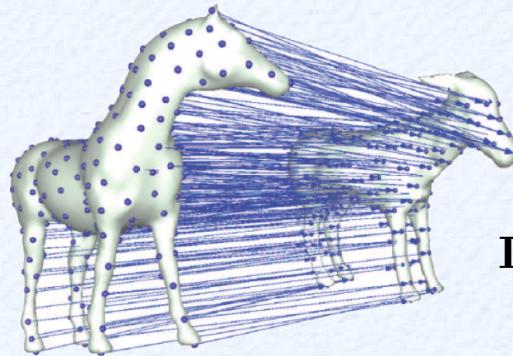
Problem Setup: A Unified Formulation



- **Information Network**
 - n vertices
 - vertex i takes value $x_i \in \{1, \dots, M\}$
- **This Talk: Pairwise Measurements**
 - measure certain pairwise relation function $f(x_i, x_j)$

Example: Pairwise Difference

- **Pairwise Difference** $x_i - x_j$
 - pairwise matching: $\Pi_j^{-1} \Pi_i$

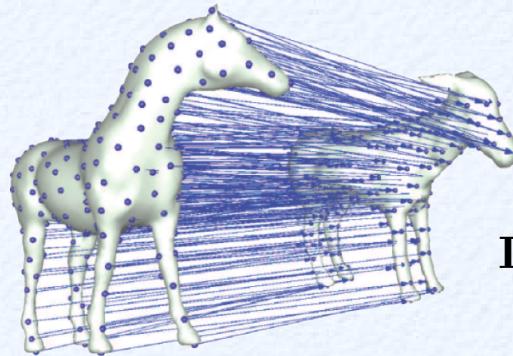


$$\Pi_i = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

Example: Pairwise Difference

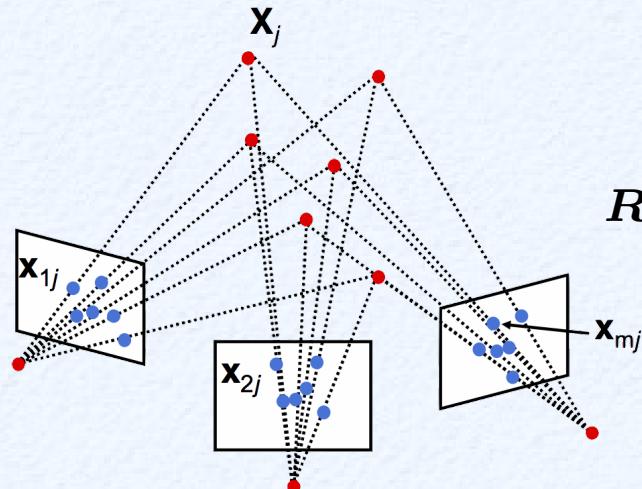
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$$\Pi_i = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

- relative rotation: $R_j^{-1} R_i$

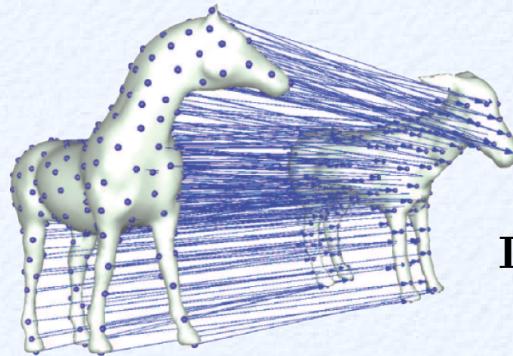


$$R_i = \begin{bmatrix} \cos \theta & \sin \theta \\ -\sin \theta & \cos \theta \end{bmatrix}$$

Example: Pairwise Difference

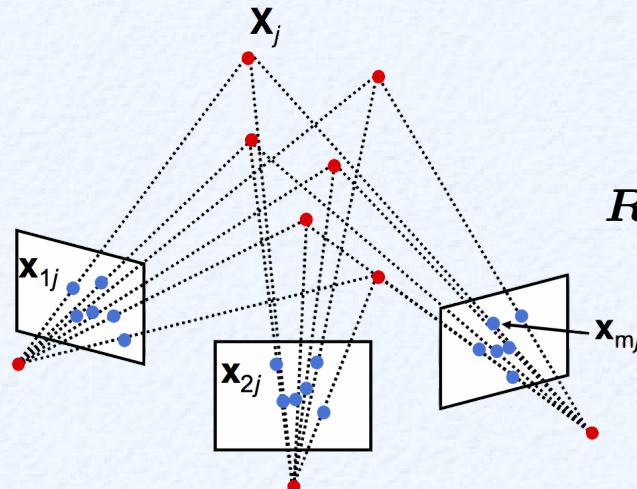
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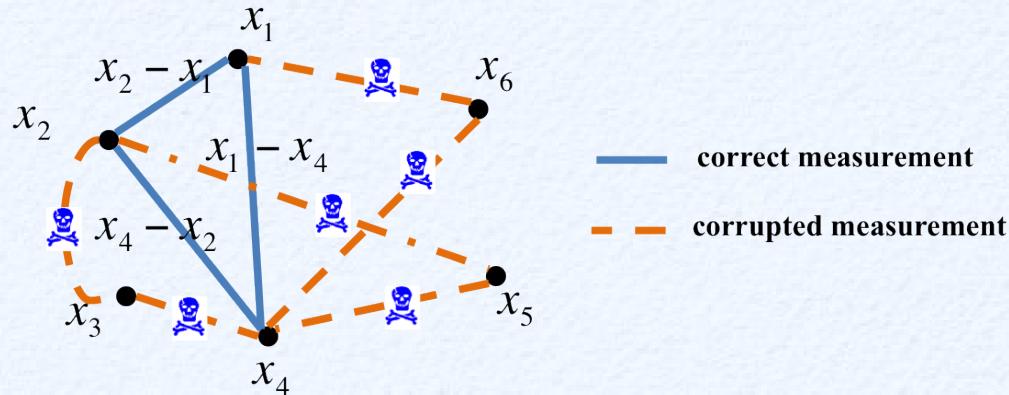
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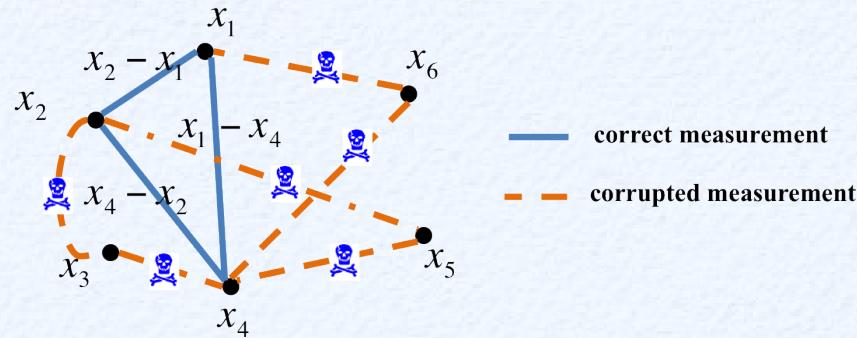
- angular difference ...

Graph-Based Measurement Model



- **Measurement Graph**
 - $\mathcal{G} = (\mathcal{V}, \mathcal{E})$
 - measurements taken over all edges $(i, j) \in \mathcal{E}$
- *Graphical channels* [AbbeMontanari'2013]
 - *community detection, graph-based coding, constrained satisfaction ...*

Challenge 1: Missing Measurements

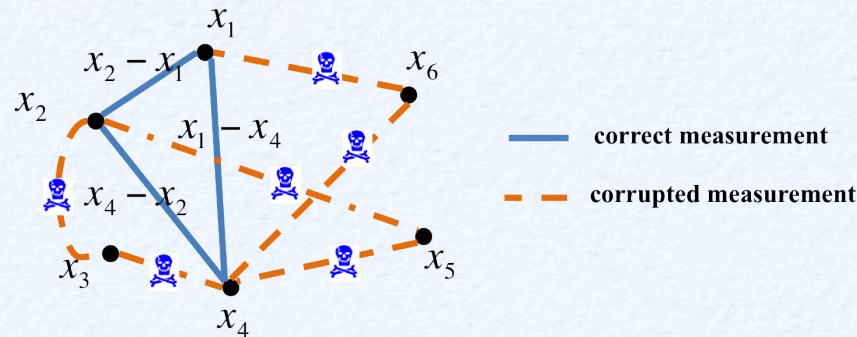


- **Incompleteness**

- measurement graph \mathcal{G} is sparse

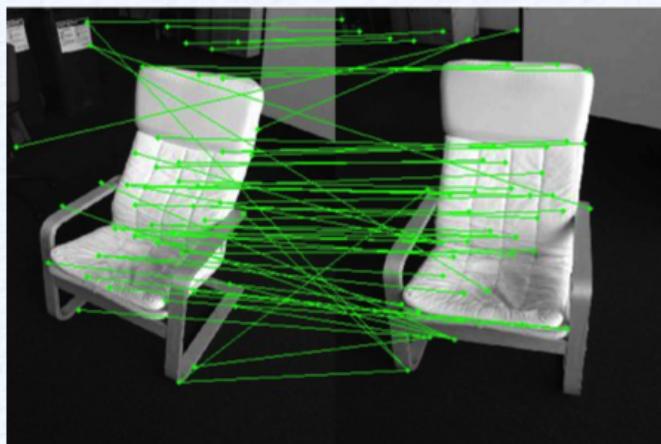


Challenge 2: Corrupted Measurements

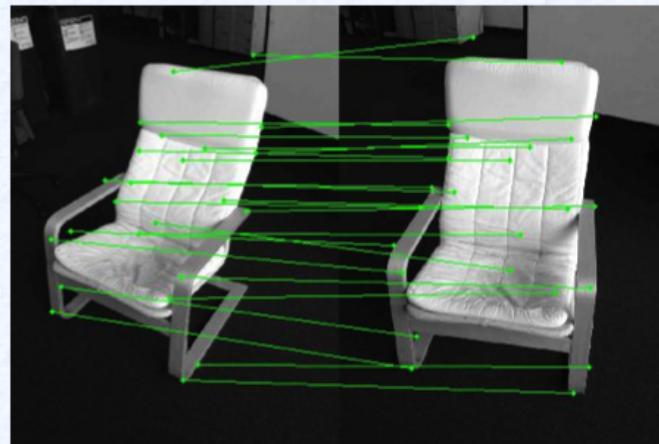


- **Dense Measurement Errors**

- A majority of measurements can be significantly corrupted



Measurements



Ground Truth

Fundamental Limit for Information Recovery

- Random Corruption Model (for Pairwise Difference)

$$\forall (i, j) \in \mathcal{G} \quad \text{measurement}(i, j) = \begin{cases} x_i - x_j, & \text{w.p. } p_{\text{true}} \\ \text{uniform } (\mathcal{M}), & \text{else} \end{cases}$$

accurate
corrupted

- \mathcal{G} : measurement graph
- \mathcal{M} : # possible values
- p_{true} : non-corruption rate

Fundamental Limit for Information Recovery

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accurate
corrupted

- \mathcal{G} : measurement graph
- M : # possible values
- p_{true} : non-corruption rate

- **Goal:** identify the threshold on p_{true} below which recovery is impossible

Fundamental Limit for Information Recovery

- Random Corruption Model (for Pairwise Difference)

$$\forall (i, j) \in \mathcal{G} \quad \text{measurement}(i, j) = \begin{cases} x_i - x_j, & \text{w.p. } p_{\text{true}} \\ \text{uniform } (\mathcal{M}), & \text{else} \end{cases} \quad \begin{matrix} & \text{accurate} \\ & \text{corrupted} \end{matrix}$$

- \mathcal{G} : measurement graph
- M : # possible values
- p_{true} : non-corruption rate

Theorem. For various homogeneous graphs, exact info recovery is possible iff

$$p_{\text{true}} \gtrsim \begin{cases} \frac{1}{\sqrt{M \cdot \text{avg-degree}}}, & \text{if } M \leq \text{avg-degree} \\ \frac{1}{\text{avg-degree}}, & \text{else} \end{cases}$$

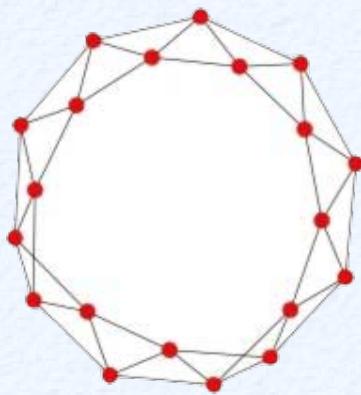
\gtrsim : order-wise bound up to a polylog factor

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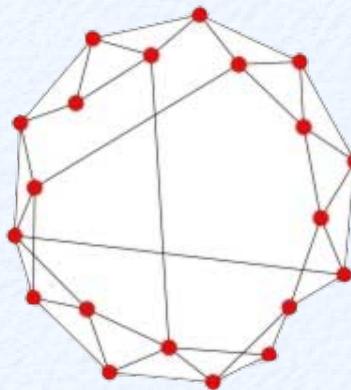
Key Metrics: Edge Sparsity

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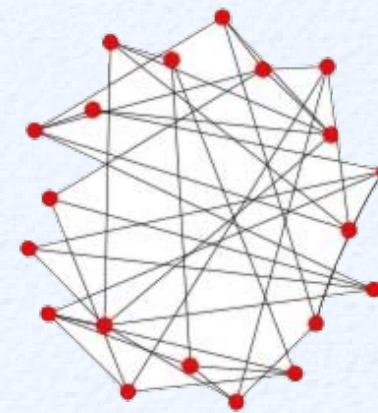
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geometric graphs



small world graphs

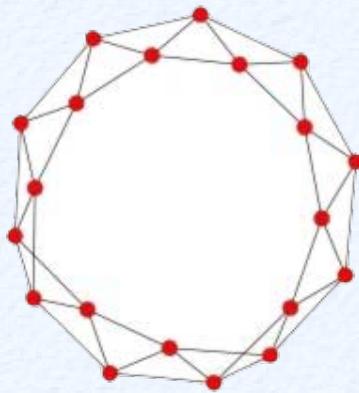


Erdos Renyi random graphs

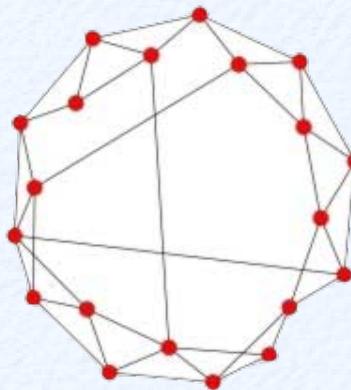
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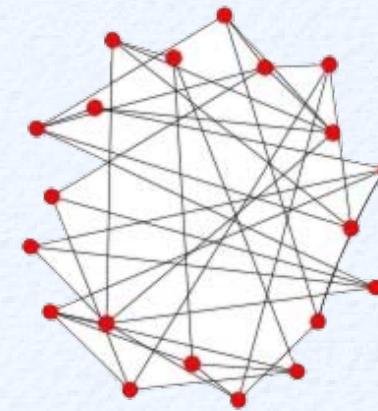
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geometric graphs



small world graphs



Erdos Renyi random graphs

- Depend almost only on edge sparsity

- irrespective of other graphical metrics (e.g. spectral gaps)

More Precisely: Which homogeneous graphs?

- Key Metric

$$\mathcal{N}_k \triangleq \left\{ \text{Vertex Set } S : \underbrace{|\partial S|}_{\# \text{ edges crossing from } S \text{ to } S^c} \leq k \right\} \quad (1)$$

- Our results work for all graphs such that

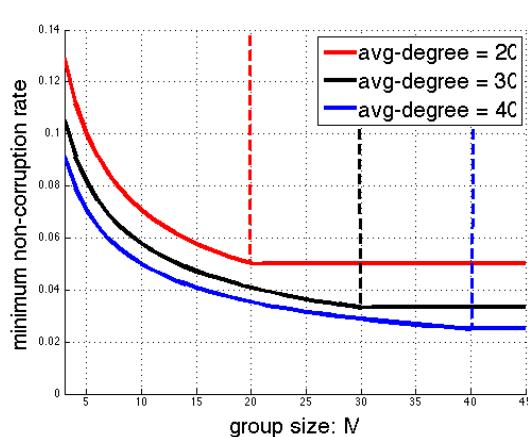
$$\forall k : \frac{\log |\mathcal{N}_k|}{k} = O(\text{poly log}(n))$$

- Erdos-Renyi Graphs
- Random Geometric Graphs
- Small-World Graphs
- Other Expander Graphs ...

Key Metrics: Universe Size

Theorem. For various homogeneous graphs, exact info recovery is possible iff

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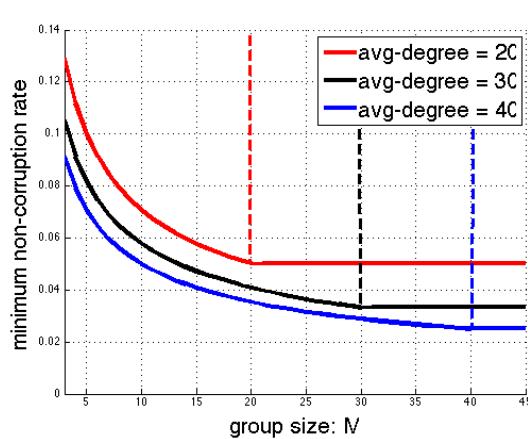


p_{true} v.s. M

Key Metrics: Universe Size

Theorem. For various homogeneous graphs, exact info recovery is possible iff

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p_{true} v.s. M

- Improve with the universe size M

- before entering a *connectivity-limited regime*

Generalization

- Same order-wise limits hold for general pairwise measurements!
 - Any relation $f(x, y)$ such that
$$\forall y, \quad f(\cdot, y) \text{ is bijective}$$
$$\forall x, \quad f(x, \cdot) \text{ is bijective}$$
 - *Pairwise Difference*
 - *Pairwise Sum*
 - *Many Other Linear Relations ...*
- Abbe et al (ISIT2014) — exact limit for Erdos-Renyi graphs when $M = 2$

Computational Feasibility?

- Is the fundamental limit achievable by tractable algorithms?

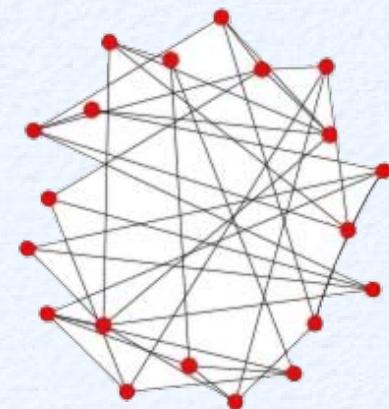
Computational Feasibility?

- Is the fundamental limit achievable by tractable algorithms?

- **Consider Convex Relaxation**

- **Applied to the following example:**

- pairwise difference: $x_i - x_j \bmod M$
 - Erdos-Renyi random graphs $\mathcal{G}(n, p_{\text{obs}})$



Erdos-Renyi random graphs

Matrix Representation

- Object Representation

- cyclic shift of I

$$x_i = k \quad \text{represented by} \quad \mathbf{X}_i = \left[\underbrace{\begin{matrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ \mathbf{1} & 0 & 0 \\ 0 & \mathbf{1} & 0 \\ 0 & 0 & \mathbf{1} \end{matrix}}_{k \text{ columns}} \quad \begin{matrix} 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{matrix} \quad \begin{matrix} 0 \\ \mathbf{1} \\ 0 \\ 0 \\ 0 \end{matrix} \right]$$

$$x_i - x_j \bmod M \quad \text{represented by} \quad \mathbf{X}_i \mathbf{X}_j^\top$$

Matrix Representation

- Object Representation

- cyclic shift of \mathbf{I}

$$x_i = k \quad \text{represented by} \quad \mathbf{X}_i = \left[\begin{array}{cccc|cc} & & & & \overbrace{\hspace{1cm}}^{k \text{ columns}} & & \\ 0 & 0 & 0 & & & \mathbf{1} & 0 \\ 0 & 0 & 0 & & & 0 & \mathbf{1} \\ \mathbf{1} & 0 & 0 & & & 0 & 0 \\ 0 & \mathbf{1} & 0 & & & 0 & 0 \\ 0 & 0 & \mathbf{1} & & & 0 & 0 \end{array} \right]$$

$$x_i - x_j \bmod M \quad \text{represented by} \quad \mathbf{X}_i \mathbf{X}_j^\top$$

- Low-Rank P.S.D. Matrix

$$\mathbf{X} := \begin{bmatrix} \mathbf{I} & \mathbf{X}_{12} & \cdots & \mathbf{X}_{1n} \\ \mathbf{X}_{21} & \mathbf{I} & \cdots & \mathbf{X}_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{X}_{n1} & \mathbf{X}_{n2} & \cdots & \mathbf{I} \end{bmatrix} = \underbrace{\begin{bmatrix} \mathbf{X}_1 \\ \mathbf{X}_2 \\ \vdots \\ \mathbf{X}_n \end{bmatrix}}_{M \text{ columns}} \begin{bmatrix} \mathbf{X}_1^\top & \mathbf{X}_2^\top & \cdots & \mathbf{X}_n^\top \end{bmatrix} \succeq 0$$

Exact Recovery via Convex Relaxation

Erdos-Renyi graph: $\mathcal{G}(n, p_{\text{obs}})$


$$\begin{array}{c} \text{ground truth} \\ = \\ \begin{array}{l} \text{minimize}_{\mathbf{X}} \quad - \langle \mathbf{X}, \mathbf{X}^{\text{input}} \rangle + \lambda \langle \mathbf{X}, \mathbf{1}\mathbf{1}^\top \rangle \\ \text{subject to} \quad \mathbf{X} \geq \mathbf{0}, \\ \quad \left[\begin{array}{cc} M & \mathbf{1}^\top \\ \mathbf{1} & \mathbf{X} \end{array} \right] \succeq \mathbf{0}, \\ \quad \mathbf{X}_{ii} = \mathbf{I}. \end{array} \end{array}$$

Theorem. Convex program with $\lambda = \sqrt{p_{\text{obs}}}$ is exact with high probability if

$$p_{\text{true}} \gtrsim \frac{\log^2(Mn)}{\sqrt{p_{\text{obs}}n}}$$

Optimality of Convex Relaxation?

Convex Relaxation v.s. Fundamental Limits.

$$p_{\text{true}} \gtrsim \frac{\log^2(Mn)}{\sqrt{p_{\text{obs}}n}} \quad \text{v.s.} \quad p_{\text{true}} \gtrsim \begin{cases} \frac{\log n}{\sqrt{Mp_{\text{obs}}n}}, & \text{if } M \leq p_{\text{obs}}n \\ \frac{\log n}{p_{\text{obs}}n}, & \text{else} \end{cases}$$

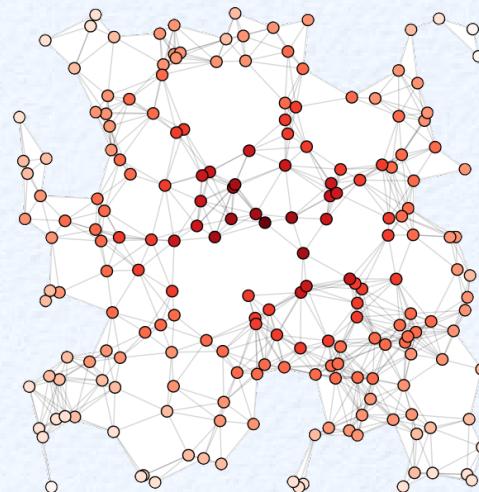
- Convex relaxation is near-optimal
 - when $M = \Theta(1)$
 - Erdos-Renyi random graph

Optimality of Convex Relaxation?

Convex Relaxation v.s. Fundamental Limits.

$$p_{\text{true}} \gtrsim \frac{\log^2(Mn)}{\sqrt{p_{\text{obs}}n}} \quad \text{v.s.} \quad p_{\text{true}} \gtrsim \begin{cases} \frac{\log n}{\sqrt{Mp_{\text{obs}}n}}, & \text{if } M \leq p_{\text{obs}}n \\ \frac{\log n}{p_{\text{obs}}n}, & \text{else} \end{cases}$$

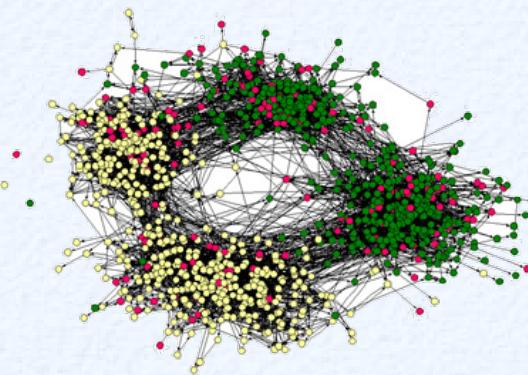
- **Convex relaxation is near-optimal**
 - when $M = \Theta(1)$
 - Erdos-Renyi random graph
- **Convex relaxation might be sub-optimal**
 - if M scales with n
 - general measurement graphs (e.g. geometric graphs)



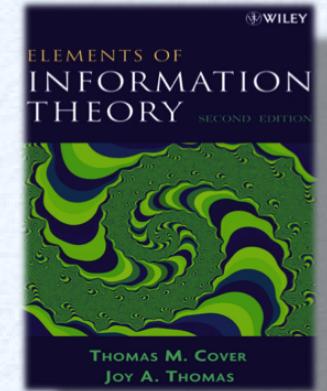
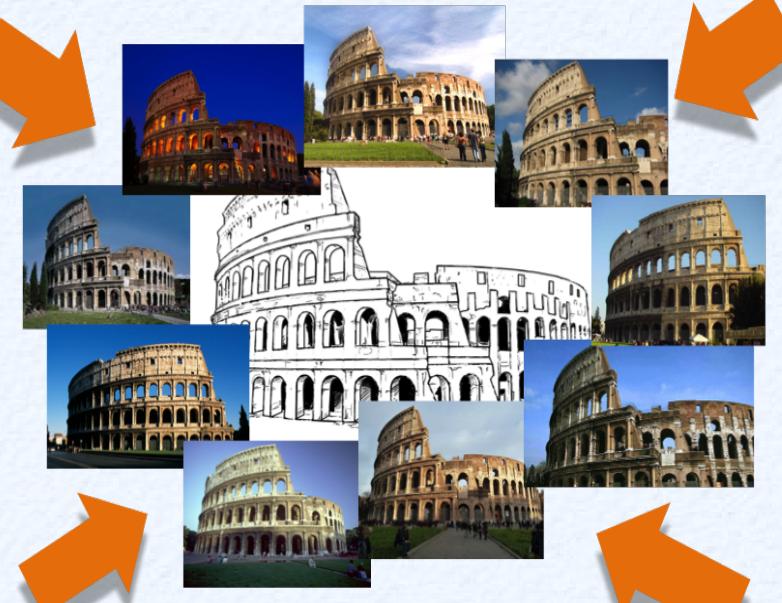
Summary: Graph-Based Joint Information Recovery

- **Fundamental Minimax Limits**
 - Depend almost only on edge sparsity
 - 2 operating regimes
- **Computational Feasibility?**
 - SDP is near-optimal in some regime
 - A large regime remains unknown (computational barrier?)
- *Thanks to A. Montanari, E. Abbe, A. Bandeira, and A. Singer*

Concluding Remarks: Learn from Data Correlation



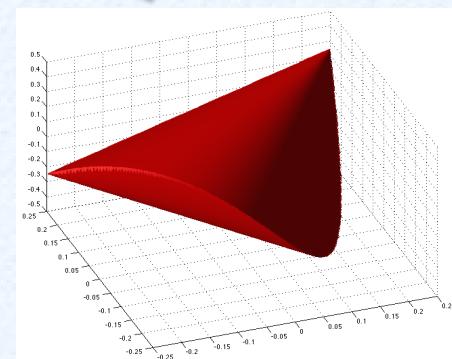
random graphs



information theory



matrix completion



convex optimization